Team Notebook

December 15, 2018

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1 01bfs

```
d.assign(n, INF);
d[s] = 0:
set<pair<int, int>> q;
q.insert({0, s});
while (!q.empty()) {
   int v = q.begin()->second;
   q.erase(q.begin());
   for (auto edge : adj[v]) {
       int u = edge.first;
       int w = edge.second;
       if (d[v] + w < d[u]) {
           q.erase({d[u], u});
           d[u] = d[v] + w;
           q.insert({d[u], u});
       }
   }
}
vector<int> d(n, INF);
d[s] = 0:
deque<int> q;
q.push_front(s);
while (!q.empty()) {
   int v = q.front();
   q.pop_front();
   for (auto edge : adj[v]) {
       int u = edge.first;
       int w = edge.second;
       if (d[v] + w < d[u]) {
           d[u] = d[v] + w;
           if (w == 1)
              q.push_back(u);
              q.push_front(u);
       }
   }
}
```

|2 2-SAT

```
int n;
vector<vector<int>> g, gt;
vector<bool> used;
vector<int> order, comp;
vector<bool> assignment;
void dfs1(int v) {
   used[v] = true:
   for (int u : g[v]) {
       if (!used[u])
           dfs1(u):
   order.push_back(v);
void dfs2(int v, int cl) {
   comp[v] = cl;
   for (int u : gt[v]) {
       if (comp[u] == -1)
           dfs2(u, cl);
   }
bool solve 2SAT() {
   used.assign(n, false);
   for (int i = 0; i < n; ++i) {</pre>
       if (!used[i])
           dfs1(i);
   comp.assign(n, -1);
   for (int i = 0, j = 0; i < n; ++i) {
       int v = order[n - i - 1];
       if (comp[v] == -1)
           dfs2(v, j++);
   assignment.assign(n / 2, false);
   for (int i = 0; i < n; i += 2) {</pre>
       if (comp[i] == comp[i + 1])
           return false:
       assignment[i / 2] = comp[i] > comp[i + 1];
   }
   return true;
```

3 2d-segtree

```
void build_y(int vx, int lx, int rx, int vy, int ly,
    int ry) {
   if (ly == ry) {
       if (lx == rx)
          t[vx][vy] = a[lx][ly];
          t[vx][vy] = t[vx*2][vy] + t[vx*2+1][vy];
   } else {
       int my = (1y + ry) / 2;
       build_y(vx, lx, rx, vy*2, ly, my);
       build_y(vx, lx, rx, vy*2+1, my+1, ry);
       t[vx][vy] = t[vx][vv*2] + t[vx][vv*2+1];
void build_x(int vx, int lx, int rx) {
   if (lx != rx) {
       int mx = (1x + rx) / 2;
       build_x(vx*2, lx, mx);
       build_x(vx*2+1, mx+1, rx);
   build_v(vx, lx, rx, 1, 0, m-1);
int sum_y(int vx, int vy, int tly, int try_, int ly,
    int ry) {
   if (lv > rv)
       return 0:
   if (ly == tly && try_ == ry)
       return t[vx][vy];
   int tmy = (tly + try_) / 2;
   return sum_y(vx, vy*2, tly, tmy, ly, min(ry, tmy))
        + sum_y(vx, vy*2+1, tmy+1, try_, max(ly, tmy
            +1), rv);
int sum_x(int vx, int tlx, int trx, int lx, int rx,
    int ly, int ry) {
   if (lx > rx)
       return 0;
```

```
if (lx == tlx && trx == rx)
       return sum_y(vx, 1, 0, m-1, ly, ry);
    int tmx = (tlx + trx) / 2;
    return sum_x(vx*2, tlx, tmx, lx, min(rx, tmx), ly,
        + sum_x(vx*2+1, tmx+1, trx, max(1x, tmx+1), rx
             , lv, rv);
}
void update_y(int vx, int lx, int rx, int vy, int ly,
    int ry, int x, int y, int new_val) {
    if (lv == rv) {
       if (lx == rx)
           t[vx][vy] = new_val;
           t[vx][vy] = t[vx*2][vy] + t[vx*2+1][vy];
   } else {
       int my = (ly + ry) / 2;
       if (y <= my)
           update_y(vx, lx, rx, vy*2, ly, my, x, y,
               new val):
       else
           update_v(vx, lx, rx, vy*2+1, my+1, ry, x, y
               , new_val);
       t[vx][vy] = t[vx][vy*2] + t[vx][vy*2+1];
   }
}
void update_x(int vx, int lx, int rx, int x, int y,
    int new_val) {
    if (lx != rx) {
       int mx = (1x + rx) / 2;
       if (x \le mx)
           update_x(vx*2, lx, mx, x, y, new_val);
       else
           update_x(vx*2+1, mx+1, rx, x, y, new_val);
    update_y(vx, lx, rx, 1, 0, m-1, x, y, new_val);
}
```

4 Articulation

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> tin, fup;
int timer;
void dfs(int v, int p = -1) {
   visited[v] = true;
   tin[v] = fup[v] = timer++;
   int children=0:
   for (int to : adj[v]) {
       if (to == p) continue;
       if (visited[to]) {
           fup[v] = min(fup[v], tin[to]);
       } else {
          dfs(to, v):
           fup[v] = min(fup[v], fup[to]);
           if (fup[to] >= tin[v] && p!=-1)
              IS_CUTPOINT(v);
           ++children:
       }
   }
   if(p == -1 \&\& children > 1)
       IS_CUTPOINT(v); // v is cut-point
void find_cutpoints() {
   timer = 0;
   visited.assign(n, false);
   tin.assign(n, -1);
   fup.assign(n, -1);
   for (int i = 0; i < n; ++i) {</pre>
       if (!visited[i])
           dfs (i):
   }
```

5 BIT

```
int BIT[1000], a[1000], n;
void update(int x, int delta){
   for(; x <= n; x += x&-x)</pre>
```

```
BIT[x] += delta;
}
int query(int x){
   int sum = 0;
   for(; x > 0; x -= x&-x)
       sum += BIT[x];
   return sum;
}
```

6 LCA

```
struct LCA {
   vector<int> height, euler, first, segtree;
   vector<bool> visited;
   int n;
   LCA(vector<vector<int>> &adj, int root = 0) {
       n = adi.size():
       height.resize(n);
       first.resize(n);
       euler.reserve(n * 2);
       visited.assign(n, false);
       dfs(adj, root);
       int m = euler.size();
       segtree.resize(m * 4);
       build(1, 0, m - 1);
   void dfs(vector<vector<int>> &adj, int node, int h
        = 0) {
       visited[node] = true;
       height[node] = h:
       first[node] = euler.size();
       euler.push_back(node);
       for (auto to : adj[node]) {
          if (!visited[to]) {
              dfs(adj, to, h + 1);
              euler.push_back(node);
       }
   void build(int node, int b, int e) {
       if (b == e) {
          segtree[node] = euler[b];
```

```
} else {
           int mid = (b + e) / 2;
           build(node << 1, b, mid);</pre>
           build(node << 1 | 1, mid + 1, e);
           int 1 = segtree[node << 1], r = segtree[</pre>
               node << 1 | 1]:
           segtree[node] = (height[1] < height[r]) ? 1</pre>
       }
   }
   int query(int node, int b, int e, int L, int R) {
       if (b > R || e < L)
           return -1;
       if (b >= L && e <= R)
           return segtree[node];
       int mid = (b + e) >> 1:
       int left = query(node << 1, b, mid, L, R);</pre>
       int right = query(node << 1 | 1, mid + 1, e, L, ]
             R):
       if (left == -1) return right;
       if (right == -1) return left;
       return height[left] < height[right] ? left :</pre>
            right;
   }
   int lca(int u, int v) {
       int left = first[u], right = first[v];
       if (left > right)
           swap(left, right);
       return query(1, 0, euler.size() - 1, left,
           right);
   }
};
```

7 LIS

```
cout << i << endl;
  return 0;
}</pre>
```

8 MO

```
11 block; //sqrt(N)
struct QUERY{
11 L,R;
};
bool compare(QUERY a,QUERY b){
if(a.L/block != b.L/block)
 return (a.L/block) < (b.L/block);</pre>
return a.R<b.R;</pre>
void mo(vector<11>a, vector<QUERY>q){
block = sqrt(a.size());
sort(q.begin(),q.end(),compare);
11 curL=0,curR=0,curSum=0;
for(int i=0;i<q.size();i++){</pre>
 11 L = q[i].L,R = q[i].R;
 while(curL<L){
  curSum-=a[curL]:
   curL++;
 while(curL>L){
  curSum+=a[curL];
  curL--;
 while(curR<=R){</pre>
  curSum+=a[curR]:
  curR++;
 while(curR>(R+1)){
  curSum-=a[curR-1];
  curR--:
 cout << curSum << "\n";</pre>
```

19 SCC

```
vector<11>adj[400005],adjr[400005];
vector<11>visited(400005,0), visitedr(400005,0);
vector<11>order,component;
void dfs1(ll src){
visited[src] = 1;
for(auto e:adj[src])
 if(!visited[e])
  dfs1(e):
order.pb(src);
void dfs2(ll src){
visitedr[src] = 1;
component.pb(src);
for(auto e:adjr[src])
 if(!visitedr[e])
  dfs2(e);
for(int i=0;i<n;i++){</pre>
cin >> a >> b;
adj[a].push_back(b);
adjr[b].push_back(a);
for(int i=0:i<n:++i)</pre>
if(!visited[i])
 dfs1(i);
for(int i=0;i<n;++i){</pre>
ll v = order[n-1-i]:
if(!visitedr[v]){
 dfs2 (v);component.clear();
```

10 Segment-CP

```
// Normal Segment tree
int n, t[4*MAXN];

void build(int a[], int v, int tl, int tr) {
   if (tl == tr) {
      t[v] = a[tl];
}
```

```
} else {
       int tm = (t1 + tr) / 2;
       build(a, v*2, t1, tm);
       build(a, v*2+1, tm+1, tr):
       t[v] = t[v*2] + t[v*2+1]:
   }
}
int sum(int v, int tl, int tr, int l, int r) {
    if (1 > r)
       return 0:
    if (1 == t1 && r == tr) {
       return t[v]:
    int tm = (tl + tr) / 2;
    return sum(v*2, tl, tm, l, min(r, tm))
          + sum(v*2+1, tm+1, tr, max(1, tm+1), r);
}
void update(int v, int tl, int tr, int pos, int
    new val) {
    if (t1 == tr) {
       t[v] = new_val;
   } else {
       int tm = (tl + tr) / 2;
       if (pos <= tm)</pre>
           update(v*2, t1, tm, pos, new_val);
           update(v*2+1, tm+1, tr, pos, new_val);
       t[v] = t[v*2] + t[v*2+1];
}
// advance version of segment tree
pair<int, int> t[4*MAXN];
pair<int, int> combine(pair<int, int> a, pair<int, int
    > b) {
   if (a.first > b.first)
       return a:
    if (b.first > a.first)
       return b:
   return make_pair(a.first, a.second + b.second);
}
```

```
void build(int a[], int v, int tl, int tr) {
   if (t] == tr) {
       t[v] = make_pair(a[tl], 1);
   } else {
       int tm = (t1 + tr) / 2:
       build(a, v*2, tl, tm);
       build(a, v*2+1, tm+1, tr):
       t[v] = combine(t[v*2], t[v*2+1]);
   }
pair<int, int> get_max(int v, int tl, int tr, int l,
    int r) {
   if (1 > r)
       return make_pair(-INF, 0);
   if (1 == t1 && r == tr)
       return t[v]:
   int tm = (tl + tr) / 2:
   return combine(get_max(v*2, tl, tm, l, min(r, tm))
                 get_max(v*2+1, tm+1, tr, max(1, tm+1))
                     , r));
void update(int v, int tl, int tr, int pos, int
    new val) {
   if (tl == tr) {
       t[v] = make pair(new val. 1):
   } else {
       int tm = (tl + tr) / 2;
       if (pos <= tm)</pre>
           update(v*2, tl, tm, pos, new_val);
           update(v*2+1, tm+1, tr, pos, new_val);
       t[v] = combine(t[v*2], t[v*2+1]);
//int find_kth(int v, int tl, int tr, int k) {
   if (k > t[v])
       return -1:
   if (tl == tr)
       return tl;
```

```
int tm = (tl + tr) / 2:
   if (t[v*2] >= k)
       return find_kth(v*2, tl, tm, k);
       return find_kth(v*2+1, tm+1, tr, k - t[v*2]);
11
struct data {
   int sum, pref, suff, ans:
};
data combine(data 1, data r) {
   data res:
   res.sum = 1.sum + r.sum:
   res.pref = max(1.pref, 1.sum + r.pref);
   res.suff = max(r.suff, r.sum + 1.suff);
   res.ans = max(max(1.ans, r.ans), 1.suff + r.pref);
   return res;
data make_data(int val) {
   data res:
   res.sum = val;
   res.pref = res.suff = res.ans = max(0, val);
   return res;
void build(int a[], int v, int tl, int tr) {
   if (t1 == tr) {
       t[v] = make_data(a[t1]);
   } else {
       int tm = (t1 + tr) / 2;
       build(a, v*2, t1, tm);
       build(a, v*2+1, tm+1, tr);
       t[v] = combine(t[v*2], t[v*2+1]):
   }
void update(int v, int tl, int tr, int pos, int
    new val) {
   if (tl == tr) {
       t[v] = make_data(new_val);
   } else {
       int tm = (t1 + tr) / 2:
```

11 bellmanford

```
struct edge
    int a, b, cost;
};
int n, m, v;
vector<edge> e;
const int INF = 1000000000;
void solve()
    vector<int> d (n, INF);
    d[v] = 0;
    for (int i=0; i<n-1; ++i)</pre>
       for (int j=0; j<m; ++j)</pre>
           if (d[e[i].a] < INF)</pre>
               d[e[j].b] = min (d[e[j].b], d[e[j].a] +
                    e[i].cost);
    // display d, for example, on the screen
}
```

```
void solve()
   vector<int> d (n, INF);
   d[v] = 0:
   for (;;)
   {
       bool any = false;
       for (int j=0; j<m; ++j)</pre>
           if (d[e[i].a] < INF)</pre>
               if (d[e[j].b] > d[e[j].a] + e[j].cost)
                  d[e[i].b] = d[e[i].a] + e[i].cost;
                  any = true;
               }
       if (!any) break;
   // display d, for example, on the screen
Retrieving Path
void solve()
   vector<int> d (n, INF);
   d[v] = 0;
   vector<int> p (n - 1);
   for (;;)
   Ł
       bool any = false;
       for (int j = 0; j < m; ++j)
           if (d[e[i].a] < INF)</pre>
               if (d[e[j].b] > d[e[j].a] + e[j].cost)
                  d[e[j].b] = d[e[j].a] + e[j].cost;
                  p[e[j].b] = e[j].a;
                  any = true;
               if (!any) break;
   }
```

```
if (d[t] == INF)
       cout << "No path from " << v << " to " << t <<
   else
   {
       vector<int> path;
       for (int cur = t; cur != -1; cur = p[cur])
           path.push_back (cur);
       reverse (path.begin(), path.end());
       cout << "Path from " << v << " to " << t << ":</pre>
       for (size_t i=0; i<path.size(); ++i)</pre>
           cout << path[i] << ' ';
   }
void solve()
   vector<int> d (n, INF);
   d[v] = 0;
   vector<int> p (n - 1);
   int x;
   for (int i=0; i<n; ++i)</pre>
       x = -1;
       for (int j=0; j<m; ++j)</pre>
           if (d[e[i].a] < INF)</pre>
               if (d[e[j].b] > d[e[j].a] + e[j].cost)
                   d[e[j].b] = max (-INF, d[e[j].a] + e
                        [j].cost);
                  p[e[i].b] = e[i].a;
                  x = e[j].b;
   }
   if (x == -1)
       cout << "No negative cycle from " << v;</pre>
   else
   {
       int y = x;
```

```
for (int i=0: i<n: ++i)</pre>
           y = p[y];
       vector<int> path;
       for (int cur=y; ; cur=p[cur])
           path.push_back (cur);
           if (cur == y && path.size() > 1)
               break;
       }
       reverse (path.begin(), path.end());
       cout << "Negative cycle: ";</pre>
       for (size_t i=0; i<path.size(); ++i)</pre>
           cout << path[i] << ' ';</pre>
    }
}
const int INF = 1000000000;
vector<vector<pair<int, int>>> adj;
bool spfa(int s, vector<int>& d) {
    int n = adj.size();
    d.assign(n, INF);
    vector<int> cnt(n, 0);
    vector<bool> inqueue(n, false);
    queue<int> q;
    d[s] = 0;
    q.push(s);
    inqueue[s] = true;
    while (!q.empty()) {
       int v = q.front();
       q.pop();
       inqueue[v] = false;
       for (auto edge : adj[v]) {
           int to = edge.first;
           int len = edge.second;
           if (d[v] + len < d[to]) {
               d[to] = d[v] + len;
               if (!inqueue[to]) {
```

```
q.push(to);
    inqueue[to] = true;
    cnt[to]++;
    if (cnt[to] > n)
        return false; // negative cycle
    }
}
}
```

12 bipart-check

```
int n:
vector<vector<int>> adj;
vector<int> side(n, -1);
bool is_bipartite = true;
queue<int> q;
for (int st = 0; st < n; ++st) {</pre>
   if (side[st] == -1) {
       q.push(st);
       side[st] = 0;
       while (!q.empty()) {
           int v = q.front();
           q.pop();
           for (int u : adj[v])
               if (side[u] == -1) {
                  side[u] = side[v] ^ 1
                  q.push(u);
              } else
                  is_bipartite &= side[u] != side[v];
       }
   }
cout << (is_bipartite ? "YES" : "NO") << endl;</pre>
```

13 bipart-match

```
#define MAX 100001
#define NIL 0
```

```
#define INF (1<<28)
vector< int > G[MAX];
int n, m, match[MAX], dist[MAX];
// n: number of nodes on left side, nodes are numbered
// m: number of nodes on right side, nodes are
    numbered n+1 to n+m
// G = NIL[0] G1[G[1--n]] G2[G[n+1--n+m]]
bool bfs() {
   int i, u, v, len;
   queue < int > Q;
   for(i=1; i<=n; i++) {</pre>
       if(match[i]==NIL) {
           dist[i] = 0:
           Q.push(i);
       else dist[i] = INF;
   dist[NIL] = INF;
   while(!Q.empty()) {
       u = Q.front(); Q.pop();
       if(u!=NIL) {
           len = G[u].size();
           for(i=0; i<len; i++) {</pre>
              v = G[u][i];
              if(dist[match[v]]==INF) {
                  dist[match[v]] = dist[u] + 1;
                  Q.push(match[v]);
   return (dist[NIL]!=INF);
bool dfs(int u) {
   int i, v, len;
   if(u!=NIL) {
       len = G[u].size();
       for(i=0; i<len; i++) {</pre>
           v = G[u][i];
           if(dist[match[v]] == dist[u] + 1) {
              if(dfs(match[v])) {
                  match[v] = u;
                  match[u] = v;
```

```
return true:
               }
           }
       }
       dist[u] = INF;
       return false:
    }
    return true:
}
int hopcroft_karp() {
    int matching = 0, i;
    // match[] is assumed NIL for all vertex in G
    while(bfs())
       for(i=1: i<=n: i++)</pre>
           if(match[i] == NIL && dfs(i))
               matching++:
    return matching;
} //calling function
```

14 bridge

```
int n: // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> tin, fup;
int timer;
void dfs(int v, int p = -1) {
   visited[v] = true;
   tin[v] = fup[v] = timer++;
   for (int to : adj[v]) {
       if (to == p) continue;
       if (visited[to]) {
          fup[v] = min(fup[v], tin[to]);
       } else {
           dfs(to, v);
          fup[v] = min(fup[v], fup[to]);
          if (fup[to] > tin[v])
              IS_BRIDGE(v, to); // bridge found from v
                  -to
       }
   }
}
```

```
void find_bridges() {
    timer = 0;
    visited.assign(n, false);
    tin.assign(n, -1);
    fup.assign(n, -1);
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs(i);
    }
}</pre>
```

15 crt

```
from functools import reduce
def chinese_remainder(n, a):
   prod = reduce(lambda a, b: a*b, n)
   for n_i, a_i in zip(n, a):
       p = prod / n_i
       sum += a_i * mul_inv(p, n_i) * p
   return sum % prod
def mul inv(a, b):
   b0 = b
   x0, x1 = 0, 1
   if b == 1: return 1
   while a > 1:
       q = a // b
       a, b = b, a\%b
       x0, x1 = x1 - q * x0, x0
   if x1 < 0: x1 += b0
   return x1
if __name__ == '__main__':
   n = [3, 5, 7] #xi=ai%ni
   a = [2, 3, 2]
   print(chinese_remainder(n, a))
```

16 digit-dp

```
/// How many numbers x are there in the range a to b,
    where the digit d occurs exactly k times in x?
vector<int> num:
int a, b, d, k;
int DP[12][12][2];
/// DP[p][c][f] = Number of valid numbers <= b from
    this state
/// p = current position from left side (zero based)
/// c = number of times we have placed the digit d so
    far
/// f = the number we are building has already become
    smaller than b? [0 = no, 1 = yes]
int call(int pos, int cnt, int f){
   if(cnt > k) return 0:
   if(pos == num.size()){
       if(cnt == k) return 1;
       return 0;}
   if(DP[pos][cnt][f] != -1) return DP[pos][cnt][f];
   int res = 0:
   int LMT;
   if(f == 0){
       /// Digits we placed so far matches with the
           prefix of b
       /// So if we place any digit > num[pos] in the
           current position, then the number will
           become greater than b
       LMT = num[pos];
   } else {
       /// The number has already become smaller than
           b. We can place any digit now.
       LMT = 9;
   }
   /// Try to place all the valid digits such that
        the number doesn't exceed b
   for(int dgt = 0; dgt<=LMT; dgt++){</pre>
       int nf = f:
       int ncnt = cnt;
       if(f == 0 \&\& dgt < LMT) nf = 1; /// The number
           is getting smaller at this position
       if(dgt == d) ncnt++;
```

```
if(ncnt <= k) res += call(pos+1, ncnt, nf);</pre>
    }
    return DP[pos][cnt][f] = res;
int solve(int b){
    num.clear():
    while(b>0){
       num.push_back(b%10);
       b/=10:
    reverse(num.begin(), num.end());
    /// Stored all the digits of b in num for
        simplicity
    memset(DP, -1, sizeof(DP));
    int res = call(0, 0, 0);
    return res;
}
int main () {
 cin >> a >> b >> d >> k:
    int res = solve(b) - solve(a-1);
    cout << res << endl;</pre>
}
```

17 diophantine

```
void shift_solution (int & x, int & y, int a, int b,
    int cnt) {
   x += cnt * b:
   v -= cnt * a:
int find_all_solutions (int a, int b, int c, int minx,
     int maxx, int miny, int maxy) { // returns no. of
     solution
   int x, y, g;
   if (! find_any_solution (a, b, c, x, y, g))
       return 0;
   a /= g; b /= g;
   int sign_a = a>0 ? +1 : -1;
   int sign_b = b>0 ? +1 : -1;
   shift_solution (x, y, a, b, (minx - x) / b);
   if (x < minx)
       shift_solution (x, y, a, b, sign_b);
   if (x > maxx)
       return 0:
   int lx1 = x;
   shift_solution (x, y, a, b, (maxx - x) / b);
   if (x > maxx)
       shift_solution (x, y, a, b, -sign_b);
   int rx1 = x;
   shift_solution (x, y, a, b, - (miny - y) / a);
   if (y < miny)</pre>
       shift_solution (x, y, a, b, -sign_a);
   if (v > maxv)
       return 0;
   int 1x2 = x;
   shift_solution (x, y, a, b, - (maxy - y) / a);
   if (y > maxy)
       shift_solution (x, y, a, b, sign_a);
   int rx2 = x:
   if (1x2 > rx2)
       swap (1x2, rx2);
   int 1x = max (1x1, 1x2);
   int rx = min (rx1, rx2);
   if (lx > rx) return 0;
   return (rx - lx) / abs(b) + 1;
```

18 dp-recurences

```
// matrix chain multiplication
// Matrix Ai has dimension p[i-1] x p[i] for i = 1..n
int MatrixChainOrder(int p[], int n)
{ /* For simplicity of the program, one extra row and
extra column are allocated in m[][]. Oth row and Oth
column of m[][] are not used */
int m[n][n]:
int i, j, k, L, q;
/* m[i,j] = Minimum number of scalar multiplications
    needed
to compute the matrix A[i]A[i+1]...A[i] = A[i..i]
dimension of A[i] is p[i-1] x p[i] */
// cost is zero when multiplying one matrix.
for (i=1: i<n: i++)</pre>
 m[i][i] = 0;
// L is chain length.
for (L=2: L<n: L++)</pre>
 for (i=1; i<n-L+1; i++) {</pre>
  j = i+L-1;
  m[i][i] = INT_MAX;
  for (k=i; k<=j-1; k++) {</pre>
   // q = cost/scalar multiplications
   q = m[i][k] + m[k+1][j] + p[i-1]*p[k]*p[j];
   if (q < m[i][j])</pre>
    m[i][i] = q; }
  return m[1][n-1];
// bell no
//Let S(n, k) be total number of partitions of n
    elements into k sets. The value of nth Bell Number
     is sum of S(n, k) for k = 1 to n.
int bellNumber(int n)
  int bell[n+1][n+1];
  bell[0][0] = 1;
```

```
for (int i=1; i<=n; i++)</pre>
     // Explicitly fill for j = 0
     bell[i][0] = bell[i-1][i-1]:
     // Fill for remaining values of j
     for (int j=1; j<=i; j++)</pre>
        bell[i][j] = bell[i-1][j-1] + bell[i][j-1];
  return bell[n][0]:
}
// subset sum
isSubsetSum(set, n, sum) = isSubsetSum(set, n-1, sum)
    -11
                         isSubsetSum(set, n-1, sum-set[
                             n-11)
Base Cases:
isSubsetSum(set, n, sum) = false, if sum > 0 and n ==
isSubsetSum(set, n, sum) = true, if sum == 0
//rod cutting
Let cutRod(n) be the required (best possible price)
    value for a rod of length n. cutRod(n) can be
    written as following.
cutRod(n) = max(price[i] + cutRod(n-i-1)) for all i in
     \{0, 1 \dots n-1\}
//LCS
/* Returns length of LCS for X[0..m-1], Y[0..n-1] */
int lcs( char *X, char *Y, int m, int n )
   if (m == 0 | | n == 0)
    return 0;
  if (X[m-1] == Y[n-1])
    return 1 + lcs(X, Y, m-1, n-1);
  else
    return max(lcs(X, Y, m, n-1), lcs(X, Y, m-1, n));
}
//
The longest common suffix has following optimal
    substructure property
  LCSuff(X, Y, m, n) = LCSuff(X, Y, m-1, n-1) + 1 if
       X[m-1] = Y[n-1]
```

```
O Otherwise (if X[m-1] != Y[n
                            -17)
The maximum length Longest Common Suffix is the
     longest common substring.
   LCSubStr(X, Y, m, n) = Max(LCSuff(X, Y, i, j))
        where 1 \le i \le m
                                                  and 1
                                                      <=
 //Kadane
    max_so_far = 0
    max ending here = 0
Loop for each element of the array
   (a) max_ending_here = max_ending_here + a[i]
   (b) if(max_ending_here < 0)</pre>
           max_ending_here = 0
  (c) if(max_so_far < max_ending_here)</pre>
            max_so_far = max_ending_here
 return max so far
 //0-1 knapsack
int knapSack(int W, int wt[], int val[], int n)
   // Base Case
   if (n == 0 || W == 0)
       return 0;
   // If weight of the nth item is more than Knapsack
        capacity W, then
   // this item cannot be included in the optimal
        solution
   if (wt[n-1] > W)
       return knapSack(W, wt, val, n-1);
   // Return the maximum of two cases:
   // (1) nth item included
   // (2) not included
   else return max( val[n-1] + knapSack(W-wt[n-1], wt,
         val, n-1),
                   knapSack(W, wt, val, n-1)
```

```
):
//Egg-Droping
k ==> Number of floors
 n ==> Number of Eggs
 eggDrop(n, k) ==> Minimum number of trials needed to
       find the critical
                  floor in worst case.
 eggDrop(n, k) = 1 + min\{max(eggDrop(n - 1, x - 1),
      eggDrop(n, k - x)):
               x in \{1, 2, ..., k\}
//Partition Problem
Let isSubsetSum(arr, n, sum/2) be the function that
    returns true if
there is a subset of arr[0..n-1] with sum equal to sum
The isSubsetSum problem can be divided into two
    subproblems
a) isSubsetSum() without considering last element
    (reducing n to n-1)
b) isSubsetSum considering the last element
    (reducing sum/2 by arr[n-1] and n to n-1)
If any of the above the above subproblems return true,
     then return true.
isSubsetSum (arr, n, sum/2) = isSubsetSum (arr, n-1,
    sum/2) | |
                           isSubsetSum (arr, n-1, sum
                               /2 - arr[n-1]
//Longest Palindromic Subsequence
// Every single character is a palindrome of length 1
L(i, i) = 1 for all indexes i in given sequence
// IF first and last characters are not same
If (X[i] != X[j]) L(i, j) = max\{L(i + 1, j), L(i, j - 1)\}
    1)}
// If there are only 2 characters and both are same
Else if (i == i + 1) L(i, j) = 2
// If there are more than two characters, and first
    and last
// characters are same
Else L(i, j) = L(i + 1, j - 1) + 2
```

```
//Coin Change
To count the total number of solutions, we can divide
    all set solutions into two sets.
1) Solutions that do not contain mth coin (or Sm).
2) Solutions that contain at least one Sm.
Let count(S[], m, n) be the function to count the
    number of solutions, then it can be written as sum
     of count(S[], m-1, n) and count(S[], m, n-Sm).
//LOngest repeating Subsequence
int findLongestRepeatingSubSeq(string X, int m, int n)
   if(dp[m][n]!=-1)
   return dp[m][n];
   // return if we have reached the end of either
        string
   if (m == 0 | | n == 0)
       return dp[m][n] = 0;
   // if characters at index m and n matches
   // and index is different
   if (X\lceil m-1 \rceil == X\lceil n-1 \rceil \&\& m != n)
       return dp[m][n] = findLongestRepeatingSubSeq(X,
                         m - 1, n - 1) + 1;
   // else if characters at index m and n don't match
   return dp[m][n] = max (findLongestRepeatingSubSeq(
        X, m, n - 1),
                        findLongestRepeatingSubSeq(X,
                             m - 1, n);
// job-scheduling
1) First sort jobs according to finish time.
2) Now apply following recursive process.
  // Here arr[] is array of n jobs
  findMaximumProfit(arr[], n)
  Ł
    a) if (n == 1) return arr[0];
    b) Return the maximum of following two profits.
        (i) Maximum profit by excluding current job, i
             .e..
           findMaximumProfit(arr, n-1)
        (ii) Maximum profit by including the current
            job
```

```
//L[0] = {iob[0]}
L[i] = {MaxSum(L[i])} + job[i] where i < i and job[i].
    finish <= job[i].start</pre>
     = job[i], if there is no such j
```

```
dsu
19
void make set(int v) {
   parent[v] = v;
int find_set(int v) {
   if (v == parent[v])
       return v:
   return find_set(parent[v]);
void union_sets(int a, int b) {
   a = find set(a):
   b = find_set(b);
   if (a != b)
       parent[b] = a;
However this implementation is inefficient. It is easy
     to construct an example, so that the trees
    degenerate into long chains. In that case each
    call find_set(v) can take O(n)
time.
This is far away from the complexity that we want to
    have (nearly constant time). Therefore we will
    consider two optimizations that will allow to
    significantly accelerate the work.
Path compression optimization
```

This optimization is designed for speeding up find_set

Union by size / rank operation. To be precise, we will change which tree gets attached to the other one. In the native implementation the second tree always got attached to the first one. In practice that can lead to trees containing chains of length O(n)

```
If we call find_set(v) for some vertex v, we actually
    find the representative p for all vertices that we
     visit on the path between v and the actual
    representative p. The trick is to make the paths
    for all those nodes shorter, by setting the parent
     of each visited vertex directly to p.
You can see the operation in the following image. On
    the left there is a tree, and on the right side
    there is the compressed tree after calling
    find_set(7), which shortens the paths for the
    visited nodes 7, 5, 3 and 2.
Path compression of call <code>find_set(7)</code>
```

The new implementation of find_set is as follows: int find set(int v) { if (v == parent[v]) return v; return parent[v] = find_set(parent[v]);

The simple implementation does what was intended: first find the representative of the set (root vertex), and the in the process of stack unwinding the visited nodes are attached directly to the representative.

This simple modification of the operation already achieves the time complexity O(logn)

per call on average (here without proof). There is a second modification, that will make it even faster

In this optimization we will change the union_set

```
. With this optimization we will avoid this by choosing very carefully which tree gets attached.
```

There are many possible heuristics that can be used.

Most popular are the following two approaches: In
the first approach we use the size of the trees as
rank, and in the second one we use the depth of
the tree (more precisely, the upper bound on the
tree depth, because the depth will get smaller
when applying path compression).

In both approaches the essence of the optimization is the same: we attach the tree with the lower rank to the one with the bigger rank.

Here is the implementation of $\underline{\text{union}}$ by size:

```
void make set(int v) {
    parent[v] = v;
    size[v] = 1;
}
void union sets(int a. int b) {
    a = find_set(a);
    b = find_set(b);
    if (a != b) {
       if (size[a] < size[b])</pre>
           swap(a. b):
       parent[b] = a;
       size[a] += size[b];
}
void make set(int v) {
    parent[v] = v;
   rank[v] = 0;
}
void union_sets(int a, int b) {
    a = find_set(a);
    b = find_set(b);
```

```
if (a != b) {
       if (rank[a] < rank[b])</pre>
           swap(a, b);
       parent[b] = a:
       if (rank[a] == rank[b])
           rank[a]++:
   }
for (int i = 0; i < L; i++) {</pre>
    make_set(i);
for (int i = m-1; i >= 0; i--) {
   int 1 = querv[i].1:
   int r = query[i].r;
    int c = query[i].c;
   for (int v = find_set(1); v <= r; v = find_set(v))</pre>
       answer[v] = c:
       parent[v] = v + 1:
   }
void make_set(int v) {
   parent[v] = make_pair(v, 0);
   rank[v] = 0;
pair<int, int> find_set(int v) {
   if (v != parent[v].first) {
       int len = parent[v].second;
       parent[v] = find_set(parent[v].first);
       parent[v].second += len;
   return parent[v];
void union_sets(int a, int b) {
   a = find_set(a).first;
   b = find_set(b).first;
```

```
if (a != b) {
       if (rank[a] < rank[b])</pre>
          swap(a, b);
       parent[b] = make_pair(a, 1);
       if (rank[a] == rank[b])
          rank[a]++:
   }
Support the parity of the path length / Checking
    bipartiteness online
We give the implementation of the DSU that supports
    parity. As in the previous section we use a pair
    to store the ancestor and the parity. In addition
    for each set we store in the array bipartite[]
    whether it is still bipartite or not.
void make_set(int v) {
   parent[v] = make_pair(v, 0);
   rank[v] = 0;
   bipartite[v] = true;
pair<int, int> find_set(int v) {
   if (v != parent[v].first) {
       int parity = parent[v].second;
       parent[v] = find_set(parent[v].first);
       parent[v].second ^= parity;
   return parent[v];
void add_edge(int a, int b) {
   pair<int, int> pa = find_set(a);
   a = pa.first;
   int x = pa.second;
   pair<int, int> pb = find_set(b);
   b = pb.first;
   int y = pb.second;
   if (a == b) {
```

```
if (x == y)
          bipartite[a] = false;
} else {
    if (rank[a] < rank[b])
        swap (a, b);
    parent[b] = make_pair(a, x^y^1);
    bipartite[a] &= bipartite[b];
    if (rank[a] == rank[b])
        ++rank[a];
}
bool is_bipartite(int v) {
    return bipartite[find_set(v).first];
}</pre>
```

20 extend-euclid

```
int gcd(int a, int b, int & x, int & y) {
   if (a == 0) {
      x = 0;y = 1;
      return b;
   }
   int x1, y1;
   int d = gcd(b % a, a, x1, y1);
   x = y1 - (b / a) * x1;
   y = x1;
   return d;
}
```

21 floydwarshall

```
for (int k = 0; k < n; ++k) {
   for (int i = 0; i < n; ++i) {
      for (int j = 0; j < n; ++j) {
         d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
      }
   }
}</pre>
```

```
for (int k = 0; k < n; ++k) {
   for (int i = 0; i < n; ++i) {
      for (int j = 0; j < n; ++j) {
        if (d[i][k] < INF && d[k][j] < INF)
            d[i][j] = min(d[i][j], d[i][k] + d[k][j
            ]);
      }
   }
}</pre>
```

22 geometry

```
// 2d point
struct point2d {
   ftype x, y;
   point2d() {}
   point2d(ftype x, ftype y): x(x), y(y) {}
   point2d& operator+=(const point2d &t) {
       x += t.x;
       y += t.y;
       return *this;
   point2d& operator-=(const point2d &t) {
       x = t.x:
       y -= t.y;
       return *this;
   point2d& operator*=(ftype t) {
       x *= t:
       y *= t;
       return *this;
   point2d& operator/=(ftype t) {
       x /= t;
       v /= t;
       return *this;
   point2d operator+(const point2d &t) const {
       return point2d(*this) += t;
   point2d operator-(const point2d &t) const {
```

```
return point2d(*this) -= t;
   point2d operator*(ftype t) const {
       return point2d(*this) *= t;
   point2d operator/(ftype t) const {
       return point2d(*this) /= t;
};
point2d operator*(ftype a, point2d b) {
   return b * a;
//3D POINT
struct point3d {
   ftype x, y, z;
   point3d() {}
   point3d(ftype x, ftype y, ftype z): x(x), y(y), z(
   point3d& operator+=(const point3d &t) {
       x += t.x;
       y += t.y;
       z += t.z:
       return *this;
   point3d& operator-=(const point3d &t) {
       x -= t.x;
       y -= t.y;
       z = t.z;
       return *this;
   point3d& operator*=(ftype t) {
       x *= t;
       v *= t;
       z *= t;
       return *this:
   point3d& operator/=(ftype t) {
       x /= t;
       v /= t;
       z /= t;
       return *this;
   point3d operator+(const point3d &t) const {
       return point3d(*this) += t;
```

```
point3d operator-(const point3d &t) const {
       return point3d(*this) -= t;
   point3d operator*(ftype t) const {
       return point3d(*this) *= t;
   point3d operator/(ftype t) const {
       return point3d(*this) /= t;
   }
};
point3d operator*(ftype a, point3d b) {
   return b * a;
}
//DOT PRODUCT
ftype dot(point2d a, point2d b) {
   return a.x * b.x + a.y * b.y;
ftype dot(point3d a, point3d b) {
   return a.x * b.x + a.y * b.y + a.z * b.z;
//PROJECTION, ANGLE,...
ftype norm(point2d a) {
   return dot(a, a);
double abs(point2d a) {
   return sqrt(norm(a));
double proj(point2d a, point2d b) {
   return dot(a, b) / abs(b);
double angle(point2d a, point2d b) {
   return acos(dot(a, b) / abs(a) / abs(b));
// CROSS PRODUCT
point3d cross(point3d a, point3d b) {
   return point3d(a.y * b.z - a.z * b.y,
                 a.z * b.x - a.x * b.z
                 a.x * b.y - a.y * b.x);
}
ftype triple(point3d a, point3d b, point3d c) {
   return dot(a, cross(b, c));
ftype cross(point2d a, point2d b) {
```

```
return a.x * b.y - a.y * b.x;
// LINE INTERSECTION
point2d intersect(point2d a1, point2d d1, point2d a2,
    point2d d2) {
   return a1 + cross(a2 - a1, d2) / cross(d1, d2) *
        d1;
// PLANE INTERSECTION
point3d intersect(point3d a1, point3d n1, point3d a2,
    point3d n2, point3d a3, point3d n3) {
   point3d x(n1.x, n2.x, n3.x);
   point3d v(n1.v, n2.v, n3.v);
   point3d z(n1.z, n2.z, n3.z);
   point3d d(dot(a1, n1), dot(a2, n2), dot(a3, n3));
   return point3d(triple(d, y, z),
                 triple(x, d, z),
                 triple(x, y, d)) / triple(n1, n2, n3)
//INTERSECTION POINTS OF LINES
struct pt {
   double x, y;
};
struct line {
   double a, b, c;
const double EPS = 1e-9;
double det(double a, double b, double c, double d) {
   return a*d - b*c;
bool intersect(line m, line n, pt & res) {
   double zn = det(m.a, m.b, n.a, n.b);
   if (abs(zn) < EPS)
       return false;
   res.x = -det(m.c, m.b, n.c, n.b) / zn;
   res.y = -det(m.a, m.c, n.a, n.c) / zn;
   return true:
```

```
bool parallel(line m, line n) {
   return abs(det(m.a, m.b, n.a, n.b)) < EPS;</pre>
bool equivalent(line m, line n) {
   return abs(det(m.a. m.b. n.a. n.b)) < EPS
       && abs(det(m.a, m.c, n.a, n.c)) < EPS
       && abs(det(m.b. m.c. n.b. n.c)) < EPS:
//CHECK IF LINE INTERSECT
struct pt {
   long long x, y;
   pt() {}
   pt(long long _x, long long _y) : x(_x), y(_y) {}
   pt operator-(const pt& p) const { return pt(x - p.
        x, y - p.y; }
   long long cross(const pt& p) const { return x * p.
        y - y * p.x; }
   long long cross(const pt& a, const pt& b) const {
        return (a - *this).cross(b - *this); }
};
int sgn(const long long & x) { return } x >= 0 ? x ? 1 :
    0 : -1:  }
bool inter1(long long a, long long b, long long c,
    long long d) {
   if (a > b)
       swap(a. b):
   if (c > d)
       swap(c, d);
   return max(a, c) <= min(b, d);</pre>
bool check_inter(const pt& a, const pt& b, const pt& c
    , const pt% d) {
   if (c.cross(a, d) == 0 \&\& c.cross(b, d) == 0)
       return inter1(a.x, b.x, c.x, d.x) && inter1(a.y
           , b.y, c.y, d.y);
   return sgn(a.cross(b, c)) != sgn(a.cross(b, d)) &&
          sgn(c.cross(d, a)) != sgn(c.cross(d, b));
//CIRCLE LINE INTERSECTION
double r, a, b, c; // given as input
```

```
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if (c*c > r*r*(a*a+b*b)+EPS)
    puts ("no points");
else if (abs (c*c - r*r*(a*a+b*b)) < EPS) {
    puts ("1 point");
    cout << x0 << ', ' << y0 << '\n';
}
else {
    double d = r*r - c*c/(a*a+b*b);
    double mult = sqrt (d / (a*a+b*b));
    double ax, ay, bx, by;
    ax = x0 + b * mult;
    bx = x0 - b * mult;
    ay = y0 - a * mult;
    by = y0 + a * mult;
    puts ("2 points");
    cout << ax << ', ' << ay << '\n' << bx << ', ' << by
         << '\n';
}
//LENGTH OF UNION OF SEGMENTS
int length_union(const vector<pair<int, int>> &a) {
    int n = a.size():
    vector<pair<int, bool>> x(n*2);
    for (int i = 0; i < n; i++) {</pre>
       x[i*2] = \{a[i].first, false\};
       x[i*2+1] = \{a[i].second, true\};
    }
    sort(x.begin(), x.end());
    int result = 0;
    int. c = 0:
    for (int i = 0; i < n * 2; i++) {
       if (i > 0 && x[i].first > x[i-1].first && c >
           result += x[i].first - x[i-1].first;
       if (x[i].second)
           c++:
       else
           --c:
    return result;
//LATTICE POINTS
```

```
int count_lattices(Fraction k, Fraction b, long long n
    ) {
    auto fk = k.floor();
    auto fb = b.floor():
    auto cnt = OLL;
    if (k >= 1 || b >= 1) {
       cnt += (fk * (n - 1) + 2 * fb) * n / 2;
       k -= fk:
       b -= fb:
    auto t = k * n + b:
    auto ft = t.floor();
   if (ft >= 1) {
       cnt += count_lattices(1 / k, (t - t.floor()) /
            k, t.floor());
    }
    return cnt;
//AREA OF POLYGON
double area(const vector<point>& fig) {
    double res = 0;
    for (unsigned i = 0; i < fig.size(); i++) {</pre>
       point p = i ? fig[i - 1] : fig.back();
       point q = fig[i];
       res += (p.x - q.x) * (p.y + q.y);
    return fabs(res) / 2:
//COMMON TANGENT OF CIRCLE
struct pt {
    double x, y;
    pt operator- (pt p) {
       pt res = \{x-p.x, y-p.y\};
       return res;
    }
};
struct circle : pt {
    double r;
};
struct line {
    double a, b, c;
};
```

```
const double EPS = 1E-9;
double sqr (double a) {
   return a * a;
void tangents (pt c, double r1, double r2, vector<line
    > & ans) {
   double r = r2 - r1:
   double z = sqr(c.x) + sqr(c.y);
   double d = z - sqr(r);
   if (d < -EPS) return;</pre>
   d = sqrt (abs (d));
   line 1;
   1.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z;
   1.c = r1;
   ans.push_back (1);
vector<line> tangents (circle a, circle b) {
   vector<line> ans;
   for (int i=-1; i<=1; i+=2)</pre>
       for (int j=-1; j<=1; j+=2)
           tangents (b-a, a.r*i, b.r*j, ans);
   for (size t i=0: i<ans.size(): ++i)</pre>
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans:
// CONVEX HULL
struct pt {
   double x, y;
};
bool cmp(pt a, pt b) {
   return a.x < b.x || (a.x == b.x && a.y < b.y);
bool cw(pt a, pt b, pt c) {
   return a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y) <
         0;
```

```
bool ccw(pt a, pt b, pt c) {
    return a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y) >
}
void convex_hull(vector<pt>& a) {
    if (a.size() == 1)
       return;
    sort(a.begin(), a.end(), &cmp);
    pt p1 = a[0], p2 = a.back();
    vector<pt> up, down;
    up.push_back(p1);
    down.push_back(p1);
    for (int i = 1; i < (int)a.size(); i++) {</pre>
       if (i == a.size() - 1 || cw(p1, a[i], p2)) {
           while (up.size() >= 2 && !cw(up[up.size()
               -2], up[up.size()-1], a[i]))
               up.pop_back();
           up.push_back(a[i]);
       }
       if (i == a.size() - 1 || ccw(p1, a[i], p2)) {
           while(down.size() >= 2 && !ccw(down[down.
               size()-2], down[down.size()-1], a[i]))
               down.pop_back();
           down.push_back(a[i]);
       }
    }
    a.clear();
    for (int i = 0; i < (int)up.size(); i++)</pre>
       a.push_back(up[i]);
    for (int i = down.size() - 2; i > 0; i--)
       a.push_back(down[i]);
}
```

23 inclu-exclu

```
int solve (int n, int r) {
  vector<int> p;
  for (int i=2; i*i<=n; ++i)</pre>
```

```
if (n % i == 0) {
       p.push_back (i);
       while (n \% i == 0)
           n /= i:
if (n > 1)
    p.push_back (n);
int sum = 0:
for (int msk=1; msk<(1<<p.size()); ++msk) {</pre>
    int mult = 1,
       bits = 0;
    for (int i=0; i<(int)p.size(); ++i)</pre>
       if (msk & (1<<i)) {</pre>
           ++bits:
           mult *= p[i];
    int cur = r / mult;
    if (bits % 2 == 1)
       sum += cur;
    else
       sum -= cur;
}
return r - sum;
```

24 kmp

```
for (int i = 0; i < n; i++)
    ans[pi[i]]++;
for (int i = n-1; i > 0; i--)
    ans[pi[i-1]] += ans[i];
for (int i = 0; i <= n; i++)
    ans[i]++;</pre>
```

25 kruskal

```
vector<int> parent, rank;
void make_set(int v) {
   parent(v) = v;
   rank[v] = 0;
int find set(int v) {
   if (v == parent[v])
       return v;
    return parent[v] = find_set(parent[v]);
void union_sets(int a, int b) {
    a = find_set(a);
    b = find_set(b);
    if (a != b) {
       if (rank[a] < rank[b])</pre>
           swap(a, b);
       parent[b] = a;
       if (rank[a] == rank[b])
           rank[a]++;
   }
struct Edge {
   int u, v, weight;
   bool operator<(Edge const& other) {</pre>
       return weight < other.weight;</pre>
   }
};
int n;
```

```
vector<Edge> edges;
int cost = 0;
vector<Edge> result;
parent.resize(n);
rank.resize(n);
for (int i = 0; i < n; i++)
    make_set(i);

sort(edges.begin(), edges.end());

for (Edge e : edges) {
    if (find_set(e.u) != find_set(e.v)) {
        cost += e.weight;
        result.push_back(e);
        union_sets(e.u, e.v);
    }
}</pre>
```

26 lazy

```
Range updates (Lazy Propagation)
Addition on segments
void build(int a[], int v, int tl, int tr) {
   if (t1 == tr) {
       t[v] = a[t1];
   } else {
       int tm = (tl + tr) / 2:
       build(a, v*2, t1, tm);
       build(a, v*2+1, tm+1, tr);
       t[v] = 0;
   }
}
void update(int v, int tl, int tr, int l, int r, int
    add) {
   if (1 > r)
       return:
   if (1 == t1 && r == tr) {
       t[v] += add;
```

```
} else {
       int tm = (t1 + tr) / 2;
       update(v*2, t1, tm, 1, min(r, tm), add);
       update(v*2+1, tm+1, tr, max(1, tm+1), r, add);
int get(int v, int tl, int tr, int pos) {
   if (tl == tr)
       return t[v]:
   int tm = (tl + tr) / 2;
   if (pos <= tm)</pre>
       return t[v] + get(v*2, t1, tm, pos);
       return t[v] + get(v*2+1, tm+1, tr, pos);
Assignment on segments
Suppose now that the modification query asks to assign
     each element of a certain segment a[lr]
to some value p. As a second query we will again
    consider reading the value of the array a[i]
void push(int v) {
   if (marked[v]) {
       t\lceil v*2\rceil = t\lceil v*2+1\rceil = t\lceil v\rceil:
       marked[v*2] = marked[v*2+1] = true;
       marked[v] = false:
   }
void update(int v, int tl, int tr, int l, int r, int
    new_val) {
   if (1 > r)
       return:
   if (1 == tl && tr == r) {
       t[v] = new val:
       marked[v] = true;
   } else {
       push(v);
       int tm = (tl + tr) / 2;
       update(v*2, tl, tm, l, min(r, tm), new_val);
```

```
update(v*2+1, tm+1, tr, max(1, tm+1), r,
           new_val);
   }
int get(int v, int tl, int tr, int pos) {
   if (t1 == tr) {
       return t[v]:
   push(v):
   int tm = (tl + tr) / 2;
   if (pos <= tm)</pre>
       return get(v*2, t1, tm, pos);
       return get(v*2+1, tm+1, tr, pos);
Adding on segments, querying for maximum
void push(int v) {
   t[v*2] += lazv[v];
   lazv[v*2] += lazv[v];
   t[v*2+1] += lazv[v]:
   lazv[v*2+1] += lazv[v];
   lazy[v] = 0;
void update(int v, int tl, int tr, int l, int r, int
    addend) {
   if (1 > r)
       return;
   if (1 == t1 && tr == r) {
       t[v] += addend:
       lazy[v] += addend;
   } else {
       push(v):
       int tm = (t1 + tr) / 2;
       update(v*2, tl, tm, l, min(r, tm), addend);
       update(v*2+1, tm+1, tr, max(1, tm+1), r, addend
           );
       t[v] = max(t[v*2], t[v*2+1]);
   }
int query(int v, int tl, int tr, int l, int r) {
```

27 linear-equation

```
int gauss (vector < vector < double> > a, vector < double>
     & ans) {
   int n = (int) a.size();
   int m = (int) a[0].size() - 1;
   vector<int> where (m. -1):
   for (int col=0, row=0; col<m && row<n; ++col) {</pre>
       int sel = row;
       for (int i=row: i<n: ++i)</pre>
           if (abs (a[i][col]) > abs (a[sel][col]))
               sel = i:
       if (abs (a[sel][col]) < EPS)</pre>
           continue:
       for (int i=col; i<=m; ++i)</pre>
           swap (a[sel][i], a[row][i]);
       where [col] = row;
       for (int i=0; i<n; ++i)</pre>
           if (i != row) {
               double c = a[i][col] / a[row][col];
               for (int j=col; j<=m; ++j)</pre>
                   a[i][j] -= a[row][j] * c;
           }
       ++row;
   ans.assign (m, 0);
   for (int i=0; i<m; ++i)</pre>
       if (where[i] != -1)
           ans[i] = a[where[i]][m] / a[where[i]][i];
   for (int i=0; i<n; ++i) {</pre>
       double sum = 0;
       for (int j=0; j<m; ++j)</pre>
```

28 matrixexpo

```
//answer in F
void power(11 F[2][2],11 n){
if(n<=1)
 return:
11 M[2][2]={{(2*f)%MOD,-1},{1,0}};
power(F,n/2);multi(F,F);
if (n%2)
 multi(F,M);
void multi(11 F[2][2],11 M[2][2]){
11 x=(F[0][0]\%MOD*M[0][0]\%MOD)\%MOD+(F[0][1]\%MOD*M
     [1] [0] %MOD) %MOD;
ll y=(F[0][0]%MOD*M[0][1]%MOD)%MOD+(F[0][1]%MOD*M
     [1] [1] %MOD) %MOD;
ll z=(F[1][0]%MOD*M[0][0]%MOD)%MOD+(F[1][1]%MOD*M
     [1] [0] %MOD) %MOD;
ll w=(F[1][0]%MOD*M[0][1]%MOD)%MOD+(F[1][1]%MOD*M
     [1] [1] %MOD) %MOD;
if(x<0)
 x=(x+MOD)%MOD;
 if(v<0)
 y=(y+MOD)%MOD;
 if(z<0)
 z=(z+MOD)%MOD:
if(x<0)
 w=(w+MOD)%MOD;
F[0][0]=x;F[0][1]=y;F[1][0]=z;F[1][1]=w;
```

29 maxflow

```
int n;
vector<vector<int>> capacity;
vector<vector<int>> adj;
int bfs(int s, int t, vector<int>& parent) {
   fill(parent.begin(), parent.end(), -1);
   parent[s] = -2;
   queue<pair<int, int>> q;
   q.push({s, INF});
   while (!q.empty()) {
       int cur = q.front().first;
       int flow = q.front().second;
       q.pop();
       for (int next : adj[cur]) {
          if (parent[next] == -1 && capacity[cur][
               nextl) {
              parent[next] = cur;
              int new_flow = min(flow, capacity[cur][
                  nextl):
              if (next == t)
                  return new_flow;
              q.push({next, new_flow});
       }
   }
   return 0:
int maxflow(int s, int t) {
   int flow = 0;
   vector<int> parent(n);
   int new_flow;
   while (new_flow = bfs(s, t, parent)) {
       flow += new_flow;
       int cur = t:
       while (cur != s) {
          int prev = parent[cur];
```

```
capacity[prev][cur] -= new_flow;
capacity[cur][prev] += new_flow;
cur = prev;
}

return flow;
}
```

30 merge-sortree

```
vector<int> t[4*MAXN];
void build(int a[], int v, int tl, int tr) {
   if (t1 == tr) {
       t[v] = vector<int>(1, a[t1]);
   } else {
       int tm = (t1 + tr) / 2:
       build(a, v*2, t1, tm);
       build(a, v*2+1, tm+1, tr);
       merge(t[v*2].begin(), t[v*2].end(), t[v*2+1].
           begin(), t[v*2+1].end(),
            back_inserter(t[v]));
}
int query(int v, int tl, int tr, int l, int r, int x)
   if (1 > r)
       return INF:
   if (1 == t1 && r == tr) {
       vector<int>::iterator pos = lower_bound(t[v].
           begin(), t[v].end(), x);
       if (pos != t[v].end())
          return *pos;
       return INF;
   }
   int tm = (tl + tr) / 2;
   return min(query(v*2, t1, tm, 1, min(r, tm), x),
             query(v*2+1, tm+1, tr, max(1, tm+1), r, x)
                 ));
```

```
void update(int v, int tl, int tr, int pos, int
    new_val) {
    t[v].erase(t[v].find(a[pos]));
    t[v].insert(new_val);
    if (tl != tr) {
        int tm = (tl + tr) / 2;
        if (pos <= tm)
            update(v*2, tl, tm, pos, new_val);
        else
            update(v*2+1, tm+1, tr, pos, new_val);
    } else {
        a[pos] = new_val;
    }
}</pre>
```

31 nge-stack

```
void printNGE(int arr[], int n) {
 stack < int > s;
 /* push the first element to stack */
 s.push(arr[0]);
 // iterate for rest of the elements
 for (int i = 1: i < n: i++) {
   if (s.empty()) {
     s.push(arr[i]);
     continue;
   /* if stack is not empty, then
      pop an element from stack.
      If the popped element is smaller
      than next, then
   a) print the pair
   b) keep popping while elements are
   smaller and stack is not empty */
   while (s.empty() == false && s.top() < arr[i])</pre>
```

```
cout << s.top() << " --> " << arr[i] << endl;
    s.pop();
}

/* push next to stack so that we can find
    next greater for it */
    s.push(arr[i]);
}

/* After iterating over the loop, the remaining
    elements in stack do not have the next greater
    element, so print -1 for them */
while (s.empty() == false) {
    cout << s.top() << " --> " << -1 << endl;
    s.pop();
}
</pre>
```

32 prim

```
int n;
vector<vector<int>> adj; // adjacency matrix of graph
const int INF = 1000000000; // weight INF means there
    is no edge
struct Edge {
   int w = INF, to = -1;
};
void prim() {
   int total_weight = 0;
   vector<bool> selected(n):
   vector<Edge> min_e(n);
   min_e[0].w = 0;
   for (int i=0; i<n; ++i) {</pre>
       int v = -1;
       for (int j = 0; j < n; ++j) {
           if (!selected[j] && (v == -1 || min_e[j].w
               < min_e[v].w))
               v = j;
       if (min_e[v].w == INF) {
           cout << "No MST!" << endl;</pre>
           exit(0);
```

33 seg-tree

```
vector<int>tree(400020),arr(100005);
int n.k:
void build(int node, int start, int end){//1,1,n
   if(start == end)
       tree[node] = A[start];
   elsef
       int mid = (start + end) / 2;
       build(2*node, start, mid);
       build(2*node+1, mid+1, end);
       tree[node] = tree[2*node] + tree[2*node+1];
   }
}
void update(int node, int start, int end, int idx, int
     val){//1,1,n,i,val i is 1 based
   if(start == end){
       arr[idx] += val;
       tree[node] += val;
   }
   else{
       int mid = (start + end) / 2:
       if(start <= idx && idx <= mid)</pre>
           update(2*node, start, mid, idx, val);
       else
           update(2*node+1, mid+1, end, idx, val);
       tree[node] = tree[2*node] + tree[2*node+1];
   }
}
```

34 segmented-sieve

```
#define MAX 46656
#define LMT 216
#define LEN 4830
#define RNG 100032
#define sq(x) ((x)*(x))
#define mset(x,v) memset(x, v , sizeof(x))
#define chkC(x,n) (x[n >> 6] & (1 << ((n >> 1) & 31)))
#define setC(x,n) (x[n >> 6] |= (1 << ((n >> 1) & 31))
using namespace std;
unsigned base[MAX/64], segment[RNG/64], primes[LEN];
* Generates all the necessary prime numbers and marks
      them in base∏
void sieve()
   unsigned i, j, k;
   for (i = 3; i < LMT; i += 2)</pre>
       if (!chkC(base, i))
          for (j = i * i, k = i << 1; j < MAX; j += k
              setC(base, j);
```

```
for (i = 3, j = 0; i < MAX; i += 2)
   {
       if (!chkC(base, i))
           primes[j++] = i;
   }
* Returns the prime-count within range [a,b] and
     marks them in segment[]
int segmented_sieve(int a, int b)
   unsigned i, j, k, cnt = (a <= 2 && 2 <=b )? 1 : 0;
   if (b < 2)
       return 0;
   if (a < 3)
       a = 3:
   if (a % 2 == 0)
       a++;
   mset (segment, 0);
   for (i = 0; sq(primes[i]) <= b; i++)</pre>
       j = primes[i] * ((a + primes[i] - 1) / primes[i
           ]);
       if (j % 2 == 0) j += primes[i];
       for (k = primes[i] << 1; j <= b; j += k)</pre>
           if (j != primes[i])
              setC(segment, (j - a));
       }
   for (i = 0; i \le b - a; i += 2)
       if (!chkC(segment, i))
           cnt++;
   return cnt;
   sieve():
   int a, b;
```

```
else {
    for (int i=l, end=(c_l+1)*len-1; i<=end; ++i)
        sum += a[i];
    for (int i=c_l+1; i<=c_r-1; ++i)
        sum += b[i];
    for (int i=c_r*len; i<=r; ++i)
        sum += a[i];
}</pre>
```

35 sqrt-decomp

```
// input data
int n;
vector<int> a (n);
// preprocessing
int len = (int) sqrt (n + .0) + 1; // size of the
    block and the number of blocks
vector<int> b (len):
for (int i=0; i<n; ++i)</pre>
    b[i / len] += a[i];
// answering the queries
for (;;) {
    int 1, r;
 // read input data for the next query
    int sum = 0;
    for (int i=1; i<=r; )</pre>
       if (i % len == 0 && i + len - 1 <= r) {// if</pre>
            the whole block starting at i belongs to [1
            : r]
           sum += b[i / len]:
           i += len:
       }
       else {
           sum += a[i];
           ++i;
       }
}
int sum = 0;
int c_1 = 1 / len, c_r = r / len;
if (c 1 == c r)
    for (int i=1; i<=r; ++i)</pre>
       sum += a[i];
```

36 string-hashing

37 topological-sort

```
int n; // number of vertices
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> ans;

void dfs(int v) {
    visited[v] = true;
    for (int u : adj[v]) {
        if (!visited[u])
            dfs(u);
    }
    ans.push_back(v);
}
```

```
void topological_sort() {
   visited.assign(n, false);
   ans.clear();
   for (int i = 0; i < n; ++i) {
       if (!visited[i])
            dfs(i);
   }
   reverse(ans.begin(), ans.end());
}</pre>
```

38 trie

```
const int ALPHABET_SIZE = 26;
struct TrieNode {
struct TrieNode *children[ALPHABET_SIZE]; //
     isEndOfWord is true if the node represents end of
bool isEndOfWord;
struct TrieNode *getNode(void) {
struct TrieNode *pNode = new TrieNode;
pNode->isEndOfWord = false;
for (int i = 0; i < ALPHABET_SIZE; i++)</pre>
 pNode->children[i] = NULL;
return pNode;
void insert(struct TrieNode *root, string key) {
struct TrieNode *pCrawl = root;
for (int i = 0; i < key.length(); i++) {</pre>
 int index = key[i] - 'a';
 if (!pCrawl->children[index])
  pCrawl->children[index] = getNode();
 pCrawl = pCrawl->children[index];
} // mark last node as leaf
pCrawl->isEndOfWord = true;
bool search(struct TrieNode *root, string key) {
struct TrieNode *pCrawl = root;
for (int i = 0; i < key.length(); i++) {</pre>
 int index = key[i] - 'a';
 if (!pCrawl->children[index])
  return false;
```

```
pCrawl = pCrawl->children[index];
}
return (pCrawl != NULL && pCrawl->isEndOfWord);
}
struct TrieNode *root = getNode();
for (int i = 0; i < n; i++) insert(root, keys[i]);
search(root, "the");</pre>
```

39 z

vector<ll>z;

```
11 k = i-1;
    if(z[k]<r-i+1)
        z.pb(z[k]);
    else
    {
        l=i;
        while(r<sz && s[r-1]==s[r])
        r++;
        z.pb(r-1);r--;
    }
}</pre>
```