

## Lecture #9

### Goals

- 1) Review MPC
- 2) Review Project
- 3) Trajectory Optimization
- 4) Solving Nonlinear Programs
- 5) Collocation
- 6) Free-Time Optimization
- 7) Combining Everything

### III. Trajectory Optimization

Def: A trajectory optimization problem is

$$\min_{u: [0,1] \rightarrow \mathbb{R}^m} \int_0^1 L(t, x(t), u(t)) dt + \phi(x(1))$$

$$\text{Subject to } \begin{aligned} \dot{x}(t) &= f(t, x(t), u(t)), \quad \forall t \in [0, 1] \\ x(0) &= x_0 \end{aligned}$$

$$g_i(x(t)) \leq 0 \quad \forall t \in [0, 1], \forall i \in \{1, \dots, n_{\text{ineq}}\}$$

where  $L: [0, 1] \times \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$  is the running cost and is assumed to be twice continuously differentiable,  $\phi: \mathbb{R}^n \rightarrow \mathbb{R}$  is the final cost and is assumed to be twice continuously differentiable,  $g_i: \mathbb{R}^n \rightarrow \mathbb{R}$  are inequality constraints and are also assumed to be twice continuously differentiable and the dynamics  $f: [0, 1] \times \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}^n$  are twice continuously differentiable.

Note: Several strategies to solve this problem:

1. indirect methods

construct a necessary condition for optimality, then discretize this condition to generate a nonlinear program.

Ex:  $\min_{x \in \mathbb{R}} J(x)$

Necessary condition for optimality is  $\nabla J(x) = 0$   
try to find zeros of  $\nabla J(x)$ .

Unfortunately finding such necessary conditions for optimal control problems w/ constraints is hard.

2. direct methods

discretize optimal control problems  
then solve using nonlinear programming.

(\*) 
$$\min_{\substack{u(\Delta T_k) \\ k=0, \dots, \frac{1}{\Delta T}-1}} \Delta T \sum_{k=0}^{\frac{1}{\Delta T}-1} L(\Delta T_k, x(\Delta T_k), u(\Delta T_k)) + \varphi(x(1))$$

$$x(\Delta T(k+1)) = x(\Delta T_k) + \Delta T f(\Delta T_k, x(\Delta T_k), u(\Delta T_k)) \\ \forall k = \{0, \dots, \frac{1}{\Delta T}-1\}$$

$$x(0) = x_0$$

$$g_i(x(\Delta T_k)) \leq 0 \quad \forall k \in \{0, \dots, \frac{1}{\Delta T}-1\} \\ \forall i \in \{0, \dots, p\}$$

How do we solve this problem?

- Use nonlinear optimization
- Need gradients of the cost and constraints to solve this problem.



## IV. Solving Nonlinear Programs

Recall

$$\begin{aligned} \min_{z \in Z} \quad & J(z) \\ \text{subject to} \quad & g_i(z) \leq 0 \quad \forall i \in \{1, \dots, p\} \\ & h_i(z) = 0 \quad \forall i \in \{1, \dots, p_e\} \end{aligned}$$

Note: 1) the problem in the earlier section (\*) can be rewritten as:

$$z = \begin{bmatrix} u(0) \\ u(\Delta T) \\ \vdots \\ u(1-\Delta T) \\ x(0) \end{bmatrix} \quad \text{so } Z = \mathbb{R}^{\frac{m}{\Delta T} \times n} \times \mathbb{R}^n$$

Notice that

$$x(\Delta T) = x(0) + \Delta T f(\Delta T, x(0), u(0)) = F_1(z)$$

Can similarly define

$$x(k\Delta T) = F_k(z)$$

So the cost function and constraints in (\*) can be written as a function of  $z$ .

2) All nonlinear optimization solvers have similar formats. We focus on `fmincon`:

$$z = \text{fmincon}(\text{cost}, z_0, A, b, A_{eq}, b_{eq}, lb, ub, \text{nonlcon}, \text{options})$$

$\swarrow$  i.c. for optimization     
  $\swarrow$   $Az \leq b$      
  $\swarrow$   $A_{eq}z = b_{eq}$      
  $\swarrow$   $lb \leq z \leq ub$

$$3) [J(z), \nabla J(z)] = \text{cost}(z)$$

$$4) [g, h, \nabla g, \nabla h] = \text{noncon}(z)$$

$$(g(i))(z) = g_i(z) \leq 0 \quad \forall i \in \{1, \dots, p_{ie}\}$$

$$(h(i))(z) = h_i(z) = 0 \quad \forall i \in \{1, \dots, p_e\}$$

$$[\nabla g]_{ij} = \frac{\partial g_i(k)}{\partial z_j}(z)$$

$$[\nabla h]_{ij} = \frac{\partial h_i(k)}{\partial z_j}(z)$$

5) The last argument to `fmincon` is a set of options! make sure to set them as follows:

options = optimoptions('fmincon', ...,  
 'SpecifyConstraintGradient', true, ...,  
 'SpecifyObjectiveGradient', true)

(a) For our optimal control problem,  
 we have to compute gradients of a  
 constraint  $g$ : assume  $p_{ie} = 1$

$$g(1) := g(F_1(z)) = g(x(0) + \Delta T f(0, x(0), u(0)))$$

$$g(2) := g(F_2(z)) = g(x(\Delta T) + \Delta T f(\Delta T, x(\Delta T), u(\Delta T)))$$

$$\text{but } x(\Delta T) = x(0) + \Delta T f(0, x(0), u(0))$$





So  $g(1)$ 's gradient is nonzero only for  $x(0)$  and  $u(0)$  term, but  $g(2)$ 's gradient is nonzero for  $x(0)$ ,  $u(0)$ , and  $u(1)$  and derivative is complex due to application of chain rule.  
Only gets worse as we keep going further along.

#### V Collocation

To avoid this problem of having to write a complex gradient, we will append the states of the system to the decision variables:

$$\min_{\substack{\{u(\Delta T k)\}_{k=0}^{\frac{1}{\Delta T}-1} \\ \{x(\Delta T k)\}_{k=0}^{\frac{1}{\Delta T}-1}}} \Delta T \sum_{k=0}^{\frac{1}{\Delta T}-1} L(\Delta T, x(\Delta T k), u(\Delta T k)) + \varphi(x(1))$$

$$\text{subject to } -x(\Delta T(k+1)) + x(\Delta T k) + \Delta T f(\Delta T, x(\Delta T k), u(\Delta T k)) = 0$$

$$g(x(\Delta T k)) \leq 0 \quad \forall k \in \{0, \dots, \frac{1}{\Delta T}-1\}$$

assume for just this presentation that it is one constraint

In this case, we treat the states and inputs as independent variables that are only linked by the equality constraint:

$$g(2) = g(x(2\Delta T)) \downarrow$$



The gradient of  $g(2)$  is treated as zero except for the term related to  $x(2\Delta T)$ .

- Note:
- 1) Larger # of decision variables
  - 2) sparser much easier to evaluate constraints
  - 3) usually much better performance.

## VI. Free-Time Optimization

What if our goal is to get the car from point 1 to point 2 as quickly as possible?

The final time should not be 1, but a decision variable for us to optimize over.

Theorem: Suppose  $\dot{x}(t) = f(t, x(t), u(t))$  with  $t \in [0, s]$  and  $x(0) = x_0$ .

Suppose we define:

$$(**) \quad \begin{bmatrix} \dot{\tau} \\ \dot{\tilde{x}} \end{bmatrix} = \begin{bmatrix} 0 \\ \tau f(\tau, \tilde{x}(\tau), u(s\tau)) \end{bmatrix}$$

$\forall \tau \in [0, 1]$  and  $\tau(0) = s$ , and  $\tilde{x}(0) = x_0$ .

a) If  $x: [0, s] \rightarrow \mathbb{R}^n$  is a solution to  $f$ , then  $\tilde{x}(\tau) = x(\tau/s) \quad \forall \tau \in [0, 1]$  is a solution to  $(**)$ .

b) Similarly if  $\tilde{x}: [0, 1] \rightarrow \mathbb{R}^n$  is a solution to  $(**)$  then  $x(s\tau) = \tilde{x}(\tau) \quad \forall \tau \in [0, 1]$  is a solution to  $f$ .

Note: 1) By optimizing over this larger state space, one can optimize over the i.c.  $\leq$  to try to find the fastest route between two points:

$$\min \left\{ u(\Delta T_k) \right\}_{k=0}^{\frac{1}{\Delta T}-1} \quad \text{Cost}$$

$$\left\{ \begin{bmatrix} \tau(\Delta T_k) \\ \tilde{x}(\Delta T_k) \end{bmatrix} \right\}_{k=0}^{\frac{1}{\Delta T}-1}$$

Subject to FWD Euler  
state constraints  
 $\tau(0) \geq 0$

2) This transformation is called the free time transformation.

3) Be careful if you have no final cost since  $\tau(0)=0$  is a good solution.

## VII Putting Everything Together.

1) Generate offline solutions using trajectory optimization.

2) Use QP-MPC in real-time to follow trajectory and avoid obstacles, and deal w/ tracking errors.

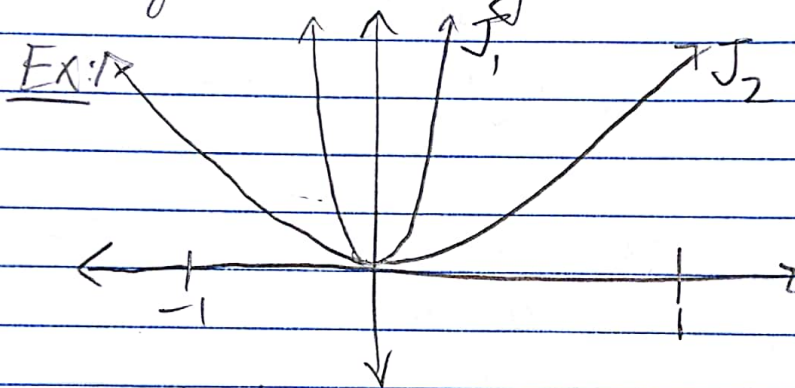


## Issues

a. nonlinear programming is prone to local minima

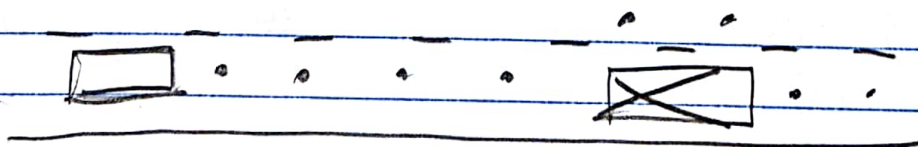
- i) initialize cleverly
- ii) try different representations of costs and constraints to avoid bad minima and slow computation.

b. success of nonlinear programming relies on gradient scaling



Small changes in  $x$  lead to large changes in  $J_1$ , when compared to  $J_2$ , but both have same minima.  $J_1$ 's gradients make it harder to use.

c. Be careful when you discretize:



Solver generates the dot, but the solution may clip the obstacle since it is continuous not discrete, be careful about using time free scaling and making  $\Delta T$  large.