# Three Address Code

**Course Name: Compiler Design** 

**Course Code: CSE331** 

Level:3, Term:3

**Department of Computer Science and Engineering** 

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## Three address code in Compiler

<u>Three address code</u> is a type of intermediate code which is easy to generate and can be easily converted to machine code. It makes use of at most three addresses and one operator to represent an expression and the value computed at each instruction is stored in temporary variable generated by compiler. The compiler decides the order of operation given by three address code.

Converting the expression a / (b + c) into three address code:

$$t_1 = b + c$$
$$t_2 = a / t_1$$

#### Intermediate Code Generation

In the analysis-synthesis model of a compiler, the front end of a compiler translates a source program into an independent intermediate code, then the back end of the compiler uses this intermediate code to generate the target code (which can be understood by the machine).

#### **EXERCISE:**

Three address code for the following expression: a + b \* c - d / (b \* c)

$$t_1 = b * c$$

$$t_2 = a + t_1$$

$$t_3 = b * c$$

$$t_4 = d / t_3$$
  
 $t_5 = t_2 - t_4$ 

$$\mathsf{t}_5 = \mathsf{t}_2 - \mathsf{t}_4$$

# Implementation of Three Address Code

There are 3 representations of three address code:

- 1. Quadruple
- 2. Triples
- 3. Indirect Triples

### For the following expression: a + b \* c - d / (b \* c)

#### Quadruples

Three Address Code:

$$t_1 = b * c$$

$$t_2 = a + t_1$$

$$t_3 = b * c$$

$$t_4 = d / t_3$$

$$\mathsf{t}_5 = \mathsf{t}_2 - \mathsf{t}_4$$

#	ор	arg1	arg2	result
0	*	b	С	t <sub>1</sub>
1	+	а	$t_1$	t <sub>2</sub>
2	*	b	С	t <sub>3</sub>
3	/	d	t <sub>3</sub>	t <sub>4</sub>
4	-	t <sub>2</sub>	t <sub>4</sub>	t <sub>5</sub>

It is structure with consist of 4 fields namely op, arg1, arg2 and result. op denotes the operator and arg1 and arg2 denotes the two operands and result is used to store the result of the expression. Easy to rearrange code for global optimization. One can quickly access value of temporary variables using symbol table. Contain lot of temporaries. Temporary variable creation increases time and space complexity.

### For the following expression: a + b \* c - d / (b \* c)

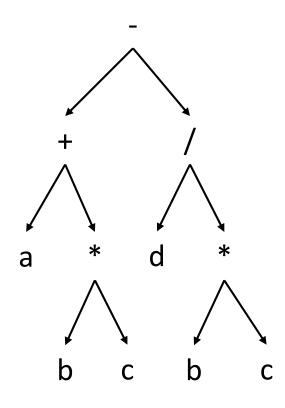
Quadruples				Triples						
	#	ор	arg1	arg2	result		#	ор	arg1	arg2
	0	*	b	С	t <sub>1</sub>		0	*	b	С
	1	+	а	$t_1$	t <sub>2</sub>		1	+	а	(0)
	2	*	b	С	t <sub>3</sub>		2	*	b	С
	3	/	d	t <sub>3</sub>	t <sub>4</sub>		3	/	d	(2)
	4	-	t <sub>2</sub>	t <sub>4</sub>	t <sub>5</sub>		4	-	(1)	(3)

This representation doesn't make use of extra temporary variable to represent a single operation instead when a reference to another triple's value is needed, a pointer to that triple is used. So, it consist of only three fields namely op, arg1 and arg2. Temporaries are implicit and difficult to rearrange code.

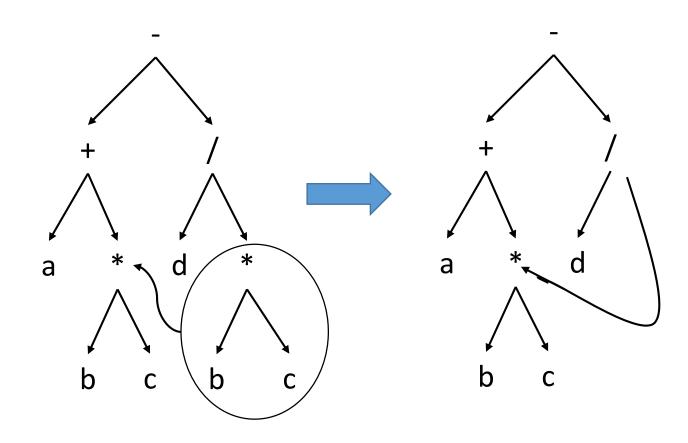
#### For the following expression: a + b \* c - d / (b \* c)

#### **Triples Indirect Triples** arg2 **Statement** arg1 op arg1 op 0 0 b 0 21 b С C (0) (0)22 + a a 2 b 2 b 23 \* С C (2) d (2) 3 d 3 24 3 (1) (3) (1) (3) 25 4 4 4

This representation makes use of pointer to the listing of all references to computations which is made separately and stored. Its similar in utility as compared to quadruple representation but requires less space than it. Temporaries are implicit and easier to rearrange code.



Syntax directed tree



Directed Acyclic Graph (DAG)

# THANK YOU