

Lossless Join and Dependency Preserving Decomposition

Decomposition of a relation is done when a relation in relational model is not in appropriate normal form. Relation R is decomposed into two or more relations if decomposition is lossless join as well as dependency preserving.

Lossless Join Decomposition

If we decompose a relation R into relations R1 and R2,

- Decomposition is lossy if $R1 \bowtie R2 \supset R$
- Decomposition is lossless if $R1 \bowtie R2 = R$

To check for lossless join decomposition using FD set, following conditions must hold:

1. Union of Attributes of R1 and R2 must be equal to attribute of R. Each attribute of R must be either in R1 or in R2.

$$\text{Att}(R1) \cup \text{Att}(R2) = \text{Att}(R)$$

2. Intersection of Attributes of R1 and R2 must not be NULL.

$$\text{Att}(R1) \cap \text{Att}(R2) \neq \emptyset$$

3. Common attribute must be a key for at least one relation (R1 or R2)

$$\text{Att}(R1) \cap \text{Att}(R2) \rightarrow \text{Att}(R1) \text{ or } \text{Att}(R1) \cap \text{Att}(R2) \rightarrow \text{Att}(R2)$$

For Example, A relation R (A, B, C, D) with FD set{A→BC} is decomposed into R1(ABC) and R2(AD) which is a lossless join decomposition as:

1. First condition holds true as $\text{Att}(R1) \cup \text{Att}(R2) = (ABC) \cup (AD) = (ABCD) = \text{Att}(R)$.

2. Second condition holds true as $\text{Att}(R1) \cap \text{Att}(R2) = (ABC) \cap (AD) \neq \Phi$
3. Third condition holds true as $\text{Att}(R1) \cap \text{Att}(R2) = A$ is a key of $R1(ABC)$ because $A \rightarrow BC$ is given.

Dependency Preserving Decomposition

If we decompose a relation R into relations $R1$ and $R2$, All dependencies of R either must be a part of $R1$ or $R2$ or must be derivable from combination of FD's of $R1$ and $R2$.

For Example, A relation $R(A, B, C, D)$ with FD set $\{A \rightarrow BC\}$ is decomposed into $R1(ABC)$ and $R2(AD)$ which is dependency preserving because FD $A \rightarrow BC$ is a part of $R1(ABC)$.

GATE Question: Consider a schema $R(A,B,C,D)$ and functional dependencies $A \rightarrow B$ and $C \rightarrow D$. Then the decomposition of R into $R1(AB)$ and $R2(CD)$ is [GATE-CS-2001]

- A. dependency preserving and lossless join
- B. lossless join but not dependency preserving
- C. dependency preserving but not lossless join
- D. not dependency preserving and not lossless join

Answer: For lossless join decomposition, these three conditions must hold true:

1. $\text{Att}(R1) \cup \text{Att}(R2) = ABCD = \text{Att}(R)$
2. $\text{Att}(R1) \cap \text{Att}(R2) = \Phi$, which violates the condition of lossless join decomposition. Hence the decomposition is not lossless.

For dependency preserving decomposition,

$A \rightarrow B$ can be ensured in $R1(AB)$ and $C \rightarrow D$ can be ensured in $R2(CD)$. Hence it is dependency preserving decomposition.

So, the correct option is C.