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CS 130

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LAB #3

A.1 [10] <COD §A.2> In addition to the basic laws we discussed in this section, there are two important theorems, called DeMorgan's theorems:

$$A + B^- = A^- \cdot B^- \text{ and } A \cdot B^- = A^- + B^-$$

Prove DeMorgan's theorems with a truth table of the form

A	B	\bar{A}	\bar{B}	$\bar{A} + \bar{B}$	$\bar{A} \cdot \bar{B}$	$\bar{A} \cdot B$	$\bar{A} + B$
0	0	1	1	1	1	1	1
0	1	1	0	0	0	1	1
1	0	0	1	0	0	1	1
1	1	0	0	0	0	0	0

A.2 [15] <COD §A.2> Prove that the two equations for E are equivalent by using DeMorgan's theorems and the axioms shown COD Section A.2 (Gates, truth tables, and logical equations).

$$E = ((A \cdot B) + (A \cdot C) + (B \cdot C)) \cdot (A \cdot B \cdot C^-)$$

$$E = (A \cdot B \cdot C^-) + (A \cdot C \cdot B^-) + (B \cdot C \cdot A^-)$$

A.3 [10] <COD §A.2> Show that there are 2^n entries in a truth table for a function with n inputs.

A.4 [10] <COD §A.2> One logic function that is used for a variety of purposes (including within adders and to compute parity) is *exclusive OR*. The output of a two-input exclusive OR function is true only if exactly one of the inputs is true. Show the truth table for a two-input exclusive OR function and implement this function using AND gates, OR gates, and inverters.

A.5 [15] <COD §A.2> Prove that the NOR gate is universal by showing how to build the AND, OR, and NOT functions using a two-input NOR gate.

A.6 [15] <COD §A.2> Prove that the NAND gate is universal by showing how to build the AND, OR, and NOT functions using a two-input NAND gate.

A. | from definition A NAND gate is equivalent to an inverter followed by an AND and a NOR gate is equivalent to inversion followed by an AND

$$\overline{AB} \equiv \overline{A} + \overline{B}$$

$$\overline{A+B} \equiv \overline{A} \overline{B}$$

from De Morgan's theorem

from true table, we compare true value between these input and the output it is same through the gate. it is true that we have the same as the definition

(A₂)

$$E = ((A.B + A.C + (B.C)) (A.B.C))'$$

$$E = (A.B.C') + (A.C.B') + (B.C.A')$$

$$E = (AB + AC + BC) (\overline{A} + \overline{B} + \overline{C})$$

$$E = (\overline{A}BC + A\overline{B}C + AB\overline{C})$$

As above

Since $A\overline{A} = 0$

$$B.\overline{B} = 0$$

$$C.\overline{C} = 0$$

(A3)

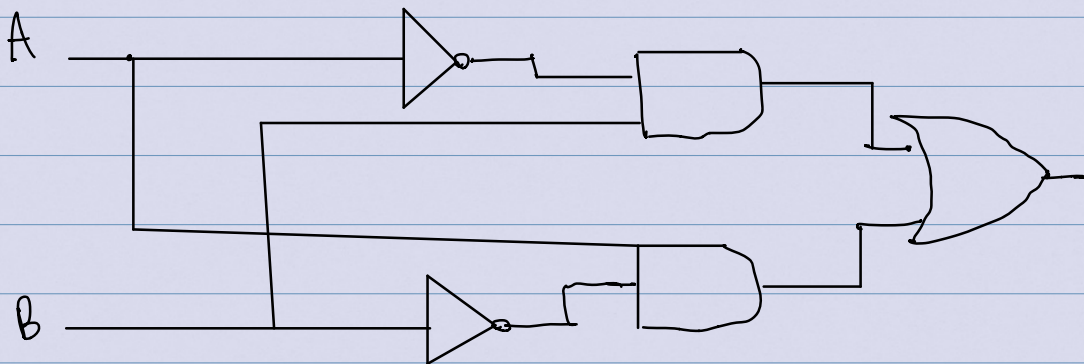
the combination of input has value of true and false, nonetheless as how many input will produce 2^n power since n equals to number input
 $2 \text{ input} = 4$ row of true table

(A.4)

A	B	$A \oplus B$
0	0	0
0	1	1
1	0	1
1	1	0

equivalent logic is equal to

$$(\bar{A} \cdot B) + (A \cdot \bar{B})$$

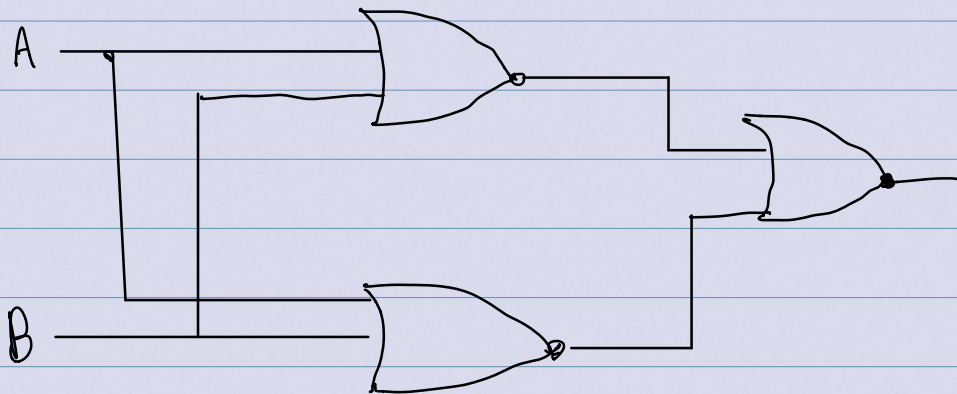


(A.5)

Notgate $\overline{A + A} = \bar{A}$

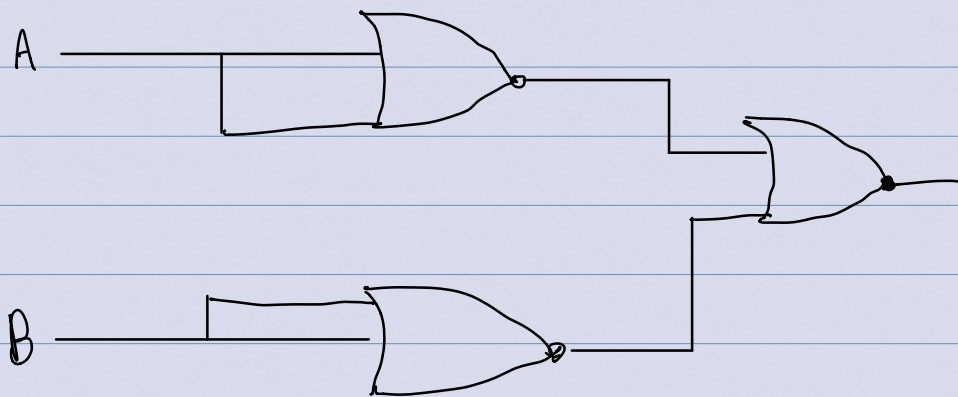
or gate

$$A + B = \overline{\overline{A + B}} = \overline{A \text{ Nor } B}$$



AND gate

$$A \cdot B = \overline{\overline{A \cdot B}} = \overline{\overline{A} + \overline{B}} \\ \bar{A} \text{ Nor } \bar{B}$$

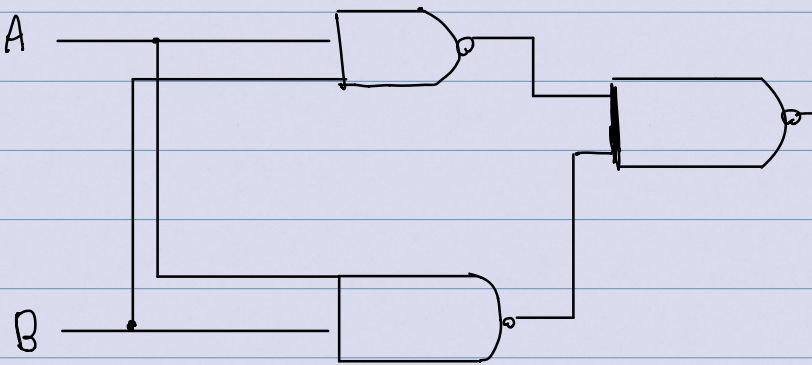


(A-b) Please NAND gate

$$\text{NOT gate} = \overline{A \cdot A} = \bar{A}$$

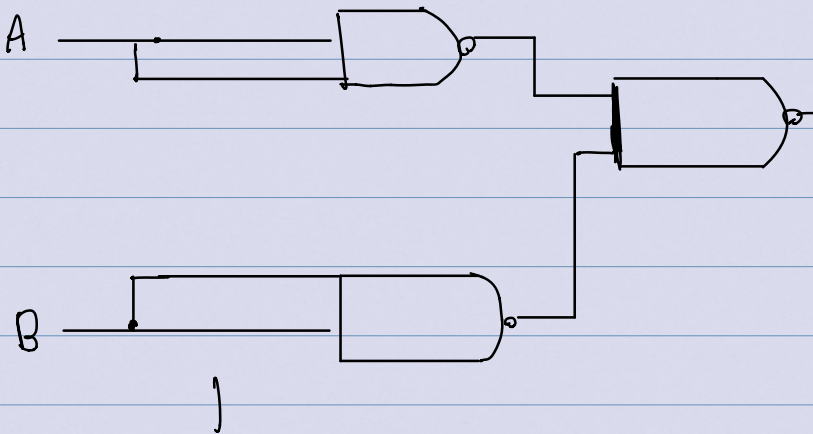
$$\text{AND gate} \quad A \cdot B = A \cdot \overline{\overline{A \cdot B}} = \overline{\overline{A \cdot B}} \text{ NAND } B$$

Therefore we must invert output of the NAND gate
so we can use three NAND gates



OR gate

$$A + B = \overline{\overline{A + B}} = \overline{\overline{A} \cdot \overline{B}} = \overline{\overline{A} \text{ NAND } \overline{B}}$$



Part A:

Methods to describe a combinational circuit

- True table
- Boolean algebraic expression
- Logic diagram

Ten properties of boolean algebra

- Commutative $x + y = y + x$, $x \cdot y = y \cdot x$
- Associative $(x + y) + z = x + (y + z)$, $(x \cdot y) \cdot z = x \cdot (y \cdot z)$
- Distributive $x + (y \cdot z) = (x + y) \cdot (x + z)$, $x \cdot (y + z) = (x \cdot y) + (x \cdot z)$
- Identity $x + 0 = x$, $x \cdot 1 = x$
- Complement $x + (x') = 1$, $x \cdot (x') = 0$

Precedence	Operator
Highest	Complement
	AND
Lowest	OR

Duality

Exchange $+$ and \cdot

Exchange 1 and 0

Zero theorem

$$x + 1 = 1$$

$$x \cdot 0 = 0$$

De Morgan's law

$$(a \cdot b)' = a' + b'$$

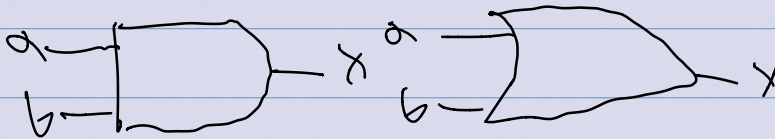
$$(a + b)' = a' \cdot b'$$

complement theorems

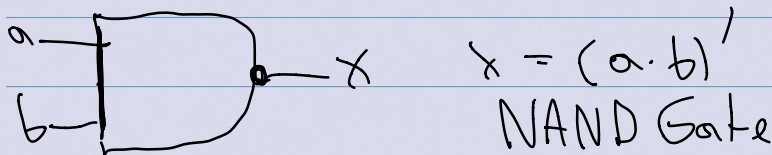
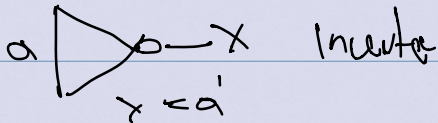
$$(x')' = x$$

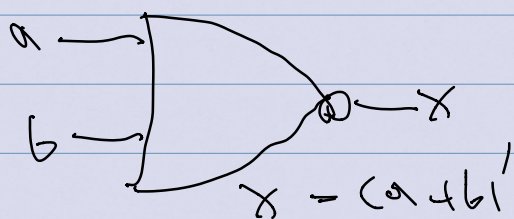
$$1' = 0$$

$$0' = 1$$

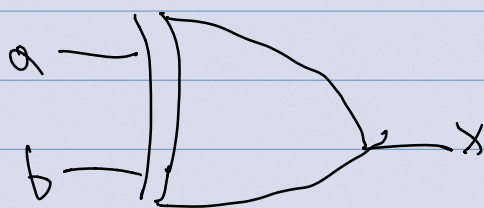


$x = a \cdot b$ AND Gate $x = a + b$ OR gate





NOR GATE



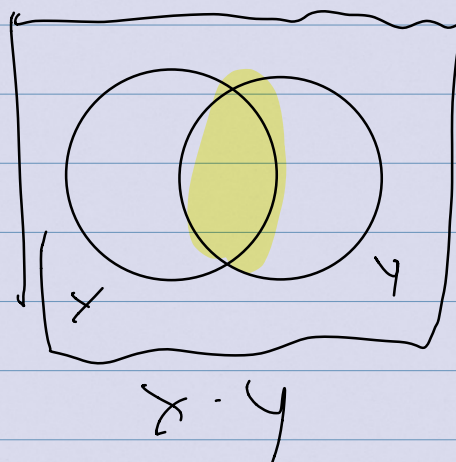
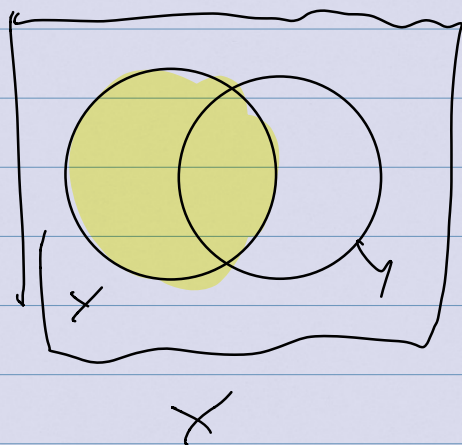
$$x = a \oplus b$$

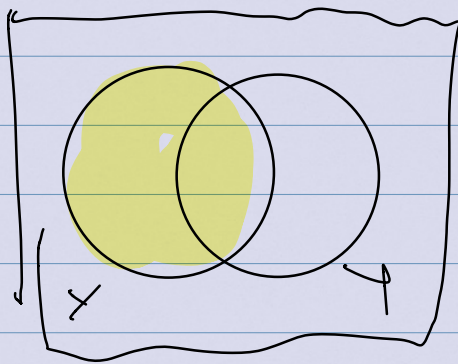
XOR GATE

OR Gate is set union

AND Gate is set intersection

INVERTER is set of complement





$$X + X \cdot Y$$