

MA3105: Numerical Analysis

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Consider the following code block, and suppose that the commented c_i 's represent the time cost of executing the corresponding statement.

```
bool unique (int a[]) {
    for (i = 0; i <= n - 2; i++)           // c_1
        for (j = i + 1; i <= n - 1; j++)   // c_2
            if (a[i] == a[j])              // c_3
                return false;
    return true;
}
```

It is clear that the first **for** statement is executed n times; the interior is executed $n - 1$ times. Similarly, the second **for** statement is executed $n - i$ times, with the interior executed $n - i - 1$ times. The **if** block is executed the full $n - i - 1$ times every loop at worst, 1 time at best. Thus, the total cost can be estimated as

$$\begin{aligned}
 T(n) &= \sum_{i=0}^{n-1} c_1 + \sum_{i=0}^{n-2} \sum_{j=i+1}^n c_2 + \sum_{i=0}^{n-2} \sum_{j=i+1}^{n-1} c_3 \\
 &= c_1 n + \sum_{i=0}^{n-2} c_2 (n - i) + \sum_{i=0}^{n-2} c_3 (n - i - 1) \\
 &= c_1 n + (c_2 + c_3) \sum_{i=0}^{n-2} (n - i) - \sum_{i=0}^{n-2} c_3 \\
 &= c_1 n + (c_2 + c_3) n(n - 1) - \frac{1}{2} (c_2 + c_3) (n - 1)(n - 2) - c_3 (n - 1).
 \end{aligned}$$

By expanding further and collecting coefficients, we have

$$\begin{aligned}
 T(n) &= c_1 n + (c_2 + c_3) n^2 - (c_2 + c_3) n - \frac{1}{2} (c_2 + c_3) n^2 + \frac{3}{2} (c_2 + c_3) n - (c_2 + c_3) - c_3 n + c_3 \\
 &= \frac{1}{2} (c_2 + c_3) n^2 + \frac{1}{2} (2c_1 + c_2 - c_3) n - c_2.
 \end{aligned}$$

As $n \rightarrow \infty$, we see that $T(n)$ grows quadratically. We write $T(n) \in O(n^2)$.

Lemma. *Given any polynomial*

$$f(n) = a_k n^k + a_{k-1} n^{k-1} + \cdots + a_1 n + a_0$$

with at least one non-zero coefficient a_i , we can write $f(n) \in O(n^k)$.

Proof. Set $M = |a_k| + |a_{k-1}| + \cdots + |a_0|$. Then for all $n \geq 1$, the triangle inequality gives

$$\begin{aligned}
 |f(n)| &= |a_k n^k + a_{k-1} n^{k-1} + \cdots + a_1 n + a_0| \\
 &\leq |a_k| n^k + |a_{k-1}| n^{k-1} + \cdots + |a_1| n + |a_0| \\
 &\leq |a_k| n^k + |a_{k-1}| n^k + \cdots + |a_1| n^k + |a_0| n^k \\
 &= M n^k.
 \end{aligned}$$

□