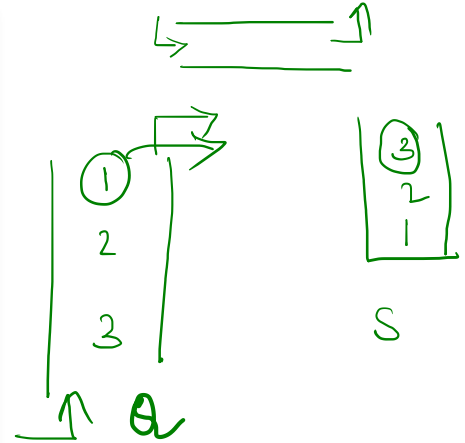


Day-7

✓ Let Q denote a queue containing sixteen numbers and S be an empty stack. Head(Q) returns the element at the head of the queue Q **without** removing it from Q. Similarly Top(S) returns the element at the top of S **without** removing it from S. Consider the algorithm given below.

```
while Q is not Empty do  
  if S is Empty OR  $\text{Top}(S) \leq \text{Head}(Q)$  then  
     $x := \text{Dequeue}(Q);$   
     $\text{Push}(S, x);$   
  else  
     $x := \text{Pop}(S);$   
     $\text{Enqueue}(Q, x);$   
  end  
end
```



The maximum possible number of iterations of the while loop in the algorithm is _____

$c = \emptyset, 1, 2, 3, 4, 5, 6, 7, 8, 9$

while Q is not Empty **do**

if S is Empty **OR** $Top(S) \leq Head(Q)$ **then**

$x := Dequeue(Q);$

$Push(S, x);$ $1 \leq 2$

else $2 \leq 3$

$x := Pop(S);$

$Enqueue(Q, x);$

end

end

entry \hookrightarrow \hookrightarrow exit
~~3~~ ~~2~~ ~~3~~ ~~1~~

Q $3 \leq 2$

$\begin{array}{|c|} \hline 3 \\ 2 \\ \cancel{3} \\ 1 \\ \hline S \end{array}$

$1 \leq 3$

$3 \leq 2$

$2 \leq 1 \times$

$2 \rightarrow 4$

$3 \rightarrow 9$

$16 \rightarrow 256$

c_1 $n = 2$ ✓

 \circ 2 1
 $\textcircled{2}$ ✓

c_2 $n = 3$

3 2 1
 $\textcircled{3}$ ✓

c_3 $n = 2$

1 2
 $\textcircled{4}$

c_4 $n = 3$ ⑨

1 2 3

① Stock-Span Problem

$p[] = \{ 100, 80, 60, 70, 60, 75, 85 \}$

$s[] = \{ 1, 1, 1, 2, 1, 4, 6 \}$

Stocks

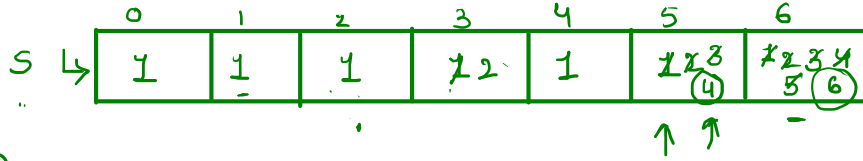
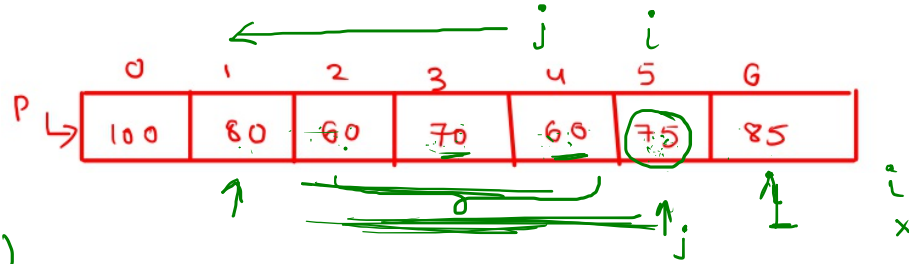
continuous

| | | | | | | | |
|-----------------|------------|------------|------------|------------|------------|------------|------------|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| P \rightarrow | 100 | 80 | 60 | 70 | 60 | 75 | 85 |
| | d_1 | d_2 | d_3 | d_4 | d_5 | d_6 | d_7 |
| | \uparrow | \uparrow | \uparrow | \uparrow | \uparrow | \uparrow | \uparrow |
| | 1 | 1 | 1 | 2 | 1 | 4 | 6 |

BF

$\rightarrow O(n^2)$ ✓

$\rightarrow O(1)/O(n)$



why (?)

Stack

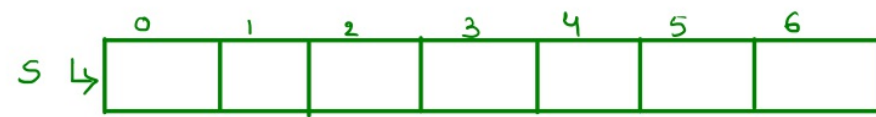
2 \rightarrow loops

loop₁ \rightarrow 0..L

i = 1 to n

loop₂ \rightarrow 1..L

2 \rightarrow nested loops



```
findSpan(price[], n, span[])
```

```
{
```

```
    ✓ Stack st
```

```
    st.push(0)
```

```
    ✓ span[0]=1
```

```
    → for(i=1; i<n; i++) →  $O(n)$ 
```

```
    {
```

```
        → while(!isEmpty() && price[st[top]] <= price[i])
```

```
        {
```

```
            ele = st.pop() ✓
```

```
        }
```

```
        span[i] = isEmpty() ? i+1 : i - st[top]
```

```
        st.push(i)
```

```
    }
```

```
}
```

→ $O(n)$ ✓

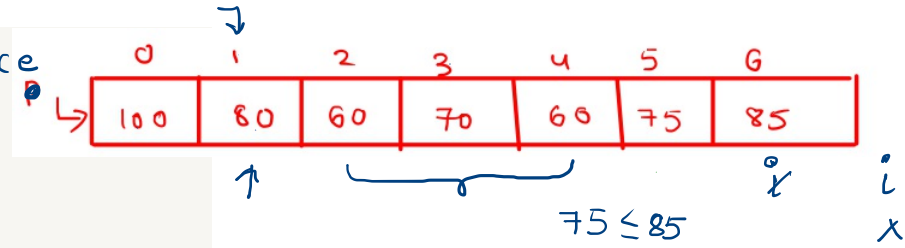
→ $O(n)$

$n \times n$

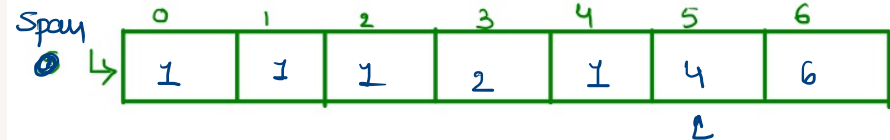
→ n^2

↑ ↑

Price



Span



st

EA \rightarrow push, pop, top

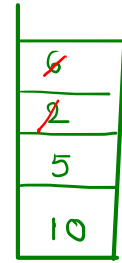
Design a stack such that getMin() is in $O(1)$, $O(1)$ Space.

10, 5, 2, 6, getMin()

\rightarrow push(10), push(5), push(2), push(6),

getMin(). pop(), pop(), getMin(),

\rightarrow 2
 $O(1)$



S_1

\rightarrow 5

AP_1

\rightarrow twice S_2
 \otimes


$$2 - (-3) = 5$$
$$-3 < 2$$

$O(1)$

$$\text{curr_min} = \text{curr_min} - \text{st.top}()$$

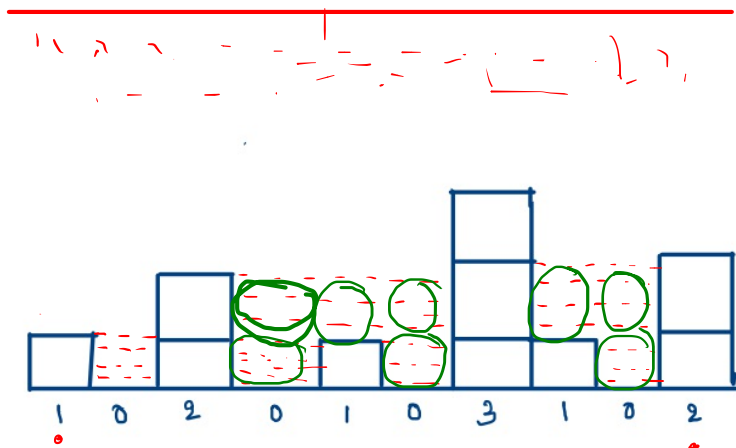
→ design a stack such that `getMin()` in $O(1)$
Let St be a global stack

```
✓ add(data)
{
    if(s.isEmpty())
    {
        s.push(data)
        curr_min=data ✓
    }
    * else
    {
        if(data < curr_min)
        {
            s.push(data - curr_min)
            curr_min=data ✓
        }
        else
            s.push(data) ✓
    }
    delete()
    {
        if(s.peak() < curr_min)
        {
            curr_min=curr_min-s.peak()
        }
        return s.pop() ✓
    }
}
```

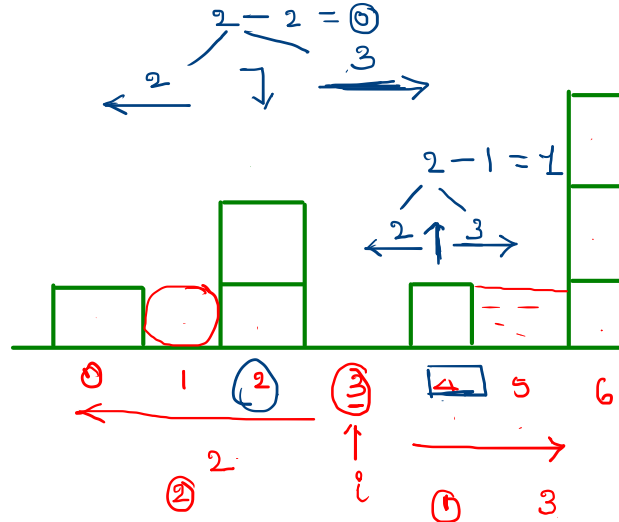


PB1 (A1)

Trapping Rain water problem



0 + 1 + 0 + 2 + 1 + 2 + 0 + 1 + 2 + 0



$$\textcircled{3} \quad \min(3, 2) = 2$$

$$\textcircled{4} \quad 2 - \text{arr}[i] = 2 - 0 = 2$$

$\Rightarrow \textcircled{9}$ ✓

```
function maxWater(arr[], n)
{
```

```
    res = 0;
```

```
    for(i = 1; i < n - 1; i++) → n
    {
```

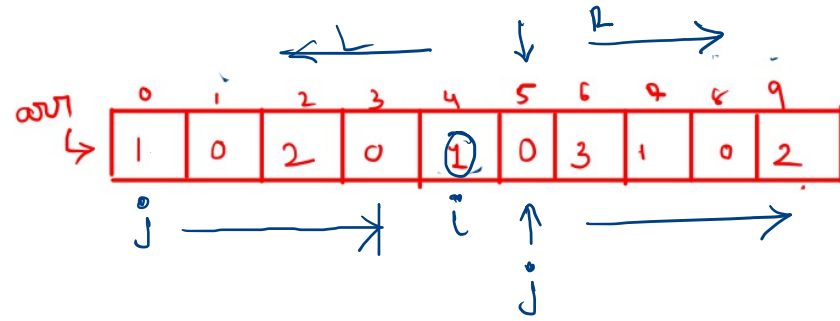
```
        left = arr[i];
        for(j = 0; j < i; j++)
        {
            left = Math.max(left, arr[j]);
        }
```

```
        right = arr[i];
        for(j = i + 1; j < n; j++)
        {
            right = Math.max(right, arr[j]);
        }
```

```
        res += Math.min(left, right) - arr[i];
```

```
    }
    return res;
```

```
}
```



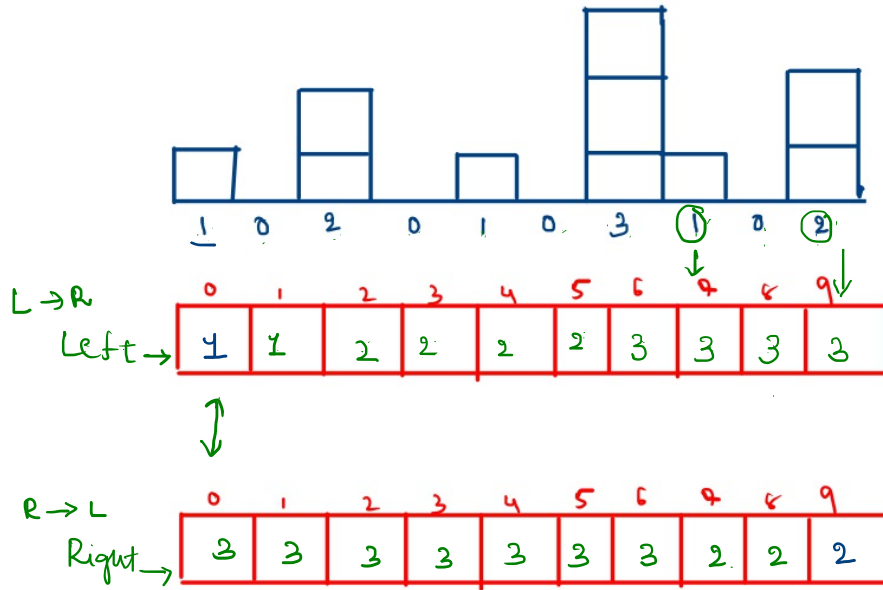
left = 2

$$n[n+n] = n^2 + n^2 = 2n^2$$

$\therefore O(n^2)$ T.C

$O(1)$ S.C

right = 3



$$T.O.C \Rightarrow n + n + n = 3n \Rightarrow O(n)$$

$$S.C \Rightarrow n + n = 2n \Rightarrow O(n)$$

→ loop₁

+

→ loop₂

+

→ loop₃

res = 0

for (i = 0; i < n; i++)

{

res = res + (min(Left[i], Right[i]) - arr[i])

}

return res

```
function findWater(arr[],n)
{
    let left[n], right[n]

    water = 0;

    left[0] = arr[0];
    for (i = 1; i < n; i++)
    {
        left[i] = Math.max(left[i - 1], arr[i]);
    }

    right[n - 1] = arr[n - 1];

    for (i = n - 2; i >= 0; i--)
    {
        right[i] = Math.max(right[i + 1], arr[i]);
    }

    for (i = 0; i < n; i++)
        water += Math.min(left[i], right[i]) - arr[i];

    return water;
}
```

```

function findWater(arr[], n)
{
    result = 0, left_max = 0, right_max = 0, lo = 0, hi = n - 1;

    while (lo <= hi)
    {
        if (arr[lo] < arr[hi])
        {
            if (arr[lo] > left_max)
                left_max = arr[lo];
            else
                result += left_max - arr[lo];
            lo++;
        }
        else
        {
            if (arr[hi] > right_max)
                right_max = arr[hi];
            else
                result += right_max - arr[hi];
            hi--;
        }
    }
    return result;
}

```

