MA3105

Numerical Analysis

Autumn 2021

Satvik Saha
19MS154

*Indian Institute of Science Education and Research, Kolkata,
Mohanpur, West Bengal, 741246, India.*

Contents

1 Time complexity

1.1 Runtime cost

When designing or implementing an algorithm, we care about its efficiency – both in terms of execution time, and the use of resources. This gives us a rough way of comparing two algorithms. However, such metrics are architecture and language dependent; different machines, or the same program implemented in different programming languages, may consume different amounts of time or resources while executing the same algorithm. Thus, we seek a way of measuring the 'cost' in time for a given algorithm.

For example, we may look at each statement in a program, and associate a cost c_i with each of them. Consider the following statements.

```
one = 1;                    // c_1
two = 2;                    // c_2
three = 3;                  // c_3
```

The total cost of running these statements can be calculated as $T = c_1 + c_2 + c_3$, simply by adding up the cost of each statement. Similarly, consider the following loop construct.

```
sum = 0;                          // c_1
for (i = 0; i < n; i++)           // c_2
    sum += a[i];                  // c_3
```

The total cost can be shown to be $T(n) = c_1 + c_2(n+1) + c_3 n$; this time, we must take into account the number of times a given statement is executed. Note that this is linear. Another example is as follows.

```
sum = 0;                          // c_1
for (i = 0; i < n; i++)           // c_2
    for (j = 0; j < n; j++)       // c_2
        sum += a[i][j];           // c_4
```

The total cost can be shown to be $T(n) = c_1 + c_2(n+1) + c_3 n(n+1) + c_4 n^2$. Note that this is quadratic. Finally, consider the following recursive call.

```
int factorial (int n) {           // c_1
    if (n == 0)                   // c_2
        return 1;                 // c_3
    return n * factorial(n - 1);  // c_4
}

f = factorial(n);                 // c_5
```

The cost can be shown to be $T(n) = c_5 + (c_1 + c_2)(n+1) + c_3 + c_4 n$. This turns out to be linear.

In all these cases, we care about our total cost as a function of the input size n. Moreover, we are interested mostly in the *growth* of our total cost; as our input size grows, the total cost can often be compared with some simple function of n. Thus, we can classify our cost functions in terms of their asymptotic growths.

1.2 Asymptotic growth

Definition 1.1. The set $O(g(n))$ denotes the class of functions f which are asymptotically bounded above by g. In other words, $f(n) \in O(g(n))$ if there exists $M > 0$ and $n_0 \in \mathbb{N}$ such that for all $n \geq n_0$,

$$|f(n)| \leq M g(n).$$

This amounts to writing

$$\limsup_{n\to\infty} \frac{|f(n)|}{g(n)} < \infty.$$

Example. Consider a function defined by $f(n) = an + b$, where $a > 0$. Then, we can write $f(n) \in O(n)$. To see why, note that for all $n \geq 1$, we have

$$|f(n)| = |an + b| \leq an + |b| \leq (a + |b|)n.$$

Thus, setting $M = a + |b| > 0$ completes the proof.

Example. Consider a polynomial function defined by

$$f(n) = a_k n^k + a_{k-1} n^{k-1} + \cdots + a_1 n + a_0,$$

with some non-zero coefficient. Then, we can write $f(n) \in O(n^k)$. Like before, note that for all $n \geq 1$, we have

$$|f(n)| \leq \sum_{i=0}^{k} |a_i| n^i \leq \sum_{i=0}^{k} |a_i| n^k = (|a_k| + |a_{k-1}| + \cdots + |a_0|) n^k.$$

Thus, setting $M = |a_k| + \cdots + |a_0| > 0$ completes the proof.

Theorem 1.1. *If $f_1(n) \in O(g_1(n))$ and $f_2(n) \in O(g_2(n))$, then*

$$f_1(n) + f_2(n) \in O(\max\{g_1(n), g_2(n)\}).$$

Definition 1.2. The set $\Omega(g(n))$ denotes the class of functions f are asymptotically bounded below by g. In other words, $f(n) \in \Omega(g(n))$ if there exists $M > 0$ and $n_0 \in \mathbb{N}$ such that for all $n \geq n_0$,

$$|f(n)| \geq M g(n).$$

This amounts to writing

$$\liminf_{n \to \infty} \frac{f(n)}{g(n)} > 0.$$

Definition 1.3. The set $\Theta(g(n))$ denotes the class of functions f which are asymptotically bounded both above and below by g. In other words, $f(n) \in \Theta(g(n))$ if there exist $M_1, M_2 > 0$ and $n_0 \in \mathbb{N}$ such that for all $n \geq n_0$,

$$M_1 g(n) \leq |f(n)| \leq M_2 g(n).$$

This amounts to writing $f(n) \in O(g(n))$ and $f(n) \in \Omega(g(n))$.

Another class of notation uses the idea of dominated growth.

Definition 1.4. The set $o(g(n))$ denotes the class of functions f which are asymptotically dominated by g. In other words, $f(n) \in o(g(n))$ if for all $M > 0$, there exists $n_0 \in \mathbb{N}$ such that for all $n \geq n_0$,

$$|f(n)| < M g(n).$$

This amounts to writing

$$\lim_{n \to \infty} \frac{|f(n)|}{g(n)} = 0.$$

Definition 1.5. The set $\omega(g(n))$ denotes the class of functions f which asymptotically dominate g. In other words, $f(n) \in \omega(g(n))$ if for all $M > 0$, there exists $n_0 \in \mathbb{N}$ such that for all $n \geq n_0$,

$$|f(n)| > Mg(n).$$

This amounts to writing

$$\lim_{n\to\infty} \frac{|f(n)|}{g(n)} = \infty.$$

Definition 1.6. We say that $f(n) \sim g(n)$ if f is asymptotically equal to g. In other words, $f(n) \sim g(n)$ if for all $\epsilon > 0$, there exists $n_0 \in \mathbb{N}$ such that for all $n \geq n_0$,

$$\left| \frac{f(n)}{g(n)} - 1 \right| < \epsilon.$$

This amounts to writing

$$\lim_{n\to\infty} \frac{f(n)}{g(n)} = 1.$$

We often abuse notation and treat the following as equivalent.

$$T(n) \in O(g(n)), \qquad T(n) = O(g(n)).$$

2 Root finding methods

Consider an equation of the form $f(x) = 0$, where $f\colon [a, b] \to \mathbb{R}$ is given. We wish to solve this equation, i.e. find the roots of f.

Note that for *arbitrary* functions, this task is impossible. To see this, consider a function f which assumes the value 1 on $[0, 1] \setminus \{\alpha\}$ and $f(\alpha) = 0$, for some $\alpha \in [0, 1]$. There is no way of pinpointing α without checking f at every point in $[0, 1]$. Besides, a computer cannot reasonably store real numbers with arbitrary precision.

Thus, we direct our attention towards *continuous* functions f. We only seek sufficiently accurate approximations of its root $\alpha \in (a, b)$.

Theorem 2.1 (Intermediate Value Theorem)**.** *Let $f\colon [a, b] \to \mathbb{R}$ be continuous. If $f(a)f(b) < 0$, then there exists $\alpha \in (a, b)$ such that $f(\alpha) = 0$.*

2.1 Tabulation method

To identify the location of a root of f on an interval $I = [a, b]$, we subdivide I into n subintervals $[x_i, x_{i+1}]$ where $x_i = a + (b - a)i/n$. Now, we simply apply the Intermediate Value Theorem to f on each of these intervals. If $f(x_i)f(x_{i+1}) < 0$, then f has a root somewhere in (x_i, x_{i+1}). Note that the error in our approximation is on the order of $|b - a|/n$. The precision of this method can be improved by increasing n.

To reach a degree of approximation ϵ, we must iterate n times, where

$$n > \frac{b - a}{\epsilon}.$$

2.2 Bisection method

Here, we first verify that $f(a)f(b) < 0$, thus ensuring that f has a root within (a, b). Now, set $x_1 = a + (b - a)/2$ and apply the Intermediate Value Theorem on the subintervals $[a, x_1]$ and $[x_1, b]$. One of these *must* contain a root of f. Note that if $f(x_1) = 0$, we are done; otherwise, let $I_1 = [a_1, b_1]$ be the subinterval containing the root. Repeat the above process, obtaining successive subintervals I_n with lengths $|b - a|/2^n$. The error in our approximation is of this order, and can be controlled by stopping at appropriately large n.

The quantity $x_{n+1} = (a_n + b_n)/2$ is a good approximation for the actual root α since we know that $x_{n+1}, \alpha \in [a_n, b_n]$, so

$$|x_{n+1} - \alpha| \leq |b_n - a_n| = 2^{-n}|b - a| \to 0.$$

To reach a degree of approximation ϵ, we must iterate n times, where

$$n > \log_2 \frac{b - a}{\epsilon}.$$

2.3 Newton-Raphson method

Assuming that f is twice differentiable, use Taylor's theorem to write

$$f(x) = f(x_0) + f'(x)(x - x_0) + \frac{1}{2} f''(c)(x - x_0)^2$$

for all $x \in [a, b]$, where c is between x and x_0. The first two terms represent the tangent line to f, drawn at $(x_0, f(x_0))$. Now, define

$$x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}.$$

Note that this is the point at which the tangent line to f at x_0 cuts the x-axis. We have implicitly assumed that $f'(x_0) \neq 0$. In this manner, create the sequence of points

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}.$$

We wish to show that $x_n \to \alpha$, under certain circumstances.

Definition 2.1 (Order of convergence)**.** Let $x_n \to \alpha$. We say that this convergence is of order $p \geq 1$ if

$$\lim_{n \to \infty} \frac{|\alpha - x_{n+1}|}{|\alpha - x_n|^p} > 0.$$

Theorem 2.2. *Let f be a real function on $[\alpha - \delta, \alpha + \delta]$ such that*

1. $f(\alpha) = 0$.
2. *f is twice differentiable, with non-zero derivatives.*
3. *f'' is continuous.*
4. $|f''(x)/f'(y)| \leq M$ *for all x, y.*

If $x_0 \in [\alpha - h, \alpha + h]$ where $h = \min\{\delta, 1/M\}$, then the Newton-Raphson sequence generated by x_0 converges to the root α quadratically.

Proof. Pick $x_n \in [\alpha - h, \alpha + h]$. Using Taylor's theorem,

$$f(\alpha) = f(x_n) + f'(x_n)(\alpha - x_n) + \frac{1}{2}f''(c)(\alpha - x_n)^2.$$

Also note that $f(\alpha) = 0$, and $x_n - x_{n+1} = f(x_n)/f'(x_n)$. Thus, dividing by $f'(x_n)$ and substituting gives

$$\alpha - x_{n+1} = -\frac{1}{2}\frac{f''(c)}{f'(x_n)}(\alpha - x_n)^2.$$

Using our estimates on $f''(c)/f'(x_n)$ and x_n along with $h \le 1/M$, we see that

$$|\alpha - x_{n+1}| \le \frac{1}{2}Mh|\alpha - x_n| \le \frac{1}{2}|\alpha - x_n|.$$

Indeed, we have shown that

$$|\alpha - x_n| \le \frac{1}{2^n}|\alpha - x_0|,$$

which directly gives the convergence $x_n \to \alpha$. Furthermore, we have

$$\frac{|\alpha - x_{n+1}|}{|\alpha - x_n|^2} = \frac{1}{2}\left|\frac{f''(c)}{f'(x_n)}\right| \le \frac{1}{2}M,$$

hence taking the limit $n \to \infty$ proves that the convergence is quadratic. □

Corollary 2.2.1. *Suppose that f satisfies the conditions of the previous theorem, along with $f' > 0$ and $f'' > 0$ on some interval $[\alpha, x]$. Then, the Newton-Raphson sequence generated by $x_0 \in [\alpha, x]$ converges to the root α quadratically.*

Remark. The convexity of f means that the tangent drawn at x_n lies below the curve, and hence cuts the x-axis between α and x_n.

Theorem 2.3. *If α is a multiple root of f such that $f(\alpha) = 0$, $f'(\alpha) = 0$, $f''(\alpha) \neq 0$, then the Newton-Raphson sequence converges to α linearly under suitable conditions.*

Proof. Use Rolle's Theorem to replace $f'(x_n) = f'(x_n) - f'(\alpha) = f''(a)(x_n - \alpha)$. □

2.4 Secant method

The chief difference between this method as Newton's method is that we approximate the tangent with a secant, i.e. perform an approximation of the derivative,

$$f'(x)h \approx f(x + h) - f(x)$$

for small h. Thus, our iterations proceed as

$$x_{n+1} = x_n - f(x_n)\frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})}.$$

Theorem 2.4. *Let f be a real function on $[a, b]$ such that*

1. *$f(\alpha) = 0$ where $\alpha \in (a, b)$.*
2. *f is continuously differentiable, with non-zero derivatives.*

Then, there exists $\delta > 0$ such that the sequence generated by the secant method converges to α when $x_0, x_1 \in (\alpha - \delta, \alpha + \delta)$.

Proof. Consider

$$\alpha - x_{n+1} = \alpha - x_n + f(x_n)\frac{x_n - x_{n-1}}{f(x_n) - x_{n-1}}.$$

Now, use the Mean Value Theorem to write $f(x_n) = f(x_n) - f(\alpha) = f'(\xi)(x_n - \alpha)$ for some ξ between α and x_n. Similarly, write $f(x_n) - f(x_{n-1}) = f'(\zeta)(x_n - x_{n-1})$ for some ζ between x_n and x_{n-1}. Thus,

$$\alpha - x_{n-1} = \alpha - x_n + \frac{f'(\xi)(x_n - \alpha)}{f'(\zeta)} = (\alpha - x_n)\left(1 - \frac{f'(\xi)}{f'(\zeta)}\right).$$

We want $|1 - f'(\xi)/f'(\zeta)| < 1$. Since $f'(\alpha) \neq 0$, there is a δ-neighbourhood of α where $3f'(\alpha)/4 < f'(x) < 5f'(\alpha)/4$ (without loss of generality) using the continuity of f'. Thus, whenever $x_0, x_1 \in (\alpha - \delta, \alpha + \delta)$, we have ξ, ζ belonging to the same neighbourhood. This gives $3/5 < f'(\zeta)/f'(\xi) < 5/3$. This gives

$$-\frac{2}{3} < 1 - \frac{f'(\xi)}{f'(\zeta)} < \frac{2}{5}.$$

In other words, $|1 - f'(\xi)/f'(\zeta)| < 2/3$, so

$$|\alpha - x_{n+1}| < \frac{2}{3}|\alpha - x_n|,$$

which directly gives $x_n \to \alpha$.

The order of convergence turns out to be $\varphi = (1 + \sqrt{5})/2$. To show this, we want

$$\lim_{n\to\infty} \frac{|\alpha - x_{n+1}|}{|\alpha - x_n|^\varphi} > 0.$$

Assume that $f'(\alpha) > 0$, $f''(\alpha) > 0$. First, we will show that

$$\lim_{n\to\infty} \frac{|\alpha - x_{n+1}|}{|\alpha - x_n||\alpha - x_{n-1}|} = \frac{f''(\alpha)}{2f'(\alpha)}.$$

Denote the quantity in the limit as $\psi(x_n, x_{n-1})$. We examine the equivalent limit

$$\lim_{x_{n-1}\to\alpha} \lim_{x_n\to\alpha} \psi(x_n, x_{n-1}).$$

Like before, write

$$\alpha - x_{n+1} = (\alpha - x_n)\left(1 - \frac{f'(\xi)(x_n - x_{n-1})}{f(x_n) - f(x_{n-1})}\right),$$

hence

$$\frac{\alpha - x_{n+1}}{(\alpha - x_n)(\alpha - x_{n-1})} = \frac{1}{\alpha - x_{n-1}}\left[1 - \frac{f'(\xi)(x_n - x_{n-1})}{f(x_n) - f(x_{n-1})}\right].$$

Thus,

$$\lim_{x_n \to \alpha} \psi(x_n, x_{n-1}) = \frac{1}{\alpha - x_{n-1}} \left[1 + \frac{f'(\alpha)(\alpha - x_{n-1})}{f(x_{n-1})} \right]$$
$$= \frac{f(x_{n-1}) + f'(\alpha)(\alpha - x_{n-1})}{f(x_{n-1})(\alpha - x_{n-1})}.$$

Use Taylor's Theorem to approximate

$$f(x_{n-1}) = f(\alpha) + f'(\alpha)(x_{n-1} - \alpha) + \frac{1}{2} f''(\eta)(x_{n-1} - \alpha)^2,$$

giving

$$\lim_{x_n \to \alpha} \psi(x_n, x_{n-1}) = \frac{f''(\eta)(\alpha - x_{n-1})^2}{2f(x_{n-1})(\alpha - x_{n-1})},$$

and use the Mean Value Theorem to write $f(x_{n-1}) = f'(\kappa)(x_{n-1} - \alpha)$ giving

$$\lim_{x_n \to \alpha} \psi(x_n, x_{n-1}) = -\frac{f''(\eta)}{2f'(\kappa)},$$

This gives

$$\lim_{x_{n-1} \to \alpha} \lim_{x_n \to \alpha} |\psi(x_n, x_{n-1})| = \frac{f''(\alpha)}{2f'(\alpha)} = C.$$

Now, suppose that

$$\lim_{n \to \infty} \frac{|\alpha - x_{n+1}|}{|\alpha - x_n|^q} = A > 0.$$

Dividing, we have

$$\lim_{n \to \infty} \frac{|\alpha - x_n|^{q-1}}{|\alpha - x_{n-1}|} = \frac{C}{A}, \qquad \lim_{n \to \infty} \frac{|\alpha - x_n|}{|\alpha - x_{n-1}|^{1/(q-1)}} = \left(\frac{C}{A} \right)^{1/(q-1)} > 0.$$

For q to be minimal, we must have $1/(q-1) = q$, or q is the golden ratio φ. □

2.5 Fixed point method

Note that a root of f is simply a fixed point of $f + x$.

Theorem 2.5. *Let $f\colon [a,b] \to [a,b]$ be continuous. Then, f has a fixed point $\beta \in [a,b]$, $f(\beta) = \beta$.*

Thus, let $f\colon [a,b] \to [a,b]$ be continuous. Define the fixed point sequence $x_{n+1} = f(x_n)$, seeded by some $x_0 \in [a,b]$. Note that if this sequence converges with $x_n \to \beta$, then β is a fixed point of f.

Definition 2.2. A function $f\colon [a,b] \to \mathbb{R}$ is said to be a contraction if there exists $L \in (0,1)$ such that $|f(x) - f(y)| \leq L|x - y|$ for all $x, y \in [a,b]$.

Remark. Note that f is Lipschitz continuous. If f is also differentiable, then $|f'| < 1$.

Theorem 2.6. *Let $f\colon [a,b] \to [a,b]$ be a contraction map. Then, any fixed point sequence converges to the unique fixed point of f.*

Proof. First, we show that f has at most one fixed point. Let β_1, β_2 be fixed points of f. Then, $|f(\beta_1) - f(\beta_2)| \leq L|\beta_1 - \beta_2|$ where $L \in (0,1)$. This forces $\beta_1 = \beta_2$. Thus, f has a unique fixed point in $[a,b]$.

Let $\{x_n\}$ be a fixed point iteration. Then,

$$|x_{n+1} - \beta| = |f(x_n) - f(\beta)| \leq L|x_n - \beta|,$$

which directly gives $x_n \to \beta$. □

3 Interpolation

3.1 Lagrange interpolation

Theorem 3.1. *Let $x_1, \ldots, x_n \in \mathbb{R}$ be distinct, and let $y_1, \ldots, y_n \in \mathbb{R}$. Then, the following polynomial of degree $n-1$ satisfies $p(x_i) = y_i$.*

$$p(x) = \sum_{i=1}^{n} \prod_{j \neq i} \frac{x - x_j}{x_i - x_j} y_i.$$

Furthermore, this choice of p is unique.

Proof. The polynomials

$$p_i(x) = \prod_{j \neq i} \frac{x - x_j}{x_i - x_j}$$

satisfy $p_i(x_j) = \delta_{ij}$. These p_i form a basis of $\mathscr{P}^{n-1}$, the space of polynomials of degree at most $n-1$. □