Divide & Conquer SOLUTIONS

Solution 1:

```
class Solution {
  public static String[] mergeSort(String[] arr, int lo, int hi) {
      String[] arr1 = mergeSort(arr, lo, mid);
      String[] arr2 = mergeSort(arr, mid + 1, hi);
      String[] arr3 = merge(arr1, arr2);
  static String[] merge(String[] arr1, String[] arr2) {
          if (isAlphabeticallySmaller(arr1[i], arr2[j])) {
```

```
arr3[idx] = arr2[j];
if (str1.compareTo(str2) < 0) {</pre>
String[] a = mergeSort(arr, 0, arr.length - 1);
```

Solution 2:

Approach 1 - Brute Force Approach

Idea : Count the number of times each number occurs in the array & find the largest count. Time complexity - $O(n^2)$

```
class Solution {
   public static int majorityElement(int[] nums) {
    int majorityCount = nums.length/2;
```

```
for (int i=0; i<nums.length; i++) {
    int count = 0;
    for (int j=0; j<nums.length; j++) {
        if (nums[j] == nums[i]) {
            count += 1;
        }
        if (count > majorityCount) {
            return nums[i];
        }
    }
    return -1;
}

public static void main(String args[]) {
    int nums[] = {2,2,1,1,1,2,2};
        System.out.println(majorityElement(nums));
}
```

Approach 2 - Divide & Conquer

Idea: If we know the majority element in the left and right halves of an array, we can determine which is the global majority element in linear time.

Here, we apply a classical divide & conquer approach that recurses on the left and right halves of an array until an answer can be trivially achieved for a length-1 array. Note that because actually passing copies of subarrays costs time and space, we instead pass lo and hi indices that describe the relevant slice of the overall array. In this case, the majority element for a length-1 slice is trivially its only element, so the recursion stops there. If the current slice is longer than length-1, we must combine the answers for the slice's left and right halves. If they agree on the majority element, then the majority element for the overall slice is obviously the same[1]. If they disagree, only one of them can be "right", so we need to count the occurrences of the left and right majority elements to determine which subslice's answer is globally correct. The overall answer for the array is thus the majority element between indices 0 and n.

Time complexity - O(n*logn)

```
class Solution {
  private static int countInRange(int[] nums, int num, int lo, int hi) {
    int count = 0;
    for (int i = lo; i <= hi; i++) {
        if (nums[i] == num) {</pre>
```

```
int mid = (hi-lo)/2 + lo;
    int right = majorityElementRec(nums, mid+1, hi);
       return left;
    int leftCount = countInRange(nums, left, lo, hi);
    int rightCount = countInRange(nums, right, lo, hi);
public static void main(String args[]) {
    int nums[] = \{2, 2, 1, 1, 1, 2, 2\};
```

Solution 3:

Approach 1 - Brute Force Approach

Idea: Traverse through the array, and for every index, find the number of smaller elements on its right side of the array. This can be done using a nested loop. Sum up the counts for all indexes in the array and print the sum.

- Traverse through the array from start to end
- For every element, find the count of elements smaller than the current number up to that index using another loop.
- Sum up the count of inversion for every index.
- Print the count of inversions.

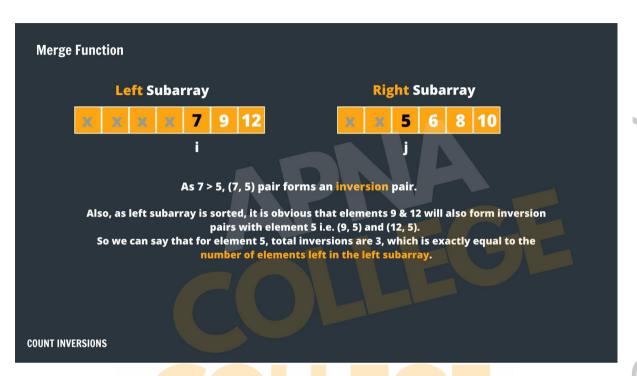
Time complexity - O(n^2)

Approach 2 - Modified Merge Sort

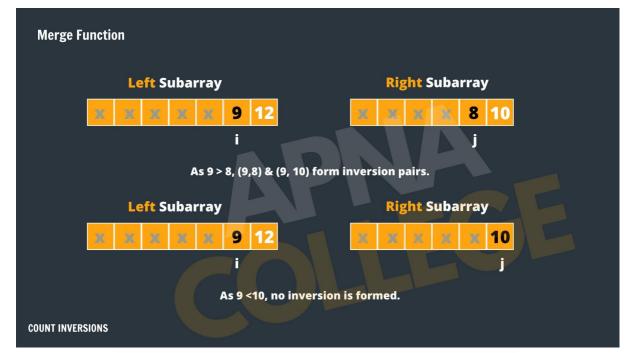
Idea: Suppose the number of inversions in the left half and right half of the array (let be inv1 and inv2); what kinds of inversions are not accounted for in Inv1 + Inv2? The answer is – the inversions that need to be counted during the merge step. Therefore, to get the total number of inversions that need to be added are the number of inversions in the left subarray, right subarray, and merge().

Basically, for each array element, count all elements more than it to its left and add the count to the output. This whole magic happens inside the merge function of merge sort.

Let's consider two subarrays involved in the merge process:









Algorithm:

- Split the given input array into two halves, left and right similar to merge sort recursively.
- Count the number of inversions in the left half and right half along with the inversions found during the merging of the two halves.
- Stop the recursion, only when 1 element is left in both halves.
- To count the number of inversions, we will use a two pointers approach. Let us consider two pointers i and j, one pointing to the left half and the other pointing towards the right half.

• While iterating through both the halves, if the current element A[i] is less than A[j], add it to the sorted list, else increment the count by mid – i.

Time complexity - O(n* logn)

```
public class Solution {
  public static int merge(int arr[], int left, int mid, int right) {
      int temp[] = new int[(right - left + 1)];
              temp[k] = arr[i];
              temp[k] = arr[j];
          temp[k] = arr[i];
          temp[k] = arr[j];
          arr[i] = temp[k];
```

```
private static int mergeSort(int arr[], int left, int right) {
        invCount = mergeSort(arr, left, mid);
        invCount += mergeSort(arr, mid + 1, right);
       invCount += merge(arr, left, mid + 1, right);
   return mergeSort(arr, 0, n - 1);
public static void main(String args[]) {
   System.out.println("Inversion Count = " + getInversions(arr));
```