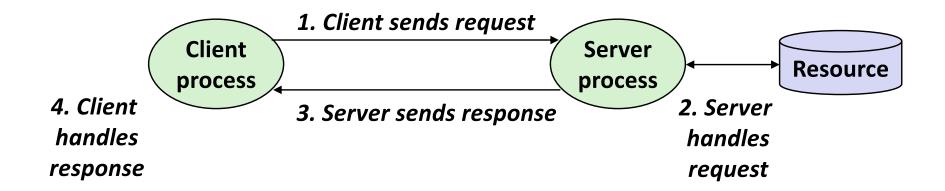
Network Programming

Computer Systems
Friday, December 01, 2023

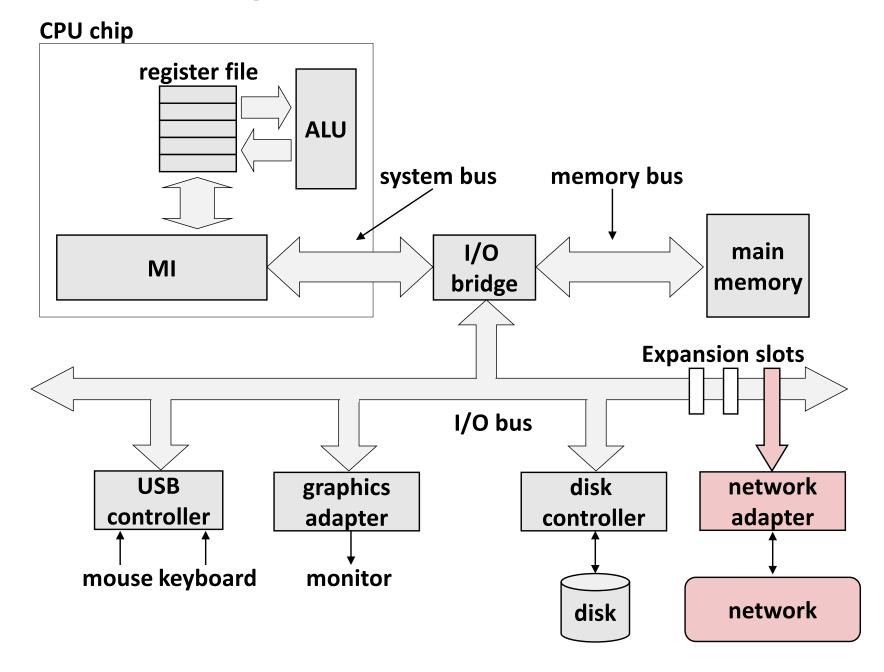
A Client-Server Transaction

- Most network applications are based on the client-server model:
 - A server process and one or more client processes
 - Server manages some resource
 - Server provides service by manipulating resource for clients
 - Server activated by request from client (vending machine analogy)



Note: clients and servers are processes running on hosts (can be the same or different hosts)

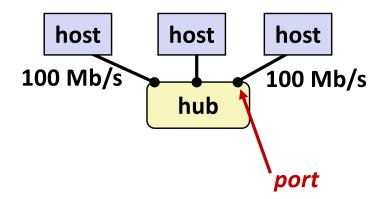
Hardware Organization of a Network Host



Computer Networks

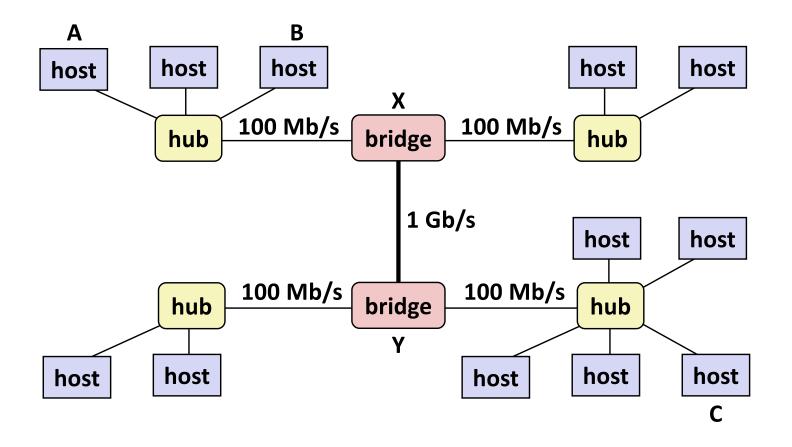
- A network is a hierarchical system of boxes and wires organized by geographical proximity
 - SAN (System Area Network) spans cluster or machine room
 - Switched Ethernet, Quadrics QSW, ...
 - LAN (Local Area Network) spans a building or campus
 - Ethernet is most prominent example
 - WAN (Wide Area Network) spans country or world
 - Typically high-speed point-to-point phone lines
- An internetwork (internet) is an interconnected set of networks
 - The Global IP Internet (uppercase "I") is the most famous example of an internet (lowercase "i")
- Let's see how an internet is built from the ground up

Lowest Level: Ethernet Segment



- Ethernet segment consists of a collection of hosts connected by wires (twisted pairs) to a hub
- Spans room or floor in a building
- Operation
 - Each Ethernet adapter has a unique 48-bit address (MAC address)
 - E.g., 00:16:ea:e3:54:e6
 - Hosts send bits to any other host in chunks called frames
 - Hub slavishly copies each bit from each port to every other port
 - Every host sees every bit
 - Note: Hubs are on their way out. Bridges (switches, routers) became cheap enough to replace them

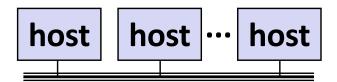
Next Level: Bridged Ethernet Segment



- Spans building or campus
- Bridges cleverly learn which hosts are reachable from which ports and then selectively copy frames from port to port

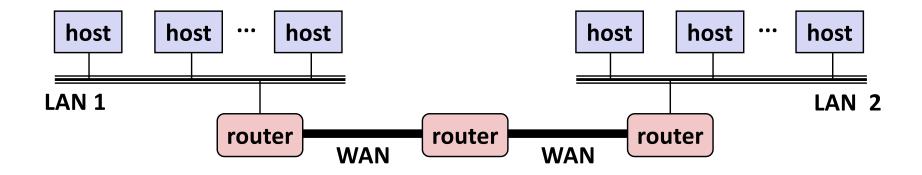
Conceptual View of LANs

For simplicity, hubs, bridges, and wires are often shown as a collection of hosts attached to a single wire:



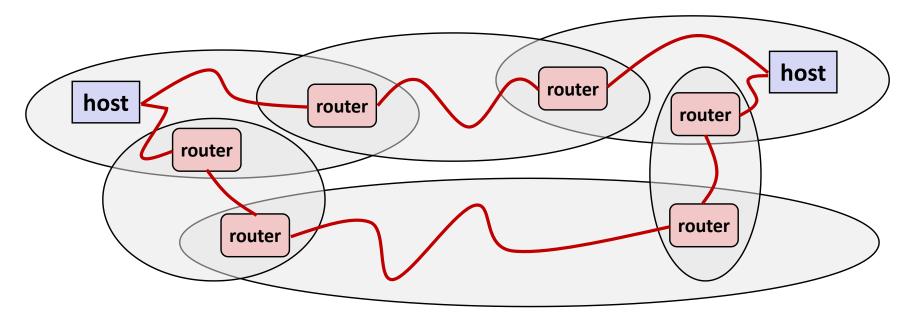
Next Level: internets

- Multiple incompatible LANs can be physically connected by specialized computers called routers
- The connected networks are called an internet (lower case)



LAN 1 and LAN 2 might be completely different, totally incompatible (e.g., Ethernet, Fibre Channel, 802.11*, T1-links, DSL, ...)

Logical Structure of an internet



- Ad hoc interconnection of networks
 - No particular topology
 - Vastly different router & link capacities
- Send packets from source to destination by hopping through networks
 - Router forms bridge from one network to another
 - Different packets may take different routes

The Notion of an internet Protocol

- How is it possible to send bits across incompatible LANs and WANs?
- Solution: protocol software running on each host and router
 - Protocol is a set of rules that governs how hosts and routers should cooperate when they transfer data from network to network.
 - Smooths out the differences between the different networks

What Does an internet Protocol Do?

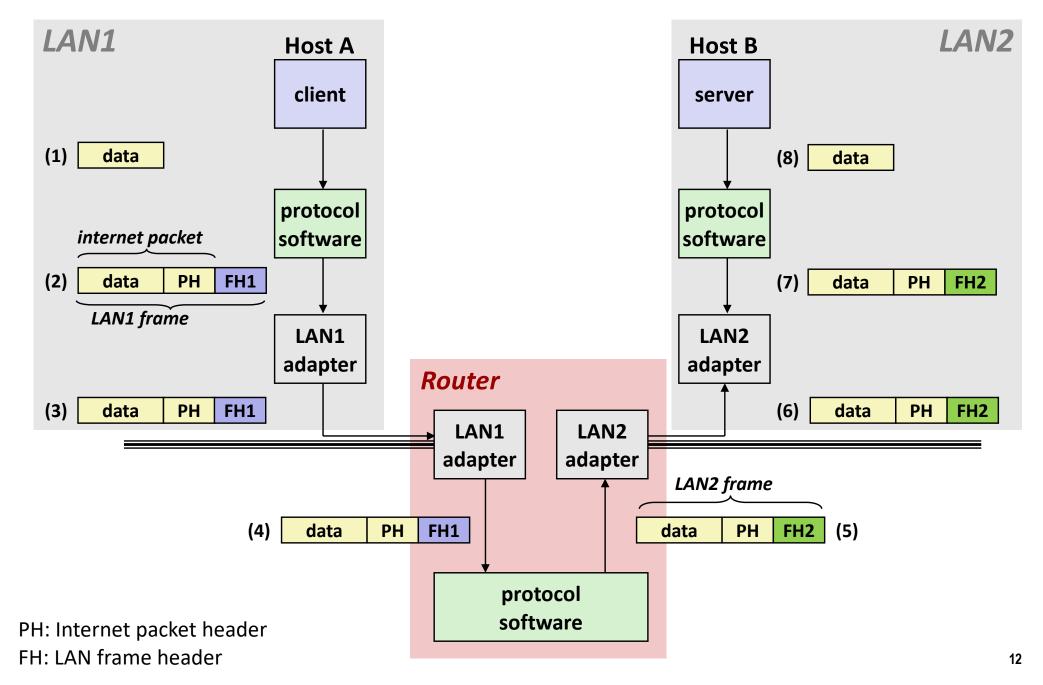
Provides a naming scheme

- An internet protocol defines a uniform format for host addresses
- Each host (and router) is assigned at least one of these internet addresses that uniquely identifies it

Provides a delivery mechanism

- An internet protocol defines a standard transfer unit (packet)
- Packet consists of *header* and *payload*
 - Header: contains info such as packet size, source and destination addresses
 - Payload: contains data bits sent from source host

Transferring internet Data Via Encapsulation



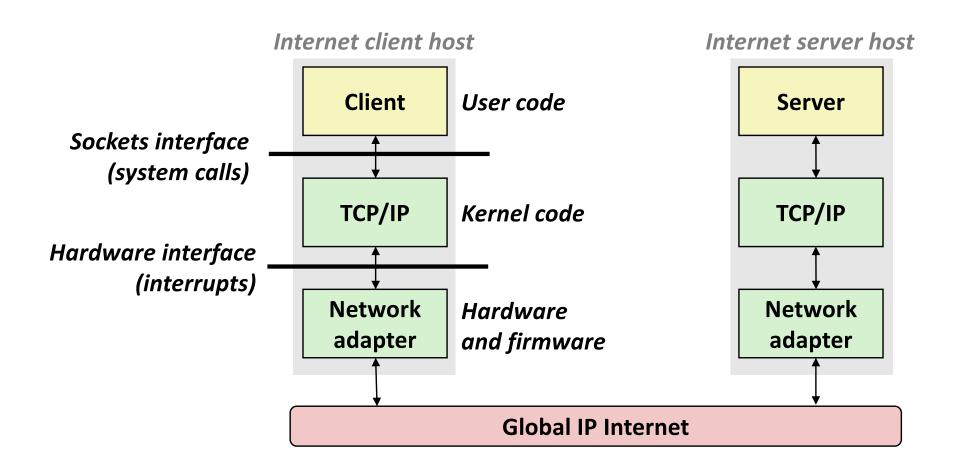
Other Issues

- We are glossing over a number of important questions:
 - What if different networks have different maximum frame sizes? (segmentation)
 - How do routers know where to forward frames?
 - How are routers informed when the network topology changes?
 - What if packets get lost?
- These (and other) questions are addressed by the area of systems known as computer networking

Global IP Internet (upper case)

- Most famous example of an internet
- Based on the TCP/IP protocol family
 - IP (Internet Protocol) :
 - Provides basic naming scheme and unreliable delivery capability of packets (datagrams) from host-to-host
 - UDP (Unreliable Datagram Protocol)
 - Uses IP to provide unreliable datagram delivery from process-to-process
 - TCP (Transmission Control Protocol)
 - Uses IP to provide reliable byte streams from process-to-process over connections
- Accessed via a mix of Unix file I/O and functions from the sockets interface

Hardware and Software Organization of an Internet Application



A Programmer's View of the Internet

- 1. Hosts are mapped to a set of 32-bit IP addresses
 - 203.246.83.225
- 2. The set of IP addresses is mapped to a set of identifiers called Internet *domain names*
 - 203.246.83.225 is mapped to seoultech.ac.kr
- 3. A process on one Internet host can communicate with a process on another Internet host over a *connection*

Aside: IPv4 and IPv6

- The original Internet Protocol, with its 32-bit addresses, is known as Internet Protocol Version 4 (IPv4)
- 1996: Internet Engineering Task Force (IETF) introduced Internet Protocol Version 6 (IPv6) with 128-bit addresses
 - Intended as the successor to IPv4
- As of 2015, vast majority of Internet traffic still carried by IPv4
 - Only 4% of users access Google services using IPv6.
- We will focus on IPv4, but will show you how to write networking code that is protocol-independent.

(1) IP Addresses

■ 32-bit IP addresses are stored in an IP address struct

- IP addresses are always stored in memory in network byte order (big-endian byte order)
- True in general for any integer transferred in a packet header from one machine to another.
 - E.g., the port number used to identify an Internet connection.

```
/* Internet address structure */
struct in_addr {
   uint32_t s_addr; /* network byte order (big-endian) */
};
```

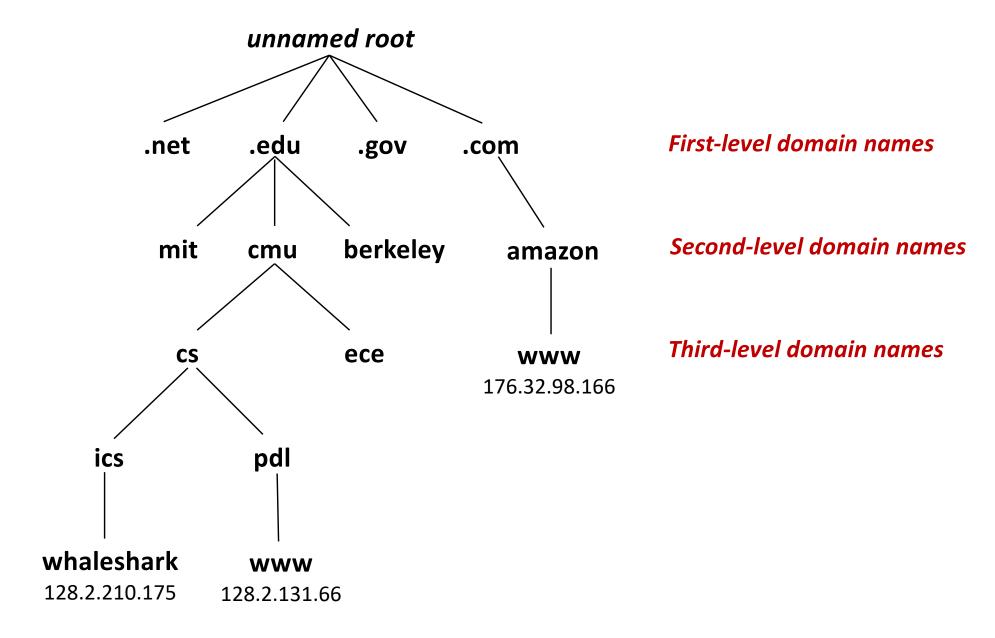
Dotted Decimal Notation

 By convention, each byte in a 32-bit IP address is represented by its decimal value and separated by a period

```
• IP address: 0x8002C2F2 = 128.2.194.242
```

 Use getaddrinfo and getnameinfo functions (described later) to convert between IP addresses and dotted decimal format.

(2) Internet Domain Names



Domain Naming System (DNS)

- The Internet maintains a mapping between IP addresses and domain names in a huge worldwide distributed database called DNS
- Conceptually, programmers can view the DNS database as a collection of millions of host entries.
 - Each host entry defines the mapping between a set of domain names and IP addresses.
 - In a mathematical sense, a host entry is an equivalence class of domain names and IP addresses.

Properties of DNS Mappings

- Can explore properties of DNS mappings using nslookup
 - Output edited for brevity

■ Each host has a locally defined domain name localhost which always maps to the *loopback address* 127.0.0.1

```
linux> nslookup localhost
Address: 127.0.0.1
```

Use hostname to determine real domain name of local host:

```
linux> hostname
whaleshark.ics.cs.cmu.edu
```

Properties of DNS Mappings (cont)

Simple case: one-to-one mapping between domain name and IP address:

```
linux> nslookup seoultech.ac.kr
Address: 203.246.83.225
```

Multiple domain names mapped to the same IP address:

```
linux> nslookup cs.mit.edu
Address: 18.62.1.6
linux> nslookup eecs.mit.edu
Address: 18.62.1.6
```

Properties of DNS Mappings (cont)

Multiple domain names mapped to multiple IP addresses:

```
linux> nslookup www.twitter.com
Address: 199.16.156.6
Address: 199.16.156.70
Address: 199.16.156.102
Address: 199.16.156.230

linux> nslookup twitter.com
Address: 199.16.156.102
Address: 199.16.156.230
Address: 199.16.156.6
Address: 199.16.156.70
```

Some valid domain names don't map to any IP address:

```
linux> nslookup ics.cs.cmu.edu
*** Can't find ics.cs.cmu.edu: No answer
```

(3) Internet Connections

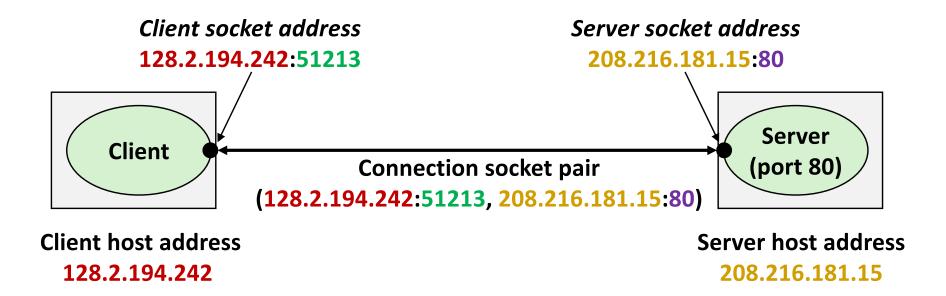
- Clients and servers communicate by sending streams of bytes over connections. Each connection is:
 - Point-to-point: connects a pair of processes.
 - Full-duplex: data can flow in both directions at the same time,
 - Reliable: stream of bytes sent by the source is eventually received by the destination in the same order it was sent.
- A socket is an endpoint of a connection
 - Socket address is an IPaddress:port pair
- A port is a 16-bit integer that identifies a process:
 - **Ephemeral port:** Assigned automatically by client kernel when client makes a connection request.
 - Well-known port: Associated with some service provided by a server (e.g., port 80 is associated with Web servers)

Well-known Ports and Service Names

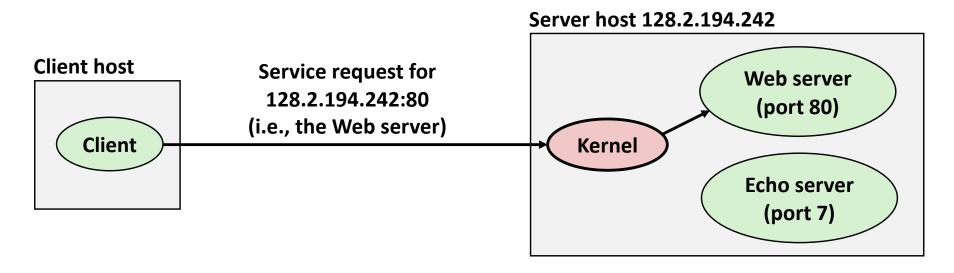
- Popular services have permanently assigned well-known ports and corresponding well-known service names:
 - echo server: 7/echo
 - ssh servers: 22/ssh
 - email server: 25/smtp
 - Web servers: 80/http
- Mappings between well-known ports and service names is contained in the file /etc/services on each Linux machine.

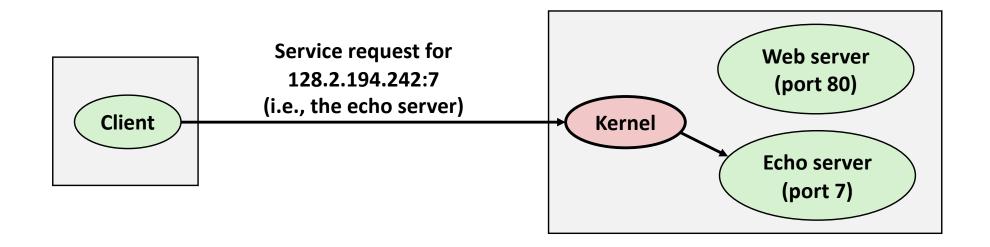
Anatomy of a Connection

- A connection is uniquely identified by the socket addresses of its endpoints (socket pair)
 - (cliaddr:cliport, servaddr:servport)



Using Ports to Identify Services



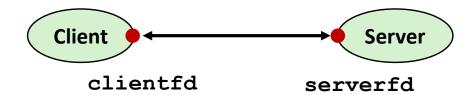


Sockets Interface

- Set of system-level functions used in conjunction with Unix I/O to build network applications.
- Created in the early 80's as part of the original Berkeley distribution of Unix that contained an early version of the Internet protocols.
- Available on all modern systems
 - Unix variants, Windows, OS X, IOS, Android, ARM

Sockets

- What is a socket?
 - To the kernel, a socket is an endpoint of communication
 - To an application, a socket is a file descriptor that lets the application read/write from/to the network
 - Remember: All Unix I/O devices, including networks, are modeled as files
- Clients and servers communicate with each other by reading from and writing to socket descriptors



The main distinction between regular file I/O and socket I/O is how the application "opens" the socket descriptors

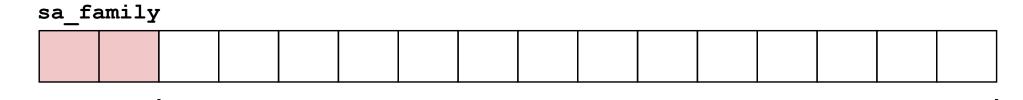
Socket Address Structures

Generic socket address:

- For address arguments to connect, bind, and accept
- Necessary only because C did not have generic (void *) pointers when the sockets interface was designed
- For casting convenience, we adopt the Stevens convention:

```
typedef struct sockaddr SA;
```

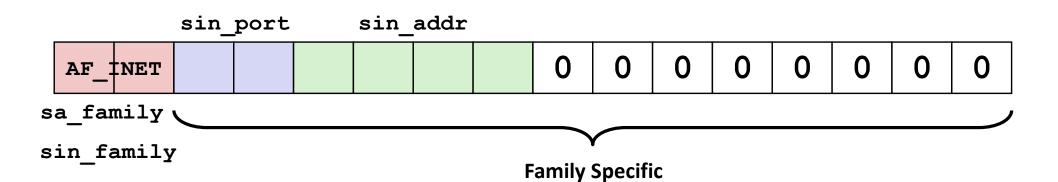
```
struct sockaddr {
  uint16_t sa_family;  /* Protocol family */
  char sa_data[14]; /* Address data. */
};
```

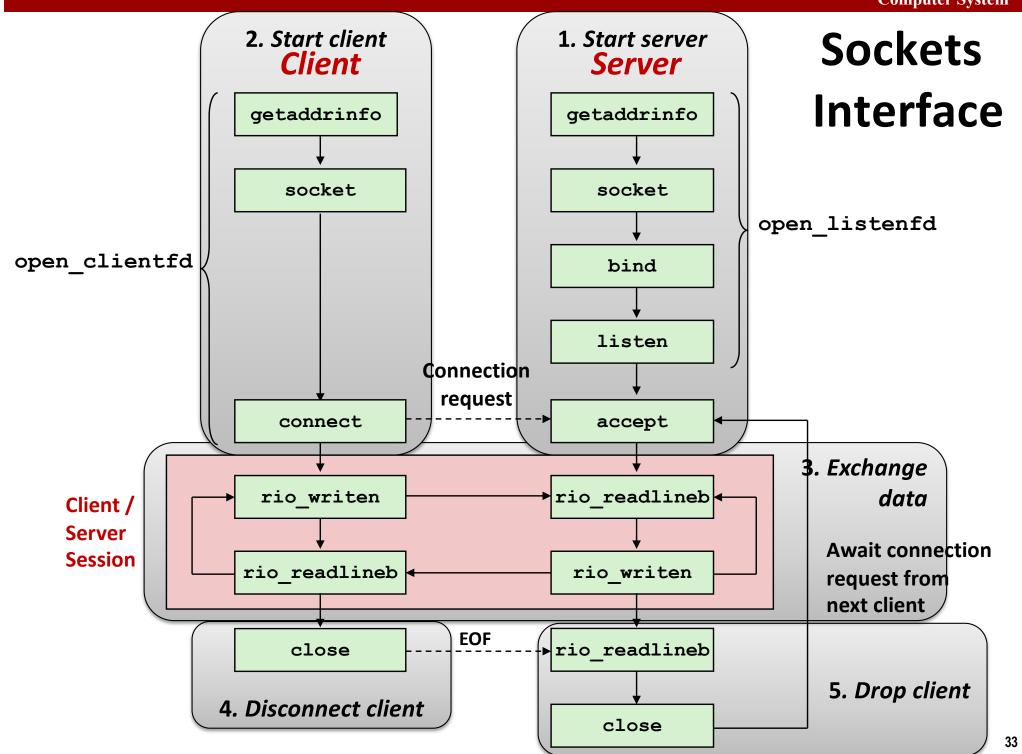


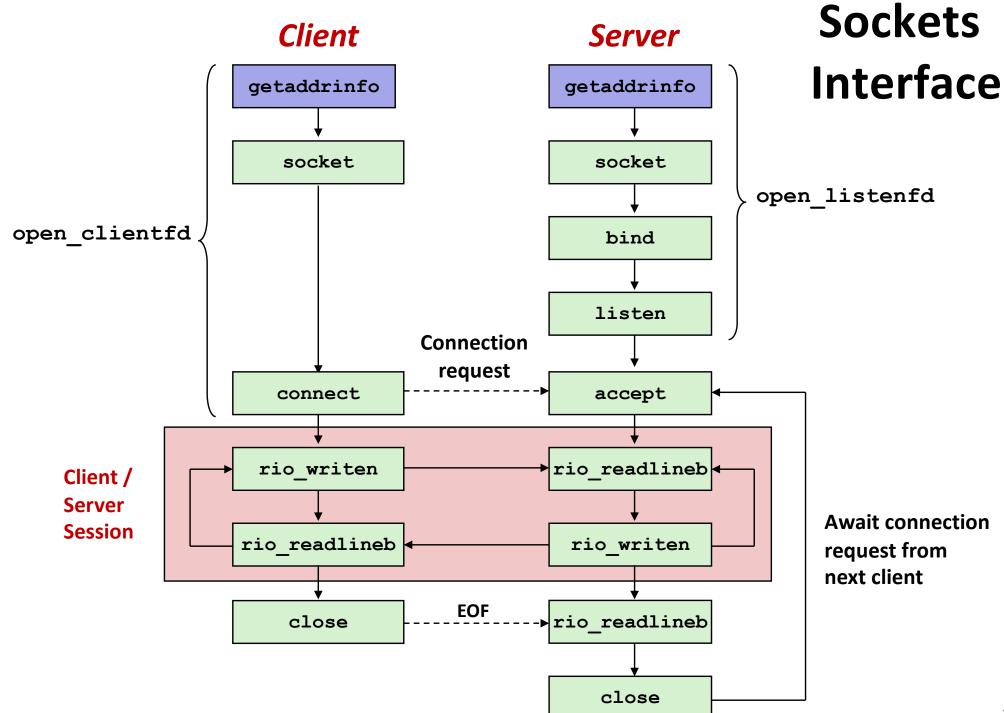
Family Specific

Socket Address Structures

- Internet-specific socket address:
 - Must cast (struct sockaddr_in *) to (struct sockaddr *) for functions that take socket address arguments.







Host and Service Conversion: getaddrinfo

- getaddrinfo is the modern way to convert string representations of hostnames, host addresses, ports, and service names to socket address structures.
 - Replaces obsolete gethostbyname and getservbyname funcs.

Advantages:

- Reentrant (can be safely used by threaded programs).
- Allows us to write portable protocol-independent code
 - Works with both IPv4 and IPv6

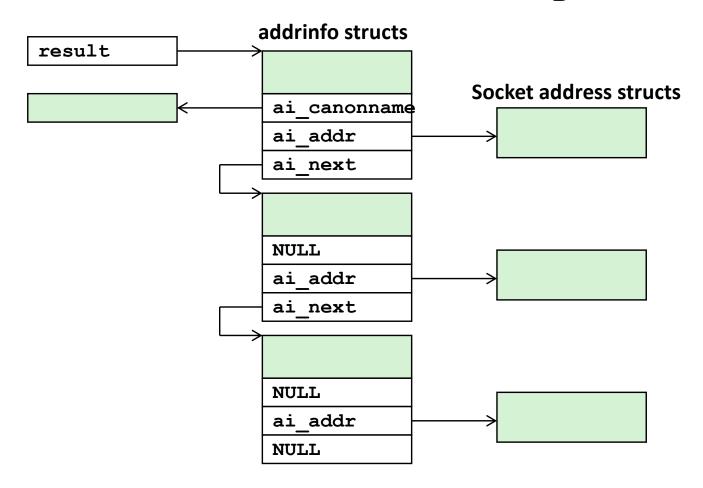
Disadvantages

- Somewhat complex
- Fortunately, a small number of usage patterns suffice in most cases.

Host and Service Conversion: getaddrinfo

- Given host and service, getaddrinfo returns result that points to a linked list of addrinfo structs, each of which points to a corresponding socket address struct, and which contains arguments for the sockets interface functions.
- Helper functions:
 - freeadderinfo frees the entire linked list.
 - gai_strerror converts error code to an error message.

Linked List Returned by getaddrinfo



- Clients: walk this list, trying each socket address in turn, until the calls to socket and connect succeed.
- Servers: walk the list until calls to socket and bind succeed.

addrinfo Struct

- Each addrinfo struct returned by getaddrinfo contains arguments that can be passed directly to socket function.
- Also points to a socket address struct that can be passed directly to connect and bind functions.

Host and Service Conversion: getnameinfo

- getnameinfo is the inverse of getaddrinfo, converting a socket address to the corresponding host and service.
 - Replaces obsolete gethostbyaddr and getservbyport funcs.
 - Reentrant and protocol independent.

Conversion Example

```
#include "csapp.h"
int main(int argc, char **argv)
   struct addrinfo *p, *listp, hints;
   char buf[MAXLINE];
   int rc, flags;
   /* Get a list of addrinfo records */
   memset(&hints, 0, sizeof(struct addrinfo));
   hints.ai family = AF INET; /* IPv4 only */
   hints.ai socktype = SOCK STREAM; /* Connections only */
   if ((rc = getaddrinfo(argv[1], NULL, &hints, &listp)) != 0) {
        fprintf(stderr, "getaddrinfo error: %s\n", gai strerror(rc));
       exit(1);
                                                              hostinfo.c
```

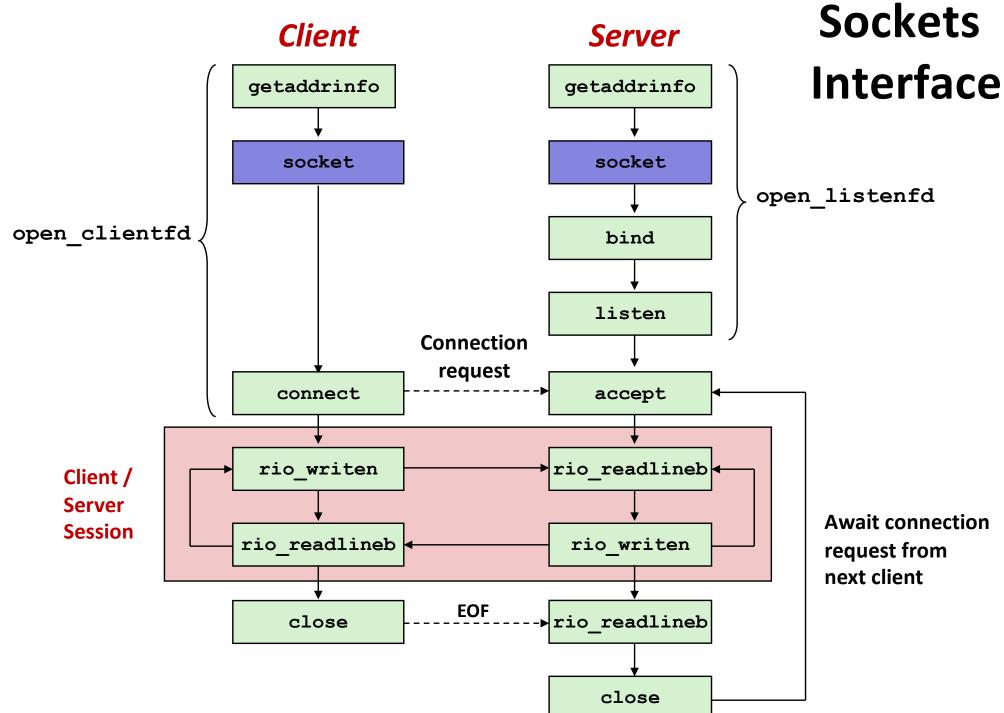
Conversion Example (cont)

Running hostinfo

```
whaleshark> ./hostinfo localhost
127.0.0.1

whaleshark> ./hostinfo whaleshark.ics.cs.cmu.edu
128.2.210.175

whaleshark> ./hostinfo twitter.com
199.16.156.230
199.16.156.38
199.16.156.102
199.16.156.198
```

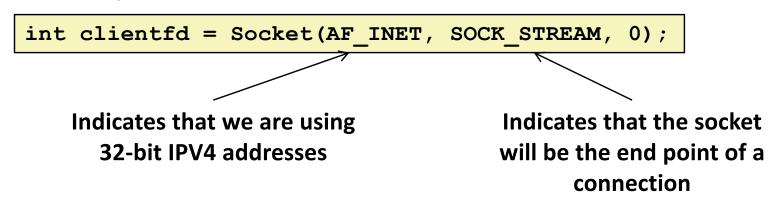


Sockets Interface: socket

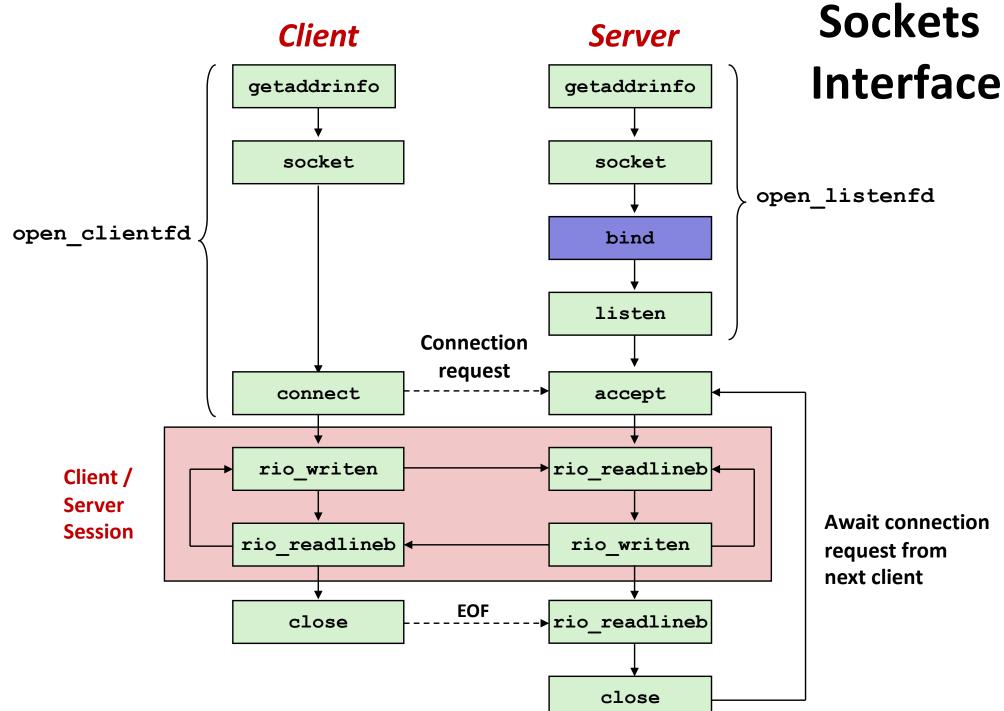
Clients and servers use the socket function to create a socket descriptor:

```
int socket(int domain, int type, int protocol)
```

Example:



Protocol specific! Best practice is to use getaddrinfo to generate the parameters automatically, so that code is protocol independent.



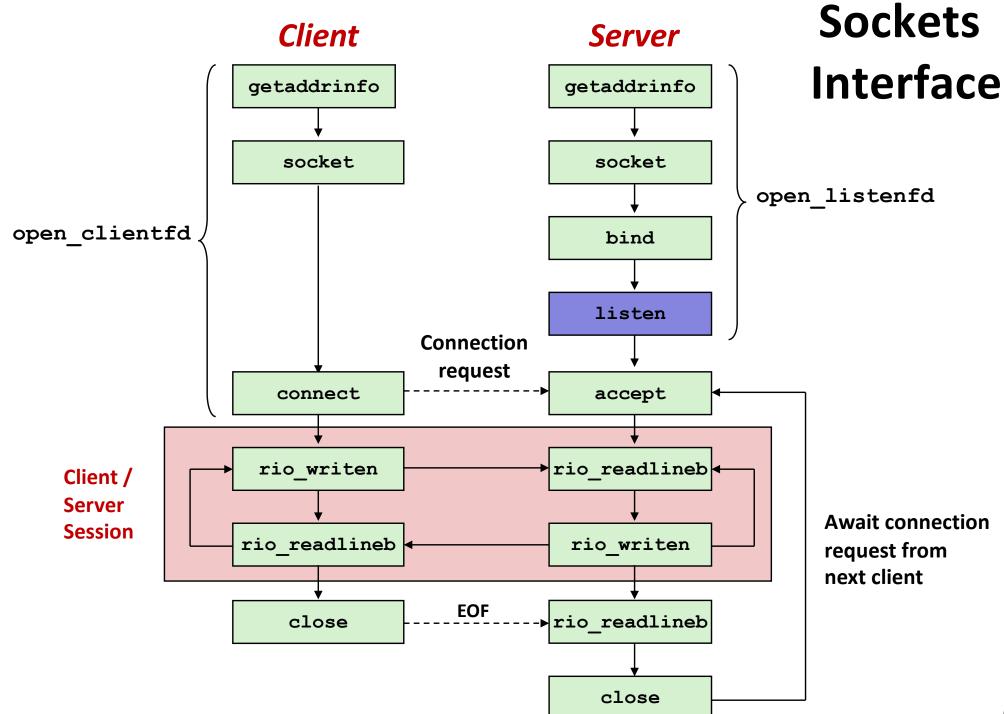
Sockets Interface: bind

A server uses bind to ask the kernel to associate the server's socket address with a socket descriptor:

```
int bind(int sockfd, SA *addr, socklen_t addrlen);
```

- The process can read bytes that arrive on the connection whose endpoint is addr by reading from descriptor sockfd.
- Similarly, writes to sockfd are transferred along connection whose endpoint is addr.

Best practice is to use getaddrinfo to supply the arguments addr and addrlen.

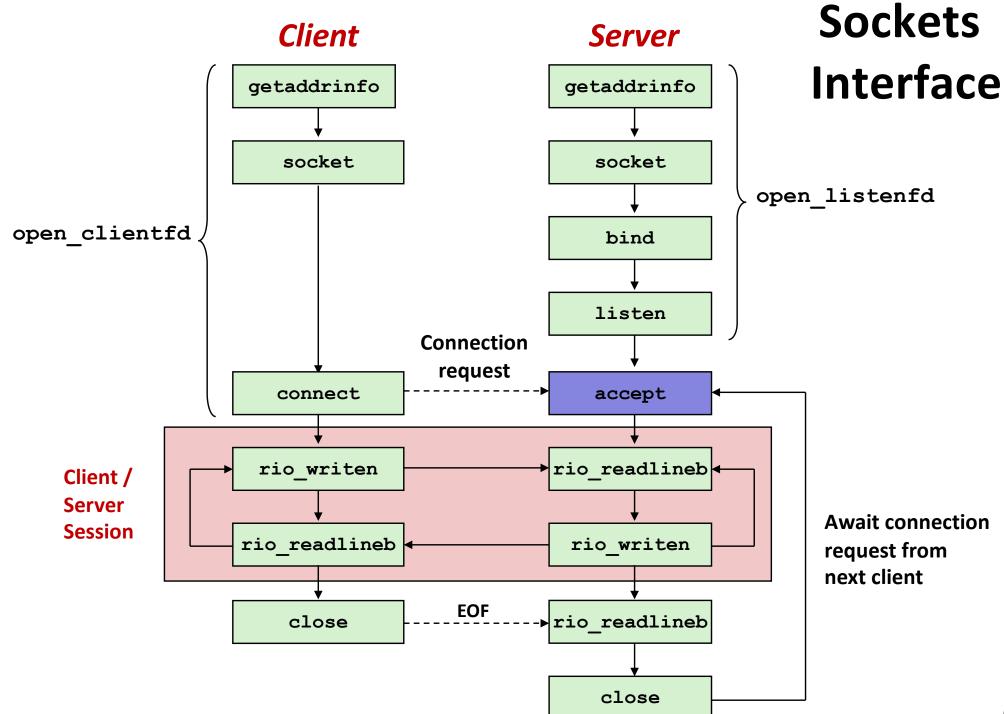


Sockets Interface: listen

- By default, kernel assumes that descriptor from socket function is an active socket that will be on the client end of a connection.
- A server calls the listen function to tell the kernel that a descriptor will be used by a server rather than a client:

```
int listen(int sockfd, int backlog);
```

- Converts sockfd from an active socket to a listening socket that can accept connection requests from clients.
- backlog is a hint about the number of outstanding connection requests that the kernel should queue up before starting to refuse requests.

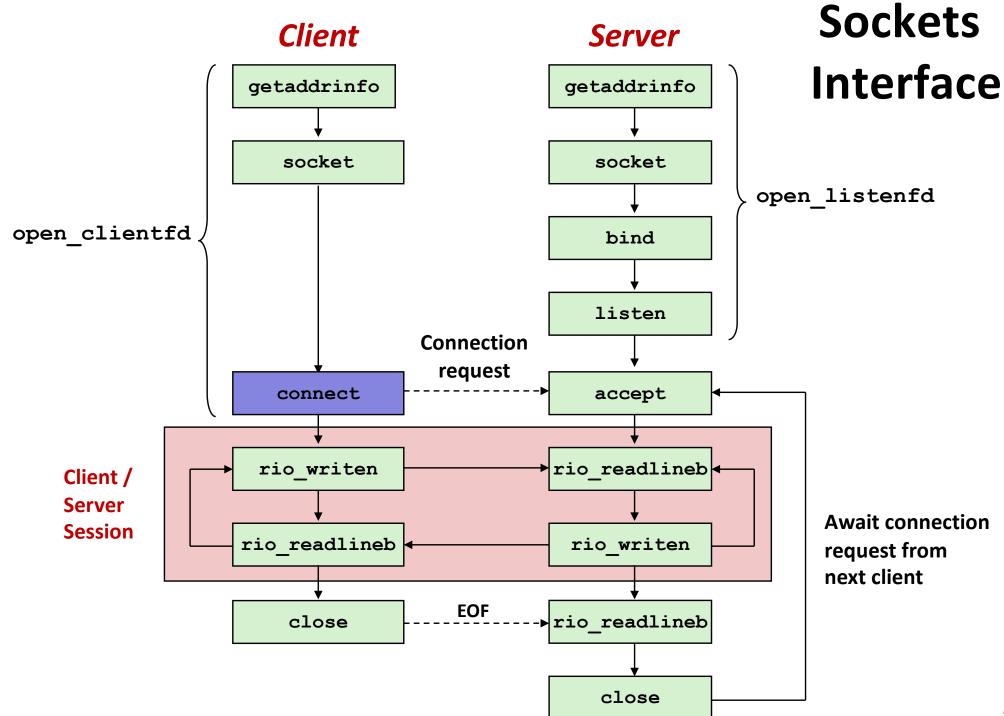


Sockets Interface: accept

Servers wait for connection requests from clients by calling accept:

```
int accept(int listenfd, SA *addr, int *addrlen);
```

- Waits for connection request to arrive on the connection bound to listenfd, then fills in client's socket address in addr and size of the socket address in addrlen.
- Returns a connected descriptor that can be used to communicate with the client via Unix I/O routines.



Sockets Interface: connect

A client establishes a connection with a server by calling connect:

```
int connect(int clientfd, SA *addr, socklen_t addrlen);
```

- Attempts to establish a connection with server at socket address addr
 - If successful, then clientfd is now ready for reading and writing.
 - Resulting connection is characterized by socket pair

```
(x:y, addr.sin_addr:addr.sin_port)
```

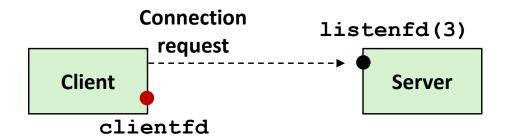
- x is client address
- y is ephemeral port that uniquely identifies client process on client host

Best practice is to use getaddrinfo to supply the arguments addr and addrlen.

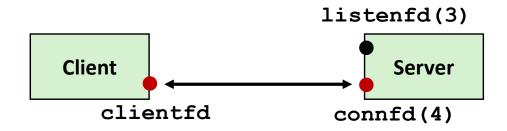
accept Illustrated



1. Server blocks in accept, waiting for connection request on listening descriptor listenfd



2. Client makes connection request by calling and blocking in connect



3. Server returns connfd from accept. Client returns from connect. Connection is now established between clientfd and connfd

Connected vs. Listening Descriptors

Listening descriptor

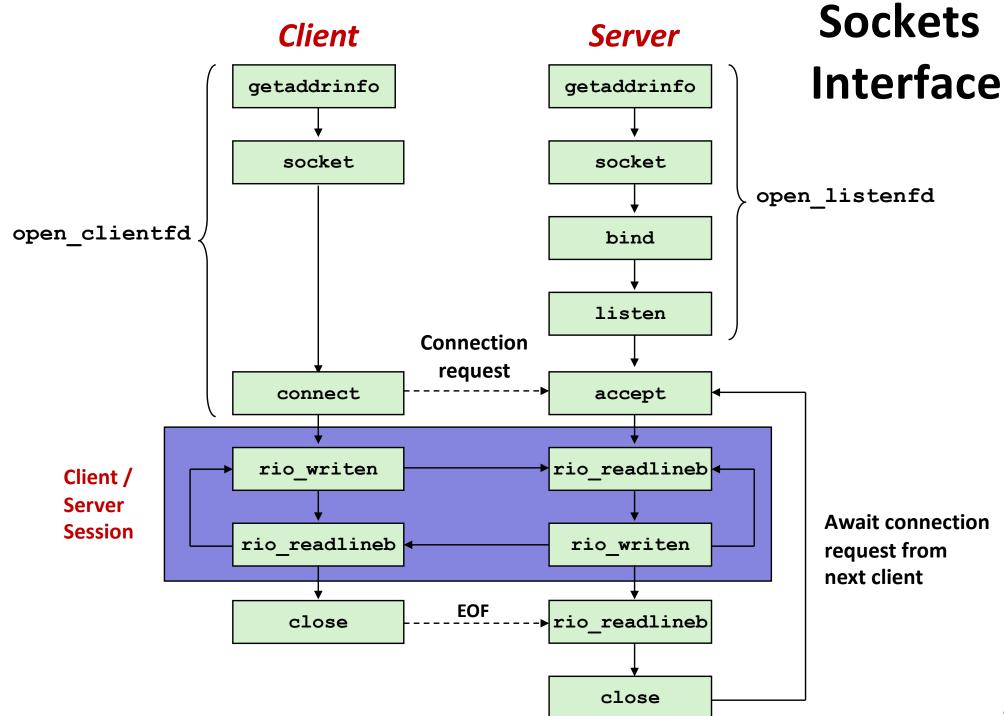
- End point for client connection requests
- Created once and exists for lifetime of the server

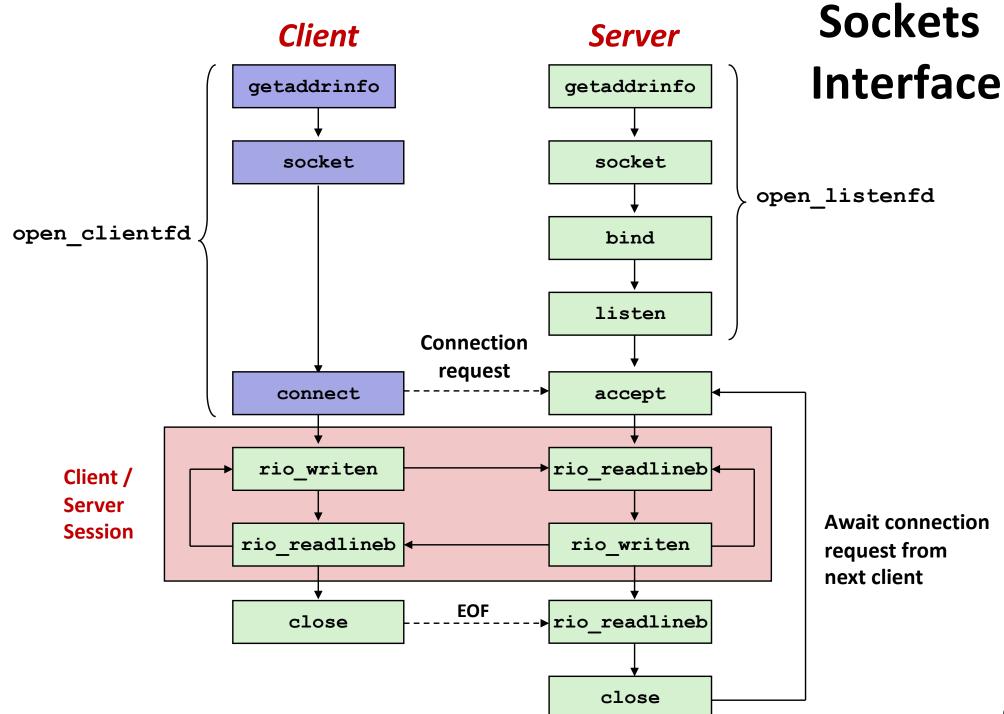
Connected descriptor

- End point of the connection between client and server
- A new descriptor is created each time the server accepts a connection request from a client
- Exists only as long as it takes to service client

Why the distinction?

- Allows for concurrent servers that can communicate over many client connections simultaneously
 - E.g., Each time we receive a new request, we fork a child to handle the request



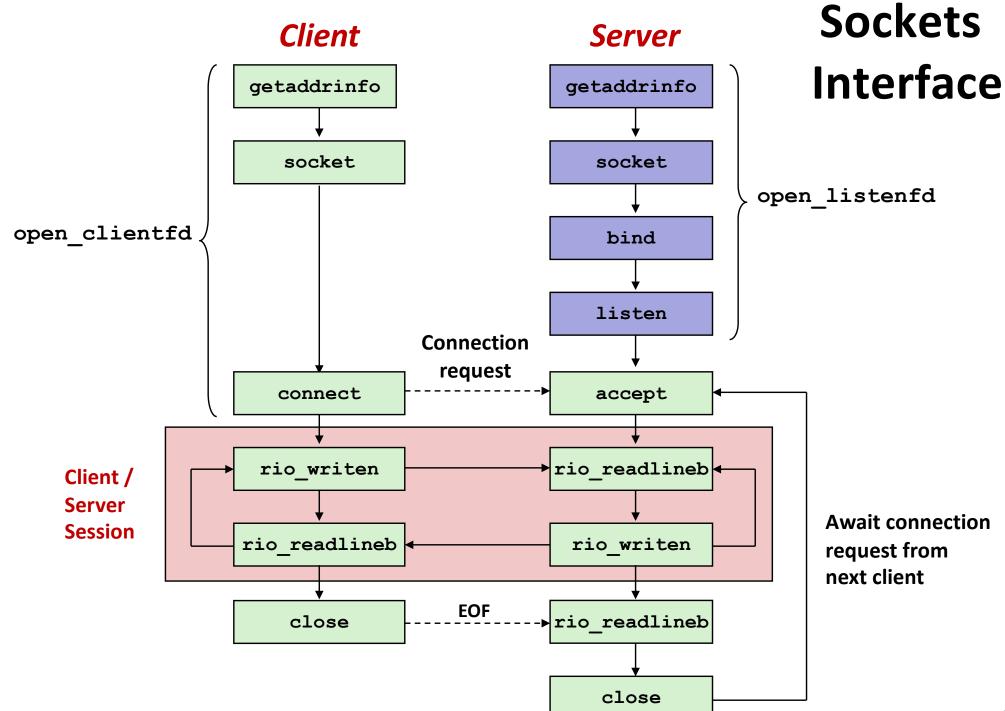


Sockets Helper: open_clientfd

Establish a connection with a server

Sockets Helper: open_clientfd (cont)

```
/* Walk the list for one that we can successfully connect to */
for (p = listp; p; p = p->ai next) {
    /* Create a socket descriptor */
    if ((clientfd = socket(p->ai family, p->ai socktype,
                           p->ai protocol)) < 0)</pre>
        continue; /* Socket failed, try the next */
    /* Connect to the server */
    if (connect(clientfd, p->ai addr, p->ai addrlen) != -1)
       break: /* Success */
    Close(clientfd); /* Connect failed, try another */
}
/* Clean up */
Freeaddrinfo(listp);
if (!p) /* All connects failed */
    return -1;
else /* The last connect succeeded */
   return clientfd;
                                                           csapp.c
```



Sockets Helper: open_listenfd

 Create a listening descriptor that can be used to accept connection requests from clients.

Sockets Helper: open_listenfd (cont)

```
/* Walk the list for one that we can bind to */
for (p = listp; p; p = p->ai next) {
   /* Create a socket descriptor */
    if ((listenfd = socket(p->ai family, p->ai socktype,
                           p->ai protocol)) < 0)
        continue; /* Socket failed, try the next */
   /* Eliminates "Address already in use" error from bind */
    Setsockopt(listenfd, SOL SOCKET, SO REUSEADDR,
               (const void *)&optval , sizeof(int));
   /* Bind the descriptor to the address */
    if (bind(listenfd, p->ai addr, p->ai addrlen) == 0)
       break; /* Success */
   Close(listenfd); /* Bind failed, try the next */
                                                         csapp.c
```

Sockets Helper: open_listenfd (cont)

```
/* Clean up */
Freeaddrinfo(listp);
if (!p) /* No address worked */
    return -1;

/* Make it a listening socket ready to accept conn. requests */
if (listen(listenfd, LISTENQ) < 0) {
    Close(listenfd);
    return -1;
}
return listenfd;
}</pre>
```

■ Key point: open_clientfd and open_listenfd are both independent of any particular version of IP.

Echo Client: Main Routine

```
#include "csapp.h"
int main(int argc, char **argv)
{
    int clientfd;
    char *host, *port, buf[MAXLINE];
    rio t rio;
    host = arqv[1];
    port = arqv[2];
    clientfd = Open clientfd(host, port);
    Rio readinitb(&rio, clientfd);
    while (Fgets(buf, MAXLINE, stdin) != NULL) {
       Rio writen(clientfd, buf, strlen(buf));
       Rio readlineb(&rio, buf, MAXLINE);
       Fputs(buf, stdout);
    Close(clientfd);
    exit(0);
                                                  echoclient.c
```

Iterative Echo Server: Main Routine

```
#include "csapp.h"
void echo(int connfd);
int main(int argc, char **argv)
    int listenfd, connfd;
    socklen t clientlen;
    struct sockaddr storage clientaddr; /* Enough room for any addr */
    char client hostname[MAXLINE], client port[MAXLINE];
    listenfd = Open listenfd(argv[1]);
    while (1) {
       clientlen = sizeof(struct sockaddr storage); /* Important! */
       connfd = Accept(listenfd, (SA *)&clientaddr, &clientlen);
       Getnameinfo((SA *) &clientaddr, clientlen,
                    client hostname, MAXLINE, client port, MAXLINE, 0);
       printf("Connected to (%s, %s)\n", client hostname, client port);
       echo(connfd);
       Close (connfd);
    exit(0);
                                                               echoserveri.c
```

Echo Server: echo function

- The server uses RIO to read and echo text lines until EOF (end-of-file) condition is encountered.
 - EOF condition caused by client calling close (clientfd)

```
void echo(int connfd)
{
    size_t n;
    char buf[MAXLINE];
    rio_t rio;

    Rio_readinitb(&rio, connfd);
    while((n = Rio_readlineb(&rio, buf, MAXLINE)) != 0) {
        printf("server received %d bytes\n", (int)n);
        Rio_writen(connfd, buf, n);
    }
}
```

Testing Servers Using telnet (or nc)

- The telnet program is invaluable for testing servers that transmit ASCII strings over Internet connections
 - Our simple echo server
 - Web servers
 - Mail servers

Usage:

- linux> telnet <host> <portnumber>
- Creates a connection with a server running on <host> and listening on port <portnumber>

Testing the Echo Server With telnet

```
server$./echoserveri 20000
Connected to (61.42.142.131, 63936)
server received 10 bytes
server received 7 bytes
moon$ telnet 122.38.251.9 20000
Trying 128.2.210.175...
Connected to 122.38.251.9.
Escape character is '^]'.
Hi there!
Hi there!
Howdy!
Howdy!
^1
telnet> quit
Connection closed.
moon$
```