

#### **UNIT-I**

Cryptography & Network Security (Jawaharlal Nehru Technological University, Hyderabad)



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#### Network Security

#### Introduction:-

computer data often travels from one computer to another, leaving the safety of its protected physical someundings. Once the data is out of hand, people with bad intension could modify or forge your data for their own benefit.

The OSI Security and tecture provides a systematic frame work for defining security attacks, mechanisms and services.

Computer Security: - generic name for the collection of tools designed to protect data and to thwart (prevent) hackers.

Network Security: - measures to protect data during their transmission. Ex. Firewalls.

Internet security: - measures to protect data during their townsmission over a collection of interconnected networks.

Ex: Bank transactions.

To askess the Security needs of an organization effectively, the manager responsible for security needs some systematic way of defining the requirements for security and characterization of approaches to satisfy. These requirements.

one approach is to consider three aspects of information security, The ITU recommendation X. 800, Security Architecture for OSI. It focuses on the following. Security attack: -

Any action that comprises the security of information owned by an organization.

security mechanism: -

A mechanism that is designed to detect, prevent or recover form a security attack.

Security Service: -

A service that embances the security of the data processing systems and the information townsfer of an organi zation.

## Basic Concepts: -

Coyptography: -

To provide the security and protect the valuable information we can use cryptography.

The aut of protecting the information by tournsfor - ming it into an unreadable format is called cryptography. (intelligible message into one that is unintelligible).

plaintext: The original intelligible message.

Ciphertext:- The towns formed me seage/encoypted text.

Cipher: - An algorithm for townsforming an intelligible message into one that is unintelligible by teamsposition

(2)

Key: - A string of bits used by a coyptographic algorithm, known only to the sender of receiver.

Encipher (encode): - The process of converting plaintent to expheritext using a cipher and a key.

Decipher (decode): - the process of converting cipher text back into plaintext using a cipher and a key.

comptarealy & s: -

The study of principles and methods of transforming an unintelligible message back into an intelligible message without knowledge of the key. Also called Code breaking. \* The basic intention of an attacker is to break a comptosystem & to-find the P.T form The C.T.

Security Attacks:-

our goals of security - confidentiality, integrity and available availability - can be threatened by security attacks. They are classify into two types. I passive attacks 2) Active attacke.

Security attacks passive & snooping attacks. E Traffic analysis.

Threat to confidentiality. ( To protect The data form unauthorised users, it is a

modification Denial of Sewice Masquerading Threat to availabil Sattacke - lity. Replaying.

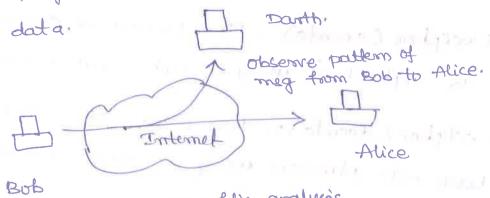
(The data, New resource Repudiation are continuously available to legitimate users).

Threat to

integrity. It means part of New security) This document is available free of charge on the studocu be done only be authorized Downloaded by Saharsh Wadekar (sorhiwadekar 2012@gmail.com) thon 3ed mechanisms).

O In general, a types of attacks threaden the confidentiality of information: snooping & Toaffic analysis

1 (a) \* Smooping refers to unauthorized access to as interception of data. Doubth.



a) Toathic analysis.

\* passive attacks are very distributed to detect there are they don't involve any attention of data.

The integraty of data can be threatened by several kinds of

attacks in of megs:

Molification of megs:

\* Abter accessing information, the attacker modifies the

\* Abter accessing information to herself. Es:- bank

information to make it beneficial to herself. Es:- bank

toansactions.

- \* Masqueroading: The attacker impersonates sombody else.

  ex: an attacker might steal the bank cound and PIN

  of a bank customer and pretend that she is that customer
- \* Replaying: The attacker obtains a copy of a meg sent by a user and later tries to replay it.
- (3) \* Denial of Service: Pt & a very common attack. It may slow down or totally intempt the Service of a system.

  Slow down or totally intempt the Service of a system.

  (Ch) Traffic analysis refers to obtaining some other type of information by monitoring online teaffic.

3	

Attacks	passive/Active	Threatening	
Snooping Trabbic Analysis.	passive	confidentiality	
Modification Masqueroading Replaying. Repudiation	Active	Integrity.	
Denial of Service	Active	Availability.	

\* In a passive attacks, the attacker's goal is fust to obtain information. i.e the attack does not modify data or harm the system.

\* An active attacks may change the data or harm the system. Active attacks are normally easier to detect than to prevent, because an attacker can launch them in a variety of ways.

| Don'th |

Bob Endernet or Alice

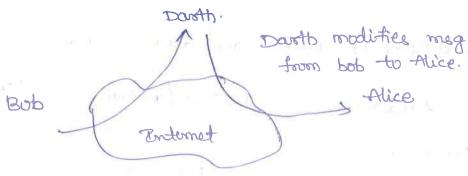
a) Masquerade.

Bob

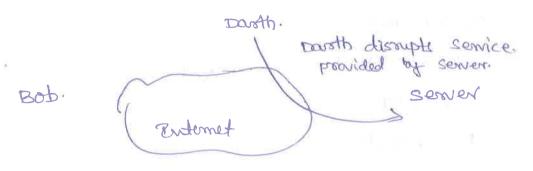
Bob

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modification of messages.



d) Denial of Service.

## Services and Mechanisms:-

The International relecommunication union-Telecommy - nication standardization Sector (ITU-T) provides some security semices and mechanisms.

## Security Services:-

The classification of Security services are as

- follows:

I confidentiality: - Ensures that the information in a computer system and transmitted information are acceptible for reading by authorised parties.

Eq: printing, displaying and other forms of disclousine

is the protection of toansmitted data from passive attacks.

Authentication: - Ensures that the origin of a message or electronic document is correctly identified, with an assurance that the identity is not false.

Data Integrity: - Ensures that only authorized parties are able to modify computer system assets and towns mitted information.

Non repudiation: Requires that neither the sender nor the receiver of a message be able to demy the toansmission.

Access control: - Requires that access to information resources may be controlled by or the target system.

Availability: - Requires that computer system assets be available to authorised paroties when needed.

# security services (X.800):-

- D Authentication: The assurance that the communicating entity is the one that it claims to be.
  - \* peers Entity Authentication: Used in association with a logical connection to provide confidence in the identity of the entities connected.
  - \* Data origin Authentication: On a connectionly transfer provides assurance that The Source of received data & as claimed.
- 2) Access control: the prevention of unauthorized use of a resource.

- 3) Data confidentiality: The protection of data from unauthorized disclosure.
  - \* connection confidentiality:
    The protection of all user data on a connection.
  - \* Connectionless confidentiality:
  - \* Selective field confidentiality!

    The confidentiality of, selected fields within the uses data on a connection or in a single data block.

# \* Traffic flow confidentiality: The protection of the information that might be derived from observation of teather flows.

The assurance that data received are exactly as sent by an authorised entity. i.e no modification, insertion, deletion or replay.

\* Connection Integrity with Recovery

provides for the integrity of all user data on a

connection and detects any modification, insertion dele

tion, or replay of any data within an entire data

sequence, with recovery otherspeed.

\* connection Enlegity without Recovery!

As above, but provides only deletion without recovery.

- \* Selective Field Connection Integrity.
- \* connectionless Integrity
- \* Selective Field Connectionless Integrity.

## 5 Nonrepudiation

provides protection against derial by one of the entities in a communication of having participated in all or part of the communication.

Nonrepudiation, origin! - proof that the meg was sent by
The specified party

Nonrepudiation, Destination: - proof that The meg was received by The specified party.

# Security Mechanisms: (X.800)

ITU-T(X.800) also recommends some security rechanisms to provide the security services defined above. (i.e OST security services).

1. Encipherment: - The use of mathematical algorithms to toans-form data into a form that is not redictly intelligible toans-formation of subsequent recovery of the data the toans-formation of subsequent recovery of the data depend on an algorithm and zero or more encryption keys.

2. Digital signature:-

A digital signature is a means by which the sender can electronically sign the data and the oeceiver can electronically verify the signature.

3. Accels control: This document is available free of charge on



Access control: - A variety of mechanisms that emforce access rights to resources.

Data Integrity! (Accuracy & consistency)

A variety of mechanisms used to assure the indegrity of a data unit or stream of data units.

Authentication Exchange: - A mechanism intended to ensure the identity of an entity by means of information exchange.

Traffic padding: The insertion of bits into gaps in a data

Routing control: - Enables selection of particular physically secure routes for certain date & allows routing changes

Notarization: The use of a tousted 1 party to assure certain properties of a data exchange.

Relation to Security Services and Mechanisms.

Delation	An Security	2001100	-43				
Service	Encipherment	Digital	Accept	70 harris 6	Authordi action exchange	postling	RC N
peers Entity Authoritied	tion Y	4			Y		
Data origin	4	4		19-			
Access control			Y				
confedentiality	<b>Y</b>					Y	
Trabbic flow	Y		n Ling			YY	
Dada Integrity	Y	Y		У			
Nonrepudiation		Y		Y			У.
Availability.				У	У		74

A mag is to be townsferred from one party to another across some sort of indemet The 2 parties, who are the principals in this transaction, must cooperate for the exchange to take place.

A logical information channel is established by defining a route through internet from Source to destination and by the cooperative use of communication protocols by and by the two principals.

Using this model requires us to:

- \* design a suitable algorithm for security transformation.
- \* generate the secret information (keys) used by the algorithm.
- \* develop methods to distribute & share the secret information
- \* specify a protocol enabling the principals to use the toans-formation and secret information for a security service.

# Network Access Security model:

Management	Info. System,
Access Chamber	Computing resource  Computing resource  Computing resource  2/0)  Data  processes  sper 8/10  unction Enternal security  controls.

- \* select appropriate gatekeeper functions to identify users.
- \* Emplement security controls to ensure only authorized users access designated information or resources.
- \* Trusted computer systems can be used to implement

  This model.

Conventional Encryption: - (Symmetric Cipher Model) The Sender & Re'r.

It is a comptographic system that uses the same key used by the sender & Re'r.

A symmetric encryption scheme has five ingredients.

plaintext: - This is the original intelligible message or data That is fed into the algorithms as imput.

Encryption algorithm: - It performs various substitutions and transformations on The plaintext.

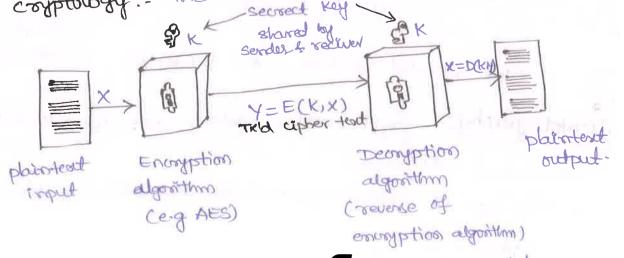
Secreet Key: - It is also imput to the encoyption algorithm. The key is a value independent of the plaintext and of the

Ciphertext: This is the scrambled may produced as ordpu It depends on the plaintext and the secret key. It is an apparently random stream of data (i.e unitellique).

Decryption algorithm: - This is essentially the encryption algorithm own in severse. It takes, expertent & the secret key and produces the original plaintent.

cryptamalysis (ande breaking):- the study of principle method of deciphering ciphers text without knowing key.

corptology: - The field of both corptography and crapptonalysis secrect key



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Two requirements of secure use of symmetric enoryption:

A strong encryption algorithm

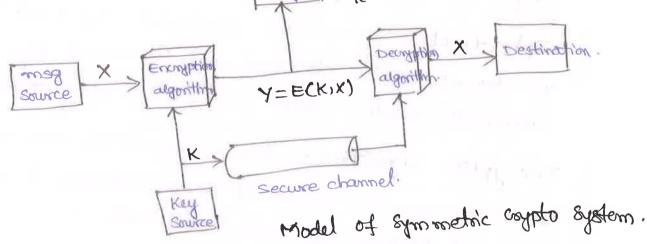
2) A secret key only known to sender Receiver.

A source produces a message in plain-terry,

X=[X1, X2, -.. XM]. The M elements of X are in some finite

alphabet.

for encryption, a key of the form K=[KI, Kz, ... KJ] is generated. Alternatively, a third party could generate The Key and securely deliver it to both source and destination enyptana > x



with the mescage X and the encryption key K as imput, The encryption algorithm forms The Ciphertext Y = [Y1, Y2, - . YN].

Advantages: .: Y = E(K, X).

The intended receiver X = D(K, Y). 2. uses fewer computers

cryptography: - cryptographic systemy are characterized along 3 independent dimensions.

) The type of operations used for toans-forming plain tent to Cipher text.

(8)

- 2) The number of kuys used.
- 3) The way in which the plaintout is processed.

# classical encryption Techniques:

There are two basic building blocks of all enoughtion techniques: 1) substitution 2) Transposition.

substitution Techniques!

In which the letters of plaintent are replaced by other letters or numbers or symbols. But The plain text is viewed as a sequence of toits, then substitution involves replacing plaintest but patterns with Cipher test but patterns.

) caleson Cipher: (shift Cipher):

which involves replacing each letter of the alphabet with the letter standing 3 places further down the alphabet.

Ex: - plain: meet me after the toga party CEPher: - PHHW PH DIWHU WKH WRID SDUMB.

Cipher: D E F G .... Z A B C.

with numerical equivalent:

c d 23 24 25 0

Then the algorithm can be expressed as

- follows

each plain-terest letter "P", substitute the This document is available free of charge on

Cipher text letter. C.

C= E(3, p) = (P+3) mod 26.

The general caesar algorithm is:

C= E(K) P) = (P+K) mod 26. K = 1 -to 25.

The decoyption algorithm is

p = DCK,c) = CC-K) mod 26.

The best known multiple-letter encoyption Cipher is the play fair, which treats diagrams in the plaintent as single units and toundates these units into ciphertext diagrams.

The playfeir algorithm is based on the use of 5x5 motions of letters, constructed using a keyword.

Ex: Let The Keywood be "MONARCHY" The matrix is constructed by filling in the letters of Kywood (minus duplicates) from lebt to right and from top to bottom, and then filling in the remainder of the matrix with the remaining Rules for E letters on apprabetical order.

Hello 

3 ball000 ba 11 00 n balk to on

94	0	N	A	R
c	H	Y	B	D
E	F	Gn	<b>z</b>  1	K
L	P	0	S	T
U	V	W	X	2

1. Digrams

2. Ropeating latters - Filter letter

3. Same Column III wrap

4. Same row 1> wrap around

I. Rectangle 1 5 Swap

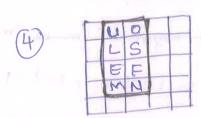
En P.T - attack Digrame: at ta CK

- 1 Repeating plaintent letters that are in the same pair are seperated with a filler letter, such as x' En:- balloon would be treated as balx loom!
- 1 Two plaintent letters that fall in the same now of the matria are each replaced by the letter to the right, with The first element of the row circularly tollowing the

Ex:- AR is enempted as! RM

Strength of playfeir Cipher: since there areas letter, 26 × 26 = 676 diagrameare possible, so identification of individual diagramues more ditti cut.

H&I are in same column, hence take little below H Them to replace: HI > OC.



M&O -> NU. (opposite corner).

Ext. plaintext: instruments. Keyword: morand After split: in st ru me nt sz (:T)- ga th mz ch rg tx.

Monoalphabetic and Polyalphabetic Cipher: -

by Et & a substitution Cipher in which for a given key, The Cipher alphabet for each plain alphabet is fixed Through - out the encryption process.

for ex:-, "et A' is enoughted as 'D': for any number of occurrence in that plaintext; A' will always get enoughed to is

tolyalphabetic Cipher is a substitution Cipher in which the Ciphers alphabet fire the plain alphabet may be different at different places during The enoughtion process. en:- play-fair & vigenere cipher are polyaphabetic.

Vigenere Cipher:

This Scheme of Cipher uses a text string (a word). as a key, which is then used for doing a number of shibts Key: deceptive deceptive do captive

on the plaintext. 3 4 2 4 15 19 8 21 4 22 4 0 17 4 3 8 18 2

P.T! - we are discovered save your selb C.T :- ZICYTWANGRZGYTW

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Ex: let us assume the key is 'point'. such alphabed **(**) of the key is converted to its respective numeric value i.e. p -> 16,0 -> 15, 1 -> 9, n -> 14, t -> 20.

Thus The Key is: 16 15 9 14 20.

C1 = CPi + K1 mod m) mod 26 Pi = (Ci-Ki mod m) mod 2

Vernam Cipher: - This works on binary data rather than letters. En this a conjetanalysis is choose a keywood that's as long as the plaintext & has no steestical relation shipton it. \* It is introduced by an ATST enginees named Gilbert Vernam 90 1918.

It can be expressed  $Ci = Pi \oplus Ki$ 

Pi = "th binary digit of plaintext

Ki = ith

ciphertext. Ci = ith

Pi = CI & Ki.

Key steers generator coxptographic bit stream (Ki) plain -> Ciphertext (Ci) teset (Pi)

## One-time Padi-

It is an unbreakable crypto system. It represents the message as a sequence of o's and 1's. This can be accomplished by writing all numbers in binary. The key is a rambom sequence of o's and 1's of same length as the mostage.

\* once a key is used, it is discarded and never used again.

each new may requires a new key of the same

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Ci = Pi +Ki

## Transposition Techniques:

Au The techniques studied so for Involve the substitution of a ciphortext symbol for a plaintext symbol.

A very different kind of mapping is achieved by persforming some sost of permutation on the plaintent letters. This is referred to as a transposition Cipher.

#### voil fence:

It is the simplest of such appear, in which the Platertost & written down as a sequence of diagonals and There read off as a sequence of rows.

Eu:- plain text: - meet me abter the toga pairty".

we write The following:

me matoht & Poy et e f e t e o a a t

The encypted mag is:-

MEMATRHT GPRY ETEFE TEOAAT

Row Transposition ciphera:

3 1 2 5 6 4 Key :naapt 4 Imput: F t suo a o M wcoixk y pe

ONLY DOWN DOWN DOWN SAMANSH WARLEY (SOM WHO DE PAXTY OKZ

A plaintext meg may be hidden in one of two ways. The methods of steganography conceal the existence of the meg. It is time-consuming to constant.

Various techniques!

- 1) character marking: selected letters of priviled or typewritten text are overwritten in pencil. The mounts are ordinality not visible unless the paper is held at angle to bright light.
- 2) Envisible Ink: no visible trace until heat or some chemical is applied to The paper
- 9 Pin punctures: Small pin punctures on selected letters
- 4) Typewrites correction vibbon: Doawbacks: - requires a lot of overhead to hide a relatively few bits of information.

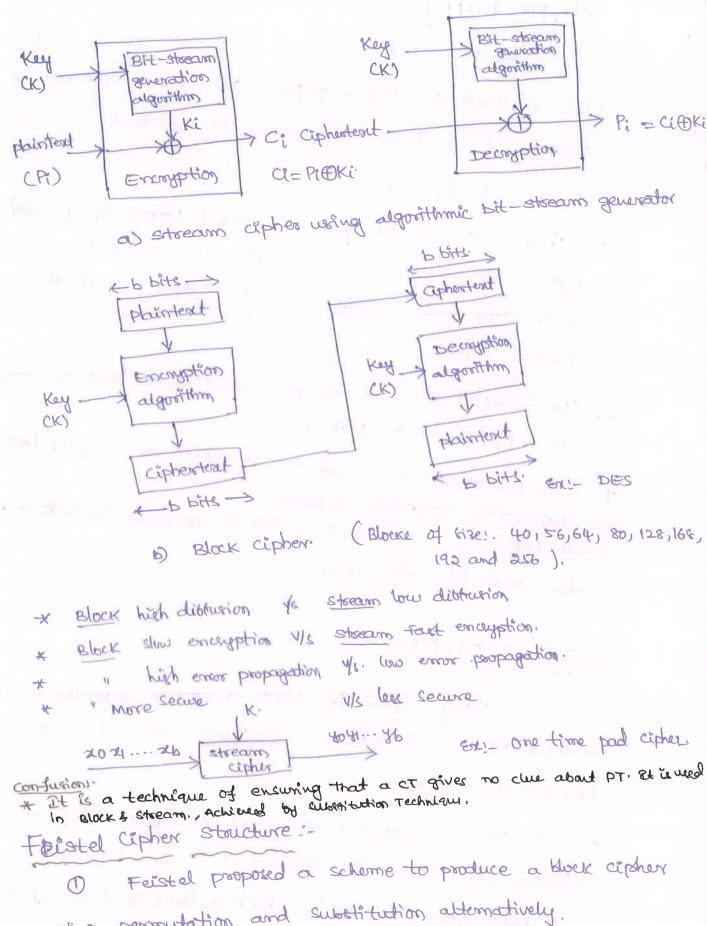
# BLOCK CIPHERS :-

Many Symmetric block encryption algorithms in current use are based on a stoucture referred to as a Feistel block Cipher.

A stream cipher is one that encrypts a digital data stream one bit us one byte at a time. Ex: streaming of date

A block cipher is one in which a block of plaintent is treated as a whole and used to produce a caphertext block of equal length. typically, a block size of 64 or 128 bits & used.

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using permutation and substitution alternatively.

permutation:

A sequence of plaintent elements is replaced by a permutation of That Sequence. (change The position of

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#### substitution:-

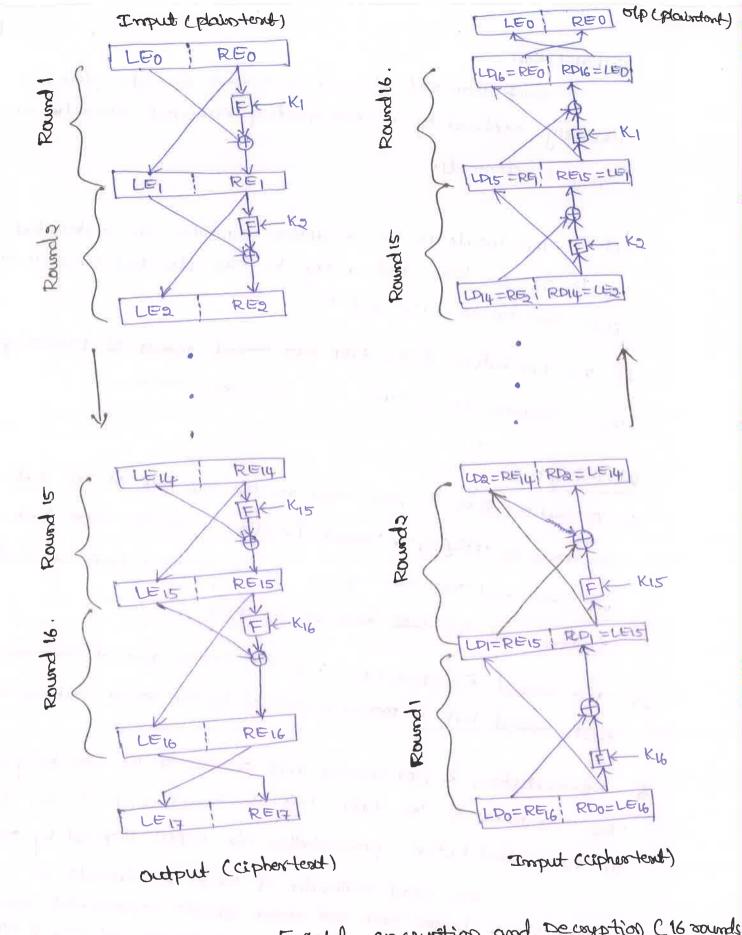
Each plaintext element or group of elements (block) is uniquely replaced by a corre spooding ciphextext elements or Sund of epowents.

- The imputs to the encryption algorithm are a plain tend bluck of length aw bits and a key Ki. The plaintent block is divided into two halves, LED and RED.
- 3) The two halves of the data pass through rounds of processing and Then combine to produce the ciphertest block.

#### Working:

- ) A substitution is performed on the lebt halt of the data. This is done by applying a sound function F to the right halb of The data and then taking the Ex-OR of the output of that function and the lebt halt of the data.
- 2) The round function (F) has the same general structure for each sound but is program eten sed by the sound Subkey Ki.
- permutation is performed that consists of the interchange of The two halves of The data. This stoucture is a particular form of the substitution - permutation New (SPN) proposed by stammon The exact realization of feisted new depends on
  - D Block Size: Larger block 8ize mean greater security, but speed V. Key 8ize: - lareges Key size means greater security, but may & E/D spec

  - 3) No. of sounds: A typical size of 16 sounds.
  - 4) Sub key generation algorithm: Greater complexity in this lead to greater distincity of cryptanalysis
  - 5) Ease of This documents available free of charge on Studocu greater complexity can make Downloaded by Saharsh Wadekar (somuwadekar 2002@gmail.com)



-fig:- Feistel encuption and Decoyption C16 sounds Dibbusion: - dissipales statistical starcture of plaintent over buck of citherten It is used in stock exphase method. It is achieved by permutation confusion: - makes relationship to ciphertout and key as complex as Increases the redundancy of the P.T by spreading it across rows and 6

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# Data Encryption standard (DES):-



- -) It is a symmetric-key algorithm using block-by-block encryption (Each block is encrypted individually and they are later changed to 64 lost later converted to -Somat foral cipher tent)
- -> Block size is 64 bits and key size 1 & 56 bits. -> No. of sub Keys -16, sub Key 8i8e \_ 48 bit.

  > It follows the Feistel Cipher structure.

- -> The algorithm towns-forms 64-bit imput in a series of steps noto a 64-bit output. The same steps, with the same key, are used to reverse the encouption.

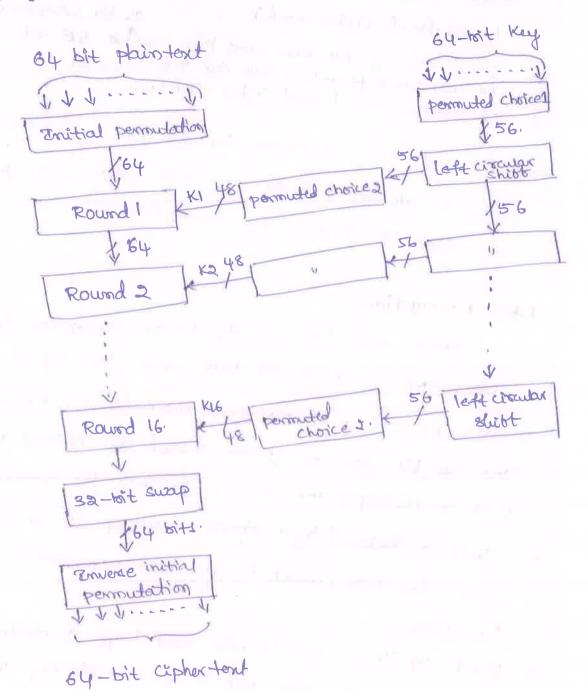
### DES Enoughtion:-

These are two imputs to the encryption function. The plaintend must be 64 bits and The Key & 56 bits in length.

- I The processing of plaintext proceeds in Three phases, first The 64 bit plaintend passes through an Emitial permutation (IP) That reamonges the bits to produce the permuted imput.
- 2) This is followed by a phase consisting of 16 rounds of the Same function, which involves both permutation and substitution
- 3) Each of These rounds will need keys. Enittally we take a 56-bit eggs cipher key but it is a single key, we pass it on to a Round-key generators, which generates 16 ditterent keys for each single round.
- 4) These keys are passed on through the rounds as 4's-tother when passing through all these rounds, we reach round 16. By The final key is passed on through the sound key generator &

we get a final permutation.

3) In the final permutation the rounds are swapped we get a final ciphertext.



General Depiction of DES Encryption algorithm

DES Decryption:-

As with any feistel cipher, decryption was the same algorithm as encryption, except took the application of the Sub Key & Teversed.

#### Future of DES:-

- Replaced by AES in 2002 as the world standard for emoryption.
- 56-bit key size easily broken by new generation computers.
- withdrawn of support for official purpose in acros.
- Triple DES still allowed for impostant data till 2030.

DES - one round. 64 bit imput Right 32 bit (data) Left 3abit 48 bot Key New Right 32 bit New left 32 bit.

Function: 32 bit imput Expansion box.

or s-bork substitution bon 132 WE permutation box

sa bit output.

Strength of DES:-

1) The use of 56-bit Keys!.

with a keylength of 56 bits, there are a 56 = 7x106

possible keys. Thus on the face of it, a boute-force appears impractical.

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## The Nedwer of DES algorithm:

- Another concern is their possibility tood cryptanalysis is possible by exploiting the characteristics of the DES algorithm.
- 2) The focus of concern Substitution tables, or S-boxes, that are used in each rteration.

# Timing attacks :-

A timing attack is one in which information about the Key or the plaintext is obtained by obsening how long it takes a given impermentation to person decryptions on various ciphertont.

\* DES's registant to these type of timing attacks.

# Block Cipher Design principles:-

## 1) Number of Rounds! -

- \* The greater the no. of sounds, the more dibbicult It is to perform cryptanalysis.
- \* more no of rounds structown the Cipher per-formance.
- \* Typically 16 sounds are used in FID.

# 2) Design of Function F:-

- \* The function F of the block clipher must be designed such that It must be impossible for any cryptanalysis to unscrambe the substitution.
- \* The criterion that strongthens the function F is it non-lineality.

- (3)
- \* more the function F & nonlinear, more it would be distinut to crock it.
- \* while designing the function F it should be confirmed that it has a good awakenche property, which stodes mad a change in one - bit of imput must reflect the change is many bits of ordput.

3) ky schedule Algorithm:

- \* The key is used to generate one subkey for each round.
  - \* It is suggested that the key schedule should confirm the strict avalanche ebbect and bit independence criterion.