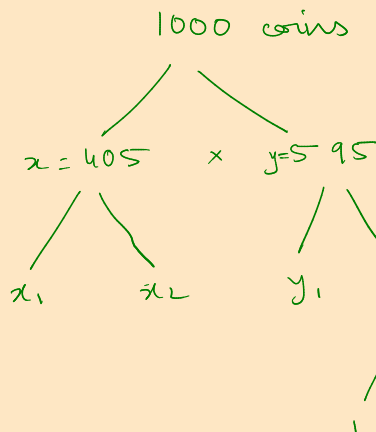
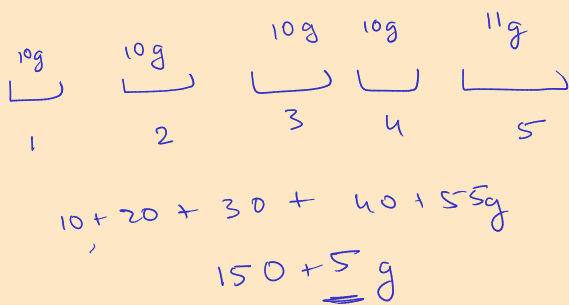


55 coins

550 g if all correct

$550 \pm i$



$$x \cdot y$$

$$+ x_1 x_2$$

$$+ y_1 y_2$$

What's the final sum?

Hint:

Start with small cases.

$$\text{split_sum}(5)$$

$$= (4 \times 1) + \text{ss}(4) + \text{ss}(1)$$

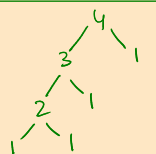
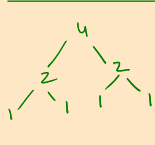
$$= (2 \times 3) + \text{ss}(2) + \text{ss}(3)$$



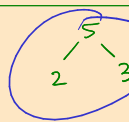
$$1 \times 1 = 1$$



$$(2 \times 1) + (1 \times 1) = 3$$



$$\begin{cases} (2 \times 2) + (1 \times 1) + (1 \times 1) = 6 \\ (3 \times 1) + (2 \times 1) + (1 \times 1) = 6 \end{cases}$$



$$\begin{cases} (2 \times 3) + 1 + 3 = ? \quad 10 \\ (4 \times 1) + \dots = ? \end{cases}$$

Notice a pattern? Can you generalize this?

1 3 6 10

$$\frac{n(n-1)}{2}$$

Base: $2: \frac{2(1)}{2} = 1$

IH: $\text{ss}(k) = \frac{k(k-1)}{2}$ for all $n \geq k \geq 2$.

IS: $\text{ss}(n+1) = \frac{(n+1)n}{2}$

Split into $x, n+1-x$.

$$\text{ss}(n+1) = \underbrace{(x)(n+1-x)} + \frac{x(x-1)}{2} + \frac{(n+1-x)(n-x)}{2}$$

$$= \frac{(n+1)n}{2}$$

$$\text{ss}(1000) = \frac{1000 \times 999}{2}$$

1 True or False ≥ 1 ≤ 1 connected \wedge no cycles

(a) Any pair of vertices in a tree are connected by exactly one path.

True: Connected $\Rightarrow \geq 1$ path.

Assume there are 2 paths between u and v .

Let w_1 be the branching pt. and w_2 be the merging pt.

Cycle involving w_1, w_2 . Contradiction.

(b) A simple graph obtained by adding an edge between two vertices of a tree creates a cycle.

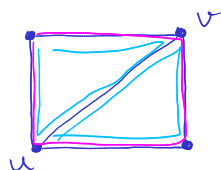
True: Let T be a tree. Pick vertices u and v . By part (a) we know

there is 1 path between u and v , $\{(u, v_1), (v_1, v_2), \dots, (v_k, v)\}$. Now we add an edge (u, v) .

Cycle: $\{(u, v), (v, v_k), \dots, (v_2, v_1), (v_1, u), (u, v)\}$.

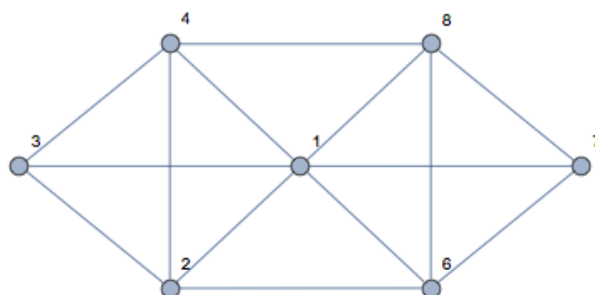
(c) Adding an edge in a connected graph creates exactly one new cycle.

False:



2 new cycles

2 Eulerian Tour and Eulerian Walk



- (a) Is there an Eulerian tour in the graph above? If no, give justification. If yes, provide an example.
- (b) Is there an Eulerian walk in the graph above? An Eulerian walk is a walk that uses each edge exactly once. If no, give justification. If yes, provide an example.
- (c) What is the condition that there is an Eulerian walk in an undirected graph? Briefly justify your answer.

3 Not everything is normal: Odd-Degree Vertices

Claim: Let $G = (V, E)$ be an undirected graph. The number of vertices of G that have odd degree is even.

Prove the claim above using:

(i) Direct proof (e.g., counting the number of edges in G). *Hint: in lecture, we proved that $\sum_{v \in V} \deg v = 2|E|$. Hint: sum of an even # of odds is even.*

(ii) Induction on $m = |E|$ (number of edges)

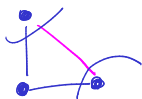
(iii) Induction on $n = |V|$ (number of vertices)

$$(i) \sum_{v \in V} \deg(v) = 2m + \sum_{v \in V_{\text{odd}}(G)} \deg(v) \quad V_{\text{odd}}(G) = \{\text{set of odd-deg vertices in } G\}$$

(ii) Base: $|E| = 0$. All vertices have $\deg 0$. 0 is even, 0 vertices w/ odd deg.

IH: $|E| = k$, even # of vertices w/ odd deg.

IS: $k+1$ $+2, 0, -2$



(iii) Base: $|V| = 1$. 0 edges. 0 vertices w/ odd deg. 0 even ✓

IH: $|V| = k$, holds.

IS:



of vertices w/ odd deg

of vertices w/ even

4 Coloring Trees

Prove that all trees with at least 2 vertices are *bipartite*: the vertices can be partitioned into two groups so that every edge goes between the two groups.

[*Hint*: Use induction on the number of vertices.]