

Goal

- 1) Motivation for Intelligence
- 2) Convexity
- 3) Convex Optimization
- 4) QP based MPC

I. Motivation for Intelligence

	Dynamics	Input/State Constraints	Real-Time	Feedback
Ackermann's Equation	LTI	no	yes	yes
LQK	LTV	no *	yes	yes
QP MPC	LTV	Yes *	yes	yes *
traj. optimization	Nonlinear	yes	no	no

Def: An **optimization problem** is:

$$\min_{z \in Z} J(z) \quad (P)$$

Subject to $g_i(z) \leq 0 \quad \forall i \in \{1, \dots, p_g\}$
 $h_i(z) = 0 \quad \forall i \in \{1, \dots, p_h\}$

where Z is the **decision space/set**,
 $J: Z \rightarrow \mathbb{R}$ is the **cost/objective function**

which is minimized over the decision variable $z \in Z$ that satisfy the inequality constraints $g_i: Z \rightarrow \mathbb{R}$ and equality constraints $h_i: Z \rightarrow \mathbb{R}$.

The set of points that satisfy the constraints $\{z \in Z \mid g_i(z) \leq 0, h_i(z) = 0, \forall i\}$ are called the feasible set.

Ex: Consider the kinematic steering model:

$$\begin{bmatrix} p_x \\ \dot{p}_x \\ p_y \\ \dot{p}_y \\ \theta \end{bmatrix} = \begin{bmatrix} v \cos \theta \\ v \sin \theta \\ w \end{bmatrix} = f(x, u)$$

↓ ↓

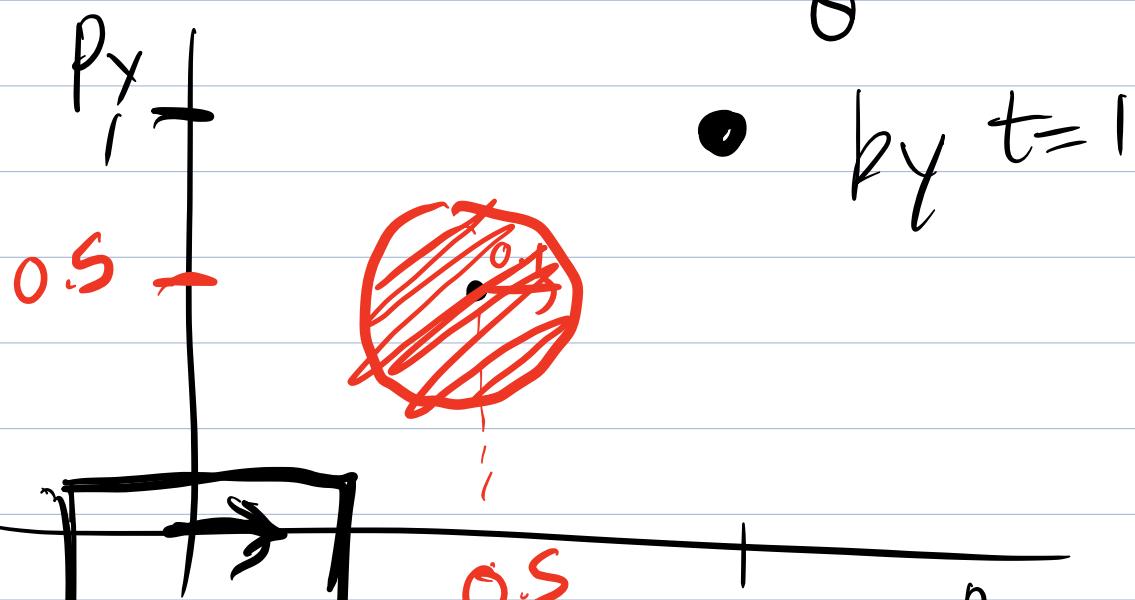
p_x v, w

\dot{p}_x

p_y

\dot{p}_y

θ



$$\begin{aligned} p_x &= 0 \\ p_y &= 0 \\ \theta &= 0 \\ t &= 1 \end{aligned}$$

$$V \in [-1, 1]$$

$$W \in [-1, 1]$$

Step 1: discretize the dynamics

use Euler Integration, choose step size $h < 1$
 $\forall h \in \mathbb{N}$, then

$$x(h(k+1)) = x(hk) + h f(x(hk), u(hk))$$

for each $k \in \{0, \dots, \frac{1}{h}-1\}$

Step 2: decision variable: $\{u(kh)\}_{k=0}^{\frac{1}{h}-1}$

$$\{x(kh)\}_{k=0}^{\frac{1}{h}}$$

Question: What is Z ? \mathbb{R}^3

$$Z = (\mathbb{R}^2)^{\frac{1}{h}-1} \times (\mathbb{R}^3)^{\frac{1}{h}}$$

Step 3:

a) Equality Constraints:

i) represent the dynamics discretized via Euler Integration

ii) initial condition

$$x(0 \cdot h) = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

$v(0)$

$w(0)$

$$-l \leq v(0) \leq l \rightarrow v(0) - l \leq 0$$

$$-(-v(0)) \leq 0$$

b) Inequality Constraints

i) ensuring input constraints are satisfied

$$-l \leq u(kh) \leq l \quad \forall k \in \{0, \dots, \frac{t}{h}-1\}$$

ii) check obstacle constraint

$$(p_x(kh) - 0.5)^2 + (p_y(kh) - 0.5)^2 - (0.1)^2 \geq 0$$

for each $k \in \{1, \dots, \frac{t}{h}\}$

Step 4: Cost Function

$$(p_x(1) - 1)^2 + (p_y(1) - 1)^2 + \sum u^2(kh)h$$

$J: \mathbb{Z} \rightarrow \mathbb{R}$

$J(p_x^{(1)}, p_y^{(1)}) \notin \mathbb{R}$

$u(0)$

$u(h)$

$u(2h)$

\vdots

$x(0)$

$x(h)$

\vdots

$Z = \begin{pmatrix} u(0) \in \mathbb{R}^2 \\ x(0) \in \mathbb{R}^3 \\ x(1) \in \mathbb{R}^3 \end{pmatrix} \leftarrow h=1$

$u(0) = \frac{v(0)}{w(0)} \times 1 \quad (1 \leftarrow)$

$x(0) = \frac{p_x(0)}{p_y(0)} \leftarrow$

$0(0) \leftarrow$

$x(1) = \frac{p_x^{(1)}}{p_y^{(1)}} \leftarrow$

$$(z(6) - 1)^2 + (z(7) - 1)^2$$

When is it easy to solve?

Our approach is to use derivatives to solve this problem.

Def: a) let J^* be the optimal value of the optimization problem (P). A **global optimizer** is a feasible z^* s.t.

$$f(z^*) = J^*$$

b) A feasible \bar{z} is a local optimizer to (P) if $J(R) \geq 0$ s.t.

$$J(\bar{z}) = \min_{z \in Z} J(z)$$

$$\text{s.t. } g_i(z) \leq 0 \quad \forall i$$

$$h_i(z) = 0 \quad \forall i$$

$$\|z - \bar{z}\| \leq R$$

When does local = global? \Rightarrow QP-MPC

objectives for today:

- 1) when is local = global?
- 2) how to formulate problems where local = global?

II Convexity

Def: 1) A set $Z \subset \mathbb{R}^n$ is convex if

$$\lambda z_1 + (1-\lambda)z_2 \in Z$$

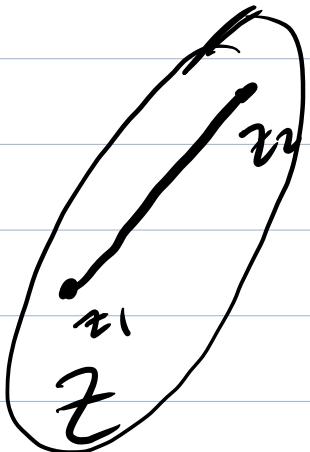
$$\forall z_1, z_2 \in Z, \lambda \in [0, 1].$$

2) A function $f: Z \rightarrow \mathbb{R}$ is convex if Z is convex and

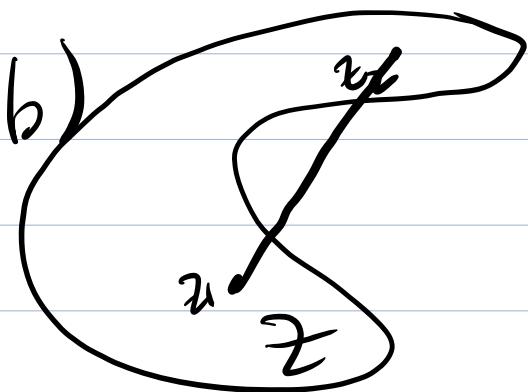
$$f(\lambda z_1 + (1-\lambda)z_2) \leq \lambda f(z_1) + (1-\lambda)f(z_2)$$

$\forall z_1, z_2 \in Z$ and $\lambda \in [0, 1]$.

Ex: a)

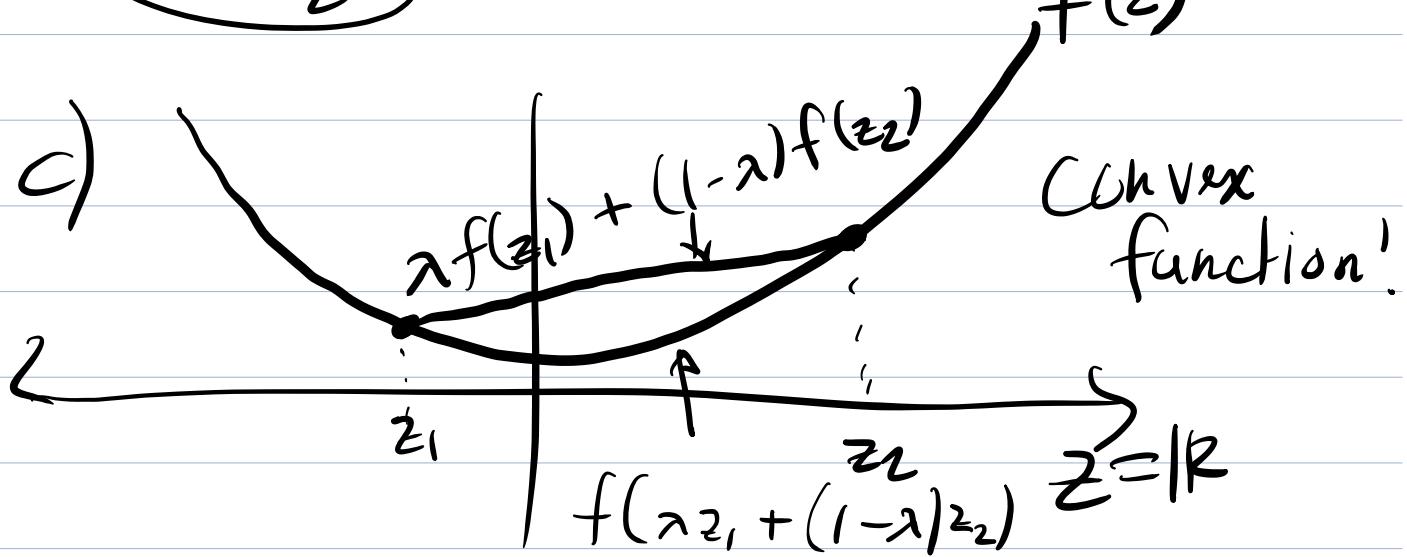


Z is convex



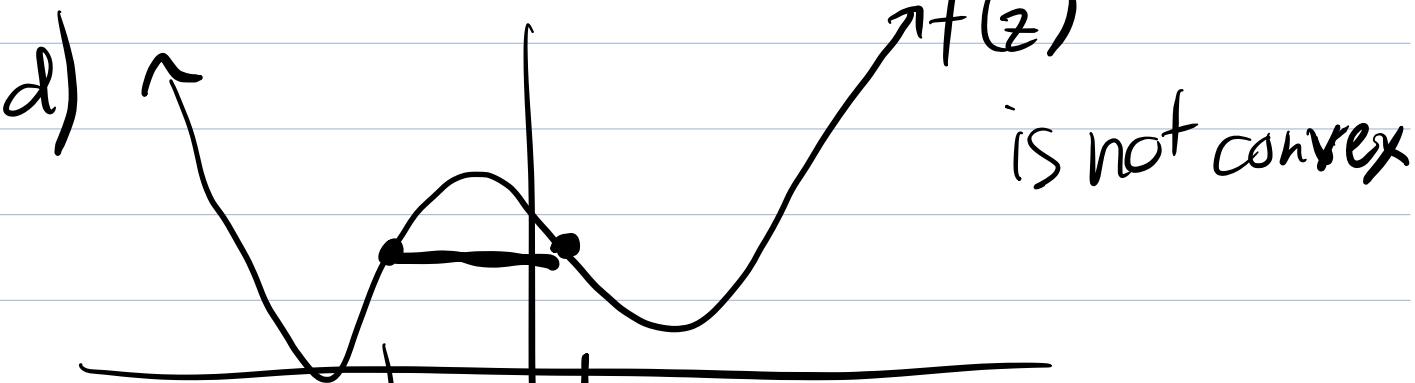
Z is not convex

c)



Convex
function!

d)



is not convex

$$z_1 \quad | \quad z_2$$

$$\mathcal{Z} = \mathbb{R}$$

e) An affine function $f(z) = c'z + d$ is convex

f) A quadratic function $f(z) = z'Qz + 2sz + r$ is convex iff $Q \geq 0$

Theorem: 1) The intersection of an arbitrary number of convex sets is convex.
 2) If f_1, \dots, f_N are convex functions, then $\sum_i \alpha_i f_i$ is a convex function for all $\alpha_i \geq 0$.

Def: A **convex polytope** $P \subset \mathbb{R}^n$ is a set defined as the intersection of a finite number of half spaces:

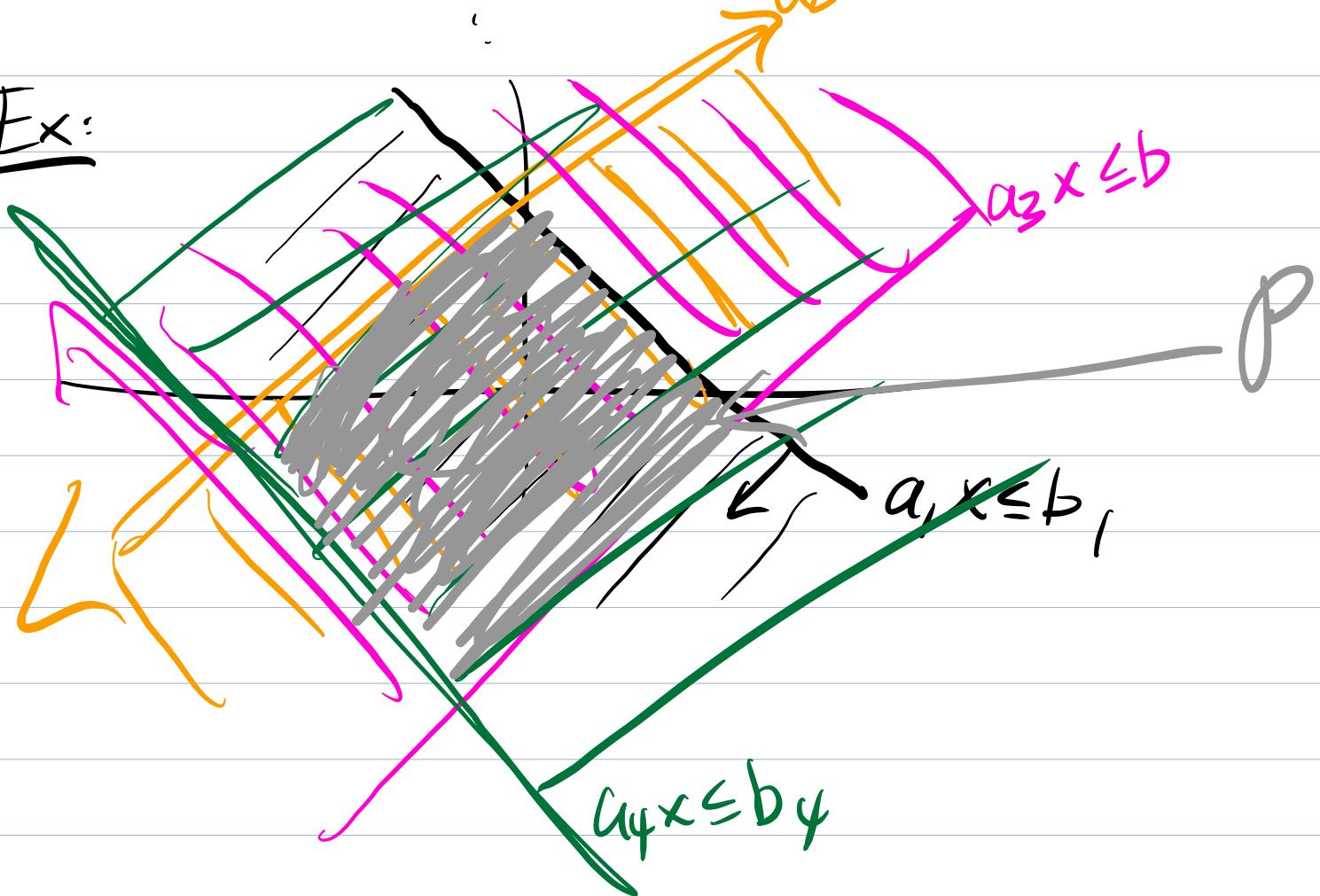
$$P = \{x \in \mathbb{R}^n \mid Ax \leq b\}$$

where $Ax \leq b$ is $a_i x \leq b_i$

$$\begin{bmatrix} - & a_1 & - \\ - & a_2 & - \\ - & a_3 & - \\ \vdots & & \end{bmatrix} \quad \underbrace{\quad}_{a_i x \leq b_i}$$

$$a_i x \leq b_i$$

Ex:



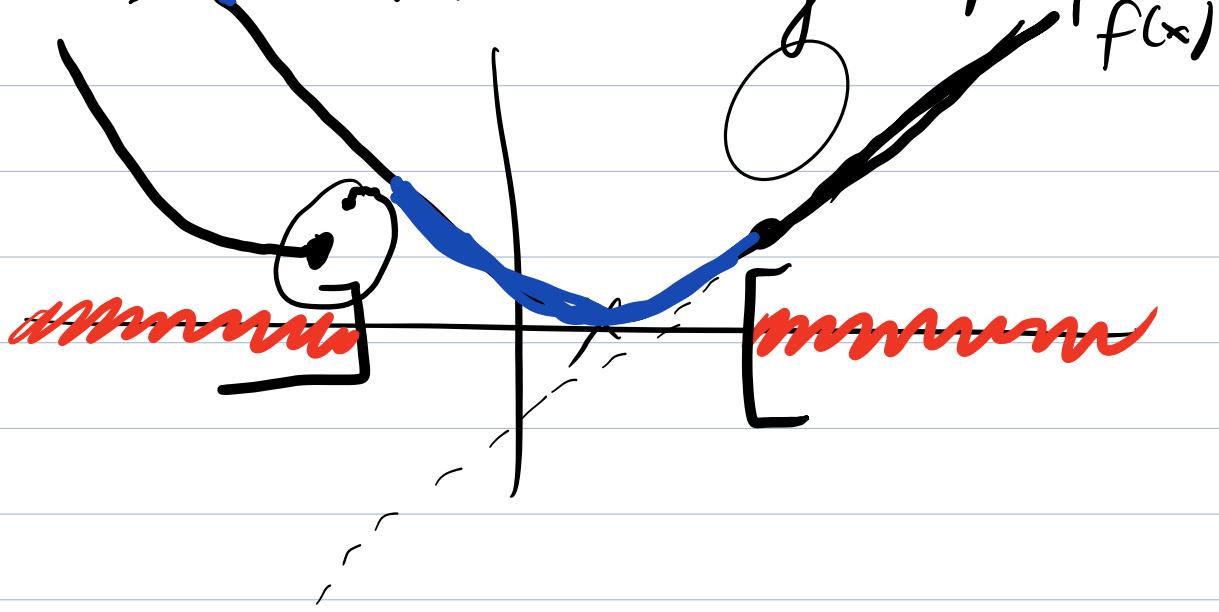
Note: Convex polytopes are convex!

\mathbb{R}^n

III. Convex Optimization

Def: An optimization problem is **Convex** if the objective function is convex and the feasible set is convex.

Thm: In convex optimization problems local minima are global optimizers.



$$f(x, y) = (x + y)^2$$

$$x = -y$$

A. Quadratic Programming

Def: A quadratic program is an optimization problem that can be written as:

$$\min_{z \in \mathbb{R}^n} \frac{1}{2} z' Q z + c' z$$

$$\text{s.t. } Az \leq b$$

where Q is symmetric.

Thm: If $Q \succeq 0$ then the quadratic program is convex.

MATLAB: quadprog

CPLEX, Gurobi, Mosek

IV QP based Model Predictive Control.

Basic Idea:

- 1) Rely on fast speed of QP solvers.
- 2) At each sampling time k , starting at current state, compute an optimal controller at a sequence of time steps over a finite horizon, $[k, k+K]$
- 3) Use the solution for the first time instance $[k, k+1]$
- 4) Repeat step 2 at the new location $k = k + 1$

Note: 1) We have LTV system discretized:

$$x(k+1) = A(k)x(k) + B(k)u(k)$$

2) Let $u_{k+k'}|_k$ be the input at time

step $k+k'$ computed at time step k .

3) Given $\{u_{k+k'}|_k\}_{k=0}^{K-1}$ let $x_{k+k'}|_k$

be the state at time step $k+k'$.

Suppose I have dynamics in ①

and I have constraints such that

$x(k) \in X$ and $u(k) \in U \forall k$

where X, U are convex polytopes

$$\min_{\{u_{k+k'}|_k\}_{k=0}^{K-1}, \{x_{k+k'}|_k\}_{k=0}^{K-1}} \quad (QP)$$

$$\sum_{k'=0}^{K-1} x_{k+k'|k}^T Q(k') x_{k+k'|k} +$$

$$+ \sum_{k'=0}^{K-1} u_{k+k'|k}^T R(k') u_{k+k'|k}$$

$$x_{k+k'+1|k} = A(k+k') x_{k+k'|k} + B(k+k') u_{k+k'|k}$$

$$\forall k' \in \{0, \dots, K-1\}$$

$$A_u u_{k+k'}|_K \leq b_u$$

$$u_{k+k'-1|_K} \in \mathcal{U} \quad \forall k'$$

$$x_{k+k'|_K} \in X \quad \forall k'$$

Algorithm for MPC

- 1) Measure the state $x_{k|_K}$ at time k
- 2) Solve the QP
- 3) Apply $u_{k|_K}$ at time step k
- 4) Wait for new sampling time $k+1$, go to step 1.