

MA3103

Introduction to Graph Theory and Combinatorics

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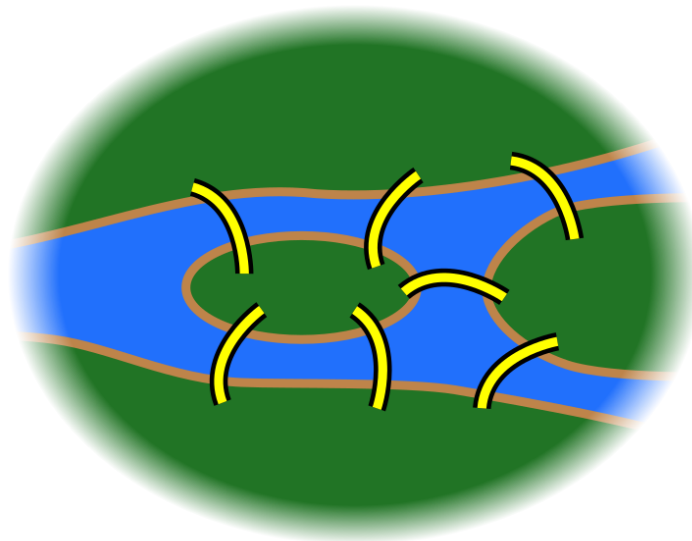
Contents

1 Introduction	1
1.1 The Seven Bridges of Königsberg	1
1.2 Basic definitions	2
1.3 Some principles	3

1 Introduction

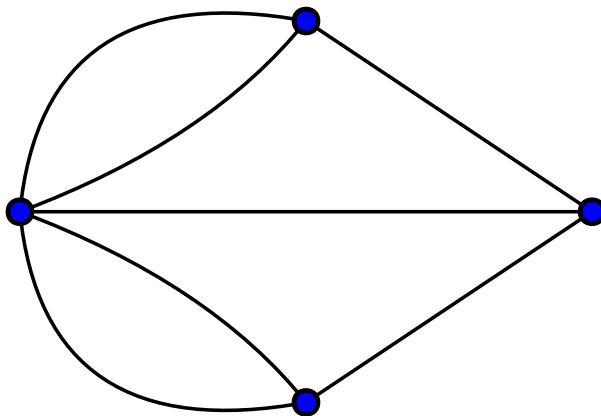
1.1 The Seven Bridges of Königsberg

The diagram below depicts a region in the city of Königsberg, Prussia. There are two islands, connected with the mainland and to each other via seven bridges. The Seven Bridges Problem is posed as follows: is it possible to walk through the entire city, visiting each one of the four landmasses by crossing each of the bridges exactly once?



Leonhard Euler showed that this is impossible; no such walk exists. The techniques he developed in doing so laid the foundations of *graph theory*.

The first thing to note is that the exact shape of the walk/trail is immaterial; all that matters is the sequence of landmasses visited and bridges crossed. Thus, each landmass can be compacted to a single point or *vertex*, and each bridge a line or *edge* connecting two such points. The resulting figure is a graph. Note that the orientations or placements of the points and lines are irrelevant, as long as the connections are undisturbed.



Now, examine a landmass which is on the trail but is neither our starting point, nor our ending point. In order to reach this landmass, we must enter via a bridge; but we cannot stay in the landmass, so we must leave via another a bridge. Thus, for each time we pass through this landmass, we can cross off two bridges joined to it. Once we are done, no bridge may remain unused; this means that we must have started with an even number of bridges joined to this landmass.

However, all four vertices in our graph connect to an odd number of edges. Since we require at least two vertices to act as intermediate points on our path, the desired walk is impossible.

1.2 Basic definitions

Definition 1.1. A graph $G(V, E)$ is an ordered pair of the set of vertices V and the set of edges E .

Definition 1.2. A simple graph is undirected, unweighted, and contains no self-loops or multiple edges joining vertices.

Definition 1.3. For a simple undirected graph, the set of edges E consists of two-element subsets of the set of vertices V .

Remark. For a directed, unweighted graph, the set of edges E consists of ordered pairs of elements from the set of vertices V .

Definition 1.4. A vertex is incident to an edge if that edge joins that vertex.

Definition 1.5. Two vertices are adjacent if there exists an edge connecting them. Two edges are adjacent if they connect to a common vertex.

Definition 1.6. The neighbours of a vertex consist of all vertices adjacent to it. The neighbours of an edge consist of all edges adjacent to it.

The number of neighbours of a vertex is called the degree of that vertex.

1.3 Some principles

Lemma 1.1 (Pigeonhole Principle). *If $n + 1$ objects are placed in n boxes, then we can find a box containing at least 2 objects.*

Proof. If every box contains at most 1 objects, then the total number of objects falls short. \square

Theorem 1.2. *There are no simple graphs where the degrees of all vertices are distinct.*

Proof. Let $G(V, E)$ be a simple graph with n vertices. The degrees of each of these vertices must be an integer among $0, 1, \dots, n - 1$. We now consider two cases.

Case I: There is a vertex of degree 0. Thus, this vertex is adjacent to no other vertex, which means that no vertex can have the full degree $n - 1$. This means that the remaining vertices have degrees among $1, 2, \dots, n - 2$, i.e. $n - 2$ choices of degree for $n - 1$ vertices.

Case II: There is no vertex of degree 0. Thus, the vertices have degrees among $1, 2, \dots, n - 1$, i.e. $n - 1$ choices of degree for n vertices.

In either case, the Pigeonhole Principle forces at least two vertices to share the same degree. \square

Lemma 1.3 (Strong Pigeonhole Principle). *Let q_1, q_2, \dots, q_n be positive integers. If*

$$N = q_1 + \dots + q_n - n + 1$$

objects are placed in n boxes, then we can find a box i containing at least q_i objects.

Proof. If every box i contains at most $q_i - 1$ objects, then the total number of objects falls short.

$$N \leq (q_1 - 1) + \dots + (q_n - 1) = q_1 + \dots + q_n - n = N - 1 \quad \square$$

Theorem 1.4. *The sum of the degrees of all vertices in a simple graph is twice the number of its edges.*

Proof. Let $G((V, E))$ be a simple graph. Define the incidence function $I: E \times V \rightarrow \{0, 1\}$, such that $I(e, v) = 1$ if e and v are incident, 0 otherwise. We perform the double counting,

$$\sum_{v \in V} \sum_{e \in E} I(e, v) = \sum_{e \in E} \sum_{v \in V} I(e, v).$$

Now, the number of edges incident to a vertex is simply its degree, so $\sum_{e \in E} I(e, v) = d(v)$. Also, every edge is incident to exactly two vertices, so $\sum_{v \in V} I(e, v) = 2$. Thus, we have

$$\sum_{v \in V} d(v) = 2|E|. \quad \square$$

Lemma 1.5 (Inclusion-Exclusion Principle). *For finite sets A_1, A_2, \dots, A_n , the number of elements in their union is given by*

$$\sum_{i=1}^n |A_i| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j| + \sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k| - \dots + (-1)^{n-1} |A_1 \cap \dots \cap A_n|.$$

Theorem 1.6. *There are $2^{\binom{n}{2}}$ simple graphs with n vertices.*

Exercise 1.1. How many simple graphs are there with n vertices and m edges?

Theorem 1.7. *Let $n, k \in \mathbb{N}$ such that $n > 3$ and $n/2 < k < n$. Let there be n points on a plane such that no three points are collinear. If every point is connected to at least k other points by segments, then there must be at least three segments forming a triangle.*

Proof. Consider a graph $G(V, E)$ with n vertices, such that every vertex has degree at least k . Pick an edge, say $\{x, y\}$, and let A be the neighbours of x apart from y , B be the neighbours of y apart from x . Note that A, B have at least $k - 1$ elements each. Suppose that $A \cap B = \emptyset$, i.e. the edge $\{x, y\}$ does not form a triangle. Thus,

$$|A \cup B| = |A| + |B| - |A \cap B| \geq 2(k - 1).$$

However, $|A \cup B| \leq n - 2$, hence $n \geq 2k$, or $k \leq n/2$. This is a contradiction. \square

Remark. We have shown that *every* segment is part of a triangle. The number of segments here is

$$|E| \geq nk > \frac{n^2}{4}.$$

Exercise 1.2. Is the condition $|E| > n^2/4$ sufficient to ensure the existence of a triangle?

Lemma 1.8 (Cauchy-Schwarz). *Let a_1, \dots, a_n and b_1, \dots, b_n be positive reals. Then,*

$$(a_1^2 + \dots + a_n^2)(b_1^2 + \dots + b_n^2) \geq (a_1b_1 + \dots + a_nb_n)^2.$$

Equality holds if and only if every $a_i = \lambda b_i$ for some fixed real λ .

Theorem 1.9 (Mantel). *In a simple graph with n vertices, the condition $|E| > n^2/4$ is sufficient to ensure the existence of a triangle.*

Proof. Let G be a simple graph with n vertices which is triangle-free. Thus, for any edge $\{x, y\} \in E$, the neighbour sets A and B of x and y intersect at no vertex. Thus, we can write

$$d(x) + d(y) = |A \cup B| \leq n.$$

Sum this over all possible edges. On the right, we have $n|E|$. On the left, we have the sum

$$\sum_{x \in V} d(x)^2 \geq \frac{1}{n} \left(\sum_{x \in V} d(x) \right)^2 \geq \frac{1}{n} \cdot 4|E|^2$$

This gives

$$\frac{4|E|^2}{n} \leq n|E|, \quad |E| \leq \frac{n^2}{4}.$$

□