# CP Template

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#### 1 Custom Codes

```
freopen("input.in", "r", stdin);
freopen("output.out", "w", stdout);
ios_base::sync_with_stdio(0), cin.tie(0);
// pragma
#pragma GCC optimize("Ofast")
#pragma GCC target("avx,avx2,fma")
#pragma GCC optimize ("03,unroll-loops")
#pragma GCC optimize ("-ffloat-store")
#pragma GCC target ("bmi,bmi2,lzcnt,popcnt")
#pragma GCC target
// random number generator
mt19937_64

→ rng(chrono::steady_clock::now().time_since_epoch().
int num = rng() % limit;
int64_t rnd(int64_t l, int64_t r) { // can handle
\hookrightarrow negative range also
 if (1 > r) swap(1, r):
  return 1 + (rng() \% (r - 1 + 1));
// ignore white space
cin.ignore(); or cin >> ws;
// STL merge function: merge two vectors or sets
vector<int> a = \{1, 2, 7\}, b = \{4, 2, 3\};
sort(a.begin(), a.end());
sort(b.begin(), b.end());
vector<int> c(a.size() + b.size());
merge(a.begin(), a.end(), b.begin(), b.end(),

    c.begin());
set < int > a = \{3, 2, 1\}, b = \{4, 5, 1\};
set<int> c:
merge(a.begin(), a.end(), b.begin(), b.end(),

    inserter(c, c.begin()));

// tuple
tuple\langle int, int, int \rangle tup = \{1, 2, 3\};
int first = get<0>(tup);
// custom compare function
struct item {
 int a, b;
bool cmp(item& a, item& b) {
  if (a.a != b.a) return a.a < b.a;
 return a.b > b.b:
bool custom(pair<int, int>& a, pair<int, int>& b) {
 if(a.first != b.first) return a.first > b.first;
  return a.second < b.second;</pre>
// custom compare in set
struct cmp{
  bool operator() (const pair<int, int>& a, const

→ pair<int, int>& b) const {
   if (a.first != b.first) return a.first > b.first;
    return a.second < b.second:
};
set<pair<int, int>, cmp> a;
```

```
// greater functions
 priority_queue<int, vector<int>, greater<int>>
 set<int, greater<int>>
 map<int, int, greater<int>>
 // fill with elements
 fill(v.begin(), v.end(), -1);
 // fill for array
 fill(\&dp[0][0][0], \&dp[0][0][0] + N * M * N, INF);
 // string to number
 stoi(num):
 // number to string
 to_string(num);
 // Binary to decimal
 int n = stoi(binary, 0, 2);
 // decimal to binary
continuity) binary = bitset<64>(n).to_string();
binary.erase(0, binary.find_first_not_of('0'));
 // count digit
 int d = log10(num) + 1;
 // set ith bit
 n \mid (1 << i)
 // unset ith bit
 n \& ~(1 << i)
 // toggle or inverse ith bit
 n^{(1)} < i
 __builtin_popcountll(x)
  __builtin_clzll(x)
  __builtin_ctzll(x)
```

# 2 Articulation Bridge

```
vector<int> adj[N];
int t = 0;
vector<int> tin(N, -1), low(N);
vector<array<int, 2>> ab;
void dfs (int u, int p) {
  tin[u] = low[u] = t++;
  for (int v: adj[u]) {
    if (v != p) {
     if (tin[v] != -1) {
        low[u] = min(low[u], tin[v]);
    } else {
        dfs(v, u);
        if (tin[u] < low[v]) {
            ab.push_back({u, v});
        }
        low[u] = min(low[u], low[v]);
    }
}</pre>
```

# 3 Articulation Point

```
vector<int> adj[N];
int t = 0;
vector<int> tin(N, -1), low(N), ap;
void dfs (int u, int p) {
  tin[u] = low[u] = t++;
  int is_ap = 0, child = 0;
```

```
for (int v: adj[u]) {
   if (v != p) {
      if (tin[v] != -1) {
        low[u] = min(low[u], tin[v]);
      } else {
        child++;
      dfs(v, u);
      if (tin[u] <= low[v]) {
            is_ap = 1;
      }
      low[u] = min(low[u], low[v]);
    }
  }
}
if ((u != p or child > 1) and is_ap) {
   ap.push_back(u);
}
```

#### 4 Bellman Ford

```
struct st {
 int a, b, cost;
} e[N];
const int INF = 2e9;
int32_t main() {
 int n, m;
  cin >> n >> m:
  for(int i = 0; i < m; i++) cin >> e[i].a >> e[i].b
  cin >> s://is there any negative cycle which is

→ reachable from s?

  vector<int> d (n, INF);//for finding any cycle(not
  → necessarily from s) set d[i] = 0 for all i
  d[s] = 0:
  vector\langle int \rangle p (n, -1);
  int x;
  for (int i=0; i<n; ++i) {
   x = -1;
   for (int j=0; j<m; ++j) {
      if (d[e[j].a] < INF) {
       if (d[e[j].b] > d[e[j].a] + e[j].cost) {
          d[e[j].b] = max (-INF, d[e[j].a] +

→ e[j].cost);//for overflow
          p[e[j].b] = e[j].a;
          x = e[j].b;
  if (x == -1) cout << "No negative cycle from "<<s;
  else {
    int y = x; //x can be on any cycle or reachable

→ from some cycle

   for (int i=0; i < n; ++i) y = p[y];
    vector<int> path;
   for (int cur=y; ; cur=p[cur]) {
      path.push_back (cur);
      if (cur == y && path.size() > 1) break;
```

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```
reverse (path.begin(), path.end());
   cout << "Negative cycle: "</pre>
   for (int i=0; i<path.size(); ++i) cout << path[i]</pre>
    return 0;
// ## Bellman-ford (V * E)
vector<int> bellman ford(int s){
 vector<int> dis(n, I);
 dis[s]=0;
 while(1){
   int any=0;
   for (auto& e: ed){
     if(dis[e.u]<I){
       if(dis[e.u]+e.cost < dis[e.v]){</pre>
          dis[e.v] = dis[e.u]+e.cost;
          any=1;
     }
   if(!any) break;
 return dis;
```

# 5 Binary & Ternary Search

```
//Binary Search (integer)
int lo = 0, hi = n - 1;
while (lo <= hi) {</pre>
 int mid = (lo + hi) / 2:
 if (f(mid)) hi = mid - 1;
 else lo = mid + 1:
//Binary Search (double)
double lo = 0, hi = inf;
int itr = 50;
while (itr--) {
 double mid = (lo + hi) / 2;
 if (f(mid)) hi = mid;
 else lo = mid;
//Ternary Search (integer, max)
int lo = 0, hi = n - 1;
while (lo < hi) {
 // int mid1 = lo + (hi - lo) / 3;
 // int mid2 = hi - (hi - lo) / 3;
 int mid = (lo + hi) / 2;
 if (f(mid) < f(mid + 1)) lo = mid + 1;
 else hi = mid;
//Ternary Search (double, max)
double lo = 0, hi = inf:
int itr = 50;
while (itr--) {
 double mid1 = lo + (hi - lo) / 3;
 double mid2 = hi - (hi - lo) / 3;
 if (f(mid1) < f(mid2)) lo = mid1;
```

```
else hi = mid2:
6 2D BIT Range Update & Range Query
const int N = 1009;
struct BIT2D {
 long long M[N][N][2], A[N][N][2];
 BIT2D() {
   memset(M, 0, sizeof M);
   memset(A, 0, sizeof A);
 void upd2(long long t[N][N][2], int x, int y, long
  → long mul, long long add) {
   for(int i = x; i < N; i += i & -i) {
     for(int j = y; j < N; j += j \& -j) {
       t[i][j][0] += mul;
       t[i][j][1] += add;
 void upd1(int x, int y1, int y2, long long mul, long
  → long add) {
   upd2(M, x, y1, mul, -mul * (y1 - 1));
   upd2(M, x, y2, -mul, mul * y2);
   upd2(A, x, y1, add, -add * (y1 - 1));
   upd2(A, x, y2, -add, add * y2);
 void upd(int x1, int y1, int x2, int y2, long long
  → val) {
   upd1(x1, y1, y2, val, -val * (x1 - 1));
   upd1(x2, y1, y2, -val, val * x2);
 long long query2(long long t[N][N][2], int x, int y)
   long long mul = 0, add = 0;
   for(int \bar{i} = y; i > 0; i -= i \& -i) {
     mul += t[x][i][0];
     add += t[x][i][1];
   return mul * y + add;
 long long query1(int x, int y) {
   long long mul = 0, add = 0;
   for(int i = x; i > 0; i = i & -i) {
     mul += query2(M, i, y);
     add += query2(A, i, y);
   return mul * x + add;
 long long query(int x1, int y1, int x2, int y2) {
   return query1(x2, y2) - query1(x1 - 1, y2) -
      query1(x2, y1 - 1) + query1(x1 - 1, y1 - 1);
};
```

# 7 Centroid Decomposition

```
struct CentroidDecomposition {
```

```
using T = vector<vector<int>>:
 int n:
 vector<int> sz, is_cen, cpar, cdep;
 CentroidDecomposition(T& adj, int root = 1) {
   n = (int)adj.size() + 1;
    sz.resize(n), is_cen.resize(n), cpar.resize(n),

    cdep.resize(n);

    Decompose(root, -1, 0, adj);
 void Cal_sz(int u, int p, T& adj) {
    sz[u] = 1:
   for (auto& v : adj[u]) {
      if (v != p && !is_cen[v]) {
       Cal_sz(v, u, adj);
       sz[u] += sz[v];
 int Get_cen(int u, int p, int csz, T& adj) {
   for (auto& v : adj[u]) {
      if (v != p \&\& !is_cen[v] \&\& (2 * sz[v] > csz)) {
       return Get_cen(v, u, csz, adj);
   }
   return u;
 void Decompose(int u, int p, int d, T& adj) {
   Cal_sz(u, p, adj);
   int c = Get_cen(u, p, sz[u], adj);
    is_cen[c] = 1, cpar[c] = p, cdep[c] = d;
   for (auto& v : adj[c]) {
      if (!is_cen[v]) {
       Decompose(v, c, d + 1, adj);
};
```

#### 8 Closest Index Where Each Element in Distinct

```
// O based, closest index where each element is

→ distinct

array<vector<int>, 2>
int n = a.size(), mxIdx = -1;
 map<int, int> mp;
 vector<int> l(n);
 for (int i = 0; i < n; i++) {
   if (mp.count(a[i])) mxIdx = max(mxIdx, mp[a[i]]);
   mp[a[i]] = i \cdot l[i] = i - mxIdx:
 mp.clear(), mxIdx = n;
 vector<int> r(n);
 for (int i = n - 1; i >= 0; i--) {
   if (mp.count(a[i])) mxIdx = min(mxIdx, mp[a[i]]);
   mp[a[i]] = i, r[i] = mxIdx - i;
 return {1, r};
```

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#### 9 Closest Min Max

```
// closest left, right index where current element is
// closest max or min element from the current element
array<vector<int>, 2> closest_min_element(vector<int>
 int n = a.size();
 vector<int> l(n), r(n);
 stack<int> st;
 st.push(-1);
 for (int i = 0; i < n; i++) {
    while (st.top() != -1 \&\& a[st.top()] > a[i])

    st.pop();
   l[i] = st.top() + 1; // closest index where
    // l[i] = st.top(); // closest min element index
   st.push(i);
 while (!st.empty()) st.pop();
 st.push(n);
 for (int i = n - 1; i \ge 0; i - -) {
    while (st.top() != n \&\& a[st.top()] >= a[i])
    \rightarrow st.pop();
   r[i] = st.top() - 1; // closest index where
    // r[i] = st.top(); // closest min element index
    st.push(i);
 return {1, r};
array<vector<int>, 2> closest_max_element(vector<int>
→ &a) {
 int n = a.size():
 vector<int> l(n), r(n);
 stack<int> st;
 st.push(-1);
 for (int i = 0; i < n; i++) {
    while (st.top() != -1 \&\& a[st.top()] < a[i])
    \hookrightarrow st.pop();
   l[i] = st.top() + 1;
   // l[i] = st.top(); // closest max element index
    st.push(i);
 while (!st.empty()) st.pop();
 st.push(n);
 for (int i = n - 1; i \ge 0; i - -) {
    while (st.top() != n && a[st.top()] <= a[i])</pre>
    \rightarrow st.pop();
   r[i] = st.top() - 1;
   // r[i] = st.top(); // closest max element index
    st.push(i);
 return {1, r};
```

# 10 Compress Array

```
void CompressArray(vector<int>& a) {
```

#### 11 Convex Hull

```
const double PI = acos((double)-1.0);
using ll = long long;
struct PT {
 11 x, y;
  bool operator < (const PT &p) const {
    return x == p.x ? y < p.y : x < p.x;
};
11 area(PT a, PT b, PT c) {
  return (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
  \rightarrow (c.x - a.x);
vector<PT> ConvexHull(vector<PT> p) {
  int n = p.size(), m = 0;
  if (n < 3) return p;
  vector<PT> hull(n + n);
  sort(p.begin(), p.end());
  for (int i = 0; i < n; ++i) {
    while (m > 1 \text{ and area}(hull[m - 2], hull[m - 1],
       p[i]) <= 0) --m;
    hull[m++] = p[i];
  for (int i = n - 2, j = m + 1; i \ge 0; --i) {
    while (m >= j and area(hull[m - 2], hull[m - 1],
    \rightarrow p[i]) <= 0) --m;
    hull[m++] = p[i];
  hull.resize(m - 1);
  return hull;
```

#### 12 DFS with LCA

```
void dfs(int u, int p, vector<vector<int>>& adj) {
    tin[u] = ++t;
    subtree_size[u] = 1;
    parent[u][0] = p;
    for (int i = 1; i < k; i++) {
      if (parent[u][i - 1] != -1) parent[u][i] =

→ parent[parent[u][i - 1]][i - 1];

      else parent[u][i] = -1;
    for (auto& v : adj[u]) {
      if (v != p) {
        depth[v] = depth[u] + 1;
        is_{leaf}[u] = 0;
        dfs(v, u, adj);
        height[u] = max(height[u], height[v] + 1);
        subtree_size[u] += subtree_size[v];
        if (heavy[u] == -1 || subtree_size[heavy[u]] <</pre>

    subtree_size[v]) {

          heavy[u] = v;
      }
    }
    tout[u] = ++t;
  bool is_ancestor(int u, int v) {
    return tin[u] <= tin[v] && tout[v] <= tout[u];
  int kth_parent(int u, int kth) {
    for (int i = k - 1; i \ge 0; i - -) {
      if (kth & (1 << i)) {
        u = parent[u][i];
        if (u == -1) return u;
    return u:
  int lca(int u, int v) {
    if (is_ancestor(u, v)) return u;
    for (int i = k - 1; i \ge 0; i--) {
      if (!is_ancestor(parent[u][i], v)) {
        u = parent[u][i];
    return parent[u][0];
  int dis(int u, int v)
    return depth[u] + depth[v] - 2 * depth[lca(u, v)];
};
13 Dijkstra
```

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# 14 Digit DP

```
#2 countWithExactDigitCount // f(0, 0, 1, 0)
int dp[11][11][2][2];
int targetDigit, targetCnt;
int f(int i, int digitCnt, int tight, int started) {
 if (i == n) {
   return targetCnt == digitCnt;
 auto& ret = dp[i][digitCnt][tight][started];
 if ("ret) return ret;
 ret = 0;
 int limit = tight ? (s[i] - '0') : 9;
 for (int d = 0; d <= limit; d++) {
   int newTight = tight && (d == limit);
   int newStarted = started || (d > 0);
   int newDigitCnt = digitCnt;
   if (newStarted) newDigitCnt += (d == targetDigit);
   ret += f(i + 1, newDigitCnt, newTight,
    → newStarted);
 return ret;
#6 countNumbersWithOnly1234 // f(0, 1, 0)
int dp[11][2][2];
int f(int i, int tight, int started) {
 if (i == n) {
   return started;
 auto& ret = dp[i][tight][started];
 if ("ret) return ret;
 ret = 0:
 int limit = tight ? (s[i] - '0') : 9;
 for (int d = 0: d <= limit: d++) {
   int newTight = tight && (d == limit);
   int newStarted = started || (d > 0);
   if (d > 4) continue;
   if (!newStarted && d == 0) ret += f(i + 1,

→ newTight, newStarted);
    else if (d >= 1 \&\& d <= 4) ret += f(i + 1,
      newTight. newStarted):
```

```
return ret:
#8 countDistinctDigitNumbers // f(0, 0, 1, 0)
int dp[11][1025][2][2];
int f(int i, int mask, int tight, int started) {
 if (i == n) return started;
  auto& ret = dp[i][mask][tight][started];
 if ("ret) return ret:
 ret = 0;
  int limit = tight ? (s[i] - '0') : 9;
 for (int d = 0; d <= limit; d++) {
   int newTight = tight && (d == limit);
    int newStarted = started || (d > 0);
    int newMask = mask:
    if (newStarted) {
     if (mask & (1 << d)) continue:
     newMask \mid = (1 \ll d);
   ret += f(i + 1, newMask, newTight, newStarted);
 return ret:
#9 countWithAnyRepeatedDigits -> n -

    f(countDistinctDigitNumbers)

#10 countNumbersWithKDistinctDigits // f(0, 0, 1, 0)
int k, dp[11][1025][2][2];
int f(int i, int mask, int tight, int started) {
 if (i == n) return started &&
  auto& ret = dp[i][mask][tight][started];
  if ("ret) return ret;
 ret = 0;
  int limit = tight ? (s[i] - '0') : 9:
 for (int d = 0; d <= limit; d++) {
   int newTight = tight && (d == limit);
    int newStarted = started || (d > 0);
    int newMask = mask;
    if (newStarted) newMask |= (1 << d);</pre>
    ret += f(i + 1, newMask, newTight, newStarted);
 return ret;
#11 count Numbers Divisible By k and Contains No Digit
\rightarrow m // f(0, 0, 1, 0)
int dp[11][100][2][2];
int k, m;
int f(int i, int rem, int tight, int started) {
 if (i == n) return started && (rem == 0);
  auto& ret = dp[i][rem][tight][started];
 if (~ret) return ret;
 ret = 0;
 int limit = tight ? (s[i] - '0') : 9;
 for (int d = 0; d <= limit; d++) {
   if (d == m) continue:
   int newTight = tight && (d == limit);
    int newStarted = started || (d > 0);
   int newRem = (rem * 10 + d) \% k;
   ret += f(i + 1, newRem, newTight, newStarted);
 return ret;
```

```
--> find sum of integers and conunt at same time
\#12 sum and count of all numbers x that has at most k
\rightarrow distinct digits // f(0, 0, 1, 0)
const int mod = 998244353:
int64_t n, k, Pow10[20];
array<int64_t, 2> dp[20][1025][2][2]; // dp[0] -> cnt,
\rightarrow dp[1] -> sum;
string s;
array<int64_t, 2> f(int i, int mask, int tight, int

    started) {

 if (i == n) {
   return {started && (__builtin_popcount(mask) <=
    \rightarrow k), 0};
  auto& ret = dp[i][mask][tight][started];
  if ("ret[0]) return ret;
  ret = \{0, 0\};
  int limit = tight ? (s[i] - '0') : 9;
  for (int d = 0; d <= limit; d++) {
    int newTight = tight && (d == limit);
    int newStarted = started || (d > 0);
    int newMask = mask;
    if (newStarted) newMask |= (1 << d):
    auto [currCnt, currSum] = f(i + 1, newMask,
    → newTight, newStarted);
    auto& [cnt, sum] = ret;
    cnt = (cnt + currCnt) % mod;
    sum = (sum + (currCnt * d \% mod) * Pow10[n - i -
    \rightarrow 1] % mod) % mod;
    sum = (sum + currSum) % mod;
 return ret;
int64_t Cnt(int64_t num) {
 if (num <= 0) return 0;
 memset(dp, -1, sizeof dp);
  s = to string(num):
  n = s.size();
  Pow10[0] = 1;
  for (int i = 1; i < 20; i++) {
   Pow10[i] = (Pow10[i - 1] * 10) \% mod;
 return f(0, 0, 1, 0)[1];
--> find digit sum upto n, this is optimized version,

    use memset just once

string s;
int64_t dp[20][2][2];
int64_t cnt[20][2][2];
pair<int64_t, int64_t> f(int i, int tight, int

    started) {
 if (i < 0) return {0, started ? 1 : 0}; // {sum,
  if (~dp[i][tight][started] && !tight) return
  int64_t totalSum = 0, totalCount = 0;
```

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```
int limit = tight ? (s[i] - '0') : 9:
 for (int d = 0; d <= limit; d++) {
    int newTight = tight && (d == limit);
    int newStarted = started || (d > 0);
    auto [nextSum, nextCount] = f(i - 1, newTight,
    → newStarted);
    if (newStarted) {
      totalSum += d * nextCount + nextSum;
      totalCount += nextCount:
   } else {
      totalSum += nextSum:
      totalCount += nextCount;
 dp[i][tight][started] = totalSum;
 cnt[i][tight][started] = totalCount;
 return {totalSum, totalCount};
int64_t sumOfAllDigits(int64_t num) {
 if (num < 0) return 0:
 s = to_string(num);
 reverse(s.begin(), s.end());
 return f(s.size() - 1, 1, 0).first;
memset(dp, -1, sizeof dp);
memset(cnt, -1, sizeof cnt);
// Optimized versions -->
## with only 1 memset, reverse the number then
\rightarrow calculate from (n - 1) \rightarrow
## so when some state is already calculated and tight
\rightarrow == 0, then return 'ret', or else again calculate
#13 count of 3 = 6 = 9, and count of 3 > 0
const int mod = 1e9 + 7, N = 17;
int dp[51][N][N][N][2]; // f(s.size() - 1, 0, 0, 0, 1)
int f(int i, int three, int six, int nine, int tight)
 if (three > 16 | | six > 16 | | nine > 16) return 0:
 if (i < 0) {
   return (three > 0) && (three == six) && (six ==
    → nine):
 auto& ret = dp[i][three][six][nine][tight];
 if (ret != -1 && !tight) return ret;
 ret = 0;
 int limit = tight ? (s[i] - '0') : 9;
 for (int d = 0; d <= limit; d++) {
   int newTight = tight && (d == limit);
   ret = (ret + f(i - 1), three + (d == 3), six + (d + 1)
    \rightarrow == 6), nine + (d == 9), newTight)) % mod;
 }
 return ret;
// optimized, without tight state
// when the state is not tight then return 'ret' also

→ memoize the dp;

int dp[51][N][N][N];
```

```
int f(int i, int three, int six, int nine, int tight)
if (three > 16 | | six > 16 | | nine > 16) return 0;
 if (i < 0) {
    return (three > 0) && (three == six) && (six ==
    \rightarrow nine);
  auto ret = dp[i][three][six][nine];
  if (ret != -1 && !tight) return ret;
  ret = 0:
  int limit = tight ? (s[i] - '0') : 9;
  for (int d = 0; d <= limit; d++) {
    int newTight = tight && (d == limit);
    ret = (\text{ret} + f(i - 1, \text{ three} + (d == 3), \text{ six} + (d
    \rightarrow == 6), nine + (d == 9), newTight)) % mod;
 if (!tight) dp[i][three][six][nine] = ret;
 return ret;
int Cnt(string a) {
 reverse(a.begin(), a.end());
 return f(s.size() - 1, 0, 0, 0, 1);
int check(string& s) {
 // problem condition
void solve() {
 string a, b;
  cin \gg a \gg b;
  cout << Cnt(b) - Cnt(a) + check(a); // check if only</pre>
  \hookrightarrow string a can satisfy the condition
// calculate between 1 and r in one function
int dp[51][2][2][18][18][18];
string 1, r;
int f(int i, int tightLower, int tightUpper, int c3,
\rightarrow int c6, int c9) {
  if (c3 >= 17 \mid | c6 >= 17 \mid | c9 >= 17) return 0;
  if (i < 0) {
    return c3 && c3 == c6 && c6 == c9;
  int& ret =
  → dp[i][tightLower][tightUpper][c3][c6][c9]:
  if (~ret && !tightLower && !tightUpper) return ret;
  int lo = tightLower ? 1[i] - '0' : 0;
  int hi = tightUpper ? r[i] - '0' : 9;
  ret = 0:
  for (int d = lo; d <= hi; d++) {
    int newTightLower = tightLower && (d == lo);
    int newTightUpper = tightUpper && (d == hi);
    ret = (ret + f(i - 1, newTightLower,
    \rightarrow newTightUpper, c3 + (d == 3), c6 + (d == 6),
       c9 + (d == 9)) % mod:
 return ret:
```

```
// optimized, without (tightLower, tightUpper) state
int f(int i, int tightLower, int tightUpper, int c3,
\rightarrow int c6, int c9) {
 if (c3 >= 17 \mid | c6 >= 17 \mid | c9 >= 17) return 0;
 if (i < 0) {
   return c3 && c3 == c6 && c6 == c9;
  int ret = dp[i][c3][c6][c9];
  if (~ret && !tightLower && !tightUpper) return ret;
  int lo = tightLower ? 1[i] - '0' : 0;
  int hi = tightUpper ? r[i] - '0' : 9;
  ret = 0:
  for (int d = lo: d <= hi: d++) {
    int newTightLower = tightLower && (d == lo);
    int newTightUpper = tightUpper && (d == hi);
    ret = (ret + f(i - 1, newTightLower,
    \rightarrow newTightUpper, c3 + (d == 3), c6 + (d == 6),
    \rightarrow c9 + (d == 9))) % mod;
 if (!tightLower && !tightUpper) dp[i][c3][c6][c9] =
 return ret;
void solve() {
  cin >> 1 >> r;
  int n = r.size();
  while (1.size() < n) 1 = "0" + 1;
  reverse(1.begin(), 1.end());
  reverse(r.begin(), r.end());
  cout << f(n - 1, 1, 1, 0, 0, 0) << '\n';
```

#### 15 Distinct Subsequence

```
const int mod = 1e9 + 7;
int distinctSubseq(string &a) {
  vector<int> last(26); // for array, use map
  int res = 1;
  for(auto& ai : a) {
    int curr = (2LL * res - last[ai - 'a']) % mod;
    if (curr < 0) curr += mod;
    last[ai - 'a'] = res;
    res = curr;
  }
  return (res - 1 + mod) % mod; // without empty set
}</pre>
```

#### 16 DSU Kruskal's Algorithm MST

```
struct DSU {
  vector<int> parent, sz;
  DSU(int n) {
    parent.resize(n + 1), sz.resize(n + 1, 1);
    iota(parent.begin(), parent.end(), 0);
}
int Find(int u) {
    if (u == parent[u]) return u;
    return parent[u] = Find(parent[u]);
}
```

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```
bool Is_connected(int u, int v) {
   return Find(u) == Find(v);
 void Merge(int u, int v) {
   u = Find(u), v = Find(v);
   if (u != v) {
     if (sz[u] < sz[v]) swap(u, v);
     parent[v] = u;
     sz[u] += sz[v];
};
```

#### 17 Exclution DP

```
vector<int> f(n), g(n);
for (int i = 1; i < N; i++) {
 for (int j = i; j < N; j += i) {
   f[i] += cnt[j];
 f[i] = nC2(f[i]);
for (int i = N - 1; i > 0; i--) {
 g[i] = f[i];
 for (int j = i * 2; j < N; j += i) {
   g[i] = g[j];
// g[i] = how many pairs have gcd i
```

### 18 Fenwick Tree 2D

```
struct FenWick2D {
 int n, m;
 vector<vector<long long>> ft;
 FenWick2D() {}
 FenWick2D(int _n, int _m) {
   n = _n, m = _m;
   ft.assign(n + 1, vector<long long> (m + 1));
 void Add(int x, int y, long long val) {
   for(int i = x; i \le n; i += i \& -i) {
     for(int j = y; j \le m; j += j \& -j) {
       ft[i][j] += val;
 long long Csum(int x, int y) {
   long long res = 0;
   for(int i = x; i > 0; i -= i & -i) {
     for(int j = y; j > 0; j -= j & -j) {
       res += ft[i][j];
   return res;
 long long Rsum(int x1, int y1, int x2, int y2) {
   return Csum(x2, y2) - Csum(x1 - 1, y2) - Csum(x2, y2)
      y1 - 1) + Csum(x1 - 1, y1 - 1);
```

# Fenwick Tree BIT

};

```
struct FenWickTree {
  int n;
  vector<long long> ft;
  FenWickTree() {}
 FenWickTree(int n) { Initial(n); }
  FenWickTree(vector<int>& a) {
    Initial((int)a.size());
    Build(a):
  void Initial(int _n) {
    ft.assign(n + 1, 0);
  void Build(vector<int>& a) {
    for (int i = 0; i < (int)a.size(); i++) {
      Add(i, i, a[i]); #change
  void Add(int idx, long long val) {
    for (int i = idx; i <= n; i += i & -i) {
     ft[i] += val:
  void Add(int 1, int r, long long val) {
    Add(1, val);
    Add(r + 1, -val);
 long long Csum(int idx) {
    long long res = 0;
    for (int i = idx; i > 0; i -= i & -i) {
     res += ft[i];
    return res;
 long long Rsum(int 1, int r) {
    return Csum(r) - Csum(l - 1);
};
20 FFT
```

```
struct cplx {
 long double a, b;
  cplx(long double a = 0, long double b = 0) : a(a),
  const cplx operator + (const cplx &c) const { return
  \rightarrow cplx(a + c.a, b + c.b); }
  const cplx operator - (const cplx &c) const { return
  \rightarrow cplx(a - c.a, b - c.b); }
  const cplx operator * (const cplx &c) const { return
  \rightarrow cplx(a * c.a - b * c.b, a * c.b + b * c.a); }
  const cplx operator / (const long double &d) const {

→ return cplx(a / d, b / d); }

const long double PI = acos(-1);
```

```
vector<int> rev:
void Preprocess(int sz) {
 if ((int)rev.size() == sz) return;
 rev.resize(sz);
 rev[0] = 0;
 int lg_n = __builtin_ctz(sz);
 for (int i = 1; i < sz; ++i) {
   rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (lg_n -
void fft(vector<cplx> &a, bool inv = 0) {
 int n = a.size();
 for (int i = 1; i < n - 1; ++i) {
   if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (int len = 2; len <= n; len <<= 1) {
   long double t = 2 * PI / len * (inv ? -1 : 1);
    cplx wlen = {cosl(t), sinl(t)};
   for (int st = 0; st < n; st += len) {
      cplx w(1);
      for (int i = 0; i < len / 2; ++i) {
       cplx ev = a[st + i], od = a[st + i + len / 2]
        → * W:
       a[st + i] = ev + od;
       a[st + i + len / 2] = ev - od;
       w = w * wlen:
 if (inv) {
   for (cplx\& z : a) z = z / n;
vector<long long> multiply(vector<int> &a, vector<int>
 int n = a.size(), m = b.size(), sz = 1;
 while(sz < n + m - 1) sz <<= 1;
 vector<cplx> x(sz), y(sz), z(sz);
 for(int i = 0; i < sz; i++) {
   x[i] = i < n ? cplx(a[i], 0) : cplx(0, 0);
   y[i] = i < m ? cplx(b[i], 0) : cplx(0, 0);
 Preprocess(sz):
 fft(x), fft(y);
 for(int i = 0; i < sz; i++) z[i] = x[i] * y[i];
 fft(z, 1);
 vector<long long> c(n + m - 1);
 for(int i = 0; i < n + m - 1; i++) c[i] =
  \rightarrow round(z[i].a);
 return c:
auto c = mutiply(a, b);
```

# 21 Floyd Warshall

void FloydWarshall(vector<vector<long long>>& adj) { int n = (int)adj.size();

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```
for (int k = 0; k < n; k++) {
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < n; j++) {
      adj[i][j] = min(adj[i][j], adj[i][k] +
      \rightarrow adi[k][i]);
      adj[i][i] = 0;
```

### 22 Geometry

```
const double PI = acos((double)-1.0);
using ld = long double;
using ll = long long;
struct point { ld x, y; };
void manhattanToChebychev(int& x, int& y) {
 int X = x + y, Y = x - y;
  x = X, y = Y;
void chebyshevToManhattan(int& x, int& y) {
 int X = (x + y) / 2, Y = (x - y) / 2;
 x = X, y = Y;
11 manhattanDistance(point& p1, point& p2) {
  return abs(p1.x - p2.x) + abs(p1.y - p2.y);
11 chebyshevDistance(point& p1, point& p2) {
 return max(abs(p1.x - p2.x), abs(p1.y - p2.y));
// determines the relative position (cross product) of
\rightarrow a point (p3)
// with respect to the line passing through (p1) and

→ (p2)

// if (d > 0): point (p3) is to the left of the line
// if (d < 0): point (p3) is to the right of the line</pre>
// if (d == 0): point (p3) lies exactly on the line.
11 determinant(point p1, point p2, point p3) {
 11 d = (p2.x - p1.x) * (p3.y - p1.y) - (p2.y - p1.y)
  \rightarrow * (p3.x - p1.x);
  return d;
bool arePointsOnSameSide(point p1, point p2, point p3,
\rightarrow point p4) {
 11 d1 = determinant(p1, p2, p3);
 11 d2 = determinant(p1, p2, p4);
 return (d1 * d2) > 0;
bool doLinesIntersect(point p1, point p2, point p3,
→ point p4) {
 11 d1 = determinant(p3, p4, p1);
  11 d2 = determinant(p3, p4, p2);
  11 d3 = determinant(p1, p2, p3);
  11 d4 = determinant(p1, p2, p4);
  return (d1 * d2 < 0) && (d3 * d4 < 0);
// distance of two points
ld distance(point& p1, point& p2) {
```

```
return sqrt((p2.x - p1.x) * (p2.x - p1.x) + (p2.y - p1.x))
  \rightarrow p1.y) * (p2.y - p1.y));
// a point inside cricle
// x1, y1 = point of circle, r = radius of circle
// x2, y2 = target point
bool isInside(int x1, int y1, int x2, int y2, int r) {
 return ((x2 - x1) * (x2 - x1)) + ((y2 - y1) * (y2 -
  \rightarrow v1)) <= r * r:
ld toDegrees(ld radians) {
 return radians * (180.0 / PI);
// circle with radius r
struct Circle {
 ld r;
  Circle(ld _r) { r = _r; }
  ld Diameter() { return 2 * r; }
 ld Circumference() { return 2 * PI * r; }
 ld Area() { return PI * r * r; }
  // if theta in radians, make it to degrees
  // all the calculations are in degree
 ld Sector_Area(ld theta) { return (theta / 360.0) *
  → Area(); }
 ld Arc_Length(ld theta) { return (theta / 360.0) *
  ld Chord_Length(ld theta) { return 2 * r *
  \rightarrow sin((theta * PI) / 360.0); }
 ld Segment_Area(ld theta) { return

    Sector_Area(theta) - (0.5 * r * r * sin((theta *
  → PI) / 180.0)); }
 ld Circumscribed_Square_Area() { return 4 * r * r; }
 ld Inscribed_Square_Area() { return 2 * r * r; }
 ld Sphere_Surface_Area() { return 4 * PI * r * r; }
 ld Sphere_Volume() { return (4.0 / 3.0) * PI * r * r
  \rightarrow * r: }
 ld Cylinder_Surface_Area(ld h) { return 2 * PI * r *
  \hookrightarrow (r + h): }
 ld Cylinder_Volume(ld h) { return PI * r * r * h; }
 ld Cone_Surface_Area(ld 1) { return PI * r * (r +
  → 1); }
 ld Cone_Volume(ld h) { return (1.0 / 3.0) * PI * r *
  \hookrightarrow r * h; }
 ld Annulus_Area(ld R) { return PI * (r * r - R * R);
  → }
// triangle with three sides: a, b, c
struct Triangle {
 ld a, b, c, s;
  Triangle(point x, point y, point z) {
    a = distance(x, y);
    b = distance(y, z);
    c = distance(z, x);
    s = (a + b + c) / 2;
 Triangle(ld _a, ld _b, ld _c) {
    a = a, b = b, c = c;
    s = (a + b + c) / 2;
```

```
ld Area() { return sqrt(s * (s - a) * (s - b) * (s -
  \rightarrow c)); }
  ld Perimeter() { return a + b + c; }
  ld Angle_A_Radians() { return acos((b * b + c * c -
  \rightarrow a * a) / (2 * b * c)); }
  ld Angle_B_Radians() { return acos((a * a + c * c -
  \rightarrow b * b) / (2 * a * c)); }
  ld Angle_C_Radians() { return acos((a * a + b * b -
  \rightarrow c * c) / (2 * a * b)); }
  ld Angle_A_Degrees() { return

    toDegrees(Angle_A_Radians()); }

  ld Angle_B_Degrees() { return

→ toDegrees(Angle_B_Radians()); }

  ld Angle_C_Degrees() { return

→ toDegrees(Angle_C_Radians()); }

  ld Inradius() { return sqrt((s - a) * (s - b) * (s -
  \rightarrow c) / s); }
  ld Circumradius() { return (a * b * c) / (4 *
  \rightarrow Area()): }
  ld Exradius_A() { return sqrt((s - a) * (s - b) * (s
  \rightarrow - c) / (s - a)); }
  ld Exradius_B() { return sqrt((s - a) * (s - b) * (s
  \rightarrow - c) / (s - b)); }
  ld Exradius_C() { return sqrt((s - a) * (s - b) * (s
  \rightarrow - c) / (s - c)); }
  ld Altitude_A() { return 2 * Area() / a; }
  ld Altitude_B() { return 2 * Area() / b; }
  ld Altitude_C() { return 2 * Area() / c; }
  ld Median_A() { return sqrt((2 * b * b + 2 * c * c -
  \rightarrow a * a) / 4): }
  ld Median_B() { return sqrt((2 * a * a + 2 * c * c -
  \rightarrow b * b) / 4); }
  ld Median_C() { return sqrt((2 * a * a + 2 * b * b -
  \hookrightarrow c * c) / 4): }
// n sided polygon (n-gon) with side length a
struct Polygon {
 int n;
 ld a;
  Polygon(int _n, ld _a) { n = _n, a = _a; }
  ld Perimeter() { return n * a; }
  ld Exterior_Angle_Degrees() { return 360.0 / n; }
  ld Exterior_Angle_Radians() { return (2 * PI) / n; }
  ld Interior_Angle_Degrees() { return (n - 2) * 180.0
  \rightarrow / n: }
  ld Interior_Angle_Radians() { return PI - (2 * PI /
  \rightarrow n); }
  ld Circumradius() { return a / (2 * sin(PI / n)); }
  ld Inradius() { return a / (2 * tan(PI / n)); }
  ld Area() { return (n * a * a) / (4 * tan(PI / n));
  ld Central_Angle_Degrees() { return 360.0 / n; }
  ld Central_Angle_Radians() { return (2 * PI) / n; }
  11 Diagonals_count() { return n * (n - 3) / 2; }
  ld Diagonal_length() { return a / sin(PI / n); }
  ld Height() {
```

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#### 23 Hash Map

```
#include <ext/pb_ds/assoc_container.hpp>
struct splitmix64_hash {
 static uint64_t splitmix64(uint64_t x) {
   x += 0x9e3779b97f4a7c15;
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
   x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
   return x ^(x >> 31);
 size_t operator()(uint64_t x) const {
   static const uint64_t FIXED_RANDOM =
    std::chrono::steady_clock::now().time_since_epoch().countMOD;[0]) % MOD[0];
   return splitmix64(x + FIXED_RANDOM);
};
template <typename K, typename V, typename Hash =

→ splitmix64 hash>

using hash_map = __gnu_pbds::gp_hash_table<K, V,</pre>
→ Hash>:
template <typename K, typename Hash = splitmix64_hash>
using hash_set = hash_map<K, __gnu_pbds::null_type,

→ Hash>;
```

# 24 Hashing with Update

```
const int MAX = int(1e6) + 9;
vector<long long> MOD = {19099999999, 1999999999};
vector<array<long long, 2>> pw(MAX), ipw(MAX);
array < long long, 2 > bs = {137, 277};
int BinExp(long long a, long long b, int mod) {
 a %= mod;
 int res = 1;
 while (b) {
   if (b \& 1) res = (res * a) % mod;
   a = (a * a) \% mod;
   b >>= 1:
 return res;
void Preprocess() {
 pw[0][0] = pw[0][1] = 1;
 for (int i = 1; i < MAX; i++) {
   pw[i][0] = (pw[i - 1][0] * bs[0]) % MOD[0];
   pw[i][1] = (pw[i - 1][1] * bs[1]) % MOD[1];
```

```
ipw[0][0] = ipw[0][1] = 1:
 long long ip1 = BinExp(bs[0], MOD[0] - 2, MOD[0]);
 long long ip2 = BinExp(bs[1], MOD[1] - 2, MOD[1]);
 for (int i = 1; i < MAX; i++) {
   ipw[i][0] = (ipw[i - 1][0] * ip1) % MOD[0];
   ipw[i][1] = (ipw[i - 1][1] * ip2) % MOD[1];
struct Hashing {
 int n:
  vector<array<long long, 2>> hs;
 Hashing(string& s) {
   if (pw[2][0] == 0) Preprocess();
   n = s.size();
   hs.resize(n + 1);
    for (int i = 0; i < n; i++) {
     hs[i + 1][0] = (hs[i][0] + (pw[i][0] * s[i]) %

    MOD[0]) % MOD[0];

     hs[i + 1][1] = (hs[i][1] + (pw[i][1] * s[i]) %
      \rightarrow MOD[1]) % MOD[1]:
 array<long long, 2> get_hash(int 1, int r) { // 0

→ based query

   l++, r++; // 1 based hashing
    long long res1 = ((hs[r][0] - hs[1 - 1][0]) +
    long long res2 = ((hs[r][1] - hs[1 - 1][1]) +
    \rightarrow MOD[1]) % MOD[1];
    res1 = (res1 * ipw[1 - 1][0]) % MOD[0];
    res2 = (res2 * ipw[1 - 1][1]) % MOD[1];
    //return res1 << 31 | res2;
    return {res1, res2};
auto Hash_Merge(array<long long, 2> left, array<long</pre>
→ long, 2> right, int left_sz) {
 for (int i = 0; i < 2; i++) {
    (right[i] *= pw[left_sz][i]) %= MOD[i];
    (left[i] += right[i]) %= MOD[i];
 return left;
// # with update, find palindrome O(nlogn)
const int N = 1e6 + 9:
int power(long long n, long long k, const int mod) {
 int ans = 1 % mod:
 n \% = mod;
 if (n < 0) n += mod;
 while (k) {
   if (k \& 1) ans = (long long) ans * n % mod;
   n = (long long) n * n % mod;
   k >>= 1;
 return ans;
using T = array<int, 2>;
const T MOD = \{127657753, 987654319\};
const T p = \{137, 277\};
```

```
T operator + (T a, int x) {return \{(a[0] + x) \%
\rightarrow MOD[0], (a[1] + x) % MOD[1]};
T operator - (T a, int x) {return \{(a[0] - x + MOD[0])\}
\rightarrow % MOD[0], (a[1] - x + MOD[1]) % MOD[1]};
T operator * (T a, int x) {return {(int)((long long))}
\rightarrow a[0] * x % MOD[0]), (int)((long long) a[1] * x %
T operator + (T a, T x) {return \{(a[0] + x[0]) \%
\rightarrow MOD[0], (a[1] + x[1]) % MOD[1]};
T operator - (T a, T x) {return \{(a[0] - x[0] +
\rightarrow MOD[0]) % MOD[0], (a[1] - x[1] + MOD[1]) %
\rightarrow MOD[1]}:}
T operator * (T a, T x) {return {(int)((long long))}
\rightarrow a[0] * x[0] % MOD[0]), (int)((long long) a[1] *
\rightarrow x[1] % MOD[1])};}
ostream& operator << (ostream& os, T hash) {return os
T pw[N], ipw[N];
void prec() {
 pw[0] = \{1, 1\};
 for (int i = 1; i < N; i++) {
    pw[i] = pw[i - 1] * p;
  ipw[0] = \{1, 1\};
  T = \{power(p[0], MOD[0] - 2, MOD[0]), power(p[1], \}\}
  \rightarrow MOD[1] - 2, MOD[1])}:
  for (int i = 1; i < N; i++) {
    ipw[i] = ipw[i - 1] * ip;
struct Hashing {
 int n;
  string s; // 1 - indexed
  vector<array<T, 2>> t; // (normal, rev) hash
  array<T, 2> merge(array<T, 2> 1, array<T, 2> r) {
   1[0] = 1[0] + r[0];
   l[1] = l[1] + r[1];
    return 1;
  void build(int node, int b, int e) {
    if (b == e) {
      t[node][0] = pw[b] * s[b];
      t[node][1] = pw[n - b + 1] * s[b];
    int mid = (b + e) >> 1, l = node << 1, r = 1 | 1;
    build(l, b, mid);
    build(r, mid + 1, e);
    t[node] = merge(t[1], t[r]);
  void upd(int node, int b, int e, int i, char x) {
    if (b > i || e < i) return;
    if (b == e && b == i) {
      t[node][0] = pw[b] * x;
      t[node][1] = pw[n - b + 1] * x;
      return;
    int mid = (b + e) >> 1, l = node << 1, r = 1 | 1;
```

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```
upd(1, b, mid, i, x);
    upd(r, mid + 1, e, i, x);
   t[node] = merge(t[1], t[r]);
 array<T, 2> query(int node, int b, int e, int i, int
  → j) {
   if (b > j \mid | e < i) return \{T(\{0, 0\}), T(\{0, 0\})\};
   if (b >= i && e <= j) return t[node];
    int mid = (b + e) >> 1, l = node << 1, r = 1 | 1;
   return merge(query(1, b, mid, i, j), query(r, mid
    \rightarrow + 1, e, i, j));
 Hashing() {}
 Hashing(string _s) {
   n = _s.size();
   s = "." + _s;
   t.resize(4 * n + 1);
    build(1, 1, n);
 void upd(int i, char c) {
    upd(1, 1, n, i, c);
    s[i] = c;
 T get_hash(int 1, int r) { // 1 - indexed
    return query(1, 1, n, l, r)[0] * ipw[1 - 1];
 T rev_hash(int 1, int r) { // 1 - indexed
   return query(1, 1, n, 1, r)[1] * ipw[n - r];
 T get_hash() return get_hash(1, n);
 bool is_palindrome(int 1, int r) return get_hash(1,
    r) == rev_hash(l, r);
};
```

#### $\operatorname{HLD}$ 25

```
struct HeavyLightDecomposition {
 int t = 0;
 vector<int> tin, depth, subtree_size, parent, heavy,
  → head, euler;
 SegmentTreeIterative<int> sg;
 HeavyLightDecomposition(int root,
  → vector<vector<int>>& adj, vector<int>& values) {
   int n = (int)adj.size() + 1;
   tin.resize(n), depth.resize(n),

    subtree_size.resize(n), parent.resize(n),
    → heavy.assign(n, -1), head.resize(n),

→ euler.resize(n);
   Dfs(root, root, adj);
   Decompose(root, root, adj, values);
    sg = SegmentTreeIterative<int> (euler);
 void Dfs(int u, int p, vector<vector<int>>& adj) {
    subtree_size[u] = 1;
   parent[u] = p;
    for (auto& v : adj[u]) {
     if (v != p) {
       depth[v] = depth[u] + 1;
       Dfs(v, u, adj);
```

```
subtree_size[u] += subtree_size[v];
        if (heavy[u] == -1 || subtree_size[heavy[u]] <</pre>

    subtree_size[v]) {

          heavy[u] = v;
  void Decompose(int u, int h, vector<vector<int>>&

→ adj, vector<int>& values) {
   tin[u] = ++t:
    euler[t] = values[u];
    head[u] = h:
    if (heavy[u] != -1) {
      Decompose(heavy[u], h, adj, values);
    for (auto& v : adj[u]) {
      if (v != parent[u] && v != heavy[u]) {
        Decompose(v, v, adj, values);
  void Update(int u, int val) {
   sg.Update(tin[u], val);
  int neutral = 0; #change
  int Merge(int a, int b) {
   return max(a, b); #change
 int PathQuery(int a, int b) {
    int res = neutral;
    for (; head[a] != head[b]; b = parent[head[b]]) {
      if (depth[head[a]] > depth[head[b]]) {
        swap(a, b);
     res = Merge(res, sg.Query(tin[head[b]],
      \rightarrow tin[b]));
    if (depth[a] > depth[b]) {
      swap(a, b);
    res = Merge(res, sg.Query(tin[a], tin[b]));
    return res;
HeavyLightDecomposition hld(1, adj, a);
// add segment tree
26 Hopcroft Karp
```

```
// O(\sqrt{V * E})
// works for only directed graph, or on two graphs
// 1 to n is left graph, n + 1 to n + m is right graph
// make directed edges for two graph
// match pairs -> if (i < match[i]) (i <-> match[i])
const int N = 2e5 + 9, INF = 1.1e9;
vector<int> adj[N];
int n, m, match[N], dist[N];
bool bfs() {
 queue<int> q;
```

```
for (int i = 1: i <= n: ++i) {
    if (!match[i]) dist[i] = 0, q.emplace(i);
    else dist[i] = INF;
 dist[0] = INF;
 while (!q.empty()) {
   int u = q.front(); q.pop();
   if (!u) continue;
   for (int v : adj[u]) {
      if (dist[match[v]] == INF) {
       dist[match[v]] = dist[u] + 1,
       q.emplace(match[v]);
   }
 return dist[0] != INF:
bool dfs (int u) {
 if (!u) return 1;
 for (int v : adj[u]) {
   if (dist[match[v]] == dist[u] + 1 and

    dfs(match[v])) {

      match[u] = v. match[v] = u:
      return 1;
 dist[u] = INF;
 return 0:
int max_matching() {
 int ret = 0;
 while (bfs()) {
   for (int i = 1; i <= n; ++i) {
      ret += !match[i] and dfs(i);
 }
 return ret;
```

## 27 Hungarian Min Assignment

```
// returns the max sum and permuation
// for minimum, convert each value to negative.
template<typename T>
pair<T, vector<int> > hungarian(const vector<vector<T>

→ > &a) {

 if (a.empty()) return {0, {}};
  int n = a.size() + 1, m = a[0].size() + 1;
  vector < int > p(m), ans(n - 1);
  vector < T > u(n), v(m);
  for (int i = 1; i < n; i++) {
    p[0] = i;
    int now = 0;
    vector\langle int \rangle pre(m, -1), vis(m + 1);
    vector<T> dis(m, numeric_limits<T>::max());
    do {
      vis[now] = true;
      int t = p[now], nxt;
```

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```
T d = numeric_limits<T>::max();
    for (int j = 1; j < m; j++)
      if (!vis[j]) {
       T cur = -a[t - 1][j - 1] - u[t] - v[j];
        if (cur < dis[j]) dis[j] = cur, pre[j] =</pre>
        if (dis[j] < d) d = dis[j], nxt = j;
    for (int j = 0; j < m; j++) {
      if (vis[j]) u[p[j]] += d, v[j] -= d;
      else dis[j] -= d;
    now = nxt;
 } while (p[now]);
  while (now) {
    int t = pre[now];
   p[now] = p[t], now = t;
for (int i = 1; i < m; i++)
 if (p[i]) ans[p[i] - 1] = i - 1;
return {v[0], ans};
```

#### 28 KMP

```
// length of longest proper prefix of P[0...i] that is
\rightarrow also a suffix of P[0...i]
// Pattern: a b a b a c a
// Index: 0 1 2 3 4 5 6
// LPS[]: 0 0 1 2 3 0 1
vector<int> get_pi(string& s) {
 int n = s.size();
 vector<int> pi(n);
 for (int i = 1, j = 0; i < n; i++) {
   if (s[i] == s[j]) pi[i] = ++j;
   else if (j == 0) pi[i] = 0;
   else j = pi[j - 1], i--;
 return pi;
// count the number of p occurs in s
string s, p;
s = p + "#" + s:
vector<int> pi = get_pi(s);
int ans = 0;
for (auto& i : pi) ans += (i == p.size());
```

# 29 Longest Path in DAG

```
vector<int> adj[N];
int dp[N], vis[N], n;
int dfs (int u) {
  vis[u] = dp[u] = 1;
  for (int v : adj[u]) {
    if (!vis[v]) dfs(v);
    dp[u] = max(dp[u], dp[v] + 1);
  } return dp[u];
}
int longestpath() {
```

```
int ans = 0;
for (int i = 1; i <= n; i++)
    if (!vis[i]) ans = max(ans, dfs(i));
    return ans;
}</pre>
```

#### 30 Manacher

```
struct Manacher { // 0 based
  int n:
  vector<int> p;
  Manacher(string& s) {
    n = s.size();
    p.resize(2 * n);
    build(s);
  void build(string& s) {
    for (int i = 0, j = 0, k; i < n * 2; i += k, j =
    \rightarrow max(i - k, 0)) {
      while (i \ge j \&\& i + j + 1 < n * 2 \&\& s[(i - j)]
      \rightarrow / 2] == s[(i + j + 1) / 2]) ++j;
      p[i] = j;
      for (k = 1; i >= k \&\& j >= k \&\& p[i - k] != j -
      \hookrightarrow k; ++k) {
        p[i + k] = min(p[i - k], j - k);
  bool is_palindrome(int 1, int r) {
    int len = (r - 1 + 1), idx = -1;
    if (len & 1) idx = (l + len / 2) * 2;
    else idx = (1 + len / 2 - 1) * 2 + 1;
    return p[idx] >= len;
 int odd_length_of_center_i(int i) {return p[i * 2];
 int even_length_of_center_i(int i) { return p[i * 2
  \leftrightarrow + 1]; }
};
```

# 31 Merge Sort Tree

```
// add ordered multiset
struct MergeSortTree {
  int size = 1;
  vector<ordered_multiset<int>> st; #change
  MergeSortTree() {}
  MergeSortTree(int n) { Initial(n); }
  MergeSortTree(vector<int>& a) {
    Initial((int)a.size() - 1);
    Build(1, 1, size, a);
  }
  void Initial(int _n) {
    size = _n;
    int tree_size = 1;
    while (tree_size < _n) tree_size *= 2;
    st.resize(tree_size * 2);
  }
  void Build(int u, int s, int e, vector<int>& a) {
```

```
if (s == e) { #change
      st[u].insert(a[s]);
    int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
    Build(v, s, m, a);
    Build(w, m + 1, e, a);
    for (int i = s; i <= e; i++) st[u].insert(a[i]);</pre>
  void Update(int u, int s, int e, int k, int prev,

    int curr) {

    if (s == e) { #change

    st[u].erase(st[u].find_by_order(st[u].order_of_key)

      st[u].insert(curr);
      return:
    int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
    if (k <= m) Update(v, s, m, k, prev, curr);</pre>
    else Update(w, m + 1, e, k, prev, curr);
    #change

    st[u].erase(st[u].find_by_order(st[u].order_of_key(p))

    st[u].insert(curr);
  void Update(int k, int prev, int curr) {
    Update(1, 1, size, k, prev, curr);
  int Query(int u, int s, int e, int 1, int r, int
  → val) {
    if (e < 1 || r < s) { #change
      return 0;
    if (1 <= s && e <= r) { #change
      return st[u].order_of_key(val);
    int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
    int lsum = Query(v, s, m, l, r, val);
    int rsum = Query(w, m + 1, e, l, r, val);
    return lsum + rsum; #change
  int Query(int 1, int r, int val) {
    return Query(1, 1, size, 1, r, val);
};
```

### 32 MEX with TRIE

```
const int N = 2e5 * 22 + 9;
int nxt[N][2], cnt[N], intCnt[N], node = 2;
void insert(int x) {
  int u = 1;
  for(int i = 20; i >= 0; i--) {
    int bit = (x >> i) & 1;
    if(!nxt[u][bit]) nxt[u][bit] = node++;
    u = nxt[u][bit];
    cnt[u]++;
}
```

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```
intCnt[u]++:
int find(int x) {
 int u = 1:
 for (int i = 20; i >= 0; i--) {
   int bit = (x >> i) & 1;
   if (!nxt[u][bit]) return 0;
   u = nxt[u][bit];
 return intCnt[u];
void erase(int x) {
 if (find(x) == 0) return;
 int u = 1:
 for (int i = 20; i \ge 0; i - -) {
   int bit = (x >> i) & 1;
   int v = nxt[u][bit];
   cnt[v]--;
   u = v;
 intCnt[u]--;
int mex() {
 int u = 1, ret = 0;
 for(int i = 20; i \ge 0; i--) {
   if(u == 0 || !nxt[u][0]) return ret;
   if(cnt[nxt[u][0]] >= (1 << i)) {
      u = nxt[u][1];
      ret |= (1 << i);
   } else {
      u = nxt[u][0];
 return ret;
```

# 33 Minimum Expression

# 34 Number Theory

```
gpf[j] = i;
      if (!lpf[j]) lpf[j] = i;
// find all factors upto N
for (int i = 2; i < N; i++) {
  int num = i:
  while (num > 1) {
    int p = lpf[num], cnt = 0;
    while (num > 1 && num % p == 0) {
      num /= p;
   factors[i].push_back({p, cnt});
int num = 10;
int total divisors = 1:
long long sum_of_divisors = 1;
for (auto& [p, c] : factors[num]) {
 total_divisors *= (c + 1);
 sum_of_divisors *= (pow(p, c + 1) - 1) / (p - 1);
// (a ^ b) % p (Binary Exponentiation)
int BinExp(long long a, long long b, int mod) {
 a %= mod;
  int res = 1:
 while (b) {
   if (b \& 1) res = (res * a) \% mod;
    a = (a * a) \% mod;
   b >>= 1;
 return res;
// (a * b) % p (Binary Multiplication)
int BinMul(long long a, long long b, int mod) {
 a %= mod;
  int res = 0;
  while (b) {
   if (b & 1) res = (res + a) % mod;
    a = (a + a) \% mod;
   b >>= 1:
 return res;
// (a ^ -1) % mod (inverse of a number)
BinExp(a, mod - 2);
// (a / b) % mod
BinMul(a, BinExp(b, mod - 2, mod), mod);
// (a ^ (b ^ c)) % mod
BinExp(a, BinExp(b, c, mod - 1), mod);
// Permutations and Combinations
struct Combinatorics {
  vector<long long> fact, inv, ifact;
  Combinatorics(int n) {
    fact.assign(n + 1, 1), inv.assign(n + 1, 1),
    \rightarrow ifact.assign(n + 1, 1);
    inv[0] = 0;
```

```
for (int i = 2: i \le n: i++) fact[i] = (fact[i -
    → 1] * i) % mod:
   for (int i = 2; i <= n; i++) inv[i] = mod - (mod /
    \rightarrow i) * inv[mod % i] % mod;
   for (int i = 2; i <= n; i++) ifact[i] = (ifact[i -

→ 1] * inv[i]) % mod;

  int nPr(int n, int r) { // Permutations
    if (n < r) return 0;
    return (fact[n] * ifact[n - r]) % mod:
 int nCr(int n, int r) { // Combinations
   return (nPr(n, r) * ifact[r]) % mod;
 int nCRr(int n, int r) { // Combinations with
   return (nPr(n + r - 1, r) * ifact[r]) \% mod;
} comb(N);
// phi of single integer
int n = 10;
long long num = n;
long long phi_of_n = n;
for (long long i = 2; i * i <= num; i++) {
 if (num % i == 0) {
    while (num \% i == 0) num /= i;
   phi_of_n -= phi_of_n / i;
if (num > 1) phi_of_n -= phi_of_n / num;
// phi upto N
vector<int> phi(N);
// initial 0 to N
iota(phi.begin(), phi.end(), 0);
for (int i = 2; i < N; i++) {
 if (phi[i] == i) {
   for (int j = i; j < N; j += i) {
      phi[j] -= phi[j] / i;
// gcd sum -> \sum gcd(i, n) for 1 <= i <= n; (n <= 1e9)
// \gcd(1, n) + \gcd(2, n) + ... + \gcd(n, n)
long long sum = 0;
for (int i = 1; 1LL * i * i <= n; i++) {
 if (n \% i == 0) {
    sum += i * phi(n / i);
    if (n / i != i) sum += (n / i) * phi(n / (n / i));
// all pair gcd sum from 1 to N (N <= 4e6)
// \sum i \text{ to n - 1}, \sum j + 1 \text{ to n } [\gcd(i, j)]
for (int i = 1; i < N; i++) {
 for (int j = i; j < N; j += i) sum[j] += i * phi[j /

    il:
```

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```
} for (int i = 1; i < N; i++) sum[i] += sum[i - 1] - i;

// lcm sum -> \( \sum \) lcm(i, n) for 1 <= i <= n; (n <= 1e6)

// lcm(1, n) + lcm(2, n) + .. + lcm(n, n)
phi[1] = 2; // phi[1] should be 2 for this algorithm
for (int i = 1; i < N; i++) {
    for (int j = i; j < N; j += i) lcm_sum[j] += 1LL * j
    \( \times \) * phi[i] * i / 2;
}

// all pair lcm sum from 1 to N (N <= 1e6)

// \( \sum \) i to n - 1, \( \sum \) j + 1 to n [lcm(i, j)]
phi[1] = 2; // phi[1] should be 2 for this algorithm
for (int i = 1; i < N; i++) {
    for (int j = i; j < N; j += i) lcm_sum[j] += (long
    \( \times \) long) j * (1LL * i * phi[i] / 2);
}
for (int i = 1; i < N; i++) lcm_sum[i] += lcm_sum[i -
    \( \times \) 1] - i;</pre>
```

#### 35 Ordered Set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template <class T> using ordered_set = tree<T,
→ null_type, less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

template <class T> using ordered_multiset = tree<T,</pre>

→ null_type, less_equal<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

template <class T, class R> using ordered_map =

    tree<T, R, less<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

// Sorting Descending (or Equal)
template <class T> using ordered_set = tree<T,</pre>

→ null_type, greater<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

template <class T> using ordered_multiset = tree<T,</pre>
→ null_type, greater_equal<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

template <class T, class R> using ordered_map =

    tree<T, R, greater<T>, rb_tree_tag,

    tree_order_statistics_node_update>;

st.order_of_key(k); // count of elements smaller than
\rightarrow k. - 0(\log n)
st.find_by_order(k); // Returns the iterator for the
\rightarrow kth element (0 based index). - O(log n)
st.size() - st.order_of_key(k + 1); // count of
\rightarrow elements greater than k. - O(\log n)
st.order_of_key(r + 1) - st.order_of_key(l); // count
\hookrightarrow of elements between 1 to r
st.erase(st.find_by_order(st.order_of_key(k))); //

→ erase in multiset
```

#### 36 SCC

```
// clear everything first, 1 based
```

```
// reverse scc for topological order
const int N = 5e5 + 10;
vector<int> adj[N], trans[N], scc[N];
int col[N], vis[N], idx = 0, n, m;
stack<int> st;
void dfs(int u) {
  vis[u] = 1;
  for (int v : adj[u]) if (!vis[v]) dfs(v);
  st.push(u);
void dfs2(int u) {
  col[u] = idx;
  scc[idx].push_back(u);
  for (int v : trans[u]) if (!col[v]) dfs2(v);
void findSCC() {
 for (int i = 1; i <= n; i++)
    if(!vis[i]) dfs(i);
  for (int u = 1; u <= n; u++)
    for (int v : adj[u])
      trans[v].push_back(u);
  while (!st.empty()) {
    int u = st.top(); st.pop();
    if (col[u]) continue;
    idx++; dfs2(u);
// take input and call findSCC();
```

# 37 Segment Tree Iterative Point Update Range Query

```
template <class T> struct SegmentTreeIterative {
 int n = 1;
 vector<T> st;
 SegmentTreeIterative() {}
 SegmentTreeIterative(int n) { Initial(n); }
 SegmentTreeIterative(vector<int>& a) {
   Initial((int)a.size() - 1);
   Build(a);
 void Initial(int _n) {
   int tree_size = 1;
   while (tree_size <= n) tree_size *= 2;</pre>
   st.resize(tree_size * 2);
 T neutral = INT_MAX; #change
 T Merge(T& a, T& b) { #change
   return min(a, b);
 void Build(vector<int>& a) {
   for (int i = 1; i <= n; ++i) {
     st[n + i] = a[i];
   for (int u = n - 1; u > 0; --u) {
     st[u] = Merge(st[u << 1], st[u << 1 | 1]);
 void Update(int idx, T val) {
```

#### 38 Segment Tree Lazy

```
const long long INF = 1.1e17;
struct node { #change
 long long sum, lazy_add;
 node() {
   sum = 0;
   lazy_add = INF;
};
struct SegmentTreeLazy {
 int size = 1;
 vector<node> st:
 SegmentTreeLazy() {}
  SegmentTreeLazy(int n) { Initial(n); }
 SegmentTreeLazy(vector<int>& a) {
   Initial((int)a.size() - 1);
    Build(1, 0, size, a);
 void Initial(int _n) {
    size = _n;
    int tree_size = 1;
    while (tree_size <= size) tree_size *= 2;</pre>
    st.resize(tree_size * 2);
 node Make_node(long long val) { #change
   node res;
   res.sum = val;
   res.lazy_add = INF;
   return res;
 node Merge(node& 1, node& r) { #change
   node res:
   res.sum = 1.sum + r.sum;
   return res;
 void Push(int u, int l, int r) { #change
   if (st[u].lazy_add == INF) return;
   if (1 != r) {
      int v = 2 * u, w = 2 * u + 1;
      if (st[v].lazy_add != INF) st[v].lazy_add +=

    st[u].lazy_add;
```

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```
else st[v].lazv add = st[u].lazv add:
    if (st[w].lazy_add != INF) st[w].lazy_add +=

    st[u].lazy_add;

    else st[w].lazy_add = st[u].lazy_add;
  st[u].sum += (r - 1 + 1) * st[u].lazy_add;
  st[u].lazy_add = INF;
void Build(int u, int s, int e, vector<int>& a) {
  if (s == e) {
    st[u] = Make_node(a[s]);
    return:
 int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
  Build(v, s, m, a);
 Build(w, m + 1, e, a);
  st[u] = Merge(st[v], st[w]);
void Update(int u, int s, int e, int l, int r, long
→ long val) {
 Push(u, s, e);
 if (e < 1 || r < s) return;
 if (1 <= s && e <= r) { #change
    if (st[u].lazy_add != INF) st[u].lazy_add +=
    \hookrightarrow val;
    else st[u].lazy_add = val;
    Push(u, s, e);
   return;
  int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
  Update(v, s, m, l, r, val);
  Update(w, m + 1, e, l, r, val);
  st[u] = Merge(st[v], st[w]);
void Update(int 1, int r, long long val) {
 Update(1, 0, size, 1, r, val);
node Query(int u, int s, int e, int l, int r) {
  Push(u, s, e);
 if (e < 1 || r < s) { #change
   return node();
 if (1 <= s && e <= r) return st[u];
  int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
 node lsum = Query(v, s, m, l, r);
 node rsum = Query(w, m + 1, e, 1, r);
 return Merge(lsum, rsum);
node Query(int 1, int r) {
 return Query(1, 0, size, 1, r);
int Idx_query(int u, int s, int e, int l, int r) {
 Push(u, s, e);
 if (e < 1 \mid | r < s) return -1;
  if (s == e) return st[u].mn == 0 ? s : -1;
  int v = u \ll 1, w = v \mid 1, m = s + e >> 1;
  int lq = Idx_query(v, s, m, l, r);
 if (lq == -1) {
    return Idx_query(w, m + 1, e, l, r);
```

```
return la:
  int Idx_query(int 1, int r) {
    return Idx_query(1, 0, size, 1, r);
};
39 Segment Tree
struct node { #change
 long long sum, pref, suff, ans;
 node() {
    sum = pref = suff = ans = 0;
};
struct SegmentTree {
  int size = 1;
  vector<node> st;
  SegmentTree() {}
  SegmentTree(int n) { Initial(n); }
  SegmentTree(vector<int>& a) {
    Initial((int)a.size() - 1);
    Build(1, 0, size, a);
  void Initial(int _n) {
    size = n:
    int tree_size = 1;
    while (tree_size <= size) tree_size *= 2;</pre>
    st.resize(tree_size * 2);
 node Make_node(int val) { #change
    node res:
    res.sum = val:
    res.pref = res.suff = res.ans = max(0, val);
    return res:
 node Merge(node& 1, node& r) { #change
    res.sum = 1.sum + r.sum;
    res.pref = max(1.pref, 1.sum + r.pref);
    res.suff = max(r.suff, r.sum + 1.suff);
    res.ans = max(max(l.ans, r.ans), l.suff + r.pref);
    return res;
  void Build(int u, int s, int e, vector<int>& a) {
    if (s == e) {
      st[u] = Make_node(a[s]);
     return:
    int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
    Build(v, s, m, a);
    Build(w, m + 1, e, a);
    st[u] = Merge(st[v], st[w]);
  void Update(int u, int s, int e, int k, long long
  → val) {
    if (s == e) { #change
      st[u] = Make_node(val);
     return;
```

```
int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2:
    if (k <= m) Update(v, s, m, k, val);</pre>
    else Update(w, m + 1, e, k, val);
    st[u] = Merge(st[v], st[w]);
  void Update(int k, long long val) {
    Update(1, 0, size, k, val);
  node Query(int u, int s, int e, int l, int r) {
    if (e < 1 || r < s) { #change
      return node();
    if (1 <= s && e <= r) return st[u];
    int v = 2 * u, w = 2 * u + 1, m = (s + e) / 2;
    node lsum = Query(v, s, m, l, r);
    node rsum = Query(w, m + 1, e, l, r);
    return Merge(lsum, rsum);
  node Query(int 1, int r) {
    return Query(1, 0, size, 1, r);
};
```

#### 40 SOS DP

```
const int mod = 1e9 + 7;
const int B = 20, N = (1 << B) + 9;
int f[N], g[N], pairF[N], pairG[N], subseqF[N],

→ subseqG[N], Pow2[N];

void sub(int& a, int& b) { // a = a - b
 a -= b;
 if (a < 0) a += mod;
int nC2(int n) {
 return (n * 1LL * (n - 1) / 2) \% mod;
void forward1() { // sum over subset
 for (int bit = 0; bit < B; bit++) {
   for (int i = 0; i < N; i++) {
      if (i & (1 << bit)) {
        f[i] += f[i ^ (1 << bit)];
 }
void backward1() { // exclude sebset sum
 for (int bit = 0; bit < B; bit++) {</pre>
   for (int i = N - 1; i \ge 0; i--) {
      if (i & (1 << bit)) {
        // f[i] -= f[i ^ (1 << bit)]:
        sub(pairF[i], pairF[i ^ (1 << bit)]);</pre>
        sub(subseqF[i], subseqF[i ^ (1 << bit)]);</pre>
void forward2() { // sum over superset
  for (int bit = 0; bit < B; bit++) {</pre>
```

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```
for (int i = N - 1; i \ge 0; i--) {
      if (i & (1 << bit)) {
       g[i ^(1 << bit)] += g[i];
 }
}
void backward2() { // exclude superset sum
 for (int bit = 0; bit < B; bit++) {
   for (int i = 0; i < N; i++) {
      if (i & (1 << bit)) {
        // g[i ^ (1 << bit)] == g[i];
        sub(pairG[i ^ (1 << bit)], pairG[i]);</pre>
        sub(subseqG[i ^ (1 << bit)], subseqG[i]);</pre>
}
void sosDP(vector<int>& a) {
 Pow2[0] = 1:
 for (int i = 1; i < N; i++) {
   Pow2[i] = (Pow2[i - 1] * 2) \% mod:
 for (auto& ai : a) {
   f[ai]++;
   g[ai]++;
 forward1();
 forward2();
 for (int i = 0; i < N; i++) {
   // all pair (i < j)
   pairF[i] = nC2(f[i]);
    pairG[i] = nC2(g[i]);
    // subsequence
    subseqF[i] = Pow2[f[i]] - 1;
    subseqG[i] = Pow2[g[i]] - 1;
 backward1();
 backward2():
int subMaskOf(int x) { // x | y = x
 return f[x];
int superMaskOf(int x) { // x & y = x
 return g[x];
int countPairsWithAnd(int x) { // y & z = x (i < j)</pre>
  return pairG[x];
int countPairsWithOr(int x) { // y | z = x (i < j)
 return pairF[x];
int countSubseqWithAnd(int x) {
 return subseqG[x];
int countSubseqWithOr(int x) {
 return subseqF[x];
sosDP(a); // clear everything
```

### 41 Sparse Table RMQ

```
const int N = (int)2e5 + 9:
int lg[N];
void Preprocess() {
 for (int i = 2; i < N; ++i) {
   lg[i] = lg[i / 2] + 1;
template <class T> struct RMQ {
 int n = 1, LOG = 1;
  vector<vector<T>> st:
 T Merge(T& a, T& b) {
    return min(a, b); #change
  RMQ() {}
  RMQ(vector<T>& a) {
    if (lg[2] == 0) Preprocess();
    n = (int)a.size(), LOG = __lg(n) + 1;
    st.assign(n, vector<T> (LOG));
    for (int j = 0; j < LOG; j++) {
      for (int i = 0; i + (1 << j) - 1 < n; i++) {
        if (j == 0) st[i][j] = a[i];
        else st[i][j] = Merge(st[i][j - 1], st[i + (1
        \leftrightarrow << (j - 1))][j - 1]);
   }
 T Query(int 1, int r) {
    // if (1 > r) return 0;
    int k = lg[r - l + 1];
    return Merge(st[1][k], st[r - (1 << k) + 1][k]);
};
```

# 42 Stress Testing

```
// rename file -> testing.sh
set -e
g++ -std=c++17 gen.cpp -o gen
g++ -std=c++17 main.cpp -o main
g++ -std=c++17 brute.cpp -o brute
for((i = 1; ; ++i)); do
    echo $i
    ./gen $i > in
    ./main < in > out
    ./brute < in > out2
    diff -w out out2 || break
done
```

# 43 Sublime Build

### 44 Suffix Array

```
for integer, just change string to vector<int> and

→ minimum value of vector must be >= 1

  for integer, lim will be the maximum value of the
 LCP of suffix (sa[i], sa[i + 1]) = lcp[i + 1]
 O(nlogn)
array<vector<int>, 2> get_sa(string& s, int lim = 128)
 int n = s.size() + 1, k = 0, a, b;
  vector < int > x(begin(s), end(s) + 1), y(n), sa(n),
  \rightarrow lcp(n), ws(max(n, lim)), rank(n);
  x.back() = 0;
  iota(begin(sa), end(sa), 0);
  for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim
  \Rightarrow = p) {
   p = \bar{j}, iota(begin(y), end(y), n - j);
   for (int i = 0; i < n; ++i) if (sa[i] >= j) y[p++]
    \rightarrow = sa[i] - j;
   fill(begin(ws), end(ws), 0);
   for (int i = 0; i < n; ++i) ws[x[i]]++;
    for (int i = 1; i < \lim; ++i) ws[i] += ws[i - 1];
   for (int i = n; i--;) sa[-ws[x[y[i]]]] = y[i];
    swap(x, y), p = 1, x[sa[0]] = 0;
   for (int i = 1; i < n; ++i) a = sa[i - 1], b =
     x[b] = (y[a] == y[b] && y[a + j] == y[b + j]) ?
      \rightarrow p - 1 : p++;
  for (int i = 1; i < n; ++i) rank[sa[i]] = i;
  for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
   for (k \&\& k--, j = sa[rank[i] - 1]; s[i + k] ==
    \rightarrow s[i + k]; k++);
```

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```
sa.erase(sa.begin()), lcp.erase(lcp.begin());
 return {sa, lcp};
 O(|S| + |alphabet|) Suffix Array
 LIM := max{s[i]} + 2
 LCP of suffix (sa[i], sa[i + 1]) = lcp[i]
void inducedSort(const vector<int>& vec, int

→ val_range, vector<int>& SA, const vector<int>& sl,

    const vector<int>& lms_idx) {

 vector<int> l(val_range, 0), r(val_range, 0);
 for (int c : vec) {
   ++r[c]; if (c + 1 < val_range) ++l[c + 1];
 partial_sum(1.begin(), 1.end(), 1.begin());
 partial_sum(r.begin(), r.end(), r.begin());
 fill(SA.begin(), SA.end(), -1);
 for (int i = lms_idx.size() - 1; i >= 0; --i)

    SA[--r[vec[lms_idx[i]]]] = lms_idx[i];

 for (int i : SA) if (i > 0 and sl[i - 1]) SA[1[vec[i
  \rightarrow - 1]]++] = i - 1;
 fill(r.begin(), r.end(), 0);
 for (int c : vec) ++r[c];
 partial_sum(r.begin(), r.end(), r.begin());
 for (int k = SA.size() - 1, i = SA[k]; k; --k, i =
  \rightarrow SA[k]) {
   if (i and !sl[i-1]) SA[--r[vec[i-1]]] = i-1;
vector<int> suffixArray(const vector<int>& vec, int
→ val range) {
 const int n = vec.size();
 vector<int> sl(n), SA(n), lms idx:
 for (int i = n - 2; i \ge 0; --i) {
   sl[i] = vec[i] > vec[i + 1] or (vec[i] == vec[i +
    \rightarrow 1] and sl[i + 1]):
   if (sl[i] and !sl[i + 1]) lms_idx.emplace_back(i +
       1):
 reverse(lms_idx.begin(), lms_idx.end());
 inducedSort(vec, val_range, SA, sl, lms_idx);
 vector<int> new_lms_idx(lms_idx.size()),

→ lms_vec(lms_idx.size());
 for (int i = 0, k = 0; i < n; ++i) {
   if (SA[i] > 0 and !sl[SA[i]] and sl[SA[i] - 1])
    \rightarrow new lms idx[k++] = SA[i]:
 int cur = 0; SA[n - 1] = 0;
 for (int k = 1; k < new_lms_idx.size(); ++k) {</pre>
   int i = new_lms_idx[k - 1], j = new_lms_idx[k];
   if (vec[i] ^ vec[i]) {
     SA[j] = ++cur; continue;
   bool flag = 0;
   for (int a = i + 1, b = j + 1; ; ++a, ++b) {
     if (vec[a] ^ vec[b]) {
       flag = 1; break;
```

```
if ((!sl[a] \text{ and } sl[a-1]) \text{ or } (!sl[b] \text{ and } sl[b-
      → 1])) {
        flag = !(!sl[a] and sl[a - 1] and !sl[b] and
        \rightarrow sl[b - 1]); break;
    SA[j] = flag ? ++cur : cur;
  for (int i = 0; i < lms_idx.size(); ++i) lms_vec[i]</pre>
  if (cur + 1 < lms_idx.size()) {</pre>
    auto lms_SA = suffixArray(lms_vec, cur + 1);
    for (int i = 0; i < lms_idx.size(); ++i)</pre>
    → new lms idx[i] = lms idx[lms SA[i]]:
  inducedSort(vec, val_range, SA, sl, new_lms_idx);

→ return SA:

vector<int> getSuffixArray(const string& s, const int
→ LIM = 128) { // change limit for integer array,
\hookrightarrow (integer > 0)
 vector<int> vec(s.size() + 1);
  copy(begin(s), end(s), begin(vec)); vec.back() =
  auto ret = suffixArrav(vec. LIM);
 ret.erase(ret.begin()); return ret;
// build RMQ on it to get LCP of any two suffix
vector<int> getLCParray(const string& s, const

    vector<int>& SA) {

  int n = s.size(), k = 0;
  vector<int> lcp(n), rank(n);
  for (int i = 0; i < n; ++i) rank[SA[i]] = i;
  for (int i = 0; i < n; ++i, k ? --k : 0) {
    if (rank[i] == n - 1) {
      k = 0; continue;
    int j = SA[rank[i] + 1];
    while (i + k < n \text{ and } j + k < n \text{ and } s[i + k] == s[j]
    \rightarrow + kl) ++k:
    lcp[rank[i]] = k;
 lcp[n - 1] = 0; return lcp;
int lower_bound(string& s, string& t, vector<int>& sa)
  int n = s.size(), m = t.size(), lo = 0, hi = n - 1;
  while (lo <= hi) {
    int mid = (lo + hi) / 2;
    if (s.substr(sa[mid], m) < t) lo = mid + 1;</pre>
    else hi = mid - 1:
 return lo;
int upper_bound(string& s, string& t, vector<int>& sa)
 int n = s.size(), m = t.size(), lo = 0, hi = n - 1;
 while (lo <= hi) {
```

```
int mid = (lo + hi) / 2;
   if (s.substr(sa[mid], m) \le t) lo = mid + 1;
   else hi = mid - 1;
 return lo;
int find_occurrence(string& s, string& t, vector<int>&
 return upper_bound(s, t, sa) - lower_bound(s, t,
const int N = 1e6 + 9:
int64_t distPref[N];
void kthSubstringDistinctPreprocess(vector<int>& sa,

    vector<int>& lcp) {

int last = 0, n = sa.size();
 for (int i = 0: i < n: i++) {
   distPref[i] = n - sa[i] - last;
   if (i) distPref[i] += distPref[i - 1]:
   last = lcp[i];
array<int, 2> kthSubstringDistinctPos(string& s,

→ int64 t k, vector<int>& sa, vector<int>& lcp) {
 int n = s.size(); k--;
 int i = upper_bound(distPref, distPref + n, k) -

→ distPref;

 int len = k - (i == 0 ? 0 : distPref[i - 1]) + (i ==
  \rightarrow 0 ? 0 : lcp[i - 1]) + 1;
 return {i. len}:
string kthSubstringDistinct(string& s, int64_t k,

    vector<int>& sa, vector<int>& lcp) {

 auto [i, len] = kthSubstringDistinctPos(s, k, sa,
  → lcp);
 return s.substr(sa[i], len);
int64 t pref[N]:
void kthSubstringPreprocess(vector<int>& sa) {
 int n = sa.size();
 for (int i = 0; i < n; i++) {
   pref[i] = n - sa[i];
   if (i) pref[i] += pref[i - 1];
string kthSubstring(string& s, int64_t k, vector<int>&
int n = s.size();
 int64_t = 1, hi = distPref[n - 1];
 while (lo <= hi) {
   int64_t mid = (lo + hi) >> 1;
   auto [i, len] = kthSubstringDistinctPos(s, mid,
    \rightarrow sa, lcp);
   int64_t totCnt = (i == 0 ? 0 : pref[i - 1]) + len;
   int mn = len;
   for (int j = i; j < n; j++) {
     mn = min(mn, lcp[j]);
```

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```
if (mm == 0) break;
   totCnt += mn;
}
if (totCnt < k) lo = mid + 1;
else hi = mid - 1;
}
return kthSubstringDistinct(s, lo, sa, lcp);
}
int main() {
   string s;
   auto sa = getSuffixArray(s);
   auto lcp = getLCParray(s, sa);
   kthSubstringDistinctPreprocess(sa, lcp);
   kthSubstringPreprocess(sa);
}</pre>
```

#### 45 Topological Sort

```
bool topological_sort(vector<vector<int>>& adj) {
 int n = adj.size();
 vector<int> in_degree(n + 1);
 for (int u = 1; u <= n; u++) {
   for (auto& v : adj[u]) in_degree[v]++;
 // without sorting the order
 queue<int> q;
 // sort with lexicographically
 // priority_queue<int, vector<int>, greater<int>> q;
 for (int u = 1; u <= n; u++) {
   if (in_degree[u] == 0) {
      q.push(\tilde{u});
 if (q.empty()) return 0;
 vector<int> order;
 while (!q.empty()) {
   # change the queue operation
   int u = q.front();
    q.pop();
   for (auto& v : adj[u]) {
     in_degree[v]--;
     if (in_degree[v] == 0) {
       q.push(v);
   order.push_back(u);
 return (int)order.size() == n:
```

#### 46 Trie

```
// iterative
const int N = 1e6 + 9, A = 26;
int nxt[N][A], pref_cnt[N], word_cnt[N], node = 2;
void insert(string& s) {
  int u = 1;
  for (auto& c : s) {
    int idx = c - 'a';
    if (!nxt[u][idx]) nxt[u][idx] = node++;
```

```
u = nxt[u][idx]:
   pref_cnt[u]++;
 word_cnt[u]++;
int countPref(string& s) {
 int u = 1;
 for (auto& c : s) {
    int idx = c - 'a';
    if (!nxt[u][idx]) return 0;
    u = nxt[u][idx];
 return pref_cnt[u];
int find(string& s) {
  int u = 1;
  for (auto& c : s) {
    int idx = c - 'a';
    if (!nxt[u][idx]) return 0;
    u = nxt[u][idx];
 return word_cnt[u];
void erase(string& s) {
 if (find(s) == 0) return:
  int u = 1:
 for (auto& c : s) {
    int idx = c - 'a':
    int v = nxt[u][idx];
    pref_cnt[v][idx]--;
    u = v;
 word_cnt[u]--;
// trie integer iterative (max xor, min xor)
const int N = 2e5 * 31 + 9, A = 2;
int nxt[N][A], pref_cnt[N], int_cnt[N], node = 2;
void insert(int num) {
  int u = 1;
  for (int bit = 30; bit >= 0; bit--) {
    int idx = (num >> bit) & 1;
    if (!nxt[u][idx]) nxt[u][idx] = node++;
    u = nxt[u][idx]:
    pref_cnt[u]++;
 int_cnt[u]++;
int find(int num) {
 int u = 1;
 for (int bit = 30; bit >= 0; bit--) {
    int idx = (num >> bit) & 1;
    if (!nxt[u][idx]) return 0;
    u = nxt[u][idx];
  return int_cnt[u];
void erase(int num) {
 if (find(num) == 0) return:
  int u = 1:
  for (int bit = 30; bit >= 0; bit--) {
```

```
int idx = (num >> bit) & 1:
    int v = nxt[u][idx];
    pref_cnt[v]--;
    \bar{u} = v:
 int_cnt[u]--;
int maxXor(int num) {
 int res = 0, u = 1;
 for (int bit = 30; bit >= 0; bit--) {
    int idx = (num >> bit) & 1;
    int flip = idx ^ 1;
    if (nxt[u][flip] && pref_cnt[nxt[u][flip]]) {
      res += (1 << bit);
      u = nxt[u][flip];
   } else {
      u = nxt[u][idx];
 return res;
int minXor(int num) {
 int res = 0, u = 1;
 for (int bit = 30; bit >= 0; bit--) {
    int idx = (num >> bit) & 1;
    if (nxt[u][idx] && pref_cnt[nxt[u][idx]]) {
      u = nxt[u][idx];
    } else {
      res += (1 << bit);
      u = nxt[u][idx ^ 1];
 }
 return res;
```

#### 47 XOR Basis

```
using ll = long long;
const int D = 60;
struct XorBasis {
 // number of subsequences of xor-sum X, for ith
  \rightarrow prefix -> 2 ^ (i - sz);
 11 basis[D] = {}:
 int sz = 0, n = 0;
 void insert(ll x) {
   for (int i = D - 1; i \ge 0; i--) {
      if (!((x >> i) & 1)) continue;
      if (!basis[i]) {
        basis[i] = x;
        sz++:
       break;
      x = basis[i];
   } n++;
 bool canRepresent(11 x) {
   for (int i = D - 1; i >= 0; i--) {
     if ((x >> i) & 1) x ^= basis[i];
```

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```
return x == 0:
11 \max Xor(11 x = 0) {
  for (int i = D - 1; i \ge 0; i - -) {
    if ((x ^ basis[i]) > x) x ^= basis[i];
11 minXor() { // except xor 0
  for (int i = 0; i < D; i++) {
    if (basis[i]) return basis[i];
 return 0;
11 maxXorWith(ll x) {
  return maxXor(x):
11 minXorWith(11 x) {
  for (int i = D - 1; i \ge 0; i--) {
    if ((x ^ basis[i]) < x) x ^= basis[i];
  return x;
11 countDistinctXors() {
  return 1LL << sz;
ll kthXor(ll k) { // returns k + 1 th xor, 1st zor
\hookrightarrow is 0
 11 \text{ res} = 0:
 11 tot = countDistinctXors();
  if (tot < k) return -1:
  for (int i = D - 1; i \ge 0; i--) {
    if (basis[i]) {
      11 low = tot / 2:
      if ((low < k && (res & 1 << i) == 0) ||
        (low >= k && (res & 1 << i) > 0)) res ^=
        → basis[i];
      if (low < k) k = low;
      tot \neq 2;
  return res;
ll kthLargestXor(ll k) {
  return kthXor(countDistinctXors() - 1 - k);
11 kthXorAllCombinations(ll k) {
  if (n - sz > 60) return kthXor(1);
 ll totComb = 1LL << n:
  11 disComb = countDistinctXors();
  11 dupPerDis = totComb / disComb:
 ll \ dis Idx = (k - 1) / dup Per Dis + 1;
  return kthXor(disIdx);
11 kthLargestXorAllCombinations(11 k) {
  if (n - sz > 60) return

    kthXor(countDistinctXors());
  11 \text{ totComb} = 1LL << n:
 11 disComb = countDistinctXors();
 11 dupPerDis = totComb / disComb;
```

```
ll disIdx = (k - 1) / dupPerDis + 1;
ll desIdx = disComb - disIdx + 1;
return kthXor(desIdx);
}
ll countSubsetsLessThan(ll x) {
    ll lo = 0, hi = countDistinctXors();
    while (lo <= hi) {
        ll mid = (lo + hi) / 2;
        if (kthXor(mid) < x) lo = mid + 1;
        else hi = mid - 1;
    }
    return hi;
}
ll countSubsetsWithXor(ll x) {
    if (!canRepresent(x)) return 0;
    return 1LL << (n - sz); // if n > 60 use mod
}
};
```

#### 48 Z Function

#### 49 Mathematical Formulas & Notes

- 1. **PI up to 31:** 3.1415926535897932384626433832795
- 2. **PI** value in **CPP**: 2 \* acos(0), 2 \* asin(1), M *PI*
- 3. Formula for angle C using the Law of Cosines  $C = \cos^{-1}\left(\frac{a^2 + b^2 c^2}{2ab}\right)$
- 4. Sum remainder: n mod 1+n mod 2+ n mod 3 +.....+n mod m: n\*m-sum of divisors form 1 to n
- 5. A number is divisible by 60 if and only if it is divisible by 3 and 20
- 6. All numbers greater than 1099 can be written as a sum of 11 and 111

- 7. Legendre's formula:  $\nu_p(n!) = \sum_{i=1}^L \left\lfloor \frac{n}{p^i} \right\rfloor$ , where  $L = \lfloor \log_p n \rfloor$
- 8.  ${}^{n}C_{r} = \binom{n}{r} = \frac{n!}{r!(n-r)!}, {}^{n}P_{r} = \frac{n!}{(n-r)!}$
- 9. (a-b)%M = (a%M b%M + M)%M
- 10.  $(a/b)\%M = (a\%M * b^{-1}\%M)\%M$
- 11.  $(a^b)\%M = ((a\%M)^b)\%M$
- 12. Euler's Totient Function (ETF): Count of numbers less than n that are co-prime to n is,  $\phi(n) = n * \prod_{p|n} (1 \frac{1}{p})$ Here, p = distinct prime factors of n
- 13. Congruence:  $a \equiv b \pmod{M}$  if a%M = b%M and  $M \mid (a b)$
- 14. Euler's Theorem:  $a^{\phi(M)} \equiv 1 \pmod{M}$ , where a and M are co-prime
- 15. Fermat's Little Theorem:  $a^{M-1} \equiv 1 \pmod{M}$ , where M is prime
- 16.  $a^b \equiv a^{b \mod \phi(M)} \pmod{M}$  or it can be written that,  $a^b \mod M = a^{b \mod \phi(M)} \mod M$  or,  $a^b \mod M = a^{b \mod M-1} \mod M$
- 17. x steps forward or backward in a circular number range:

newPos = l + ((pos - l + x)%N) + N)%NWhere N = l - r + 1 (total numbers in range)

- 18. Stars & Bars:  $x_1 + x_2 + \cdots + x_k = n$  with  $x_i \ge 0$  has  $\binom{n+k-1}{n}$  solutions and with  $x_i > 0$  has  $\binom{n-1}{k-1}$  solutions.
- 19. Number of subsequences of length k from an array of size n such that x appears at least once in the subsequence, where that x appears c times in the array:  ${}^{n}C_{k} {}^{n-c}C_{k}$
- 20. Hockey-stick Identity:  $n > r, \sum_{i=r}^{n} {i \choose r} = {n+1 \choose r+1}$

21. 
$$a^k - b^k = (a - b) \cdot (a^{k-1}b^0 + a^{k-2}b^1 + \dots + a^0b^{k-1})$$

22.  $ab \mod ac = a(b \mod c)$ 

23. 
$$|a-b|+|b-c|+|c-a|=2(\max(a,b,c)-\min(a,b,c))$$

24. if 
$$a \cdot b \leq c$$
 then  $a \leq \left| \frac{c}{b} \right|$  Same for  $<, \leq, >, \geq$ 

25. For positive integer 
$$n$$
 & arbitrary real numbers  $m, x$ ,  $\left\lfloor \frac{\lfloor x/m \rfloor}{n} \right\rfloor = \left\lfloor \frac{x}{mn} \right\rfloor$  and  $\left\lceil \frac{\lceil x/m \rceil}{n} \right\rceil = \left\lceil \frac{x}{mn} \right\rceil$ 

# Simple Formulas

1. 
$$(a \pm b)^3 = a^3 \pm 3a^2b + 3ab^2 \pm b^3$$

2. 
$$a^3 \pm b^3 = (a+b)(a^2 \mp ab + b^2)$$

3. 
$$\sin(A \pm B) = \sin A \cos B \pm \cos A \sin B$$

4. 
$$\cos(A \pm B) = \cos A \cos B \mp \sin A \sin B$$

5. 
$$\tan(A \pm B) = \frac{\tan A \pm \tan B}{1 \pm \tan A \tan B}$$

5. 
$$\tan(A \pm B) = \frac{\tan A \pm \tan B}{1 \mp \tan A \tan B}$$
  
6.  $\sin A \pm \sin B = 2\sin\left(\frac{A \pm B}{2}\right)\cos\left(\frac{A \mp B}{2}\right)$ 

7. 
$$\cos A + \cos B = 2\cos\left(\frac{A+B}{2}\right)\cos\left(\frac{A-B}{2}\right)$$

8. 
$$\cos A - \cos B = 2\sin\left(\frac{A+B}{2}\right)\sin\left(\frac{B-A}{2}\right)$$

9. 
$$\sin A \sin B = -\frac{1}{2} [\cos(A+B) - \cos(A-B)]$$

10. 
$$\cos A \cos B = \frac{1}{2} [\cos(A+B) + \cos(A-B)]$$

11. 
$$\sin A \cos B = \frac{1}{2} [\sin(A+B) + \sin(A-B)]$$

$$12.\sin 2\theta = 2\sin \theta\cos \theta = \frac{2\tan \theta}{1+\tan^2 \theta}$$

13. 
$$\cos 2\theta = \cos^2 \theta - \sin^2 \theta = 2\cos^2 \theta - 1$$
  
=  $1 - 2\sin^2 \theta = \frac{1 - \tan^2 \theta}{1 + \tan^2 \theta}$ 

$$14. \tan 2\theta = \frac{2 \tan \theta}{1 - \tan^2 \theta}$$

$$15.\sin 3\theta = 3\sin \theta - 4\sin^3 \theta$$

$$16.\cos 3\theta = 4\cos^3 \theta - 3\cos \theta$$

$$17. \tan 3\theta = \frac{3 \tan \theta - \tan^3 \theta}{1 - 3 \tan^2 \theta}$$

$$18.1 + \cos 2\theta = 2\cos^2 \theta$$

$$19.1 - \cos 2\theta = 2\sin^2 \theta$$

$$20.1 \pm \sin 2\theta = (\cos \theta \pm \sin \theta)^2$$

$$21.\sin^{-1} x = \cos^{-1} \sqrt{1 - x^2}$$

$$22.2\sin^{-1}x = \sin^{-1}(2x\sqrt{1-x^2})$$

$$23.2\cos^{-1}x = \cos^{-1}(2x^2 - 1)$$

$$24.2 \tan^{-1} x = \tan^{-1} \left( \frac{2x}{1-x^2} \right)$$

$$25.2\tan^{-1}x = \sin^{-1}\left(\frac{2x}{1+x^2}\right)$$

$$26.2 \tan^{-1} x = \cos^{-1} \left( \frac{1 - x^2}{1 + x^2} \right)$$

$$27.3\sin^{-1}x = \sin^{-1}(3x - 4x^3)$$

$$28.3\cos^{-1}x = \cos^{-1}(4x^3 - 3x)$$

$$29.3 \tan^{-1} x = \tan^{-1} \left( \frac{3x - x^3}{1 - 3x^2} \right)$$

$$30.\sin^{-1} x \pm \sin^{-1} y = \sin^{-1} \left[ x\sqrt{1 - y^2} \pm y\sqrt{1 - x^2} \right]$$

$$31.\cos^{-1}x \pm \cos^{-1}y = \cos^{-1}\left[xy \mp \sqrt{1-x^2}\sqrt{1-y^2}\right]$$

$$32. \tan^{-1} x \pm \tan^{-1} y = \tan^{-1} \left( \frac{x \pm y}{1 \mp xy} \right)$$

33. 
$$\tan^{-1} x + \tan^{-1} y + \tan^{-1} z = \tan^{-1} \left( \frac{x+y+z-xyz}{1-xy-yz-zx} \right)$$

#### **Short Formulas**

- 1.  $\sin(-\theta) = -\sin\theta$
- 2.  $\cos(-\theta) = \cos\theta$
- 3.  $tan(-\theta) = -tan \theta$
- 4.  $\sin\left(\frac{\pi}{2} \theta\right) = \cos\theta$
- 5.  $\cos\left(\frac{\pi}{2} \theta\right) = \sin\theta$
- 6.  $\tan\left(\frac{\pi}{2} \theta\right) = \cot\theta$
- 7.  $\cot\left(\frac{\pi}{2} \theta\right) = \tan\theta$ 9.  $\sec\left(\frac{\pi}{2} - \theta\right) = \csc\theta$
- 8.  $\csc\left(\frac{\pi}{2}-\theta\right)=\sec\theta$  $10.\sin\left(\frac{\pi}{2} + \theta\right) = \cos\theta$
- $11.\cos\left(\frac{\pi}{2}+\theta\right)=-\sin\theta$
- $12.\tan\left(\frac{\pi}{4}+\theta\right)=\frac{1+\tan\theta}{1-\tan\theta}$
- $13.\tan\left(\frac{\pi}{4}-\theta\right)=\frac{1-\tan\theta}{1+\tan\theta}$
- $14.\sin(\pi+\theta) = -\sin\theta$
- $15.\cos(\pi+\theta) = -\cos\theta$
- $16.\sin(\pi-\theta)=\sin\theta$
- $17.\cos(\pi-\theta) = -\cos\theta$
- $18.\log_a a = 1$

 $19. a^{\log_a b} = b$ 

- $20.\log_b a = \frac{\log a}{\log b}$
- $21.\log a^m = m\log a$