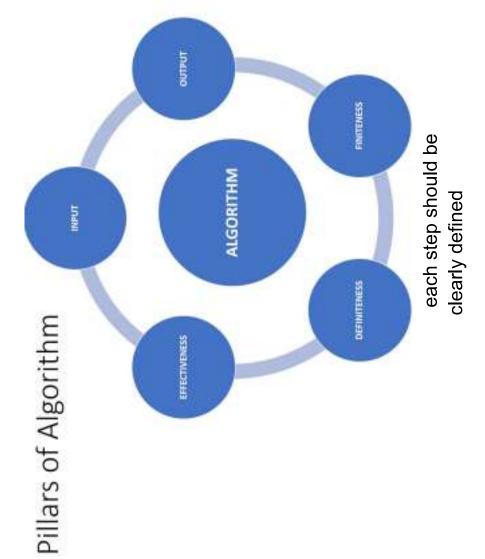
What is Algorithm?

 Set of finite steps if followed will solve given problem in finite steps and finite time.



12/1/22, 5:08 PM

12/1/22, 5:08 PM

Analysis ?

one.

Why?
 So one can predict time and space complexity and then select tl

Analysis?

- When ?
- Priori
- Done before implementation
- Using pen paper
- Posteriori
- Done after implementation
- Using software by professionals

Analysis ?

- What we get?Time Complexity

Space Complexity

https://docs.google.com/document/d/1Wyh7M9U_LmeJqSGCXQlAjJ5lpVCKaxkkPMSUu2hhst8/edit

Analysis ? On What? Best Case Worst Case

https://docs.google.com/document/d/1Wyh7M9U_LmeJqSGCXQIAjJ5lpVCKaxkkPMSUu2hhst8/edit

Time Complexity Comments -0 Operation---1

https://docs.google.com/document/d/1Wyh7M9U_LmeJqSGCXQIAjJ5lpVCKaxkkPMSUu2hhst8/edit

Time Complexity Comments -0 Operation---1

Problem

 An algorithm takes 0.5milli seconds for 100 inputs calculate time if input is increased to 500 and if complexities are

Linear

Quadratic

·Cubic

Log based

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Asymptotic Notation

 Asymptotic Notation is used to describe the running time of an algorithm - how much time an algorithm takes with a given inp

Algorithmic Common Runtimes

fastest to slowest are:

constant: 0(1)

logarithmic: O(log N)

linear: O(N)

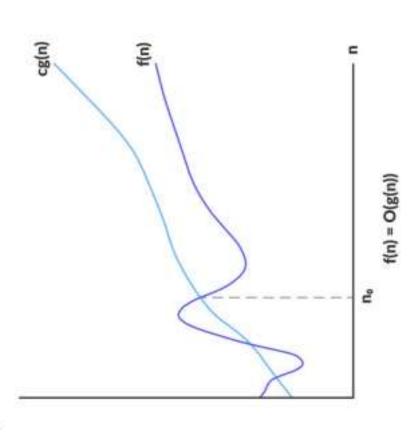
polynomial: O(N^2)

exponential: O(2^N)

factorial: O(N!)

Big-O Notation (O-notation)

 Big-O notation represents the upper bound of the running time of an algorith Thus, it gives the worst-case complexity of an algorithm.



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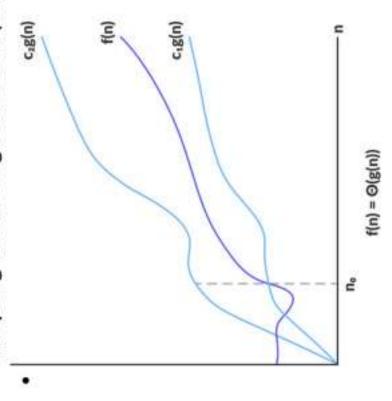
Big-O Notation (O-notation)

Rules

- Constant has no value
- In addition use max(worst only)
- In nesting multiply and then use addition rule

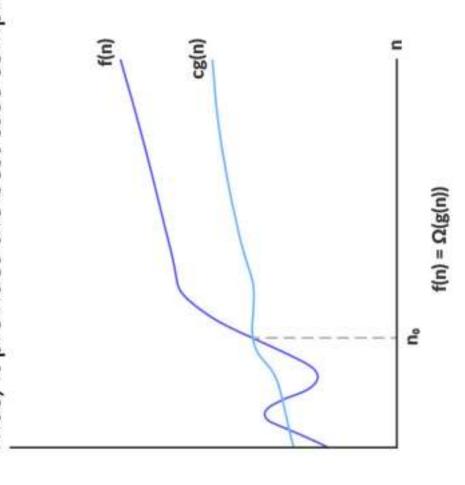
Theta Notation (O-notation)

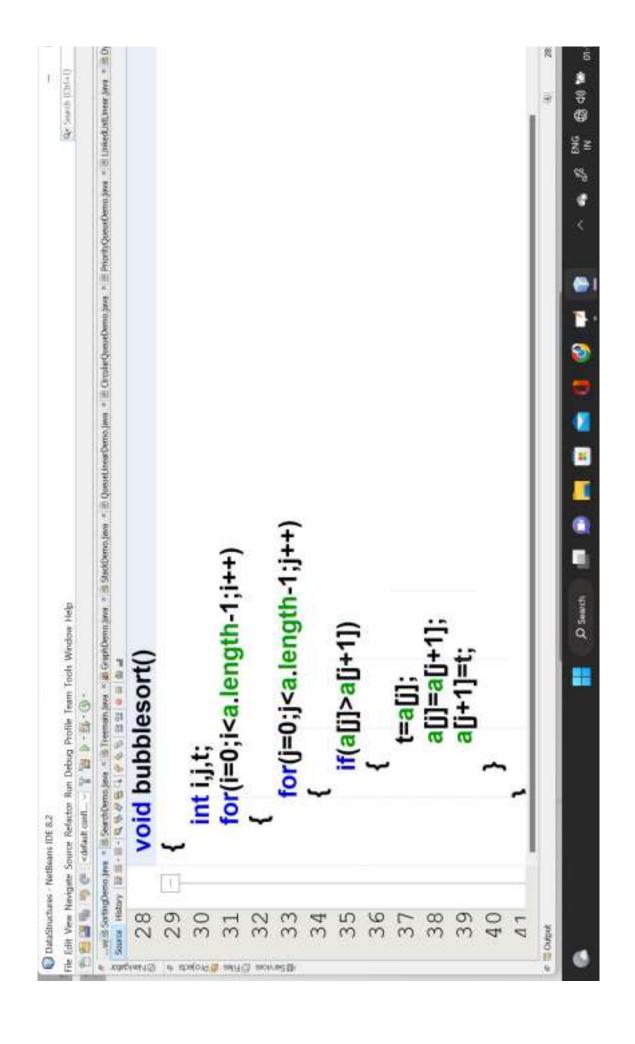
the upper and the lower bound of the running time of an algorithm, it is used for Theta notation encloses the function from above and below. Since it represents analyzing the average-case complexity of an algorithm.



Omega Notation (Ω-notation)

 Omega notation represents the lower bound of the running time of an algo Thus, it provides the best case complexity of an algorithm.





```
while (j>0 && a[j-1]>newelement)
                    { int i,j,newelement; for(i=0;i<a.length-1;i++)
                                                                                    newelement=a[i+1];
                                                                                                                                                                                                                                           a[j]=newelement;
void insertionsort()
                                                                                                                                                                         a[j]=a[j-1];
                                                                                                                                                                                                                                                                                                                               void quicksort(int start, int end)
                                                                                                           j=i+1;
                                                                                                                                                                                               <u>;</u>
```

https://docs.google.com/document/d/1Wyh7M9U_LmeJqSGCXQIAjJ5lpVCKaxkkPMSUu2hhst8/edit

```
int i,j.pivot;
i=start; j=end;pivot=start;
while(i<j)
{
    while(a[i]<a[pivot])
    i++;
    while(a[i]>a[pivot])
    i--;
    if(i<j)
    {
        int t=a[i];
        a[i]=a[i];
    int i,j.pivot;
    int i,j.pivot;
    int i,j.pivot;
    i=start; j=end;pivot=start;
    while(i<j)</pre>
```

while(a[i]<a[pivot])//i will work if pivot is at end

```
mergesort(mid+1,end);
merger(start,mid,end);
}

void merger(int start,int mid,int
end) {
    int temp[]=new
    int[a.length]; int i,j,index;
    i=index=start;
    j=mid+1;
    while(i<=mid && j<=end)
    {
        if(a[i]<a[j])
        temp[index++]=a[i++];
        else
        temp[index++]=a[i++];
    }
    while(i<=mid)
    temp[index++]=a[i++];
    while(j<=end)
    temp[index++]=a[i++];
    while(j<=end)
    temp[index++]=a[i++];
    while(j<=end)
    temp[index++]=a[i++];</pre>
```

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for (i=start;i<=end;i++) a[i]=temp[i];

data=index->a[data%10]

12/1/22, 5:08 PM