

# Concepts Notes

@saffron

## Logic

### logical operators

- $\neg, \wedge, \vee, \Rightarrow, \Leftrightarrow$

### quantifiers

- $\forall, \exists$

### helpful equivalences

- $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$
- $p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$
- $\neg(p \wedge q) \equiv \neg p \vee \neg q$
- $\neg(p \vee q) \equiv \neg p \wedge \neg q$
- $\neg(p \Rightarrow q) \equiv p \wedge \neg q$
- $p \Rightarrow q \equiv \neg q \Rightarrow \neg p$  (“contrapositive form”)
- $\neg(\forall x \in S, p(x)) \equiv \exists x \in S, \neg p(x)$
- $\neg(\exists x \in S, p(x)) \equiv \forall x \in S, \neg p(x)$

### definitions

truth tables can prove equivalences

tautology: always true / contradiction: always false

divisibility: Let  $a, b \in \mathbb{Z}$ . Then  $a \mid b \iff (\exists k \in \mathbb{Z}, ak = b)$

parity: an integer  $a$  is even iff  $2 \mid a$  and odd otherwise (i.e.  $2 \mid a + 1$ ).

## Proof Strategies

TODO AFSOC, contraposition, etc etc from your old notes.

- $(p \Rightarrow q)$  assume  $p$ , prove  $q$
- $(p \vee q \Rightarrow r)$  by cases: assume  $p$ , prove  $r$ . assume  $q$ , prove  $r$
- we find a contradiction

## Sets

- list notation:  $A = \{1, 2, 3, \dots\}$
- set builder notation:  $B = \{n : n \text{ is prime}\}$
- number sets
  - naturals  $\mathbb{N} = \{0, 1, 2, \dots\}$
  - integers  $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$
  - rationals  $\mathbb{Q} = \left\{ \frac{a}{b} : a, b \in \mathbb{Z} \wedge b \neq 0 \right\}$
  - reals  $\mathbb{R} = \{\text{the number line}\}$
  - closure properties
    - $+/\times$  of two  $\mathbb{N}$  is still  $\mathbb{N}$
    - $+/-/\times$  of two  $\mathbb{Z}$  is still  $\mathbb{Z}$

- $+ / - / \times / \div$  (except 0) of two  $\mathbb{Q}/\mathbb{R}$  is still  $\mathbb{Q}/\mathbb{R}$

definitions

- $x \notin X$  iff  $\neg(x \in X)$
- $X = \emptyset$  iff  $\forall x, x \notin X$
- $A \subseteq B$  iff  $\forall x \in A, x \in B$
- $X \cap Y = \{a : a \in X \wedge a \in Y\}$
- $X \cup Y = \{a : a \in X \vee a \in Y\}$
- $X \setminus Y = \{a : a \in X \wedge a \notin Y\}$
- $X \times Y = \{(x, y) : x \in X \wedge y \in Y\}$
- $X = Y$  iff  $X \subseteq Y \wedge Y \subseteq X$
- $\bigcap_{i \in I} X_i = \{a : \forall i \in I, a \in X_i\}$
- $\bigcup_{i \in I} X_i = \{a : \exists i \in I, a \in X_i\}$
- $\mathcal{P}(A) = \{X : X \subseteq A\}$

de Morgan's for sets

- $X \setminus (Y \cap Z) = (X \setminus Y) \cup (X \setminus Z)$
- $X \setminus (Y \cup Z) = (X \setminus Y) \cap (X \setminus Z)$
- $X \setminus \left(\bigcap_{i \in I} Y_i\right) = \bigcup_{i \in I} (X \setminus Y_i)$
- $X \setminus \left(\bigcup_{i \in I} Y_i\right) = \bigcap_{i \in I} (X \setminus Y_i)$

## Functions

definitions

- a function  $f : X \rightarrow Y$  is an assignment that satisfies
  - totality:  $\forall x \in X, f(x)$  is defined (i.e. domain =  $X$ )
  - existence:  $\forall x \in X, f(x) \in Y$  (i.e. codomain =  $Y$ )
  - uniqueness:  $\forall x, x' \in X, (x = x' \Rightarrow f(x) = f(x'))$  (i.e. VLT)
- $\text{Gr}(f) = \{(x, y) \in X \times Y : f(x) = y\}$
- let  $f : X \rightarrow Y, g : Y \rightarrow Z$ . then  $g \circ f : X \rightarrow Z = g(f(x))$
- $f = g$  iff same domain, same codomain, and  $\forall x \in X, f(x) = g(x)$
- $\text{id}_X : X \rightarrow X$  is  $\text{id}_X(x) = x$
- characteristic function of  $U$  in  $X$  is  $\chi_U(a)$  is 1 if  $a \in U$ , 0 if not.
- image of  $U$  under  $f$  is  $f[U] = \{y \in Y : \exists x \in U, y = f(x)\}$
- preimage of  $V$  under  $f$  is  $f^{-1}[V] = \{x \in X : f(x) \in V\}$
- $f$  is injective iff  $\forall a, b \in X, (f(a) = f(b) \Rightarrow a = b)$
- $f$  is surjective iff  $\forall y \in Y, \exists x \in X, f(x) = y$
- $f$  is bijective iff  $f$  is both injective and surjective
- a left inverse for  $f$  is a function  $g : Y \rightarrow X$  s.t.  $g \circ f = \text{id}_X$ 
  - this implies  $f$  is injective
- a right inverse for  $f$  is a function  $g : Y \rightarrow X$  s.t.  $f \circ g = \text{id}_Y$ 
  - this implies  $f$  is surjective
- an inverse (two-sided) for  $f$  is a function  $g$  that is both left and right inverse, also denoted  $f^{-1}$ 
  - this implies  $f$  is bijective

## Induction

The Fundamental Theorem of Arithmetic (Existence)

- Every natural number  $n \geq 2$  can be expressed as a product of primes

recursive function definitions

$$\sum_{k=1}^n a_k = \begin{cases} \sum_{k=1}^0 a_k = 0 \\ \sum_{k=1}^{m+1} a_k = a_{m+1} + \sum_{k=1}^m a_k \text{ if } n = m + 1 \end{cases}$$

$$\prod_{k=1}^n a_k = \begin{cases} \prod_{k=1}^0 a_k = 1 \\ \prod_{k=1}^{m+1} a_k = a_{m+1} \cdot \prod_{k=1}^m a_k \text{ if } n = m + 1 \end{cases}$$

$$n! = \sum_{k=1}^n k$$

weak induction formula

- to prove  $\forall n \in \mathbb{N}$  with  $n \geq n_0$ ,  $p(n)$ :
  - (Base Case) Verify  $p(n_0)$  holds
  - (Inductive Step) Verify  $\forall n \geq n_0, p(n) \Rightarrow p(n + 1)$

strong induction formula

- to prove  $\forall n \in \mathbb{N}$  with  $n \geq n_0$ ,  $p(n)$ :
  - (Base Case) Verify  $p(n_0)$  holds
  - (Inductive Step) Verify  $\forall n \geq n_0, (p(i) \text{ is true for all } i \text{ such that } n_0 \leq i \leq n) \Rightarrow p(n + 1)$

strong induction, multiple base cases formula

- to prove  $\forall n \in \mathbb{N}$  with  $n \geq n_0$ ,  $p(n)$  with base cases from  $n_0$  to  $n_1$ :
  - (Base Case) Verify all base cases hold
  - (Inductive Step) Verify  $\forall n \geq n_1, (p(i) \text{ is true for all } i \text{ such that } n_0 \leq i \leq n) \Rightarrow p(n + 1)$

## Relations

relation equivalence

- relations  $R = S$  iff same domain, codomain, and  $\forall x, y \in X \times Y, xRy \Leftrightarrow xSy$

properties of relations (let  $R$  be a relation on  $X$ )

- $R$  is reflexive iff  $\forall x \in X, xRx$
- $R$  is symmetric iff  $\forall x, y \in X, xRy \Rightarrow yRx$
- $R$  is antisymmetric iff  $\forall x, y \in X, xRy \wedge yRx \Rightarrow y = x$
- $R$  is transitive iff  $\forall x, y, z \in X, xRy \wedge yRz \Rightarrow xRz$
- $R$  is an equivalence relation iff  $R$  is reflexive, symmetric, and transitive
- $R$  is a partial order iff  $R$  is reflexive, antisymmetric, and transitive (may be denoted  $\preceq$ )
  - if  $\preceq$  is a partial order of  $X$ , we say  $(X, \preceq)$  is a poset

- $R$  is total iff  $\forall x, y \in X, (x \neq y \Rightarrow xRy \vee yRx)$  (every two distinct elements are relatable)
- $R$  is a total order iff  $R$  is a partial order and total

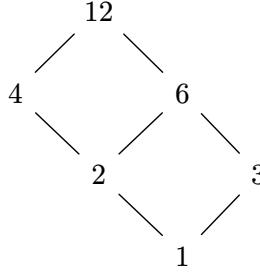


Figure 1: Hasse diagram for the poset  $(\{1, 2, 3, 4, 6, 12\}, |)$

modular congruence

- let  $n$  be positive  $\in \mathbb{Z}$ ,  $a, b \in \mathbb{Z}$ . then  $a \equiv b \pmod{n}$  aka  $a \equiv_n b$  iff  $n \mid (b - a)$

equivalence class of  $a$  under an equivalence relation  $\sim$  on  $X$

- $[a]_\sim = \{b \in X : a \sim b\}$
- the set of items related to  $a$
- $[a]_\sim = [b]_\sim$  iff  $a \sim b$

the quotient (letting  $\sim$  be an equivalence relation on  $X$ )

- $X / \sim = \{[a]_\sim : a \in X\}$
- the collection of all equivalences classes in  $X$  under  $\sim$

$P$  is a partition of  $X$  iff

- $\emptyset \notin P$
- $\forall U, V \in P, (U \neq V \Rightarrow U \cap V = \emptyset)$  i.e. different sets have no elements in common
- $\bigcup_{U \in P} U = X$  i.e. each partition together becomes  $P$

theorem for equivalence relations/partitions, with a nonempty set  $X$

- let  $\sim$  be an equivalence relation on  $X$ . then,  $X / \sim$  is a partition for  $X$
- let  $P$  be a partition for  $X$ . then there exists a unique equivalence relation on  $P$  whose quotient is  $P$ .

## Cardinality

- $[n] = \{1, 2, 3, \dots, n\}$
- infinite sets
  - $X$  is finite iff there exists a bijection  $f : [n] \rightarrow X$  for some natural  $n$
  - otherwise,  $X$  is infinite
- for finite  $X$ , then  $|X|$  is the natural  $n$  where there exists bijective  $f : [n] \rightarrow X$
- let  $X, Y$  be finite sets
  - $\exists$  an injection  $f : X \rightarrow Y$ , then  $|X| \leq |Y|$
  - $\exists$  a surjection  $f : X \rightarrow Y$ , then  $|X| \geq |Y|$
  - $\exists$  a bijection  $f : X \rightarrow Y$ , then  $|X| = |Y|$
  - if  $A \subseteq X$ , then  $A$  finite and  $|A| \leq |X|$
  - if  $A \subseteq X$ , then  $A$  finite and  $|A| \leq |X|$
  - $X \cap Y$  is finite and  $|X \cap Y| \leq \min\{|X|, |Y|\}$
  - $X \cup Y$  is finite and  $|X \cup Y| = |X| + |Y| - |X \cap Y|$

- $X \times Y$  is finite and  $|X \times Y| = |X| \cdot |Y|$
- $X$  is infinitely countable if  $\exists$  a bijection  $f : \mathbb{N} \rightarrow X$
- $X$  is countable if  $X$  is either finite or infinitely countable
- $X$  is uncountable if  $X$  is not countable (infinitely or otherwise)
- if  $X$  and  $Y$  are countable, then  $X \times Y$  is countable
- let  $X$  be countable
  - if  $\exists$  an injection  $f : Y \rightarrow X$ , then  $Y$  is countable
  - if  $\exists$  a surjection  $f : X \rightarrow Y$ , then  $Y$  is countable
  - all  $A \subseteq X$  are countable
- diagonalization (prove  $X$  uncountable)
  - show that every  $f : \mathbb{N} \rightarrow X$  is not surjective by finding a  $b \in X$  that disagrees with  $f(n)$  with something involving  $n$
- let  $X$  be any set
  - $|X| = |Y|$  iff  $\exists$  a bijection  $f : X \rightarrow Y$
  - $|X| \leq |Y|$  iff  $\exists$  an injection  $f : X \rightarrow Y$
  - $|X| < |Y|$  iff  $\exists$  an injection  $f : X \rightarrow Y \wedge \nexists$  a surjection  $g : X \rightarrow Y$
- Cantor's:  $|X| < |\mathcal{P}(X)|$

## Number Theory

- Well Ordering Principle: Any nonempty subset  $X \subseteq \mathbb{N}$  has the least element
  - use in “let  $n$  be the least element, AFSOC something, wait a minute  $m < n$ ” proofs
- division theorem: Let  $a, b \in \mathbb{Z}, b \neq 0$ . There is a quotient and a remainder.
  - (59 and 8?  $59 = 8 \cdot 7 + 3$ )
- let  $a, b \in \mathbb{Z}$ 
  - $c \in \mathbb{Z}$  is a common divisor of  $a$  and  $b$  if  $c$  divides both  $a$  and  $b$
  - $d \in \mathbb{Z}$  is a greatest common divisor of  $a$  and  $b$  if
    1.  $d$  divides  $a$  and  $b$
    2. if  $c$  is a common divisor of  $a$  and  $b$ , then  $c | d$
  - $a$  and  $b$  are coprime iff  $\gcd(a, b) = 1$
- Euclidean Algorithm
  - lemma:  $a = qb + r \Rightarrow \gcd(a, b) = \gcd(b, r)$
  - let us use an example. find  $\gcd(148, 40)$ .

$$\begin{aligned} 148 &= 40 \cdot 3 + 28 \\ 40 &= 28 \cdot 1 + 12 \\ 28 &= 12 \cdot 2 + 4 \\ 12 &= 4 \cdot 3 + 0 \end{aligned}$$

the last nonzero remainder, 4, is the gcd.

- Linear Diophantine Equations
  - Bézout’s Lemma: let  $a, b, c \in \mathbb{Z}$ . the equation  $ax + by = c$  has integer solutions for  $x, y$  iff  $\gcd(a, b) | c$
  - let us use an example, via the backwards Euclidean Algorithm. find  $252x + 198y = 36$ .

we first find  $\gcd(a, b) = \gcd(252, 198)$ .

$$\begin{aligned} 252 &= 198 \cdot 1 + 54 & (1) \\ 198 &= 54 \cdot 3 + 36 & (2) \\ 54 &= 36 \cdot 1 + 18 & (3) \\ 36 &= 18 \cdot 2 + 0 \end{aligned}$$

Thus,  $\gcd(252, 198) = 18$ . Since  $18 | 36$ , by Bézout’s, we have an integer solution. We proceed by backtracking (starting from the most recent iterations).

$$\begin{aligned} 18 &= 54 - 36 \cdot 1 & \text{(from 3)} \\ 18 &= 54 - (198 - 54 \cdot 3) & \text{(from 2)} \\ 18 &= (252 - 198) - (198 - (252 - 198) \cdot 3) & \text{(from 1)} \\ 18 &= 252 \cdot 4 - 198 \cdot 5 \\ 36 &= 252 \cdot 8 - 198 \cdot 10 \end{aligned}$$

Thus,  $x = 8$  and  $y = -10$ .

- Multiplicative Inverses Modulo  $n$ : let  $a, n \in \mathbb{Z}$  and  $n \geq 1$ .
  - we say  $a$  has a multiplicative inverse for  $a$  modulo  $n$  if  $ax \equiv_n 1$  has an integer solution for  $x$ . then  $x$  is that multiplicative inverse

- $a$  has a multiplicative inverse modulo  $n$  iff  $\gcd(a, n) = 1$
- find this inverse similarly to a Linear Dio. Euclidean to get to  $\gcd = 1$  and backsolve.
- (a normal multiplicative inverse is  $ax = 1$ )
- Divisibility Properties
  - Euclid's Lemma: Let  $a, b \in \mathbb{Z}$  and  $p$  prime. If  $p \mid ab$ , then  $p \mid a$  or  $p \mid b$ .
  - $a, b, c \in \mathbb{Z}$ . If  $c \mid ab$  and  $\gcd(a, c) = 1$ , then  $c \mid b$ .
  - every integer  $n \geq 2$  can be expressed as a unique product of positive primes
  - there are infinitely many primes
  - Fermat's little theorem: let  $a, p \in \mathbb{N}$  with  $p$  prime. Then  $a^p \equiv a \pmod{p}$ 
    - Corollary: let  $a \in \mathbb{N}$ ,  $p$  prime with  $p \nmid a$ . Then  $a^{p-1} \equiv 1 \pmod{p}$
- Bonkers Remainders When Divided by Primes
  - use Fermat's Little Corollary
  - let us use an example. find the remainder of  $3^{244886}$  when divided by 13.

by Fermat's, since  $13 \nmid 3$ , we have  $3^{12} \equiv_{13} 1$ . let us use the power of 12 to simplify  $3^{244886}$ . via long division, we have  $244886 = 12 \cdot 20407 + 2$ . We proceed.

$$3^{244886} \equiv_{13} 3^{12 \cdot 20407 + 2} \equiv_{13} 3^{12 \cdot 20407} \cdot 3^2 \equiv_{13} 1 \cdot 3^2 \equiv_{13} 9$$

Thus, the remainder is 9.