

```
1 String joinWords(String[] words) {
2     String sentence = "";
3     for (String w : words) {
4         sentence = sentence + w;
5     }
6     return sentence;
7 }
```

On each concatenation, a new copy of the string is created, and the two strings are copied over, character by character. The first iteration requires us to copy x characters. The second iteration requires copying $2x$ characters. The third iteration requires $3x$, and so on. The total time therefore is $O(x + 2x + \dots + nx)$. This reduces to $O(xn^2)$.

Why is it $O(xn^2)$? Because $1 + 2 + \dots + n$ equals $n(n+1)/2$, or $O(n^2)$.

`StringBuilder` can help you avoid this problem. `StringBuilder` simply creates a resizable array of all the strings, copying them back to a string only when necessary.

```
1 String joinWords(String[] words) {
2     StringBuilder sentence = new StringBuilder();
3     for (String w : words) {
4         sentence.append(w);
5     }
6     return sentence.toString();
7 }
```

A good exercise to practice strings, arrays, and general data structures is to implement your own version of `StringBuilder`, `HashTable` and `ArrayList`.

Additional Reading: Hash Table Collision Resolution (pg 636), Rabin-Karp Substring Search (pg 636).

Interview Questions

- 1.1 Is Unique:** Implement an algorithm to determine if a string has all unique characters. What if you cannot use additional data structures?

Hints: #44, #117, #132

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- 1.2 Check Permutation:** Given two strings, write a method to decide if one is a permutation of the other.

Hints: #1, #84, #122, #131

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- 1.3 URLify:** Write a method to replace all spaces in a string with '%20'. You may assume that the string has sufficient space at the end to hold the additional characters, and that you are given the "true" length of the string. (Note: If implementing in Java, please use a character array so that you can perform this operation in place.)

EXAMPLE

Input: "Mr John Smith ", 13

Output: "Mr%20John%20Smith"

Hints: #53, #118

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- 1.4 Palindrome Permutation:** Given a string, write a function to check if it is a permutation of a palindrome. A palindrome is a word or phrase that is the same forwards and backwards. A permutation is a rearrangement of letters. The palindrome does not need to be limited to just dictionary words.

EXAMPLE

Input: Tact Coa

Output: True (permutations: "taco cat", "atco cta", etc.)

Hints: #106, #121, #134, #136

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- 1.5 One Away:** There are three types of edits that can be performed on strings: insert a character, remove a character, or replace a character. Given two strings, write a function to check if they are one edit (or zero edits) away.

EXAMPLE

pale, ple -> true

pales, pale -> true

pale, bale -> true

pale, bake -> false

Hints: #23, #97, #130

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- 1.6 String Compression:** Implement a method to perform basic string compression using the counts of repeated characters. For example, the string aabcccccaaa would become a2b1c5a3. If the "compressed" string would not become smaller than the original string, your method should return the original string. You can assume the string has only uppercase and lowercase letters (a - z).

Hints: #92, #110

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- 1.7 Rotate Matrix:** Given an image represented by an NxN matrix, where each pixel in the image is 4 bytes, write a method to rotate the image by 90 degrees. Can you do this in place?

Hints: #51, #100

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- 1.8 Zero Matrix:** Write an algorithm such that if an element in an MxN matrix is 0, its entire row and column are set to 0.

Hints: #17, #74, #102

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- 1.9 String Rotation:** Assume you have a method isSubstring which checks if one word is a substring of another. Given two strings, s1 and s2, write code to check if s2 is a rotation of s1 using only one call to isSubstring (e.g., "waterbottle" is a rotation of "erbottlewat").

Hints: #34, #88, #104

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Additional Questions: Object-Oriented Design (#7.12), Recursion (#8.3), Sorting and Searching (#10.9), C++ (#12.11), Moderate Problems (#16.8, #16.17, #16.22), Hard Problems (#17.4, #17.7, #17.13, #17.22, #17.26).

Hints start on page 653.

2

Linked Lists

A linked list is a data structure that represents a sequence of nodes. In a singly linked list, each node points to the next node in the linked list. A doubly linked list gives each node pointers to both the next node and the previous node.

The following diagram depicts a doubly linked list:



Unlike an array, a linked list does not provide constant time access to a particular “index” within the list. This means that if you’d like to find the Kth element in the list, you will need to iterate through K elements.

The benefit of a linked list is that you can add and remove items from the beginning of the list in constant time. For specific applications, this can be useful.

► Creating a Linked List

The code below implements a very basic singly linked list.

```
1 class Node {
2     Node next = null;
3     int data;
4
5     public Node(int d) {
6         data = d;
7     }
8
9     void appendToTail(int d) {
10        Node end = new Node(d);
11        Node n = this;
12        while (n.next != null) {
13            n = n.next;
14        }
15        n.next = end;
16    }
17 }
```

In this implementation, we don’t have a `LinkedList` data structure. We access the linked list through a reference to the head `Node` of the linked list. When you implement the linked list this way, you need to be a bit careful. What if multiple objects need a reference to the linked list, and then the head of the linked list changes? Some objects might still be pointing to the old head.

We could, if we chose, implement a `LinkedList` class that wraps the `Node` class. This would essentially just have a single member variable: the head `Node`. This would largely resolve the earlier issue.

Remember that when you're discussing a linked list in an interview, you must understand whether it is a singly linked list or a doubly linked list.

► Deleting a Node from a Singly Linked List

Deleting a node from a linked list is fairly straightforward. Given a node `n`, we find the previous node `prev` and set `prev.next` equal to `n.next`. If the list is doubly linked, we must also update `n.next` to set `n.next.prev` equal to `n.prev`. The important things to remember are (1) to check for the null pointer and (2) to update the head or tail pointer as necessary.

Additionally, if you implement this code in C, C++ or another language that requires the developer to do memory management, you should consider if the removed node should be deallocated.

```

1  Node deleteNode(Node head, int d) {
2      Node n = head;
3
4      if (n.data == d) {
5          return head.next; /*moved head */
6      }
7
8      while (n.next != null) {
9          if (n.next.data == d) {
10             n.next = n.next.next;
11             return head; /*head didn't change */
12         }
13         n = n.next;
14     }
15     return head;
16 }
```

► The “Runner” Technique

The “runner” (or second pointer) technique is used in many linked list problems. The runner technique means that you iterate through the linked list with two pointers simultaneously, with one ahead of the other. The “fast” node might be ahead by a fixed amount, or it might be hopping multiple nodes for each one node that the “slow” node iterates through.

For example, suppose you had a linked list $a_1 \rightarrow a_2 \rightarrow \dots \rightarrow a_n \rightarrow b_1 \rightarrow b_2 \rightarrow \dots \rightarrow b_n$ and you wanted to rearrange it into $a_1 \rightarrow b_1 \rightarrow a_2 \rightarrow b_2 \rightarrow \dots \rightarrow a_n \rightarrow b_n$. You do not know the length of the linked list (but you do know that the length is an even number).

You could have one pointer `p1` (the fast pointer) move every two elements for every one move that `p2` makes. When `p1` hits the end of the linked list, `p2` will be at the midpoint. Then, move `p1` back to the front and begin “weaving” the elements. On each iteration, `p2` selects an element and inserts it after `p1`.

► Recursive Problems

A number of linked list problems rely on recursion. If you're having trouble solving a linked list problem, you should explore if a recursive approach will work. We won't go into depth on recursion here, since a later chapter is devoted to it.

However, you should remember that recursive algorithms take at least $O(n)$ space, where n is the depth of the recursive call. All recursive algorithms *can* be implemented iteratively, although they may be much more complex.

Interview Questions

2.1 **Remove Dups:** Write code to remove duplicates from an unsorted linked list.

FOLLOW UP

How would you solve this problem if a temporary buffer is not allowed?

Hints: #9, #40

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2.2 **Return Kth to Last:** Implement an algorithm to find the k th to last element of a singly linked list.

Hints: #8, #25, #41, #67, #126

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2.3 **Delete Middle Node:** Implement an algorithm to delete a node in the middle (i.e., any node but the first and last node, not necessarily the exact middle) of a singly linked list, given only access to that node.

EXAMPLE

Input: the node c from the linked list $a \rightarrow b \rightarrow c \rightarrow d \rightarrow e \rightarrow f$

Result: nothing is returned, but the new linked list looks like $a \rightarrow b \rightarrow d \rightarrow e \rightarrow f$

Hints: #72

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2.4 **Partition:** Write code to partition a linked list around a value x , such that all nodes less than x come before all nodes greater than or equal to x . If x is contained within the list, the values of x only need to be after the elements less than x (see below). The partition element x can appear anywhere in the "right partition"; it does not need to appear between the left and right partitions.

EXAMPLE

Input: 3 -> 5 -> 8 -> 5 -> 10 -> 2 -> 1 [partition = 5]

Output: 3 -> 1 -> 2 -> 10 -> 5 -> 5 -> 8

Hints: #3, #24

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- 2.5 Sum Lists:** You have two numbers represented by a linked list, where each node contains a single digit. The digits are stored in *reverse* order, such that the 1's digit is at the head of the list. Write a function that adds the two numbers and returns the sum as a linked list.

EXAMPLE

Input: (7 -> 1 -> 6) + (5 -> 9 -> 2). That is, 617 + 295.

Output: 2 -> 1 -> 9. That is, 912.

FOLLOW UP

Suppose the digits are stored in forward order. Repeat the above problem.

EXAMPLE

Input: (6 -> 1 -> 7) + (2 -> 9 -> 5). That is, 617 + 295.

Output: 9 -> 1 -> 2. That is, 912.

Hints: #7, #30, #71, #95, #109

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- 2.6 Palindrome:** Implement a function to check if a linked list is a palindrome.

Hints: #5, #13, #29, #61, #101

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- 2.7 Intersection:** Given two (singly) linked lists, determine if the two lists intersect. Return the intersecting node. Note that the intersection is defined based on reference, not value. That is, if the *k*th node of the first linked list is the exact same node (by reference) as the *j*th node of the second linked list, then they are intersecting.

Hints: #20, #45, #55, #65, #76, #93, #111, #120, #129

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- 2.8 Loop Detection:** Given a circular linked list, implement an algorithm that returns the node at the beginning of the loop.

DEFINITION

Circular linked list: A (corrupt) linked list in which a node's next pointer points to an earlier node, so as to make a loop in the linked list.

EXAMPLE

Input: A -> B -> C -> D -> E -> C [the same C as earlier]

Output: C

Hints: #50, #69, #83, #90

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Additional Questions: Trees and Graphs (#4.3), Object-Oriented Design (#7.12), System Design and Scalability (#9.5), Moderate Problems (#16.25), Hard Problems (#17.12).

Hints start on page 653.

3

Stacks and Queues

Questions on stacks and queues will be much easier to handle if you are comfortable with the ins and outs of the data structure. The problems can be quite tricky, though. While some problems may be slight modifications on the original data structure, others have much more complex challenges.

► Implementing a Stack

The stack data structure is precisely what it sounds like: a stack of data. In certain types of problems, it can be favorable to store data in a stack rather than in an array.

A stack uses LIFO (last-in first-out) ordering. That is, as in a stack of dinner plates, the most recent item added to the stack is the first item to be removed.

It uses the following operations:

- `pop()`: Remove the top item from the stack.
- `push(item)`: Add an item to the top of the stack.
- `peek()`: Return the top of the stack.
- `isEmpty()`: Return true if and only if the stack is empty.

Unlike an array, a stack does not offer constant-time access to the *i*th item. However, it does allow constant-time adds and removes, as it doesn't require shifting elements around.

We have provided simple sample code to implement a stack. Note that a stack can also be implemented using a linked list, if items were added and removed from the same side.

```
1 public class MyStack<T> {
2     private static class StackNode<T> {
3         private T data;
4         private StackNode<T> next;
5
6         public StackNode(T data) {
7             this.data = data;
8         }
9     }
10
11     private StackNode<T> top;
12
13     public T pop() {
14         if (top == null) throw new EmptyStackException();
15         T item = top.data;
```

```

16     top = top.next;
17     return item;
18 }
19
20 public void push(T item) {
21     StackNode<T> t = new StackNode<T>(item);
22     t.next = top;
23     top = t;
24 }
25
26 public T peek() {
27     if (top == null) throw new EmptyStackException();
28     return top.data;
29 }
30
31 public boolean isEmpty() {
32     return top == null;
33 }
34 }

```

One case where stacks are often useful is in certain recursive algorithms. Sometimes you need to push temporary data onto a stack as you recurse, but then remove them as you backtrack (for example, because the recursive check failed). A stack offers an intuitive way to do this.

A stack can also be used to implement a recursive algorithm iteratively. (This is a good exercise! Take a simple recursive algorithm and implement it iteratively.)

► Implementing a Queue

A queue implements FIFO (first-in first-out) ordering. As in a line or queue at a ticket stand, items are removed from the data structure in the same order that they are added.

It uses the operations:

- `add(item)`: Add an item to the end of the list.
- `remove()`: Remove the first item in the list.
- `peek()`: Return the top of the queue.
- `isEmpty()`: Return true if and only if the queue is empty.

A queue can also be implemented with a linked list. In fact, they are essentially the same thing, as long as items are added and removed from opposite sides.

```

1  public class MyQueue<T> {
2      private static class QueueNode<T> {
3          private T data;
4          private QueueNode<T> next;
5
6          public QueueNode(T data) {
7              this.data = data;
8          }
9      }
10
11     private QueueNode<T> first;
12     private QueueNode<T> last;
13
14     public void add(T item) {

```



```
15     QueueNode<T> t = new QueueNode<T>(item);
16     if (last != null) {
17         last.next = t;
18     }
19     last = t;
20     if (first == null) {
21         first = last;
22     }
23 }
24
25 public T remove() {
26     if (first == null) throw new NoSuchElementException();
27     T data = first.data;
28     first = first.next;
29     if (first == null) {
30         last = null;
31     }
32     return data;
33 }
34
35 public T peek() {
36     if (first == null) throw new NoSuchElementException();
37     return first.data;
38 }
39
40 public boolean isEmpty() {
41     return first == null;
42 }
43 }
```

It is especially easy to mess up the updating of the first and last nodes in a queue. Be sure to double check this.

One place where queues are often used is in breadth-first search or in implementing a cache.

In breadth-first search, for example, we used a queue to store a list of the nodes that we need to process. Each time we process a node, we add its adjacent nodes to the back of the queue. This allows us to process nodes in the order in which they are viewed.

Interview Questions

3.1 Three in One: Describe how you could use a single array to implement three stacks.

Hints: #2, #12, #38, #58

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3.2 Stack Min: How would you design a stack which, in addition to push and pop, has a function `min` which returns the minimum element? Push, pop and `min` should all operate in $O(1)$ time.

Hints: #27, #59, #78

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- 3.3 Stack of Plates:** Imagine a (literal) stack of plates. If the stack gets too high, it might topple. Therefore, in real life, we would likely start a new stack when the previous stack exceeds some *threshold*. Implement a data structure *SetOfStacks* that mimics this. *SetOfStacks* should be composed of several stacks and should create a new stack once the previous one exceeds capacity. *SetOfStacks.push()* and *SetOfStacks.pop()* should behave identically to a single stack (that is, *pop()* should return the same values as it would if there were just a single stack).

FOLLOW UP

Implement a function *popAt (int index)* which performs a pop operation on a specific sub-stack.

Hints: #64, #81

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- 3.4 Queue via Stacks:** Implement a *MyQueue* class which implements a queue using two stacks.

Hints: #98, #114

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- 3.5 Sort Stack:** Write a program to sort a stack such that the smallest items are on the top. You can use an additional temporary stack, but you may not copy the elements into any other data structure (such as an array). The stack supports the following operations: *push*, *pop*, *peek*, and *isEmpty*.

Hints: #15, #32, #43

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- 3.6 Animal Shelter:** An animal shelter, which holds only dogs and cats, operates on a strictly “first in, first out” basis. People must adopt either the “oldest” (based on arrival time) of all animals at the shelter, or they can select whether they would prefer a dog or a cat (and will receive the oldest animal of that type). They cannot select which specific animal they would like. Create the data structures to maintain this system and implement operations such as *enqueue*, *dequeueAny*, *dequeueDog*, and *dequeueCat*. You may use the built-in *LinkedList* data structure.

Hints: #22, #56, #63

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Additional Questions: Linked Lists (#2.6), Moderate Problems (#16.26), Hard Problems (#17.9).

Hints start on page 653.