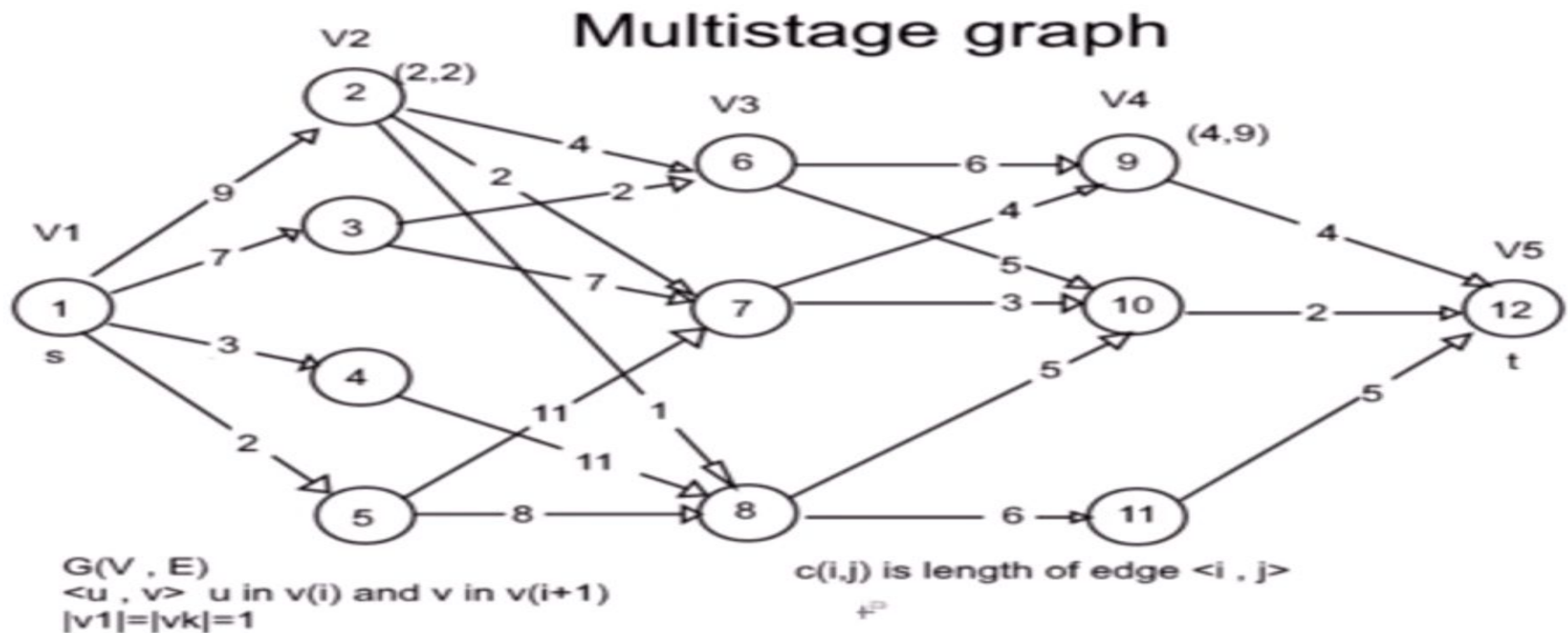


Finding Shortest Path in Multistage Graph using Dynamic Programming

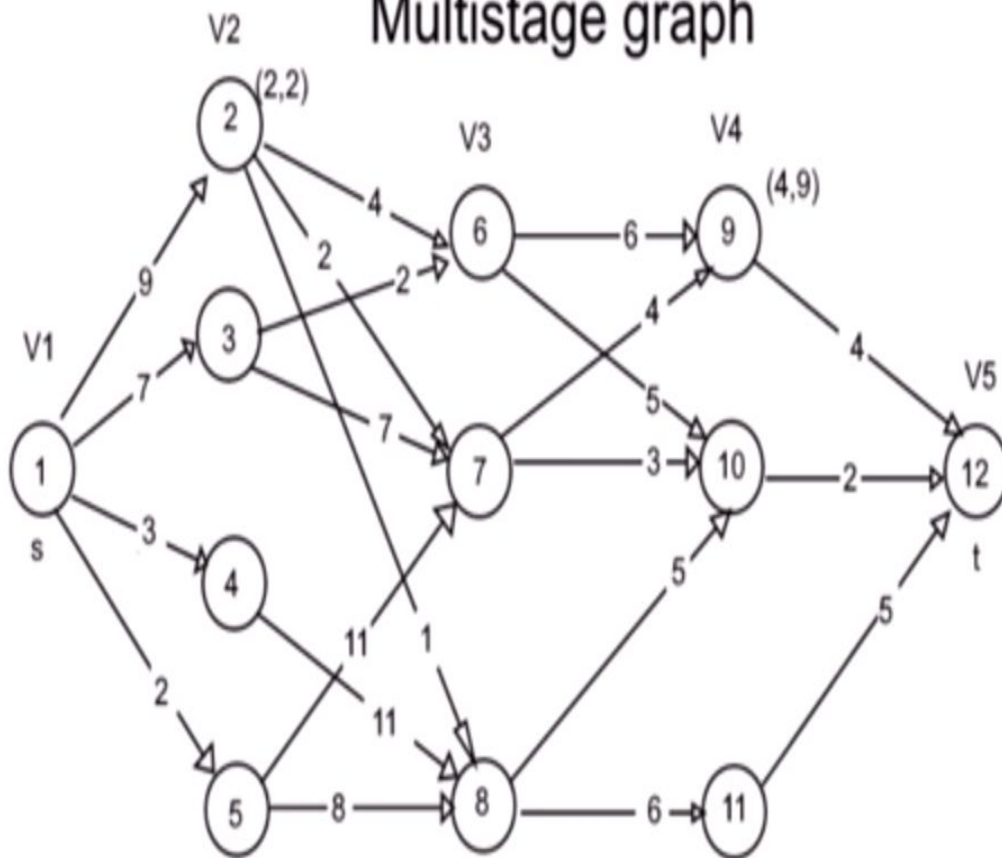
Multistage Graph

- To find the shortest path between the source vertex s and the destination vertex t .
- A multistage graph is a directed graph which is divided into stages V_1, V_2, \dots
- Vertices from one stage are connected to vertices of next stage (no edges between vertices of the same stage and from a vertex of current stage to previous stage).
- The first and the last stage have single vertex.



Applying Greedy approach to solve

Multistage graph



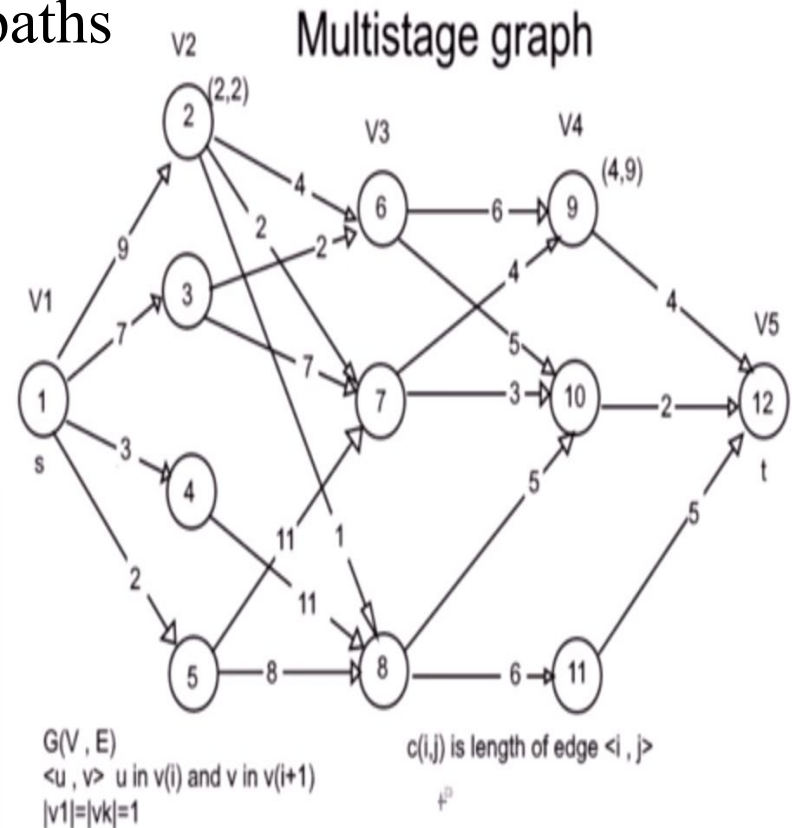
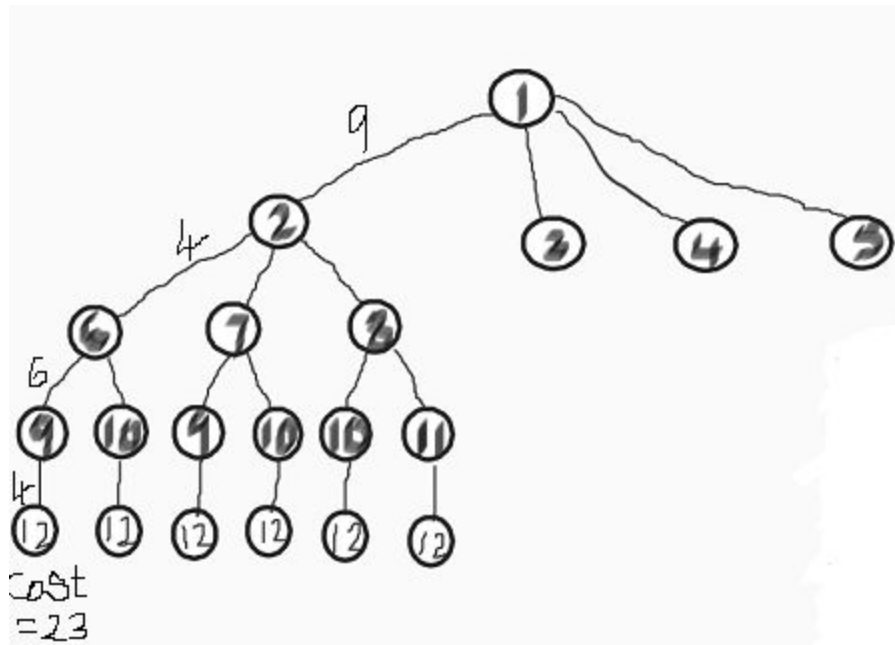
$G(V, E)$
 $\langle u, v \rangle$ u in $v(i)$ and v in $v(i+1)$
 $|v1|=|vk|=1$

$c(i,j)$ is length of edge $\langle i, j \rangle$
 \dagger^D

- Greedy Choice 1:
- Edge: (1,5) (5,8) (8,10) (10,12)
- Cost: $2 + 8 + 5 + 2 = 17$
- Choice 2:
- Edge: (1,2) (2,7) (7,10) (10,12)
- Cost: $9 + 2 + 3 + 2 = 16$
- Greedy choice fails

Applying Brute force to solve

Brute Force: Enumerate all possible paths
And find the minimum cost path.

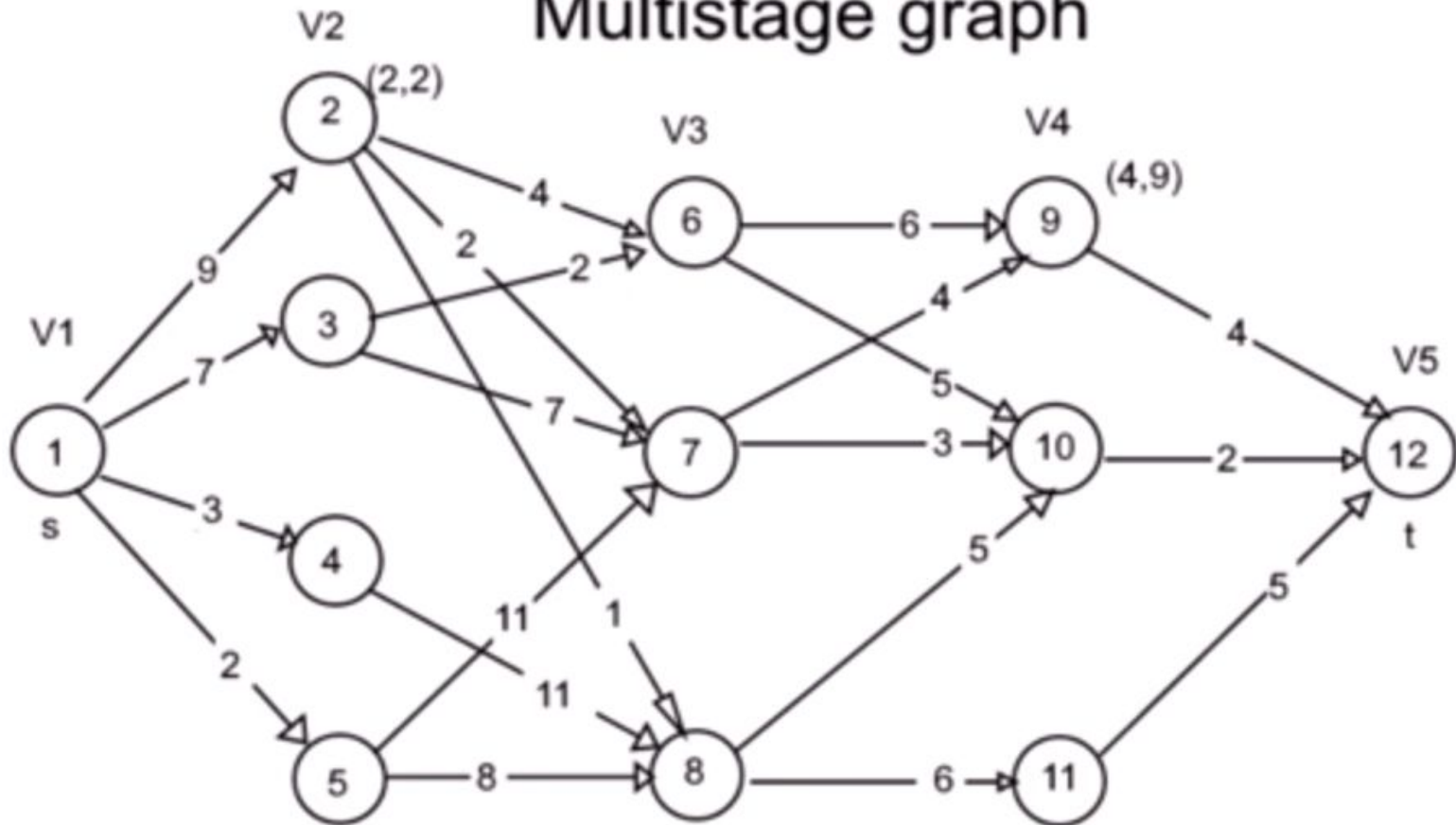


Solving using Dynamic Programming

- Forward approach
- Backward approach

Solving using Dynamic Programming

Multistage graph



$G(V, E)$

$\langle u, v \rangle$ u in $v(i)$ and v in $v(i+1)$

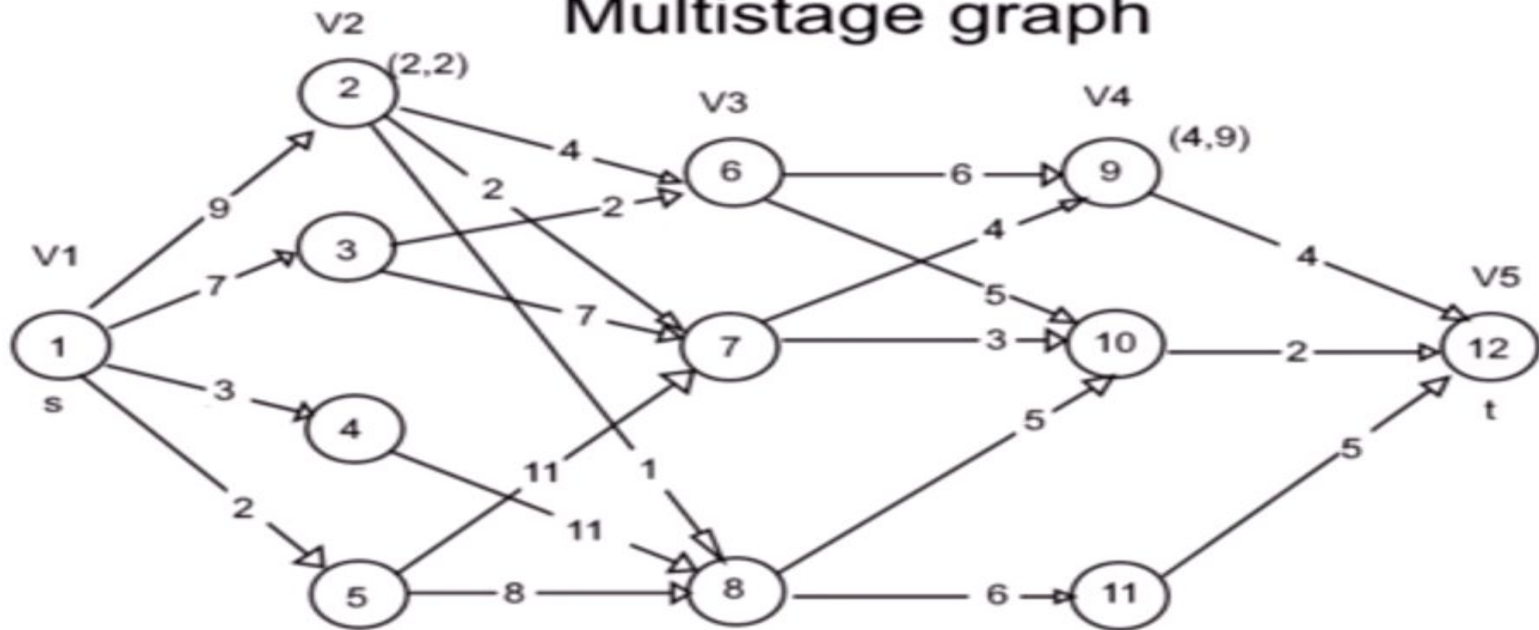
$|v1|=|vk|=1$

$c(i,j)$ is length of edge $\langle i, j \rangle$

t^2

Solving using Dynamic Programming

Multistage graph



$G(V, E)$
 $\langle u, v \rangle$ u in $v(i)$ and v in $v(i+1)$
 $|v_1| = |v_k| = 1$

$c(i, j)$ is length of edge $\langle i, j \rangle$
 ∞

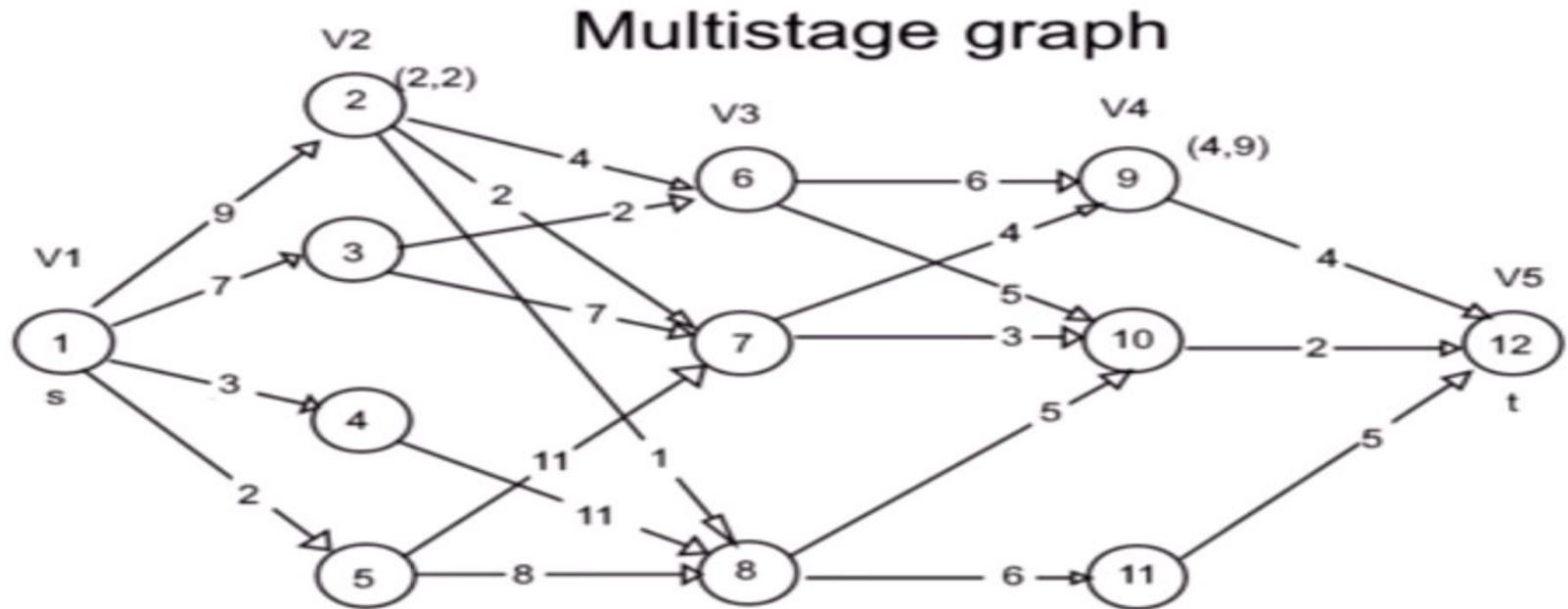
$$\text{cost}(i, j) = \min\{ c(j, l) + \text{cost}(i+1, l) \}$$

stage

node

length of edge $\langle j, l \rangle$
 if there is no edge infinity

Solving using Dynamic Programming: Forward approach

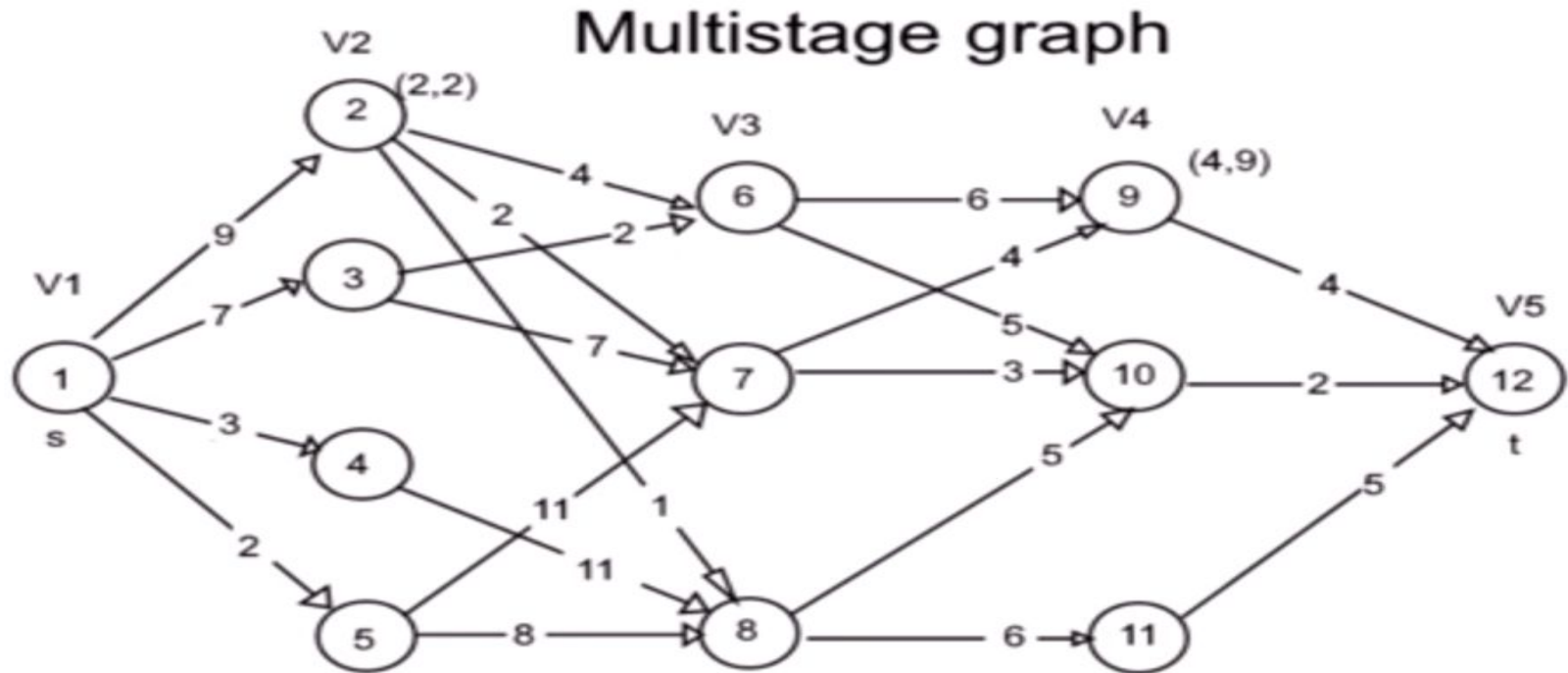


$$\begin{aligned} \text{cost}(4,9) &= c(9,12) = 4 \\ \text{cost}(4,10) &= c(10,12) = 2 \\ \text{cost}(4,11) &= c(11,12) = 5 \end{aligned}$$

path : 1 -> 2 -> 7 -> 10 -> 12
 1 -> 3 -> 6 -> 10 -> 12

$$\begin{aligned} \text{cost}(3,6) &= \min \{c(6,9) + \text{cost}(4,9), c(6,10) + \text{cost}(4,10)\} = \min \{6 + 4, 5 + 2\} = 7 \\ \text{cost}(3,7) &= \min \{c(7,9) + \text{cost}(4,9), c(7,10) + \text{cost}(4,10)\} = \min \{4 + 4, 3 + 2\} = 5 \\ \text{cost}(3,8) &= \min \{c(8,10) + \text{cost}(4,10), c(8,11) + \text{cost}(4,11)\} = \min \{5 + 2, 6 + 5\} = 7 \end{aligned}$$

Solving using Dynamic Programming: **Forward approach**



$$\text{cost}(2,2) = \min\{c(2,6) + \text{cost}(3,6), c(2,7) + \text{cost}(3,7), c(2,8) + \text{cost}(3,8)\} = \min\{4+7, 2+5, 1+7\} = 7$$

$$\text{cost}(2,3) = \min\{c(3,6) + \text{cost}(3,6), c(3,7) + \text{cost}(3,7)\} = \min\{2+7, 7+5\} = 9$$

$$\text{cost}(2,4) = \min\{c(4,8) + \text{cost}(3,8)\} = 11+7 = 18$$

$$\text{cost}(2,5) = \min\{c(5,7) + \text{cost}(3,7), c(5,8) + \text{cost}(3,8)\} = \min\{11+5, 8+7\} = 15$$

$$\begin{aligned} \text{cost}(4,9) &= c(9,12) = 4 \\ \text{cost}(4,10) &= c(10,12) = 2 \\ \text{cost}(4,11) &= c(11,12) = 5 \end{aligned}$$

path : 1 -> 2 -> 7 -> 10 -> 12
1 -> 3 -> 6 -> 10 -> 12

$$\begin{aligned} \text{cost}(3,6) &= \min \{c(6,9) + \text{cost}(4,9), c(6,10) + \text{cost}(4,10)\} = \min \{6 + 4, 5 + 2\} = 7 \\ \text{cost}(3,7) &= \min \{c(7,9) + \text{cost}(4,9), c(7,10) + \text{cost}(4,10)\} = \min \{4 + 4, 3 + 2\} = 5 \\ \text{cost}(3,8) &= \min \{c(8,10) + \text{cost}(4,10), c(8,11) + \text{cost}(4,11)\} = \min \{5 + 2, 6 + 5\} = 7 \end{aligned}$$

$$\begin{aligned} \text{cost}(2,2) &= \min \{c(2,6) + \text{cost}(3,6), c(2,7) + \text{cost}(3,7), c(2,8) + \text{cost}(3,8)\} = \min \{4 + 7, 2 + 5, 1 + 7\} = 7 \\ \text{cost}(2,3) &= \min \{c(3,6) + \text{cost}(3,6), c(3,7) + \text{cost}(3,7)\} = \min \{2 + 7, 7 + 5\} = 9 \\ \text{cost}(2,4) &= \min \{c(4,8) + \text{cost}(3,8)\} = 11 + 7 = 18 \\ \text{cost}(2,5) &= \min \{c(5,7) + \text{cost}(3,7), c(5,8) + \text{cost}(3,8)\} = \min \{11 + 5, 8 + 7\} = 15 \end{aligned}$$

verte x	1	2	3	4	5	6	7	8	9	10	11	12
Cost	16	7	9	18	15	7	5	7	4	2	5	0
d-dest inatio n	2/3	7	6	8	8	10	10	10	12	12	12	12

Solving using Dynamic Programming: **Forward approach**

$$\text{cost}(1,1) = \min\{ c(1,2) + \text{cost}(2,2) , c(1,3) + \text{cost}(2,3) \\ c(1,4) + \text{cost}(2,4), c(1,5) + \text{cost}(2,5) \}$$

$$= \min\{9 + 7 , 7 + 9, 3 + 18, 2 + 15\} = 16 \quad d(1,1) = 2,3$$

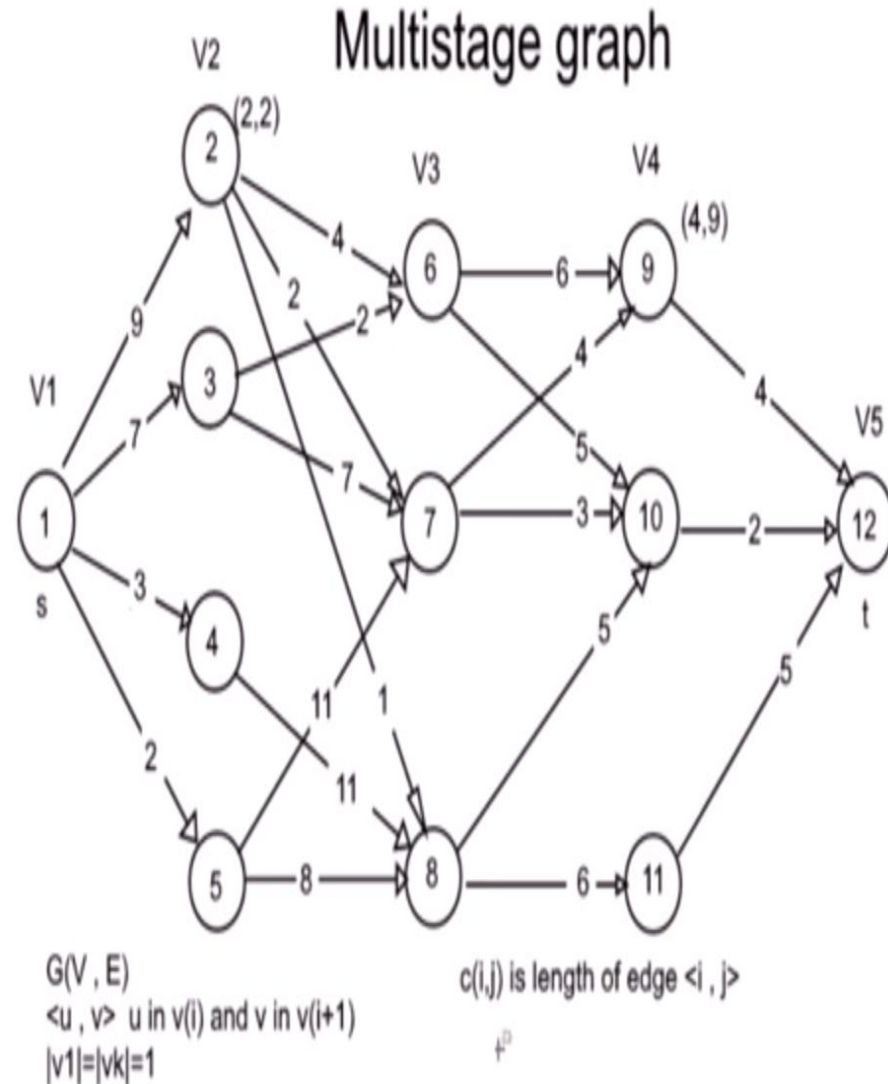
verte x	1	2	3	4	5	6	7	8	9	10	11	12
Cost	16	7	9	18	15	7	5	7	4	2	5	0
d-dest inatio n	2/3	7	6	8	8	10	10	10	12	12	12	12

Multistage Graph pseudo code : **forward approach**

```
1  Algorithm FGraph( $G, k, n, p$ )
2  // The input is a  $k$ -stage graph  $G = (V, E)$  with  $n$  vertices
3  // indexed in order of stages.  $E$  is a set of edges and  $c[i, j]$ 
4  // is the cost of  $\langle i, j \rangle$ .  $p[1 : k]$  is a minimum-cost path.
5  {
6       $cost[n] := 0.0$ ;
7      for  $j := n - 1$  to  $1$  step  $-1$  do
8      { // Compute  $cost[j]$ .
9          Let  $r$  be a vertex such that  $\langle j, r \rangle$  is an edge
10         of  $G$  and  $c[j, r] + cost[r]$  is minimum;
11          $cost[j] := c[j, r] + cost[r]$ ;
12          $d[j] := r$ ;
13     }
14     // Find a minimum-cost path.
15      $p[1] := 1$ ;  $p[k] := n$ ;
16     for  $j := 2$  to  $k - 1$  do  $p[j] := d[p[j - 1]]$ ;
17 }
```

Solving using Dynamic Programming: Backward approach

- $\text{bcost}(i,j)$: Minimum cost path from vertex s to vertex j in V_i .
- $\text{bcost}(i,j) = \min \{ \text{bcost}(i-1, k) + c(k,j) \}$
- $k \in V_{i-1}$
- $\text{bcost}(2,2) = 9$
- $\text{bcost}(2,3) = 7$
- $\text{bcost}(2,4) = 3$
- $\text{bcost}(2,5) = 2$
- $\text{bcost}(3,6) = \min(\text{bcost}(2,2) + c(2,6),$
- $\text{bcost}(2,3) + c(3,6))$
- $= \min \{ 9 + 4, 7 + 2 \} = 9$
- $\text{bcost}(3,7) = \min(\text{bcost}(2,2) + c(2,7),$
- $\text{bcost}(2,3) + c(3,7)$
- $\text{bcost}(2,5) + c(2,7))$
- $= \min \{ 9 + 2, 7 + 7, 2 + 11 \} = 11$



Solving using Dynamic Programming: **Backward approach**

$$bcost(3, 7) = 11$$

$$bcost(3, 8) = 10$$

$$bcost(4, 9) = 15$$

$$bcost(4, 10) = 14$$

$$bcost(4, 11) = 16$$

$$bcost(5, 12) = 16$$

Multistage Graph pseudo code: **backward approach**

```
1  Algorithm BGraph( $G, k, n, p$ )
2  // Same function as FGraph
3  {
4       $bcost[1] := 0.0;$ 
5      for  $j := 2$  to  $n$  do
6          { // Compute  $bcost[j]$ .
7              Let  $r$  be such that  $\langle r, j \rangle$  is an edge of
8               $G$  and  $bcost[r] + c[r, j]$  is minimum;
9               $bcost[j] := bcost[r] + c[r, j];$ 
10              $d[j] := r;$ 
11         }
12     // Find a minimum-cost path.
13      $p[1] := 1; p[k] := n;$ 
14     for  $j := k - 1$  to  $2$  do  $p[j] := d[p[j + 1]];$ 
15 }
```

Solve

Find minimum cost path from s to t in the multistage graph given below using:

a. Forward approach

b. Backward approach

