

# Overview

EE450: Introduction to Computer Networks

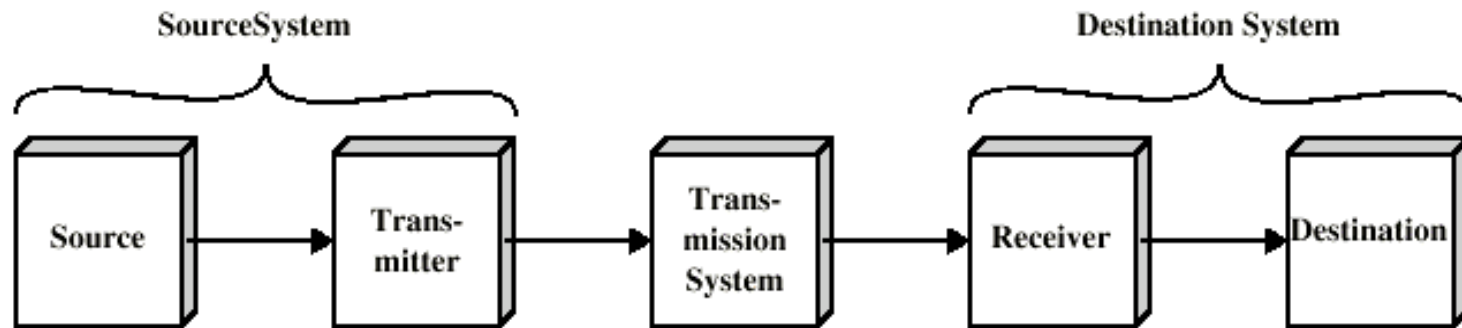
Professor A. Zahid

# Course Overview

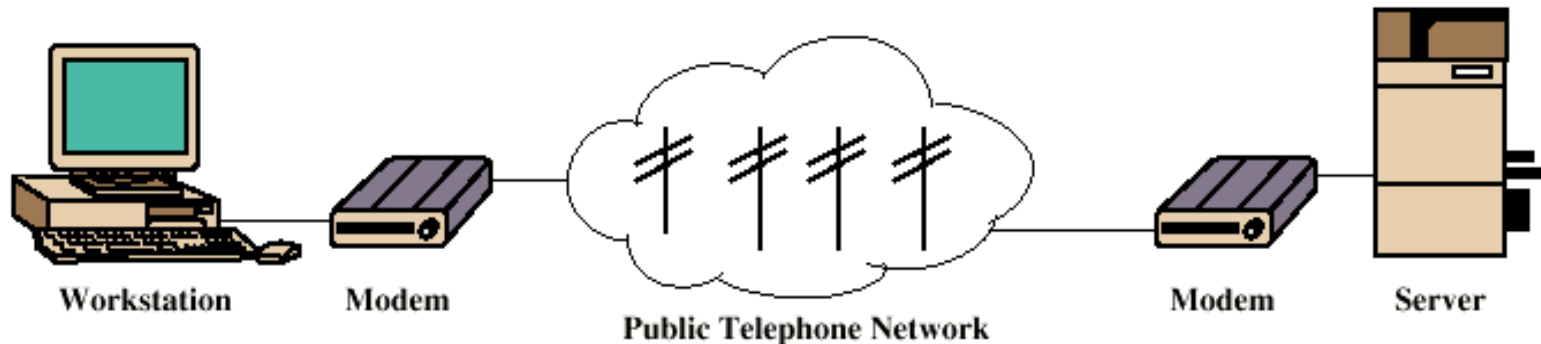
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- Part 1: Data Communications & Networking
- Part 2: Computer Networking Protocols (TCP/IP)
- Part 3: Wide Area Networks (WANs)
- Part 4: Local Area Networks (LANs)
- Part 5: Internetworking
- Part 6: Transport Layer Protocols
- Part 7: Network Applications

# Simple Data Communications Model



(a) General block diagram



(b) Example

# Key Data Communications Tasks

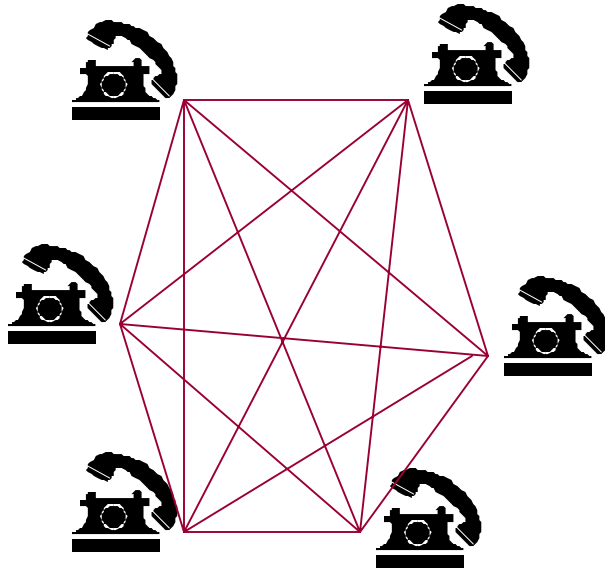
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- Interfacing
- Signal Generation
- Synchronization
- Exchange Management
- Error detection and correction
- Addressing and routing
- Recovery
- Message formatting
- Security
- Network Management
- Many more...

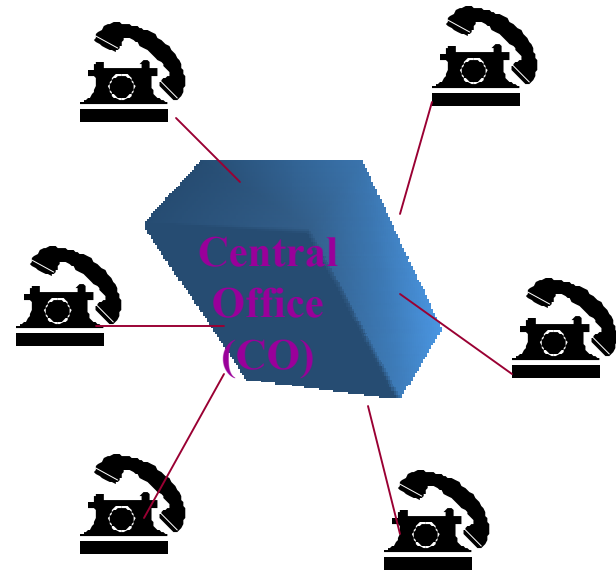
# Computer Network

- A Computer Network is a set of nodes such as routers, switches, hosts, etc.. interconnected via transmission facilities such as copper, cable, fiber, satellite, radio, microwave, etc.. for the purpose of providing *services* to end systems/users
- Do we need networking? Yes we do!
- Point-to-point communication is not practical!
  - Devices are too far apart
  - Large set of devices would need impractical number of connections. See illustration next chart

# Example: Telephone Network



Fully-Connected Mesh  
# of FDX links =  $N(N-1)/2$   
e.g.,  $N=6$ ;  $6(5)/2=15$  links  
Total # ports =  $N(N-1)$   
e.g.,  $N=6$ ;  $6(5)=30$  ports



With Central Office  
# of FDX links =  $N$   
e.g.,  $N=6$ ; 6 links  
Total # of ports =  $N$   
e.g.  $N=6$ , 6 ports

# Multimedia Convergence

## Major Industries

Broadcast TV  
Film

## Data

- e-mail
- files
  - executables
  - source code
  - data
  - html
  - image

## Computer Software

Financial, e-Commerce, etc.

Recording  
Broadcast  
Radio

## Video

- pre-recorded / on-demand (e.g., MPEG, Real Net)
- live (video phone, video teleconference)

## Voice / Audio

- pre-recorded / on-demand (streaming or file [mp3])
- live (Real Net, VoIP)

## Wireless



Network

Telephone companies (Telco)  
Internet Service Providers (ISP)



# Clients, Servers and Peers

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- A network computer can either provide service or request service
- A *server* is a service provider, providing access to network resources
- A *Client* is a service requester
- A *Peer-to-Peer* network does not have a dedicated server. All computers are equal and they both provide and request services.



# Server Roles

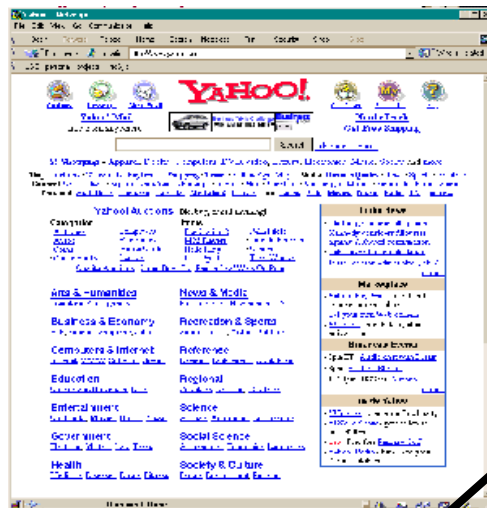
- Servers can assume several roles and a single server could also have several roles
- Examples of Servers include:
  - File Servers: Manages user access to shared files
  - Print Servers: Manages user access to print resources
  - Application Servers: Similar to FS with some processing
  - Mail Servers: Manages electronic messages between users
  - Communications (Remote Access) Servers: Manages data flow and e-messages from one network to another
  - Web Servers: Runs WWW and FTP servers for access via the Internet/Intranet
  - Directory (DNS) Servers: Locates information about networks such as domains.

# Client/Server Model

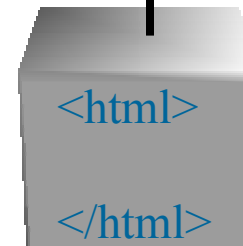
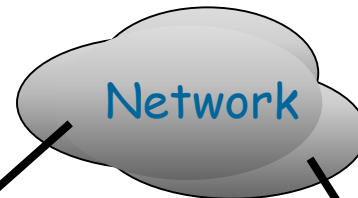
Example: World Wide Web

Client  
(web browser)  
-source-

Server  
(web server)  
-destination-



Request  
(GET index.html)



html file



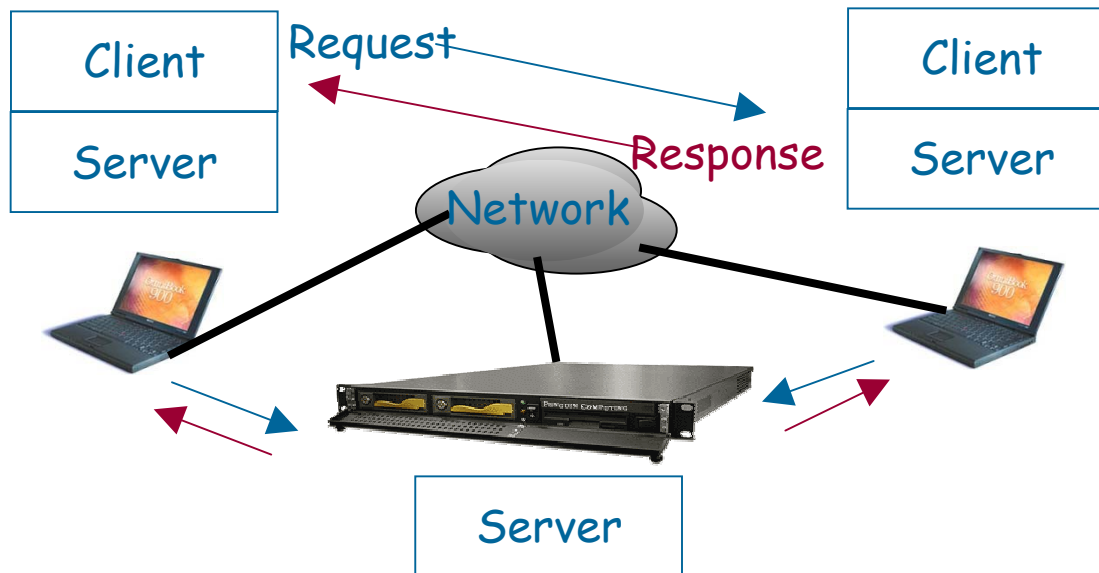
image file

Response  
(HTML file  
gif, jpeg, ..)



# Peer-to-Peer Model

- Peer-to-Peer
  - Each host has both client and server functionalities
  - CPU cycle sharing
  - Example: Gnutella, KaZaA, Skype, etc...



# Network Software (I)

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- NOS include special functions for connecting hosts into a network
- NOS manages network resources and services
- NOS provide network security for multiple users
- Most common Client/Server NOS include:
  - UNIX
  - Microsoft NT/Windows 2000
  - Novell Netware
  - LINUX
  - OS/2
  - Others

# Network Software (II)

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- Network hosts communicate through the use of client software called "Shells, Redirectors, Requesters"
- Network Protocols (such as TCP/IP, SPX/IPX, NETBEUI, etc..) enables data transmission across the network
- Client software resides on top of the network protocols.

# Network Hardware

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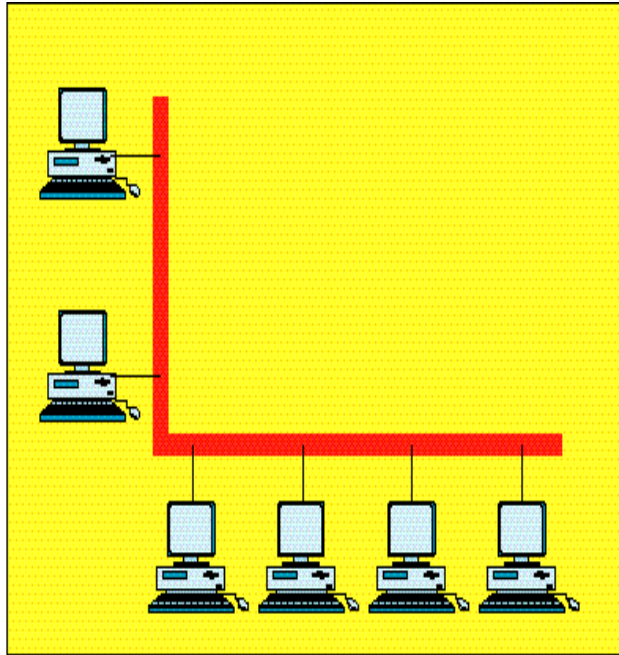
- Users accessing network resources must have a Pathway to those resources.
- Host connect to networks using expansion cards known as Network Interface Cards (NICs), a.k.a. Adapter Cards.
- Network cards communicate by sending signals through the medium (Twisted pair, Coax, Fiber, Radio, etc..)

# Network Classifications

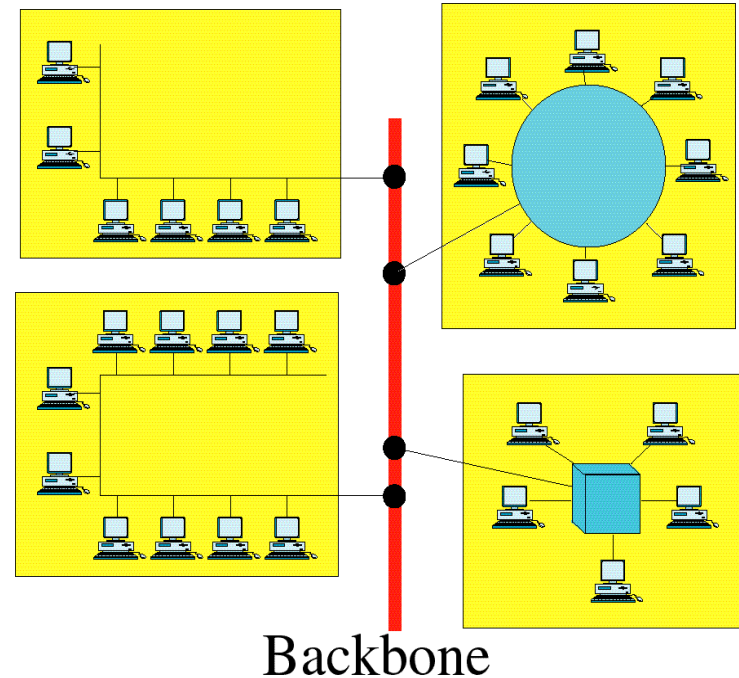
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- Networks can be classified based on Coverage into
  - LANs: Local Area Networks
  - WANs: Wide Area Networks
  - Others including MAN (Metropolitan Area Networks), PAN (Personal Area Networks), Home Networks, etc...
- Networks could also be classified as Switched or Shared (Broadcast) networks

# Local Area Networks (I)



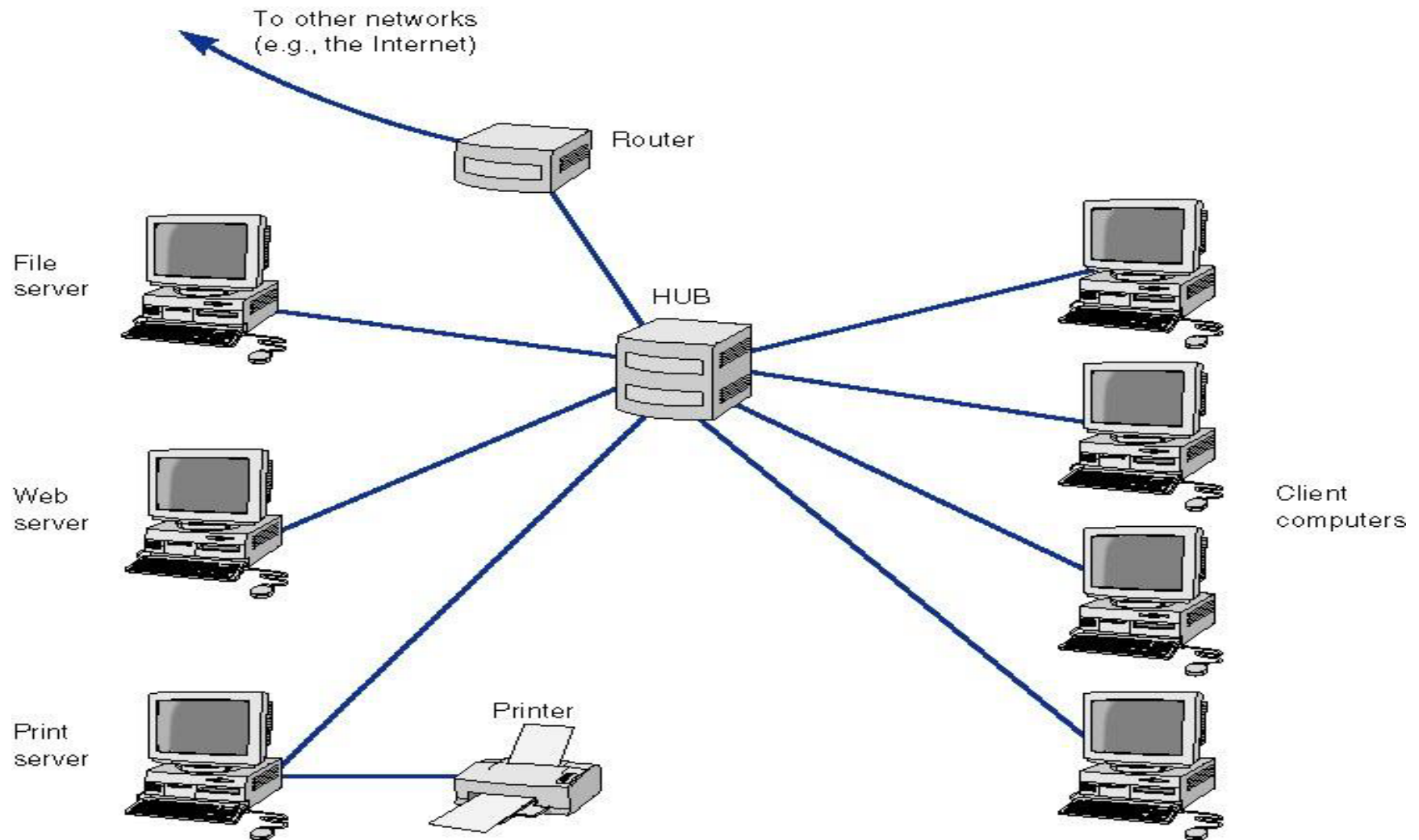
Single building LAN



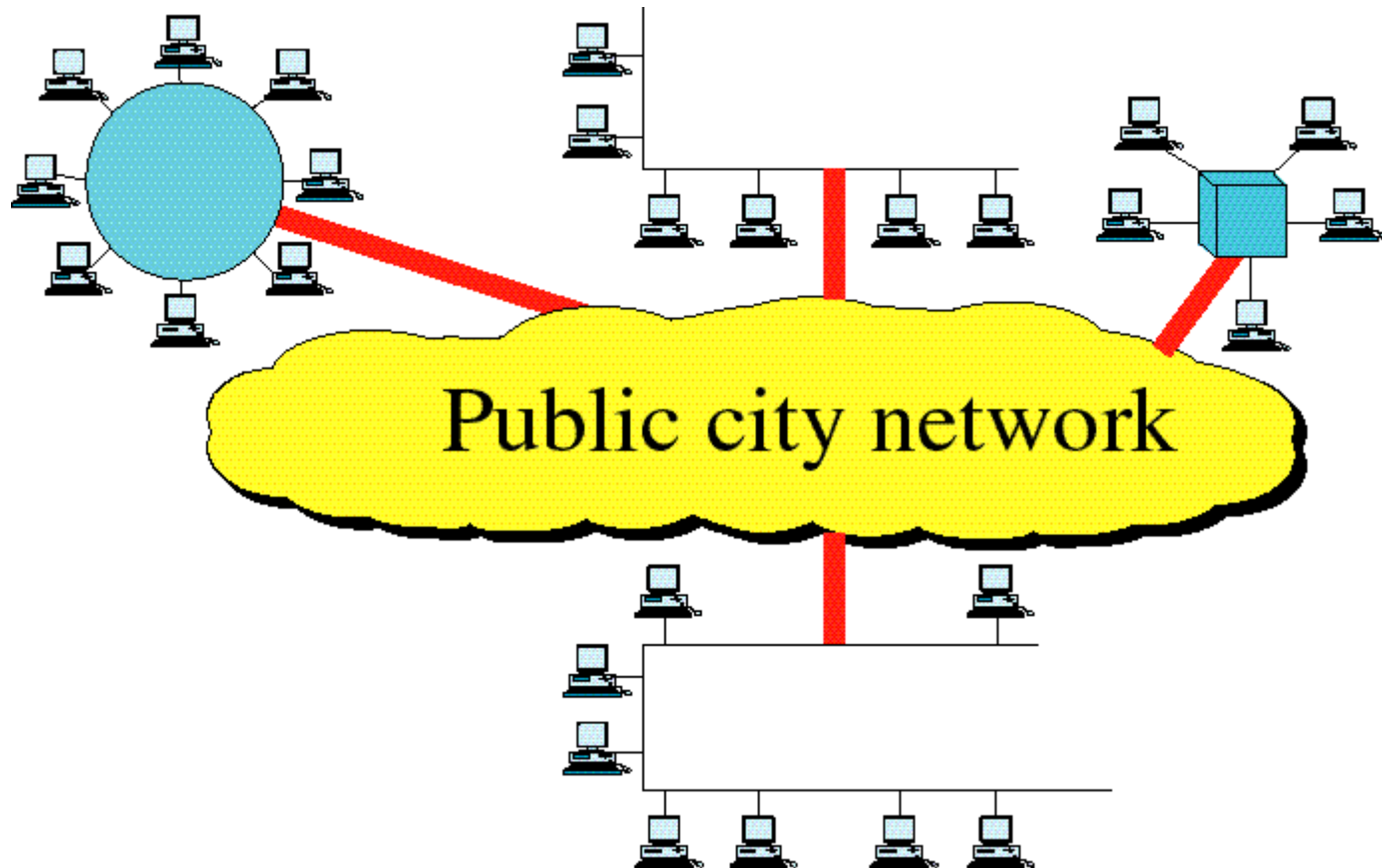
Multiple building LAN



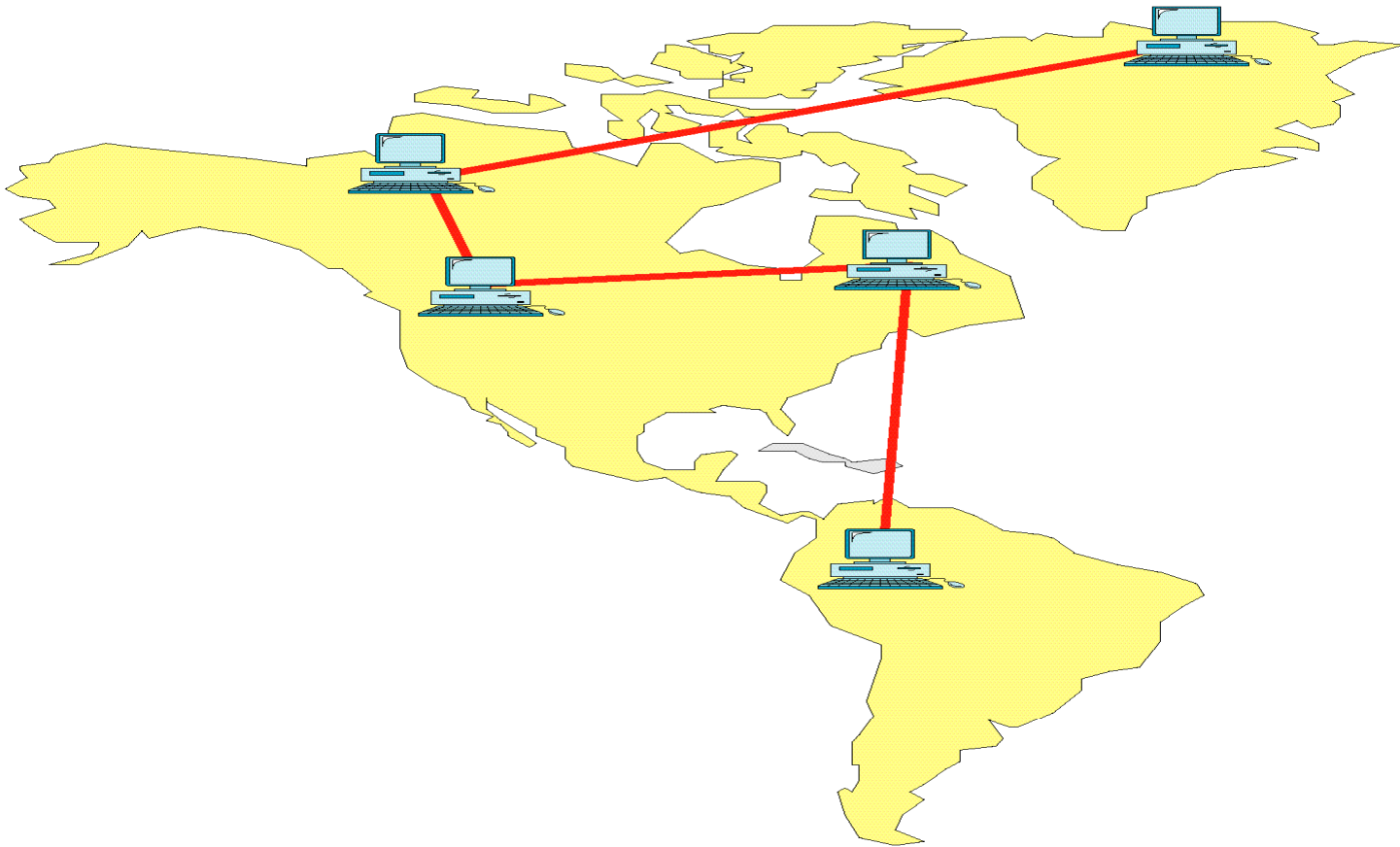
# Local Area Networks (II)



# Metropolitan Area Network



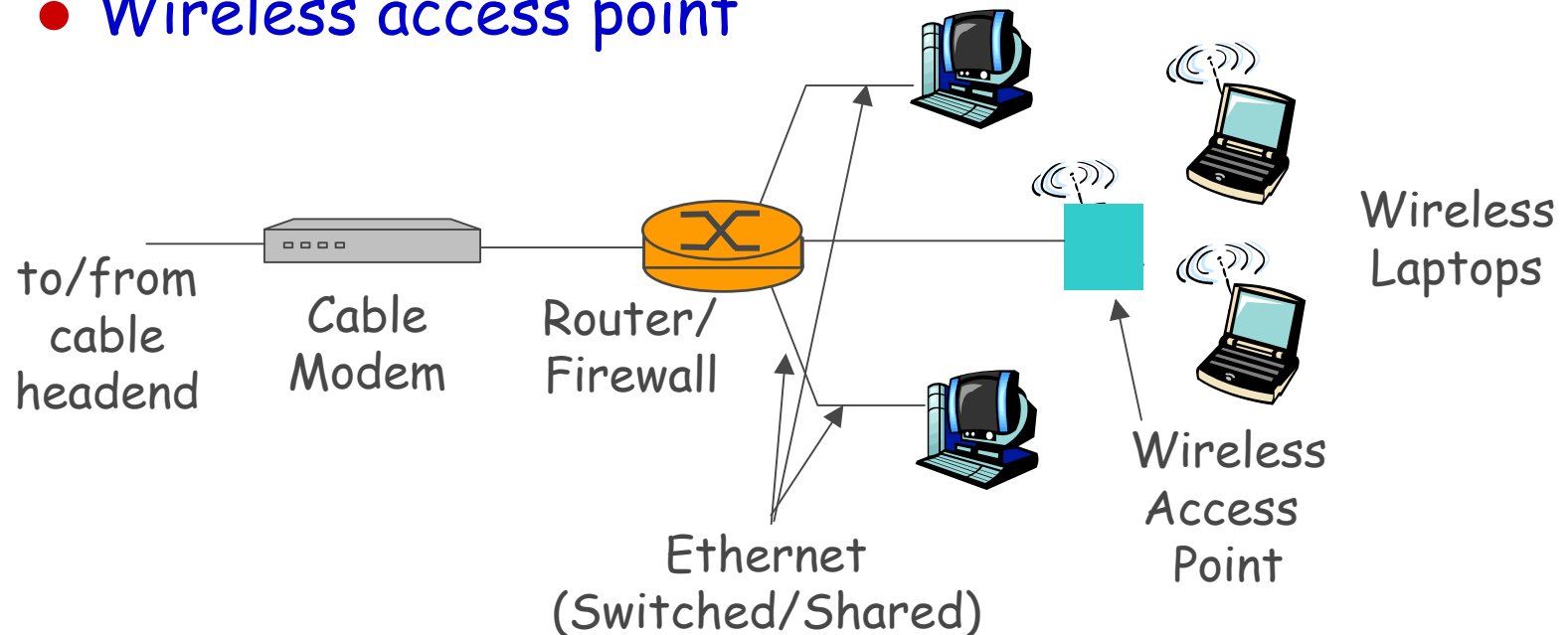
# Wide Area Networks



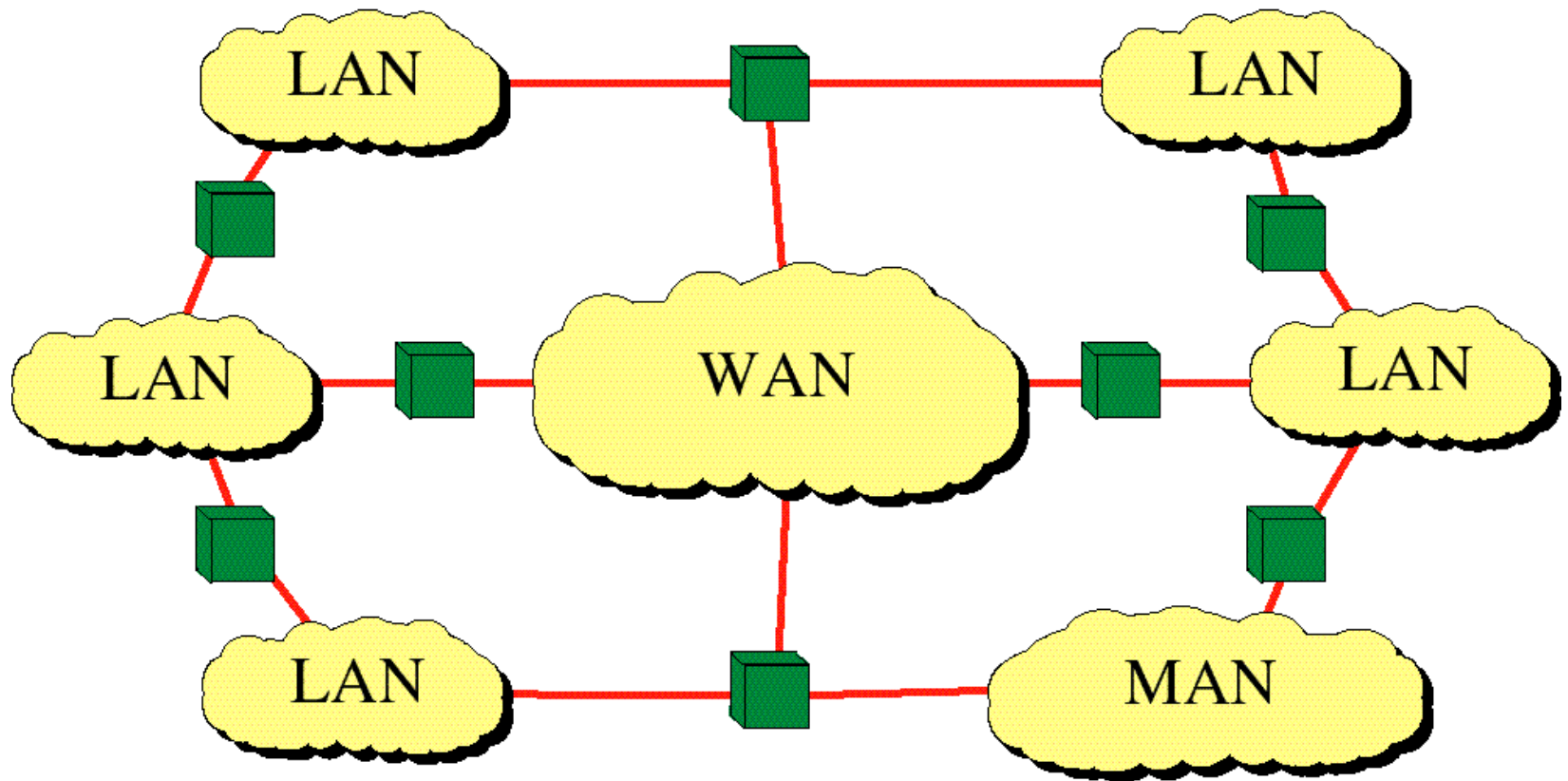
# Home Networks

## Typical home network components

- ADSL or cable modem
- Router/firewall/NAT
- Ethernet
- Wireless access point



# Internetworking

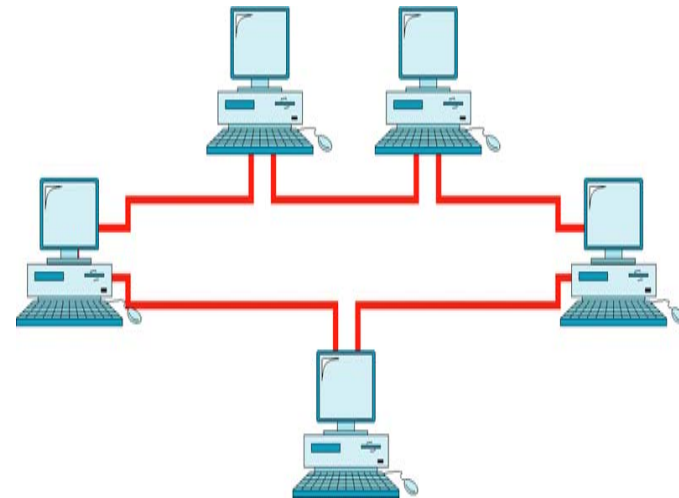
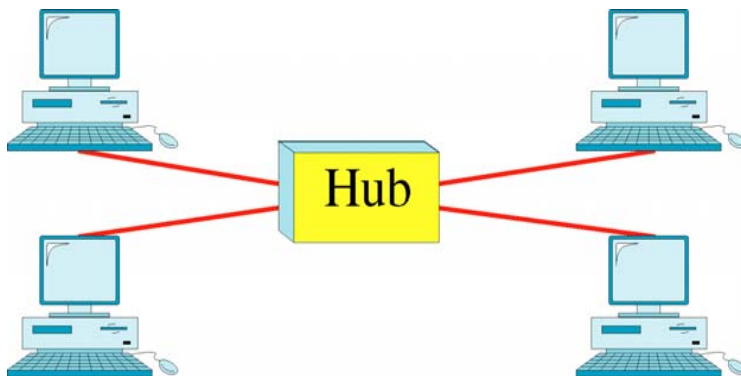
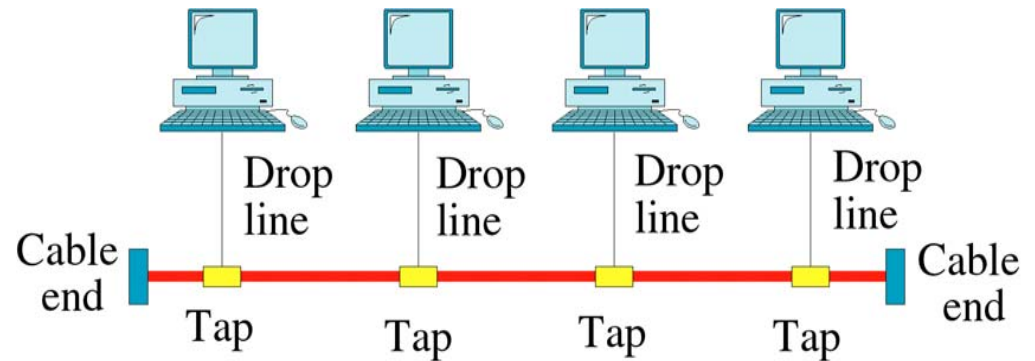
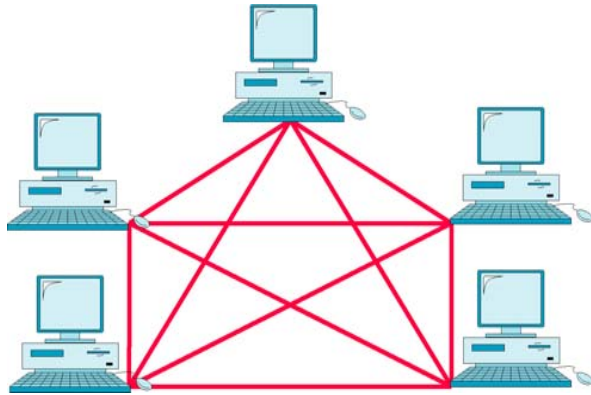


# Network Topologies

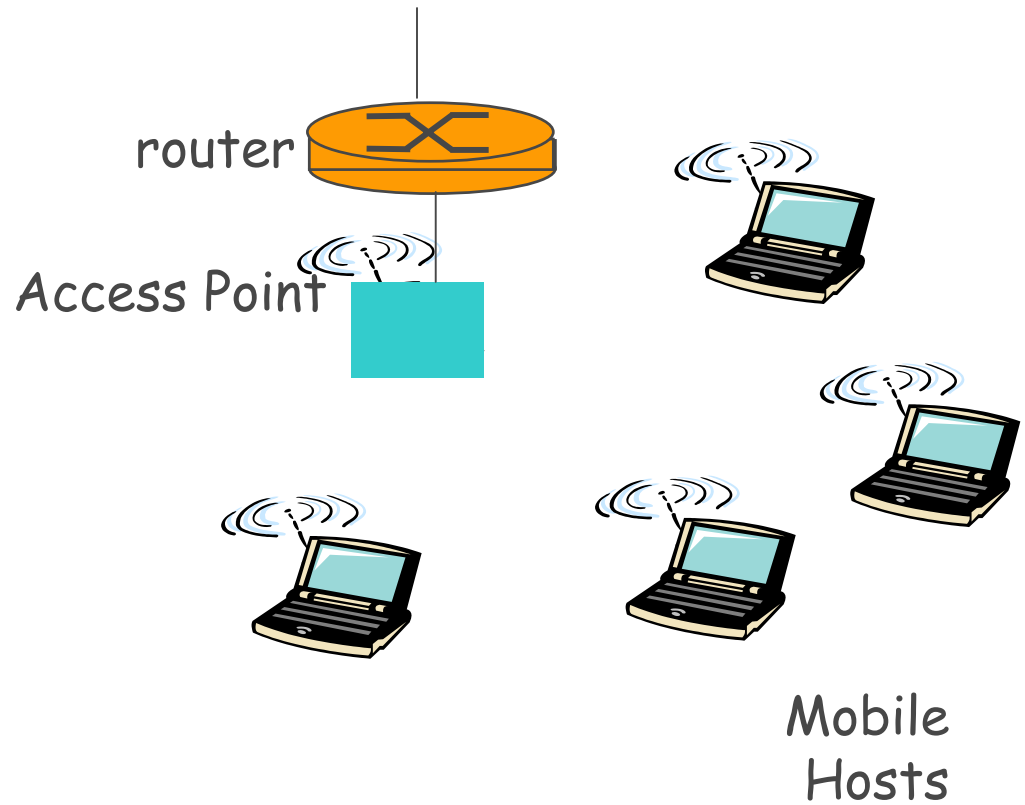
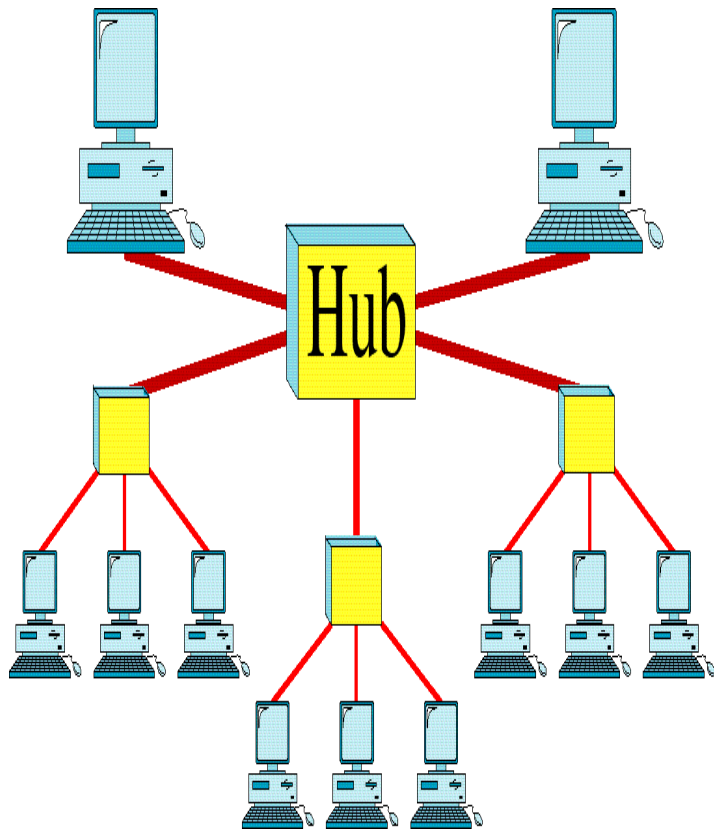
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- Network topology is the physical arrangement of the network nodes and the links interconnecting them
  - Mesh topology
  - Star/Hub topology
  - Bus topology
  - Tree Topology
  - Ring topology
- A fully connected network is one in which every node is connected to every other node

# Mesh, Hub, Bus and Rings (I)



# Tree and Wireless (II)





# Link Topologies

- Point-to-point

- Direct link
- Only 2 devices share link



- Multipoint

- More than two devices share the link



# Link Duplicity

- Simplex

- One direction
- e.g. Radio/Television broadcasting



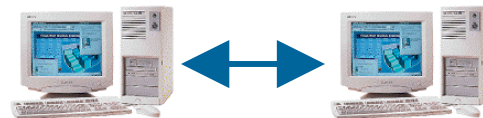
- Half duplex (HDX)

- Either direction, but only one way at a time
- e.g. Police radio

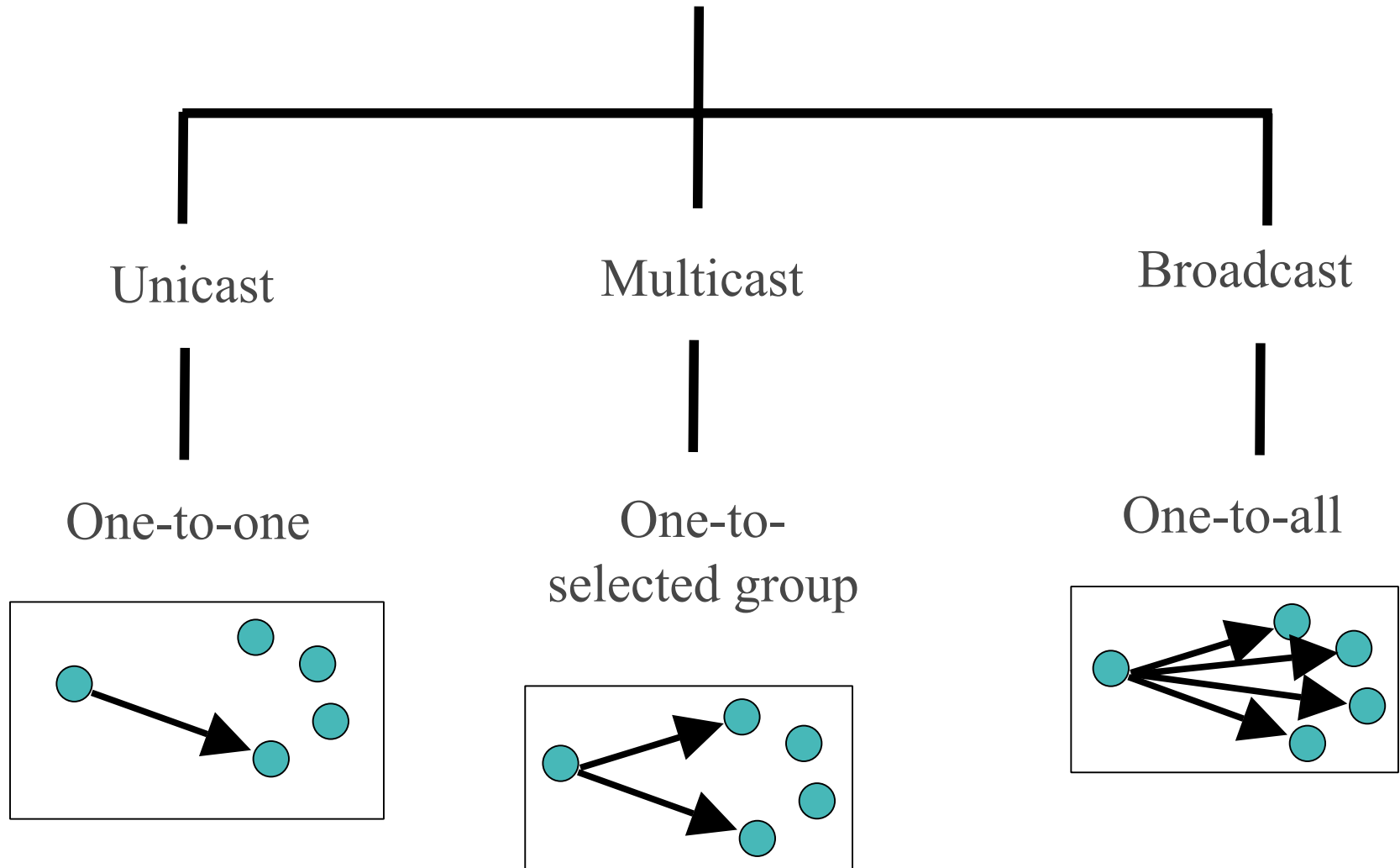


- Full duplex (FDX)

- Both directions at the same time
- e.g. Telephony



# Transmission Modes



# Physical Media (I): Copper

- **Bit:** propagates between transmitter/rcvr pairs
- **Physical link:** what lies between transmitter & receiver
- **Guided media:**
  - signals propagate in solid media: copper, fiber, coax
- **Unguided media:**
  - signals propagate freely, e.g., radio

## Twisted Pair (TP)

- two insulated copper wires
  - Category 3: traditional phone wires, 10 Mbps Ethernet
  - Category 5: 100Mbps Ethernet



# Physical Media (II): Cable

## Coaxial Cable:

- Two concentric copper conductors
- Bi-directional
- Baseband:
  - single channel on cable
  - legacy Ethernet
- Broadband:
  - Multiple channels on cable
  - CATV, Cable Access



## Fiber Optic Cable:

- Glass fiber carrying light pulses, each pulse a bit
- High-speed operation:
  - high-speed point-to-point transmission (e.g., 10's-100's Gps)
- Low error rate: repeaters spaced far apart ; immune to electromagnetic noise



# Physical Media (III): Radio

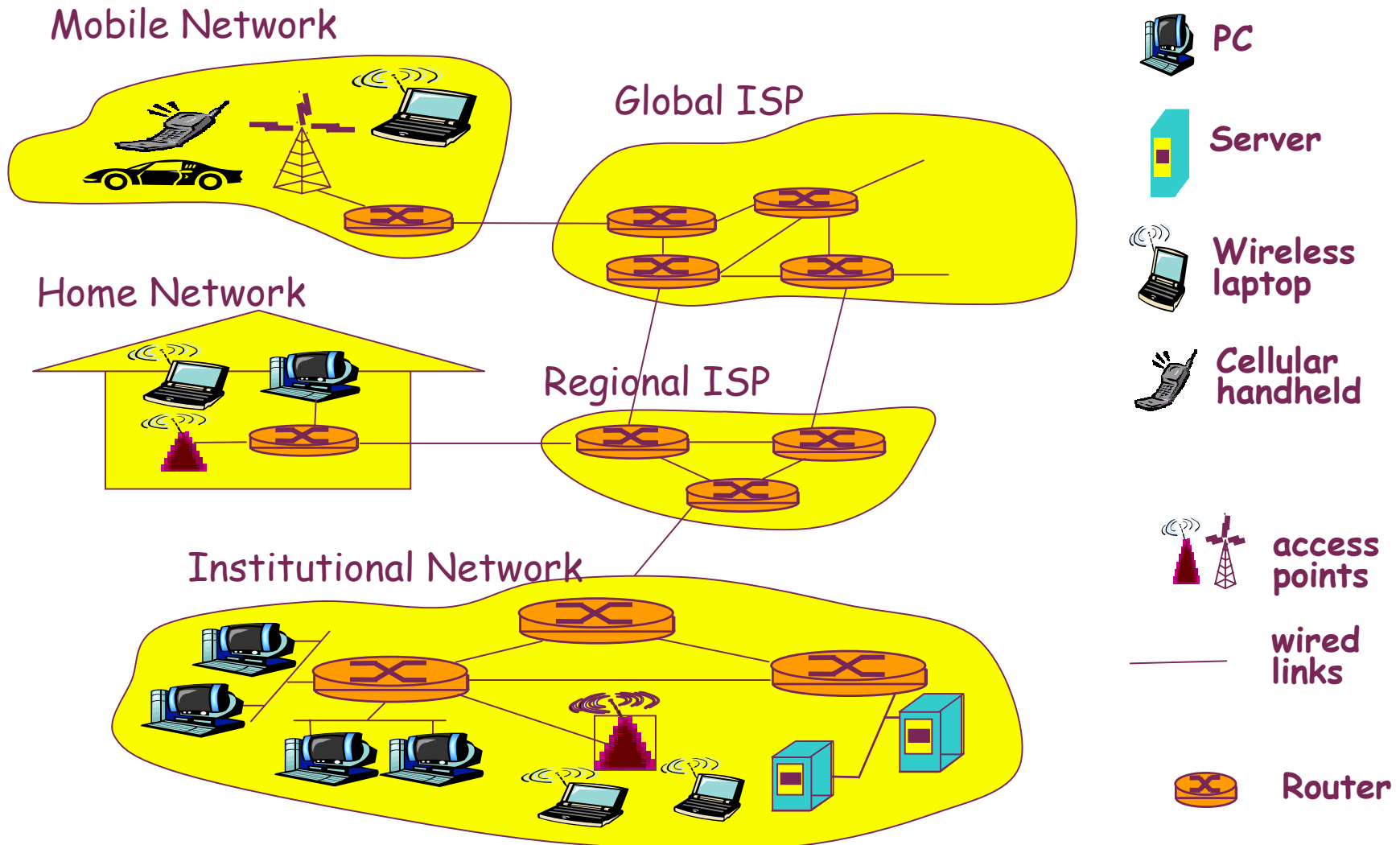
- Signal carried in electromagnetic spectrum
- No physical “wire”
- Bi-directional
- Propagation environment effects:
  - Reflection
  - Obstruction by objects
  - Interference

## Radio link types:

- **Terrestrial Microwave**
  - e.g. up to 45 Mbps channels
- **WLAN** (e.g., Wi-Fi)
  - 2Mbps, 11Mbps, 54 Mbps
- **Wide-area** (e.g., cellular)
  - e.g. 3G: hundreds of kbps
- **Satellite**
  - Kbps to 45Mbps channel (or multiple smaller channels)
  - 270 msec end-end delay
  - GEO/LEO

# What is the Internet?

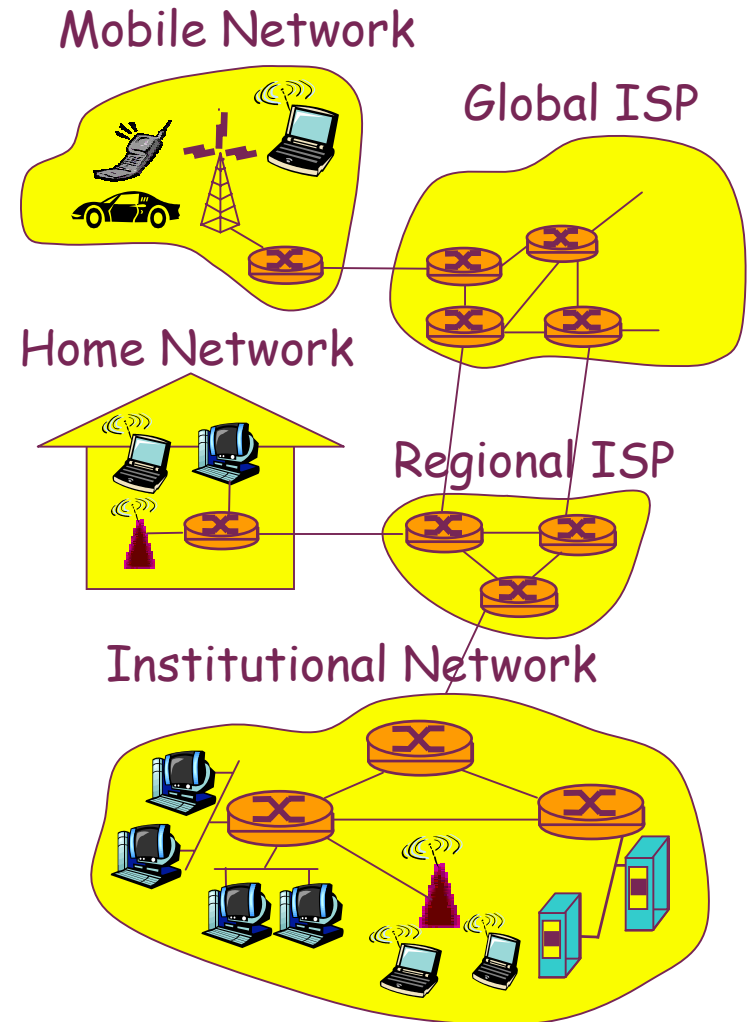
## Nuts and Bolts View



# What is the Internet?

## Nuts and Bolts View

- Millions of connected computing devices
  - Hosts = end systems, run Network Applications
- Communication links
  - Fiber, Copper, Radio, Cable, Satellite, etc...
  - Provides Bandwidth
- Routers:
  - Forward Packets (chunks of data) from source to destination

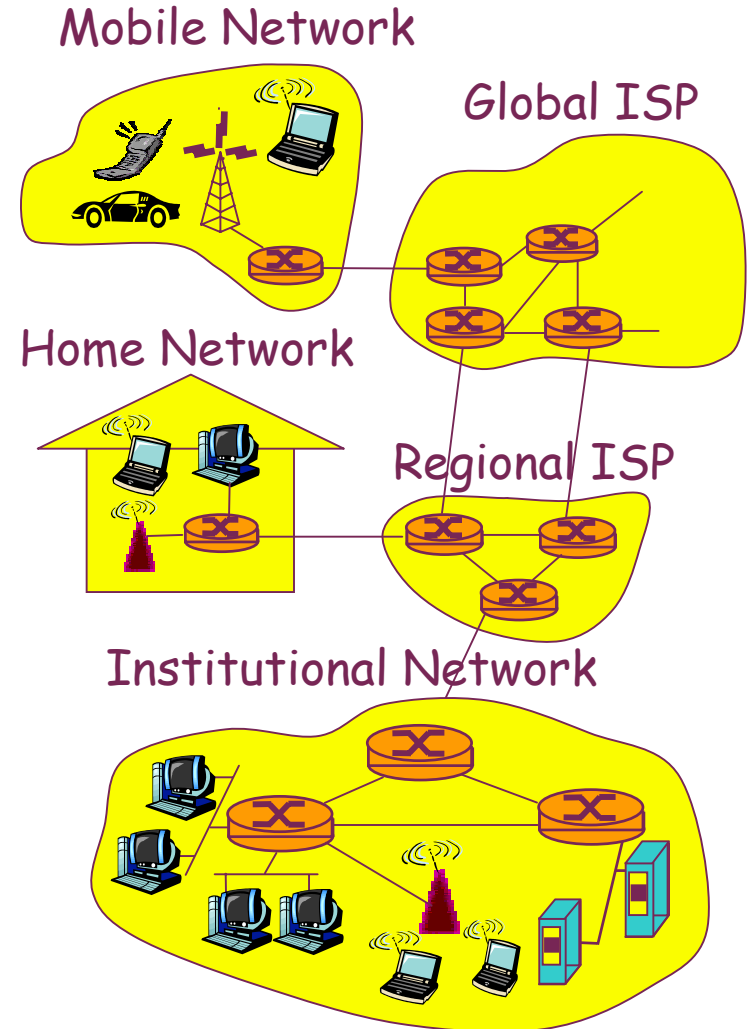




# What is the Internet?

## Nuts and Bolts View (Continued)

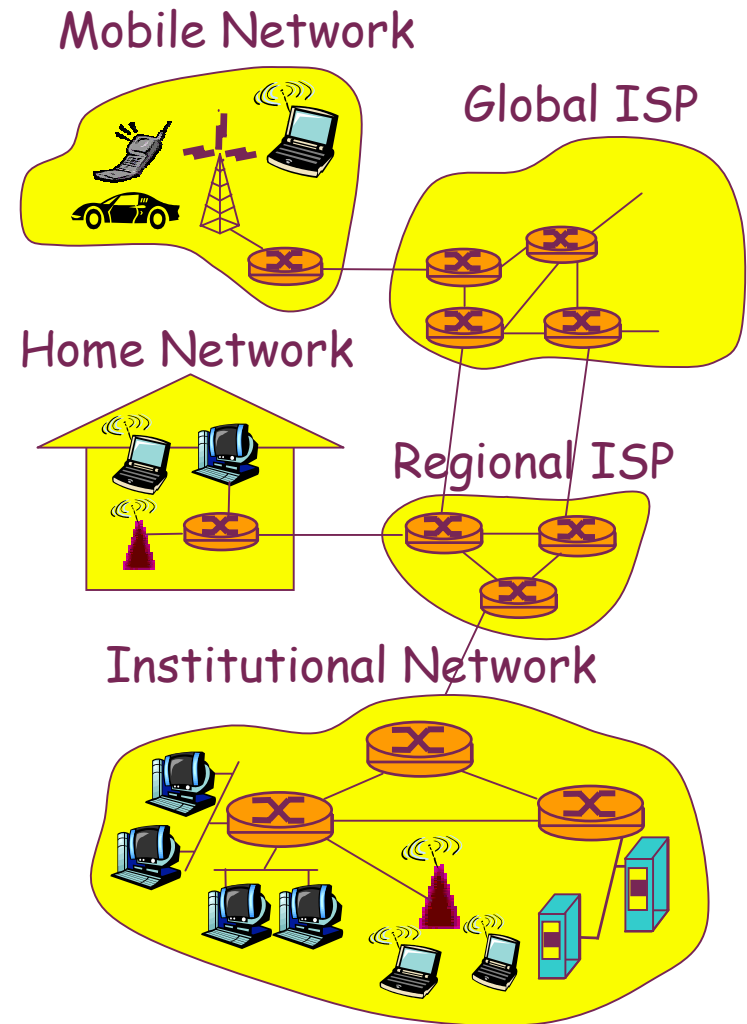
- **Protocols** control sending, receiving of msgs
  - TCP, IP, HTTP, Ethernet
- **Internet:**
  - "Network of networks"
  - Loosely hierarchical
  - Public Internet v.s. Private Intranet
- **Internet standards**
  - RFC: Request for comments
  - IETF: Internet Engineering Task Force



# What is the Internet?

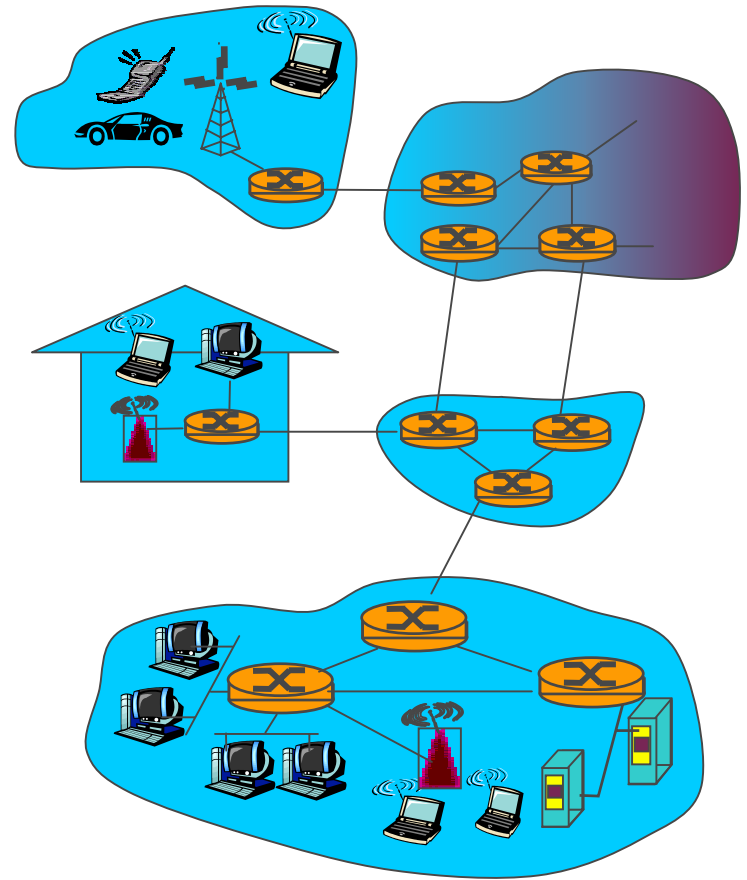
## A Service View

- **Communication Infrastructure** enables distributed applications:
  - Web, VoIP, email, games, e-commerce, file sharing
- **Communication services provided to applications include**
  - Reliable data delivery from source to destination
  - "Best effort" (unreliable) data delivery



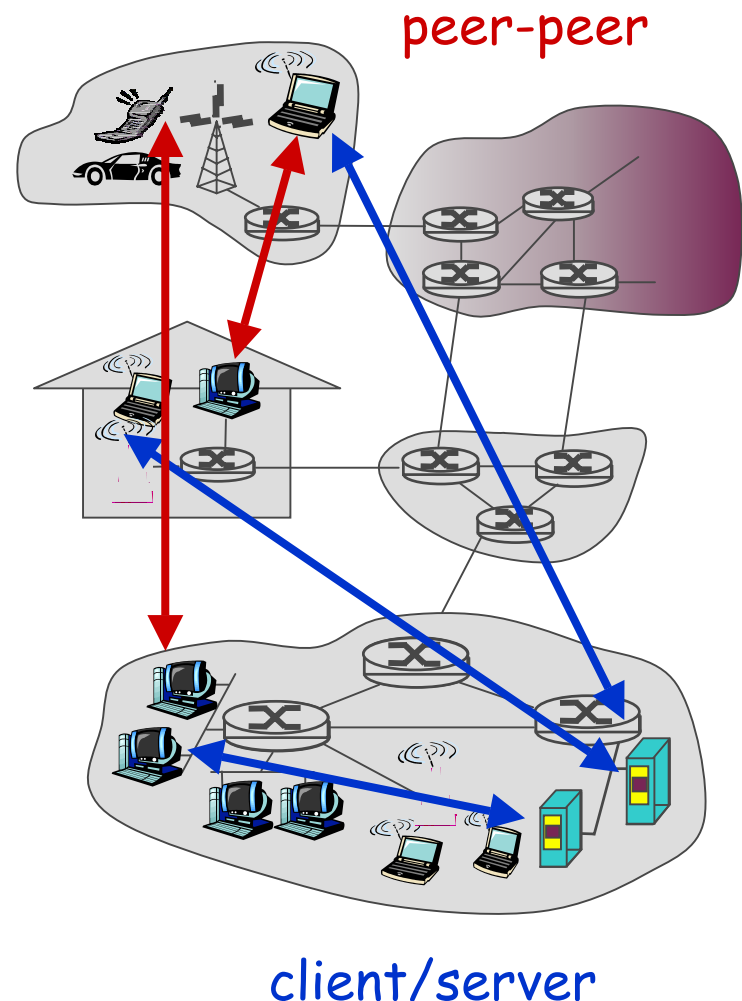
# Network Structure

- Network edge:  
applications and hosts
- Access Networks,  
physical media: wired,  
wireless communication  
links
- Network core:
  - interconnected routers
  - Network of networks



# Network Edge

- **End systems (hosts):**
  - run application programs
  - e.g. Web, email
  - at “edge of network”
- **Client-Server model**
  - client host requests, receives service from always-on server
  - e.g. Web browser/server; email client/server
- **Peer-to-peer model:**
  - Minimal (or no) use of dedicated servers
  - e.g. Kazaa, Gnutella, Skype, Napster,...



# Network Edge Services (I)

## "Reliable Service"

- Goal: data transfer  
between end systems
- handshaking: setup  
(prepare for) data  
transfer ahead of time
    - Hello, initial  
establishment
    - **set up "state"** in two  
communicating hosts
  - TCP - Transmission  
Control Protocol
    - Internet's reliable  
data transfer service

### TCP service

- reliable, in-order byte-  
stream data transfer
  - loss:  
acknowledgements  
and retransmissions
- flow control:
  - sender won't  
overwhelm receiver
- congestion control:
  - senders "slow down  
sending rate" when  
network congested

# Network Edge Services (II)

## Best Effort “Unreliable” Service

Goal: data transfer

between end systems

- same as before!

- **UDP** - User Datagram

Protocol:

- connectionless
- unreliable data transfer
- no flow control
- no congestion control

App's using TCP:

- HTTP (Web), FTP (file transfer), Telnet (remote login), SMTP (email)

App's using UDP:

- streaming media, teleconferencing, DNS, Internet telephony

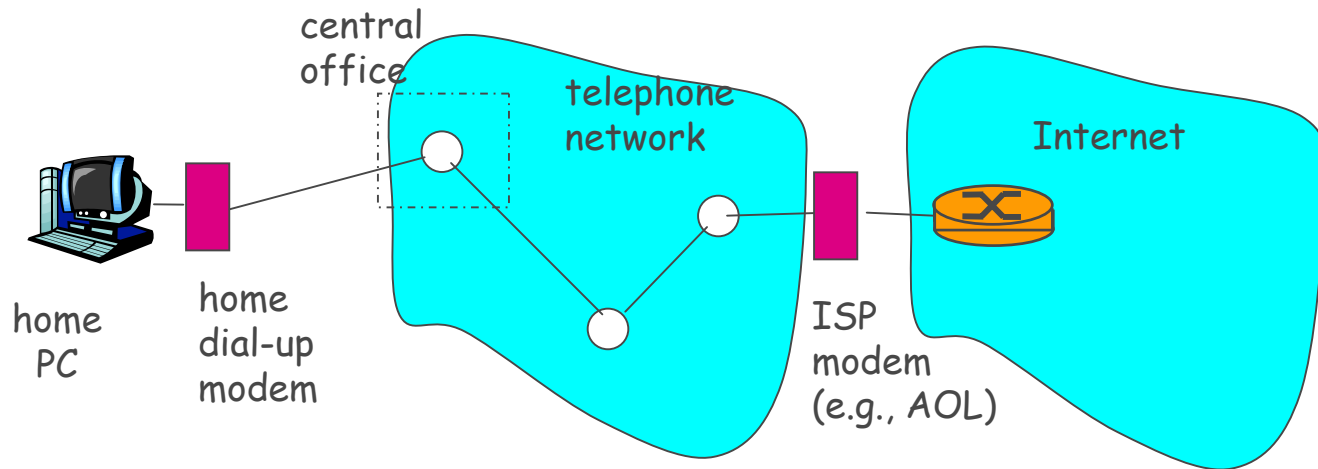
# Access Networks

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Q: How to connect end systems to edge router?

- Residential Access Networks
- Institutional access networks (school, company)
- Mobile access networks
  - Keep in mind:
  - bandwidth (bits per second) of access network?
  - shared or dedicated?

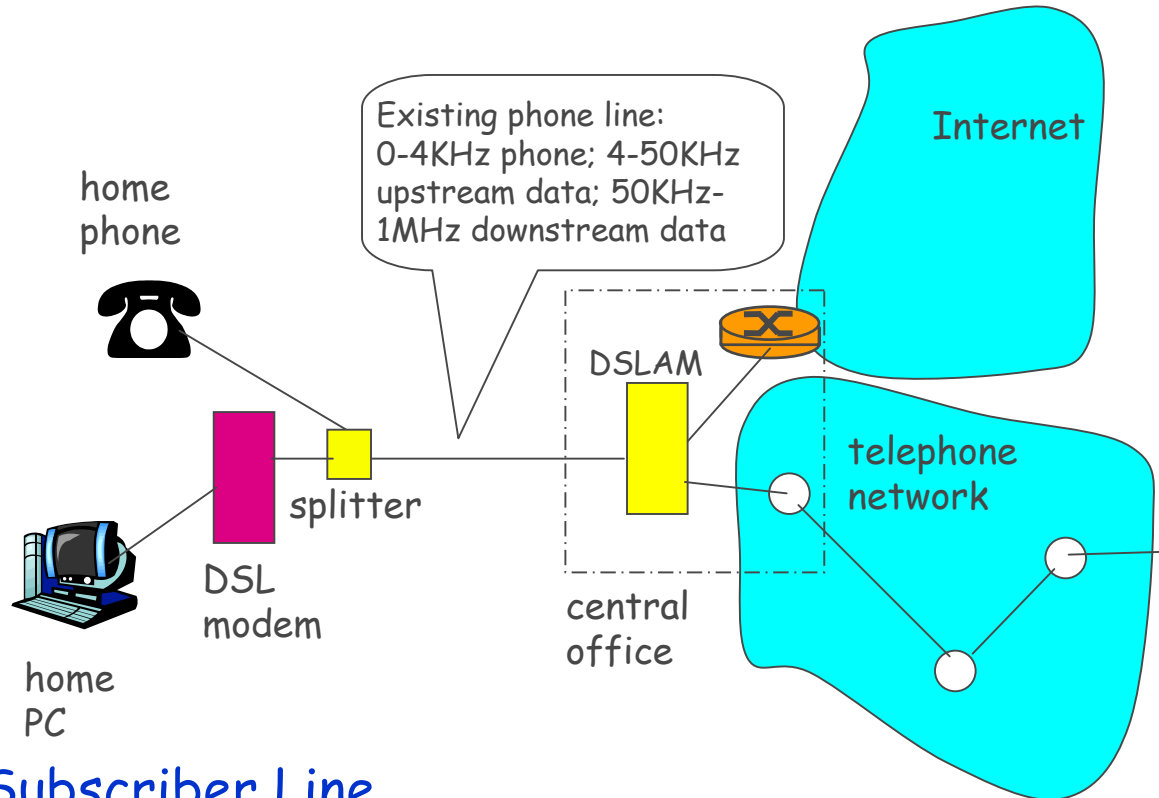
# Dial-up Connection



- ❖ Uses existing telephony infrastructure
- ❖ Home is connected to central office
- ❖ up to 56Kbps direct access to router (often less)
- ❖ Can't surf and phone at same time: not "always on"

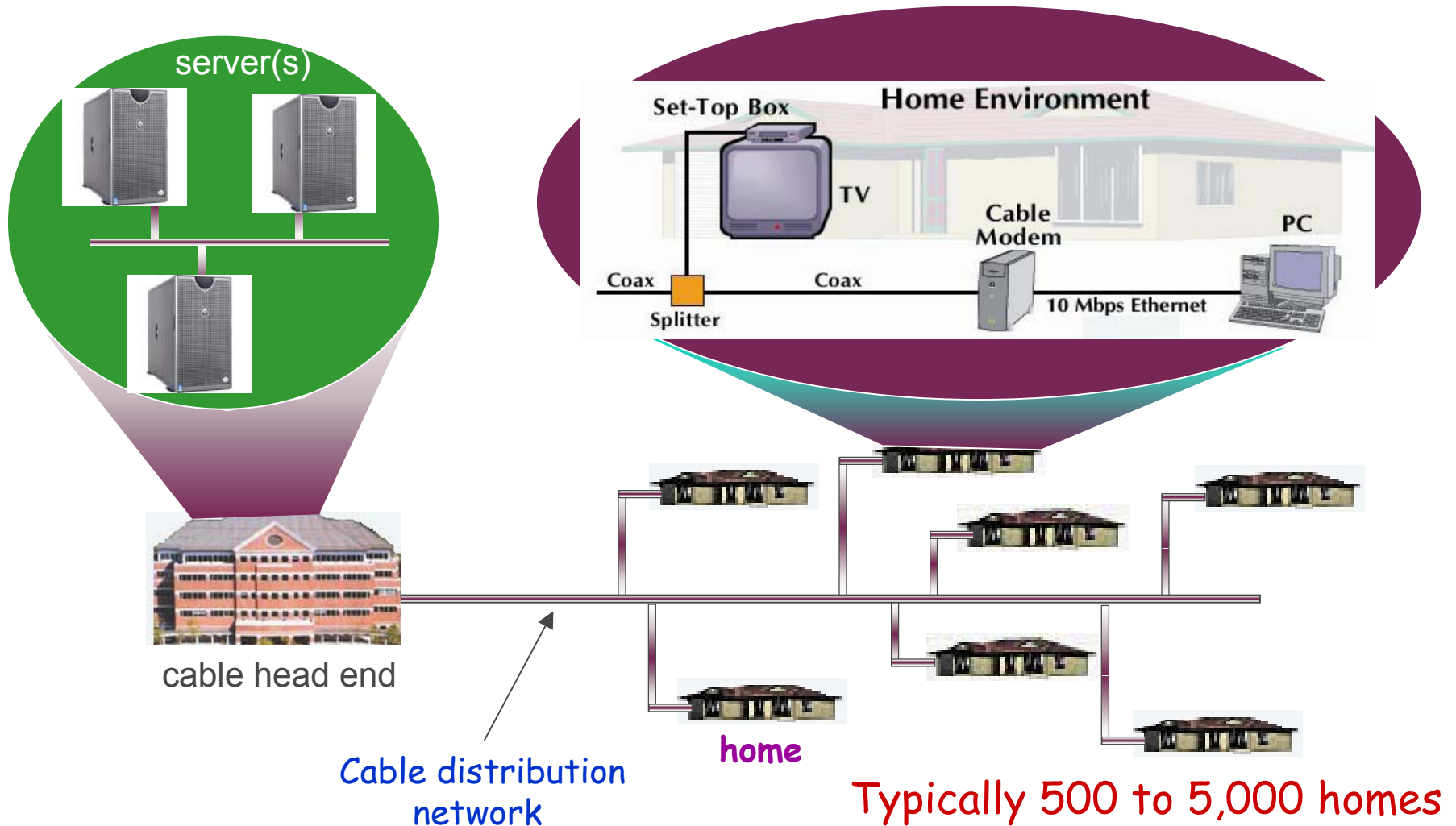


# DSL Connection

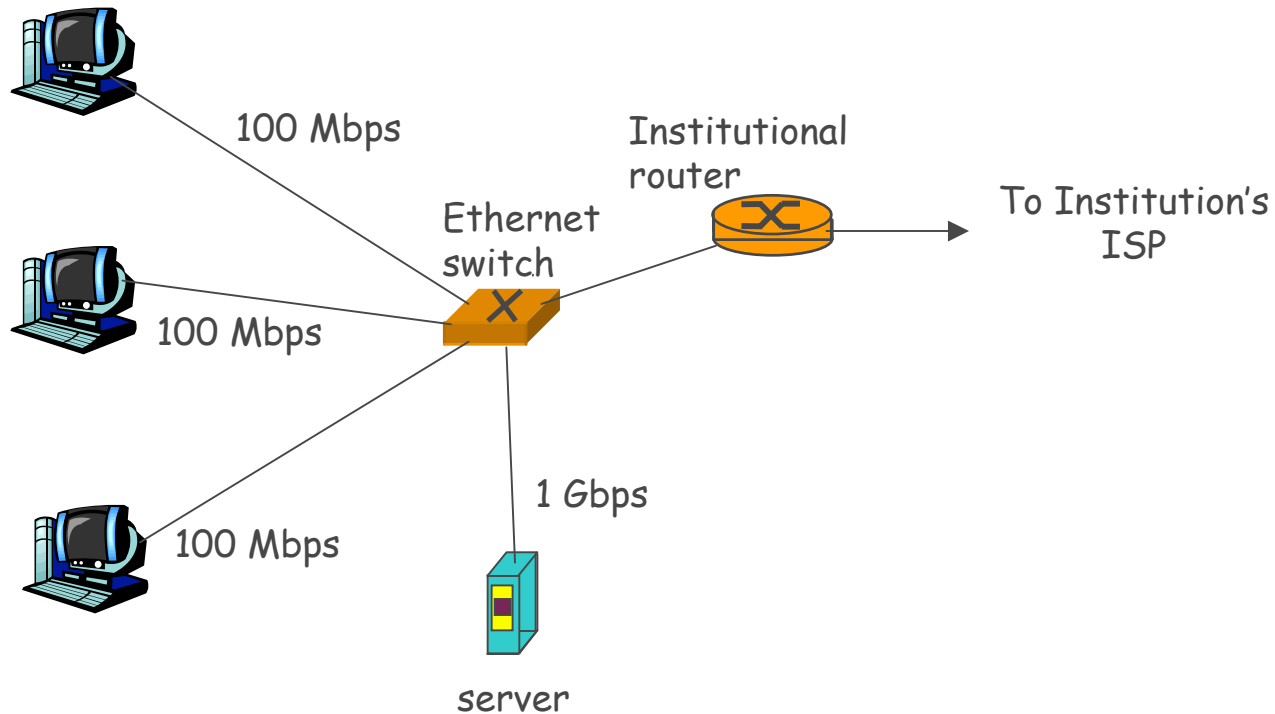


- Digital Subscriber Line
  - deployment: telephone company (typically)
  - up to 1 Mbps upstream (today typically < 256 kbps)
  - up to 8 Mbps downstream (today typically ~ 1 Mbps)
  - Dedicated physical line to telephone central office

# Cable Residential Access



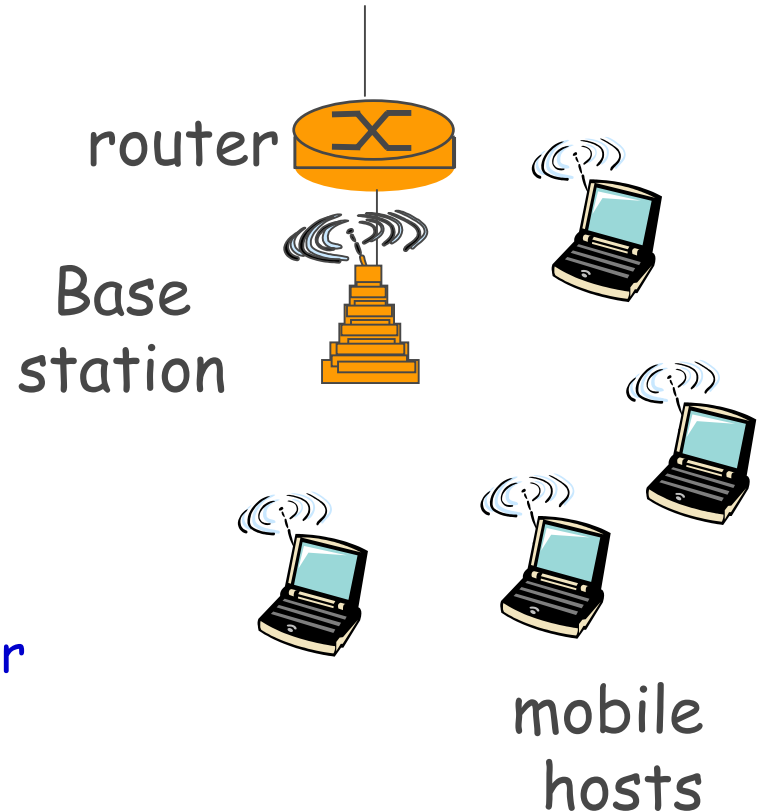
# Ethernet Access



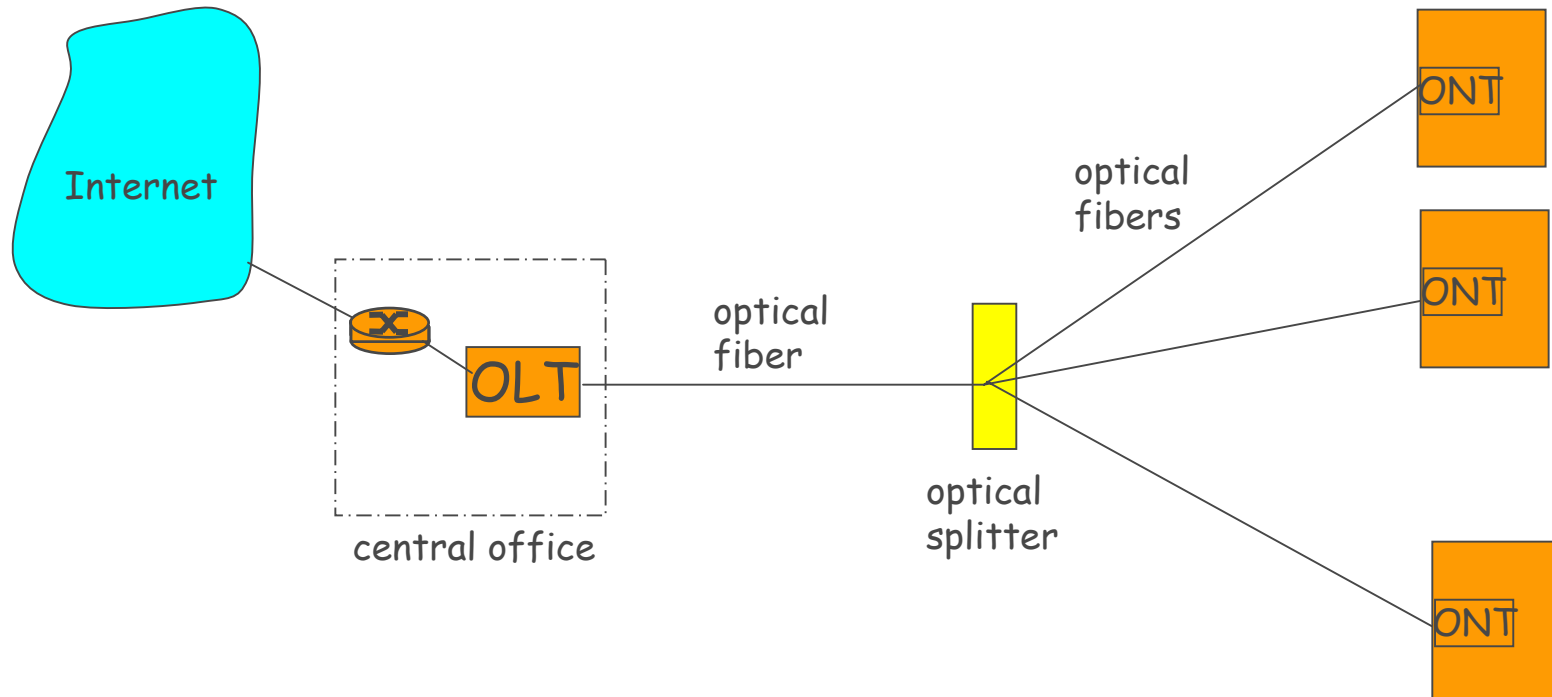
- ❑ Typically used in companies, universities, etc
- ❑ 10 Mbps, 100Mbps, 1Gbps, 10Gbps Ethernet
- ❑ Today, end systems typically connect into Ethernet switch

# Wireless Access

- shared wireless access network connects end system to router
  - via base station aka "access point"
- **wireless LANs:**
  - 802.11ab/g (Wifi): 11/54 Mbps
- **wider-area wireless access**
  - provided by Telco operator
  - ~1Mbps over cellular system (EVDO, HSDPA)
  - Next up: WiMax (10's Mbps) over wide area



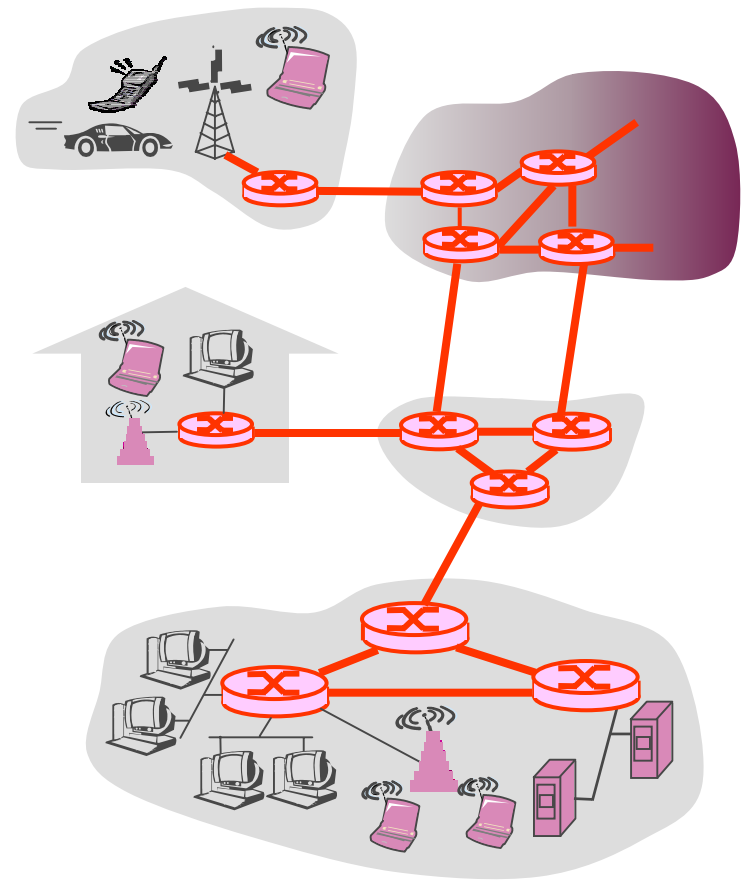
# Fiber to the Home (FTTH)



Much higher Internet rates; fiber also carries television and phone services

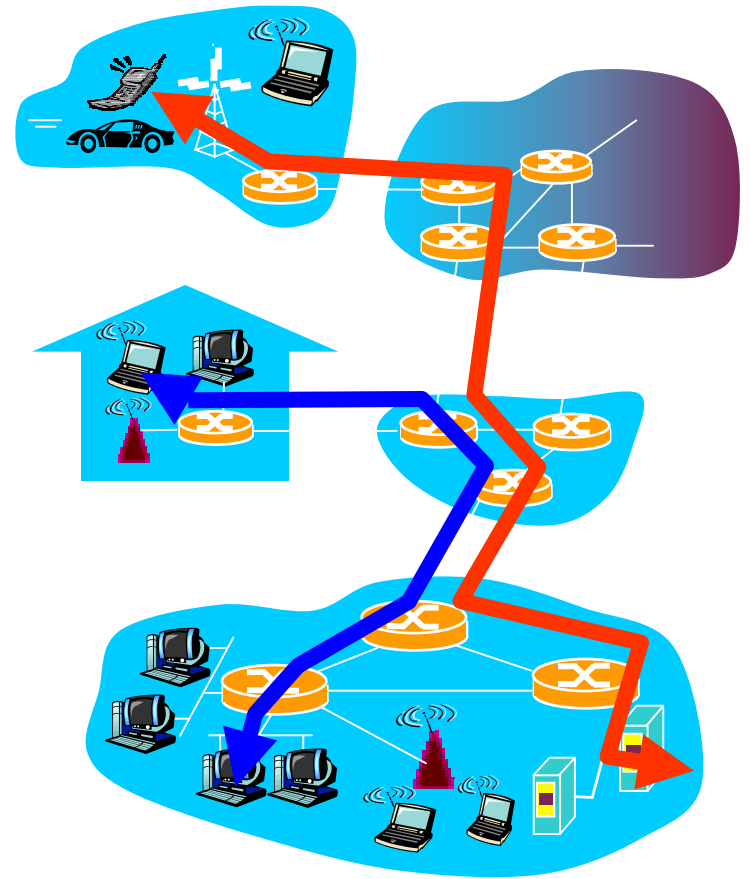
# The Core Network

- Mesh of interconnected routers
- The fundamental question: how is data transferred through net?
  - Circuit Switching: dedicated circuit per call: telephone network (PSTN)
  - Packet Switching: data sent thru net in discrete "chunks": Internet




# Circuit Switching

- End-end resources reserved for duration of call
- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required
- re-establish call upon failure
- Example: PSTN

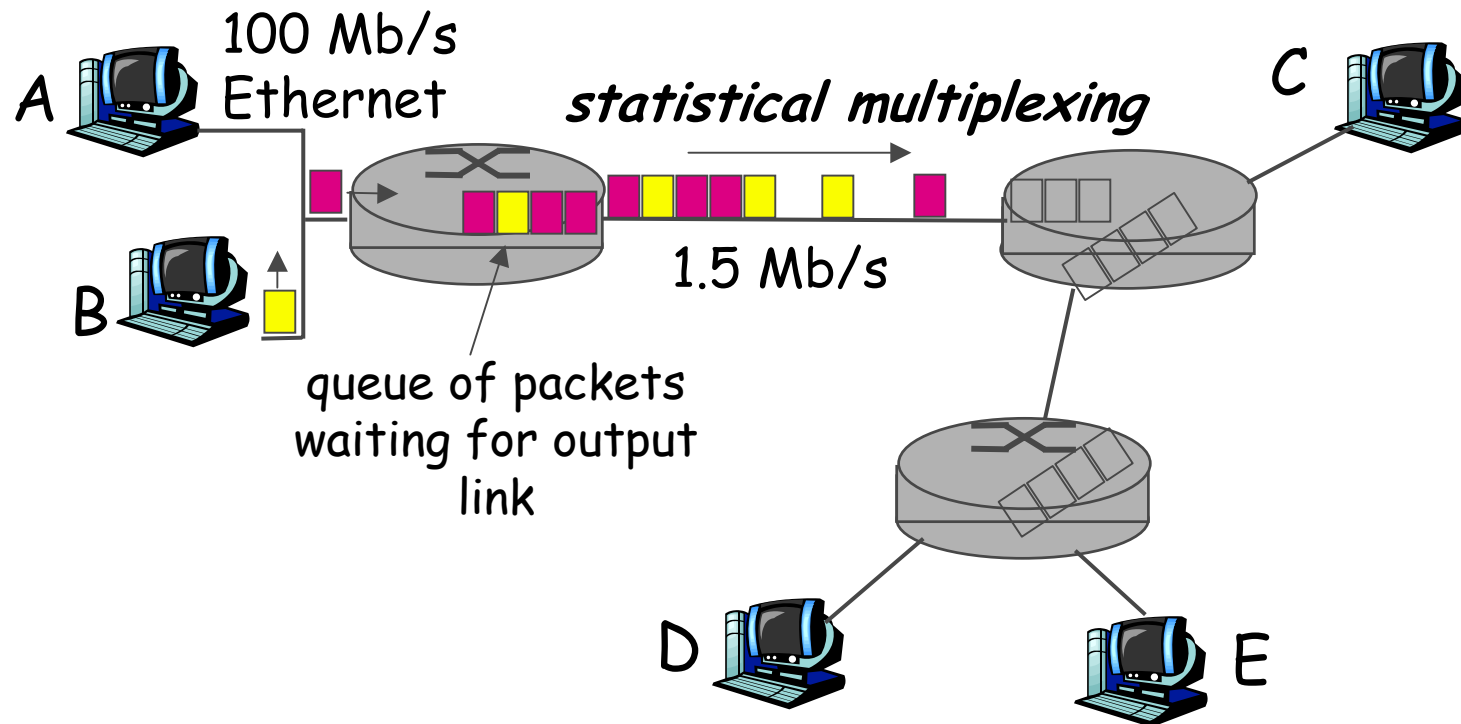


# Packet Switching

- each end-end data stream divided into **packets**
  - user A, B packets **share** network resources
  - Each packet uses full link bandwidth
  - Resources used *as needed*
- Bandwidth division into "pieces"  
Dedicated allocation  
Resource reservation
- 
- **Resource contention:**
    - aggregate resource demand can exceed amount available
    - congestion: packets queue, wait for link use
    - store and forward: packets move one hop at a time Node receives complete packet before forwarding



# Packet Switching (Continued)



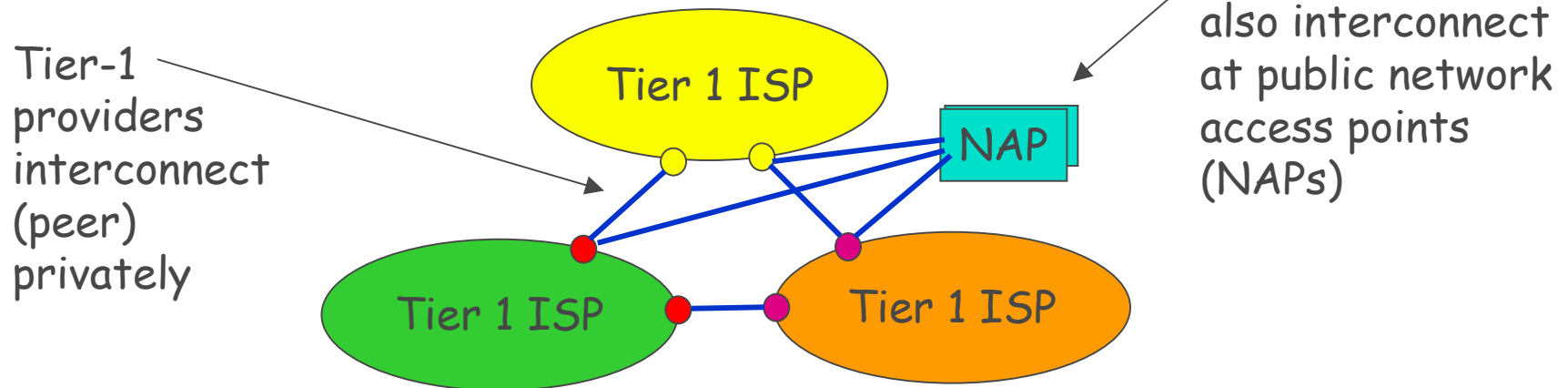
# Packet vs. Circuit Switching

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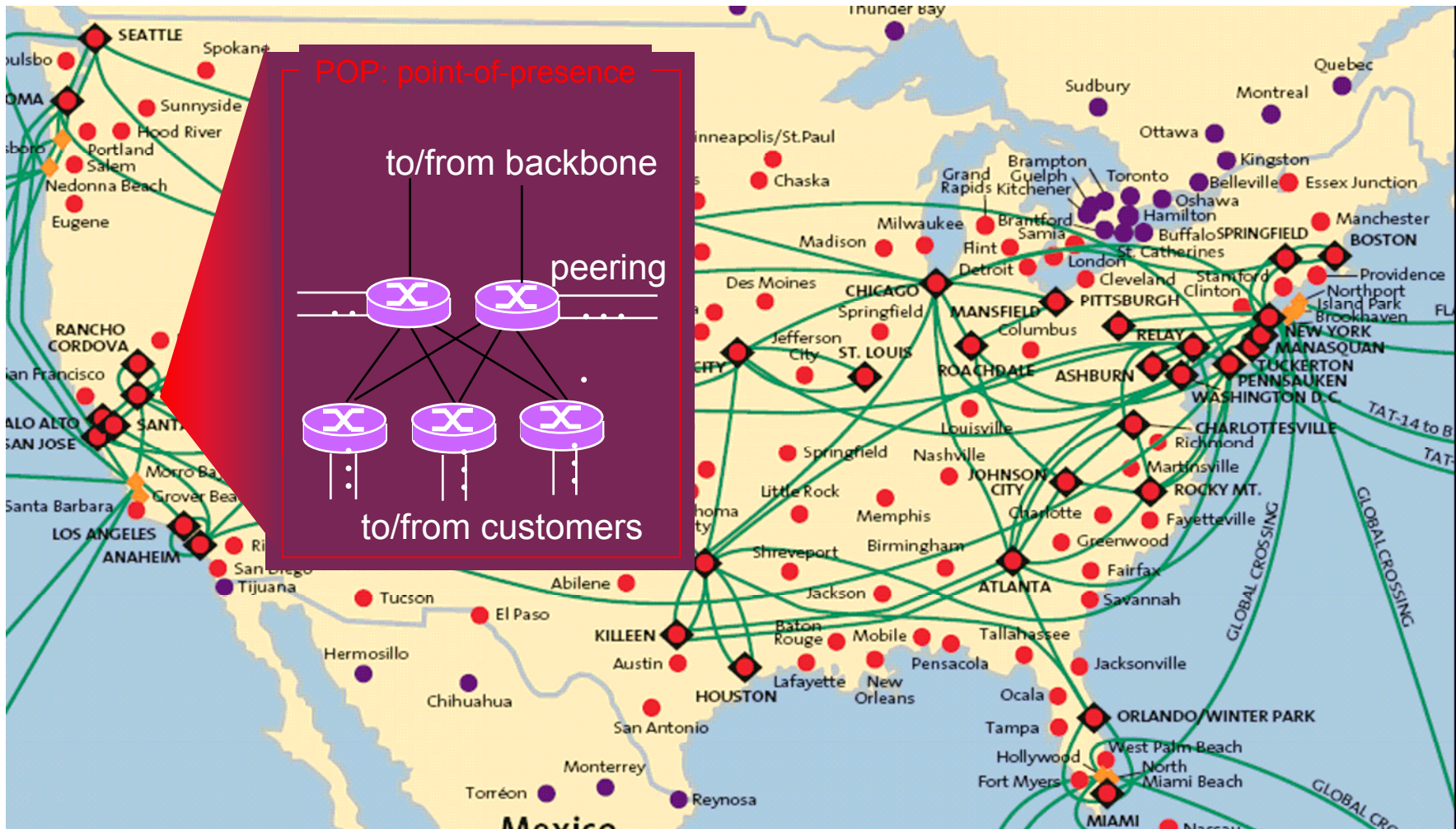
- PS great for bursty data
  - resource sharing (scalable!)
  - simpler, no call setup, more robust (re-routing)
- excessive congestion: packet delay and loss
  - Without admission control: protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
  - Bandwidth guarantees needed for audio/video apps
  - Possible solution: Virtual circuit

# Internet Structure (Tier 1)

- Roughly hierarchical
- at center: "tier-1" ISPs (e.g., Verizon, Sprint, AT&T, Cable and Wireless), national/international coverage
- treat each other as equals



# Tier 1 ISP: Sprint

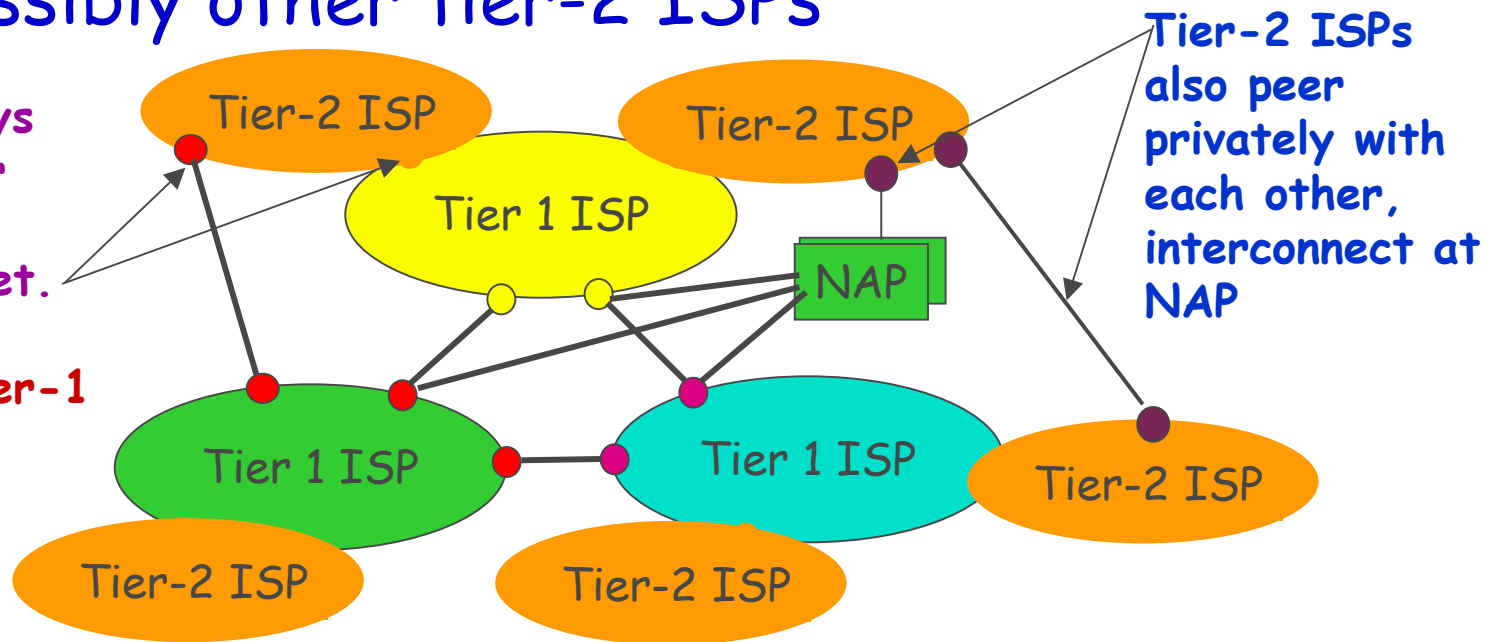


# Internet Structure (Tier 2)

- "Tier-2" ISPs: smaller (often regional) ISPs

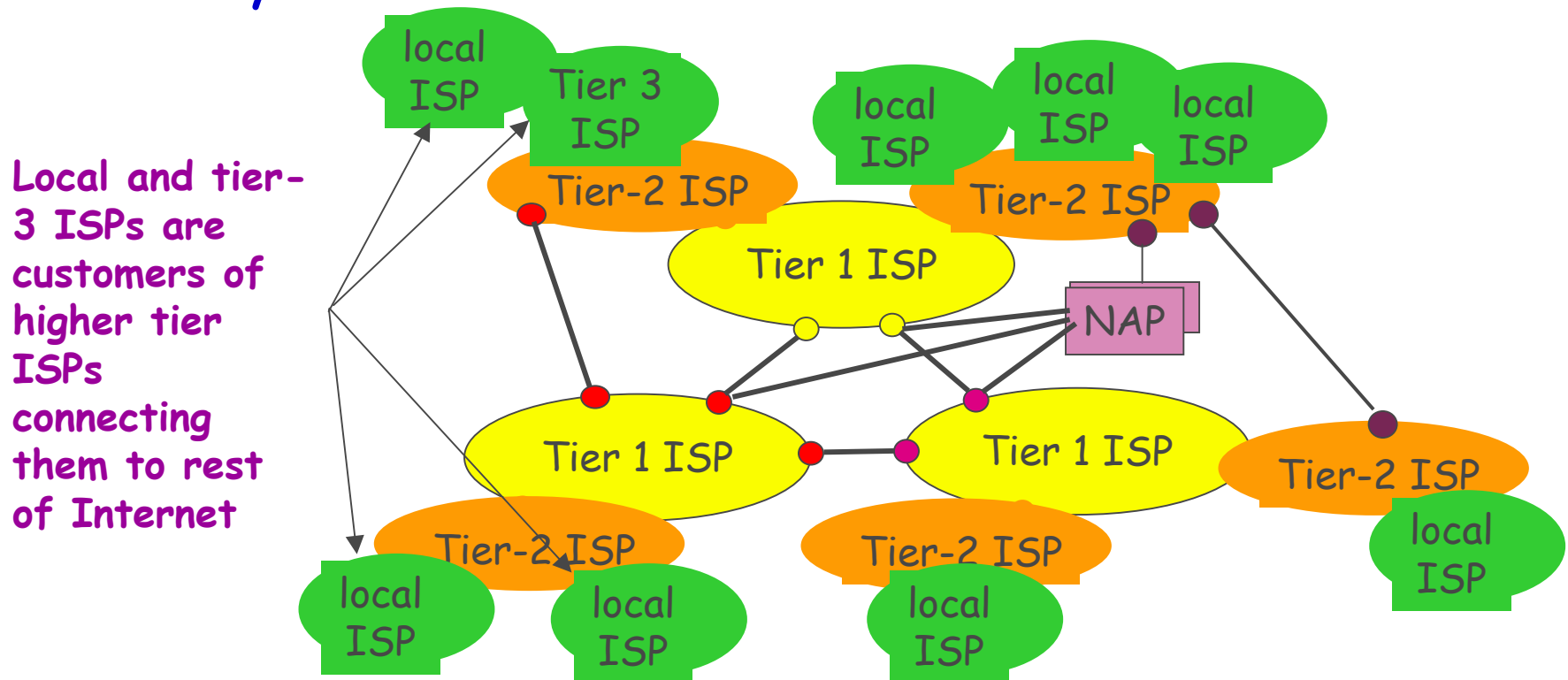
- Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs

Tier-2 ISP pays Tier-1 ISP for connectivity to rest of Internet. Tier-2 ISP is customer of tier-1 provider

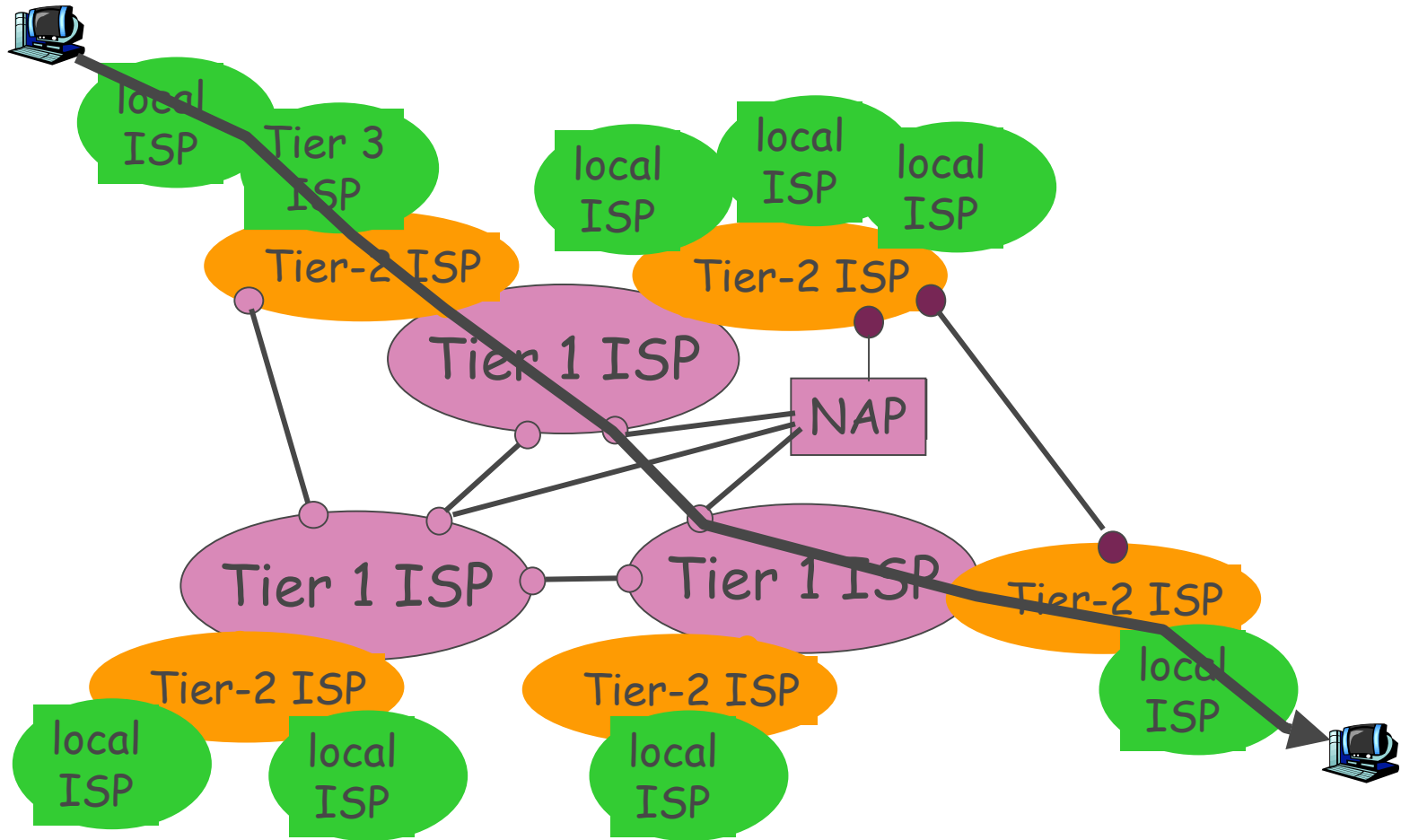


# Internet Structure (Tier 3)

- "Tier-3" ISPs and local ISPs
  - last hop, access network, closest to end systems



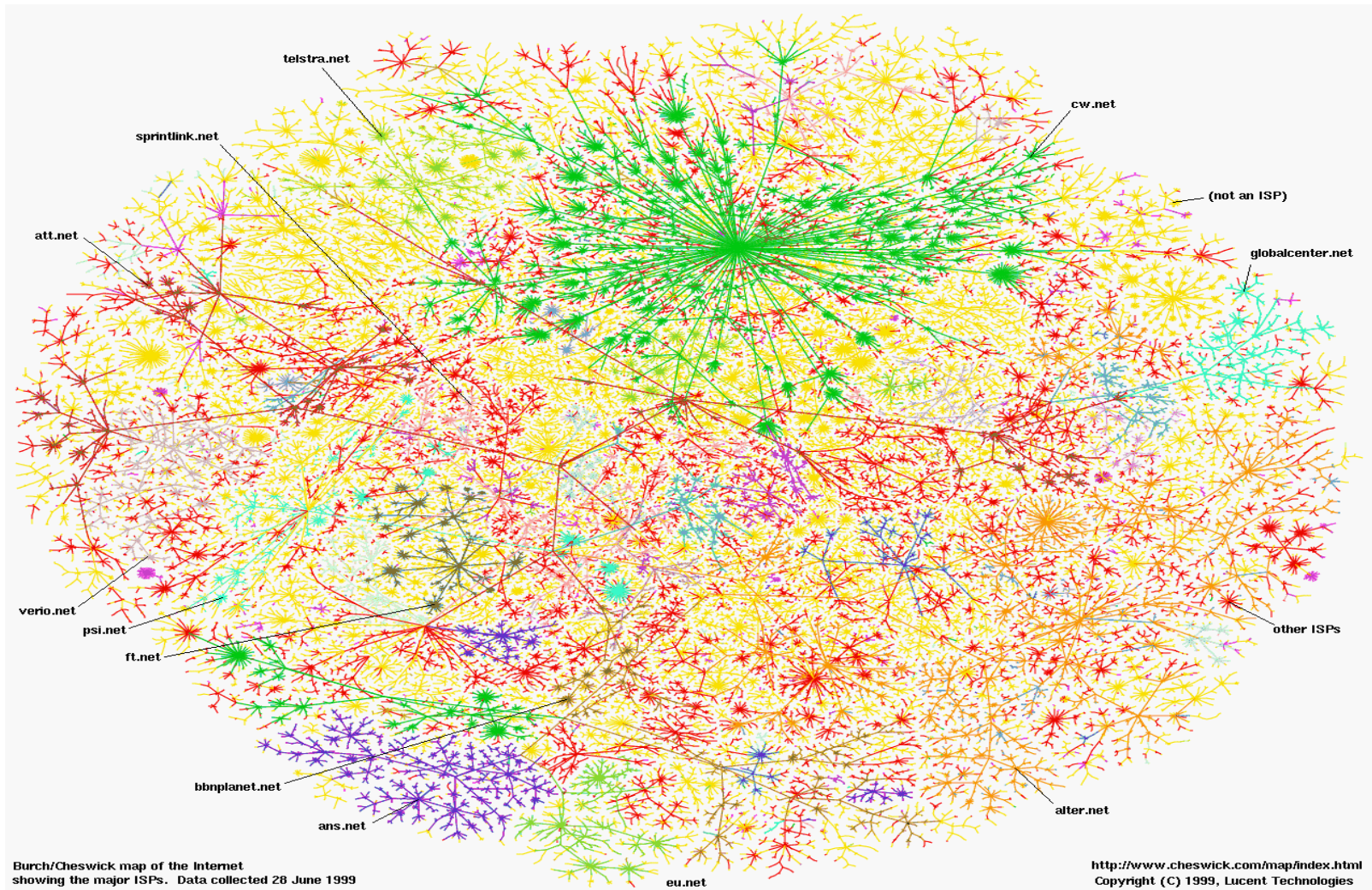
# Internet Structure (Summary)



a Packet passes through many networks!

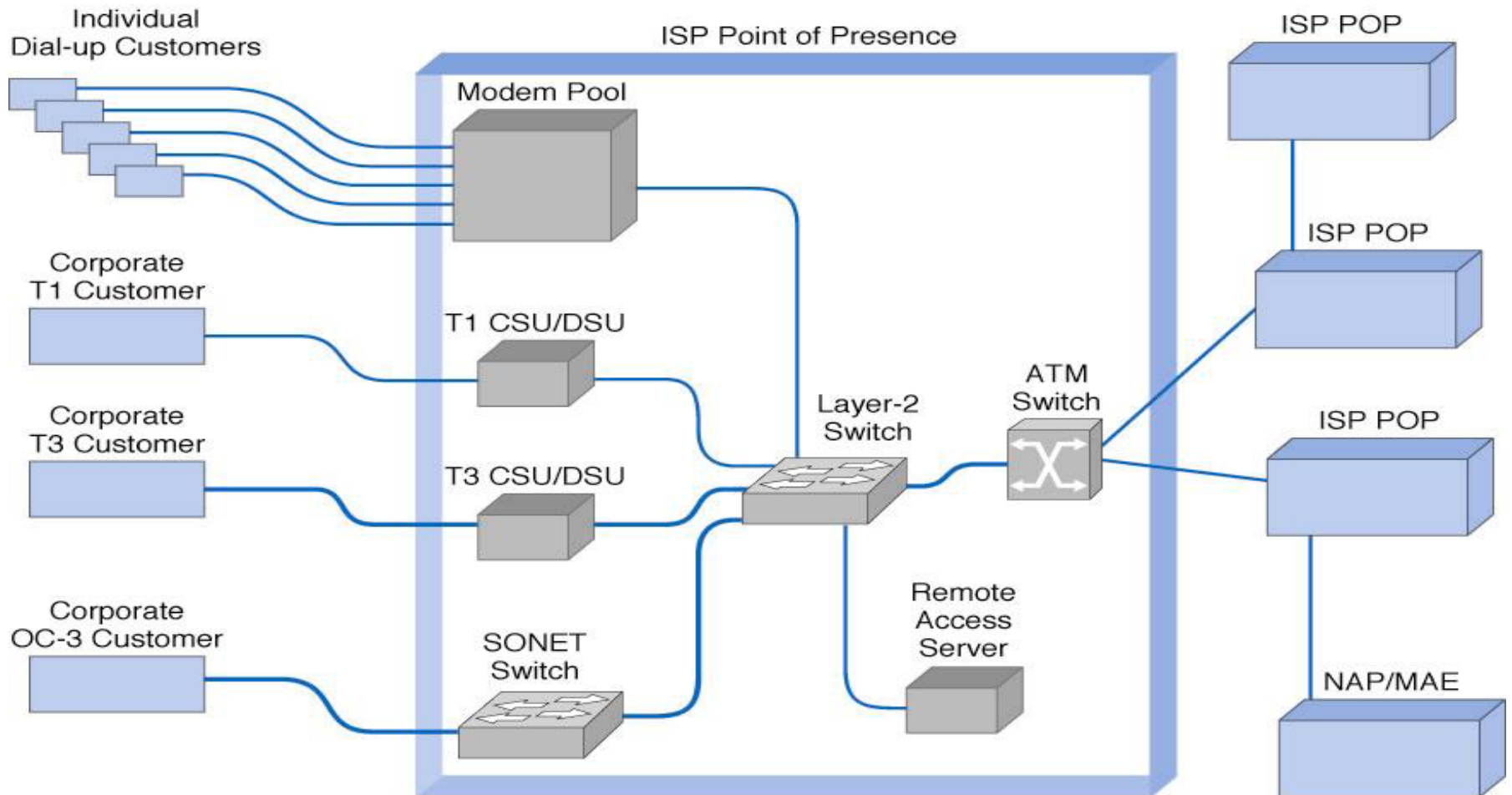


# The Internet ?





# Anatomy of an ISP



# Network Performance Measures

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- The two most important network performance measures are Delay/Latency & Throughput
- End-to-end delay consists of several components
  - Transmission time
  - Propagation delay
  - Nodal processing
  - Queuing delay (Random, depends on network loading, link capacities, disciplines, etc..)

# Transmission Time

- Transmission Time ( $t_{\text{trans}}$ )
  - The time it takes to transmit a group of bits (e.g., a Message/Package/Frame) of bits into a network

$$t_{\text{tran}} = \frac{\text{Number of message bits}}{\text{Data rate [bps]}}$$

# Propagation Delay

- Propagation time ( $t_{\text{prop}}$ )
  - The time it takes for a bit to traverse the link

$$t_{\text{prop}} = \frac{\text{link length}[m]}{v_{\text{prop}} [m/s]}$$

- Example propagation velocities:
  - Air/Free space:  $c = 3 \times 10^8$  meters/sec
  - Cat 5 UTP:  $2 \sim 2.5 \times 10^8$  meters/sec
  - Optical Fiber:  $2 \sim 2.5 \times 10^8$  meters/sec

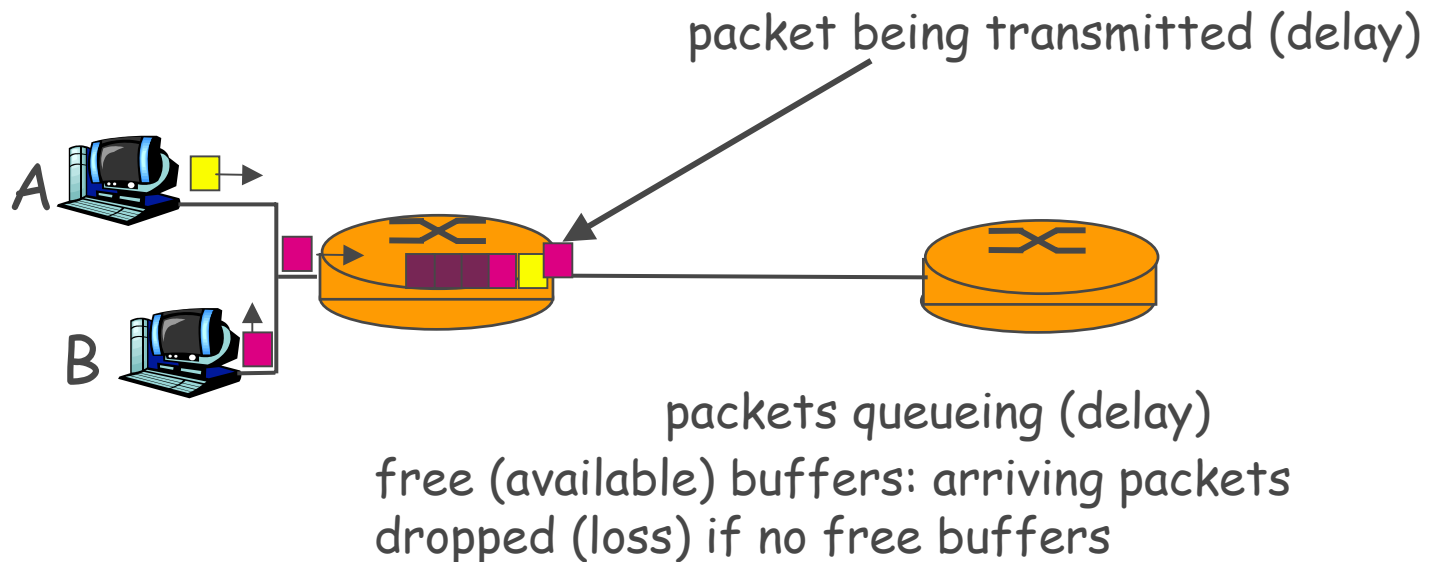
# Nodal Processing/Queueing

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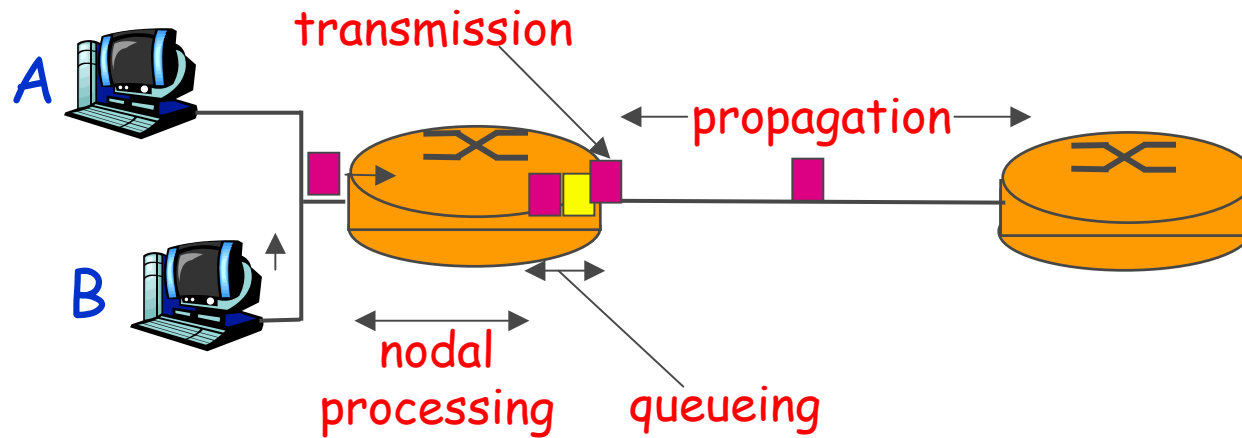
- Nodal processing:
  - Check bit errors
  - Determine output link (Routing decision)
- Queueing
  - Time waiting at output link for transmission
  - Depends on congestion level of router

# Loss and Delay: Why?

- packets *queue* in router buffers
- packet arrival rate to link exceeds output link capacity
- packets queue, wait for turn



# Summary of Delay Components



# Message Transfer Time

- Message Transfer Time ( $t_{xfr}$ ) = Message latency
  - Time for sender to transmit message to the receiver and for the receiver to receive the entire message. Also known as the end-to-end delay

$$t_{xfr} = t_{trans} + t_{prop} + t_{queuing/processing}$$

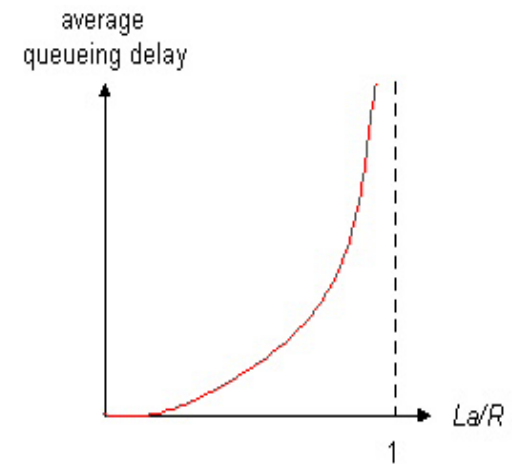


# More on Queuing Delay

- $R$  = link bandwidth (bps)
- $L$  = packet length (bits)
- $\lambda$  = average packet arrival rate (Packets/sec)

**Traffic intensity =  $\lambda L/R$**

- $\lambda L/R \sim 0$ : average queueing delay
- $\lambda L/R \rightarrow 1$ : delays become large
- $\lambda L/R > 1$ : more “work” arriving than can be serviced, average delay infinite!

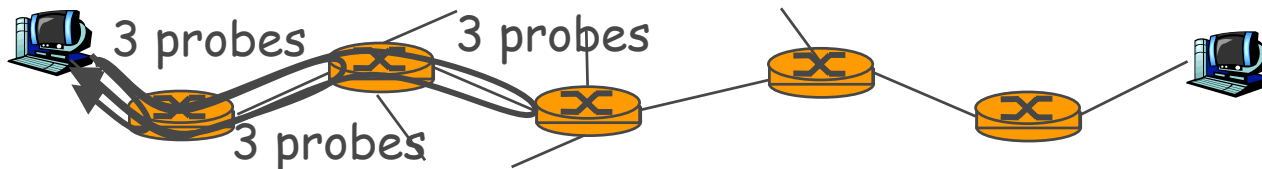


# Round Trip Time (RTT)

- **Round Trip Time:** The time to send a message from a sender to the receiver and receive a response back
- RTT depends on message size, length of link, direction of propagation, propagation velocity, network node processing, network loading, etc...
- For simplicity, RTT is normally assumed to be twice the end-to-end propagation delay although this might not be true if the message and the response traverses different links

# Real Internet Delays & Routes

- What do “real” Internet delay & loss look like?
- Traceroute program: provides delay measurement from source to router along end-end Internet path towards destination. For all  $i$ :
  - sends three packets that will reach router  $i$  on path towards destination
  - router  $i$  will return packets to sender
  - sender times interval between transmission and reply.

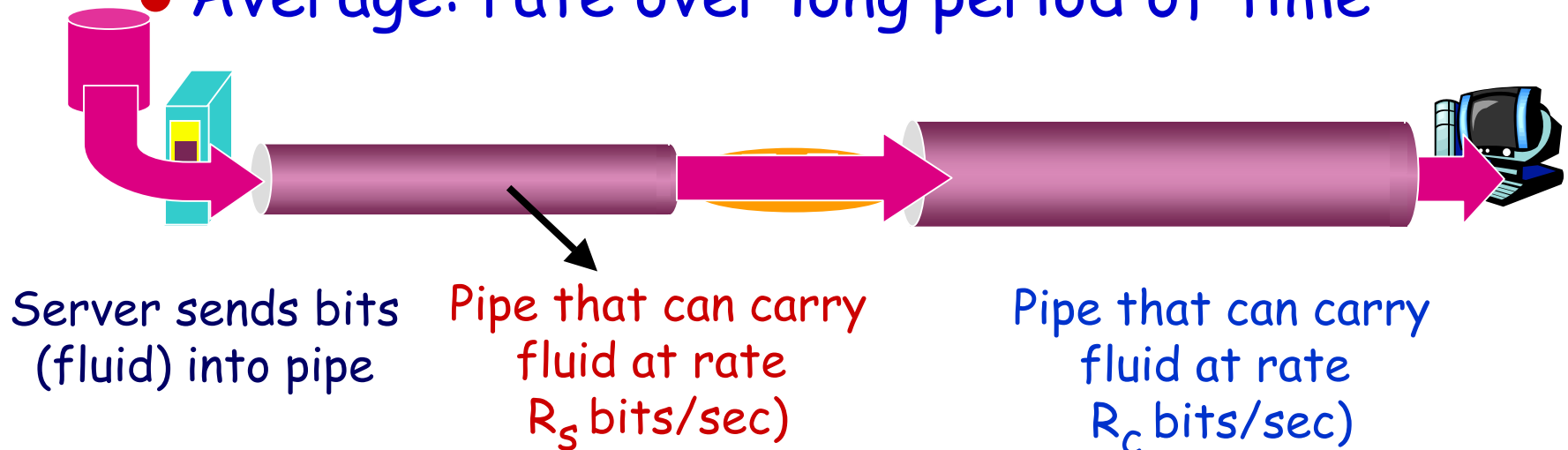


# Throughput

- The **Throughput** is defined as the number of information bits that can be transferred reliably over a certain period of time. It is measured in "bps"
- The throughput is the carried load and it is not equal to the offered load
- Protocols add overhead bits and time delays in addition to the transmission time of the actual information bits. That would result in reduced throughput.
- Link errors are result in reduced throughput

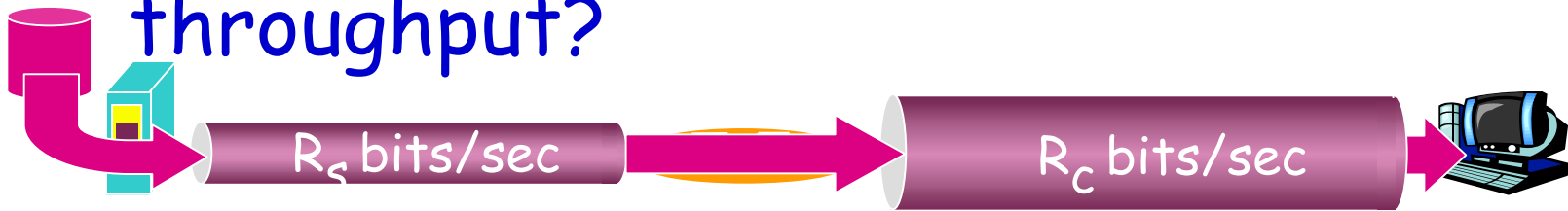
# Instantaneous vs. Average Throughput

- Throughput: rate (bits/time unit) at which bits transferred between sender/receiver
  - Instantaneous: rate at given point in time
  - Average: rate over long period of time

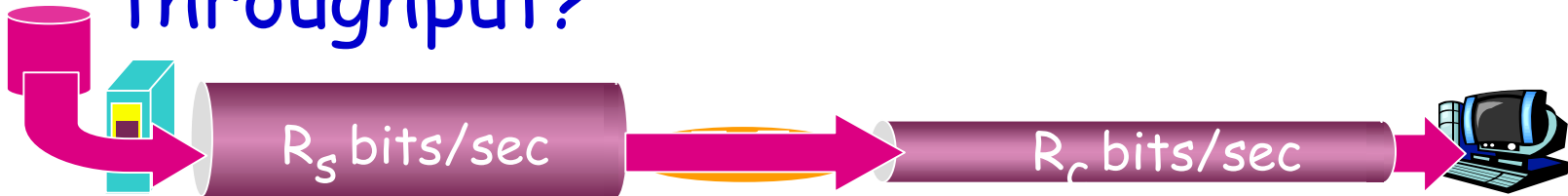


# More on Throughput

- $R_s < R_c$  What is average end-end throughput?



- $R_s > R_c$  What is average end-end throughput?

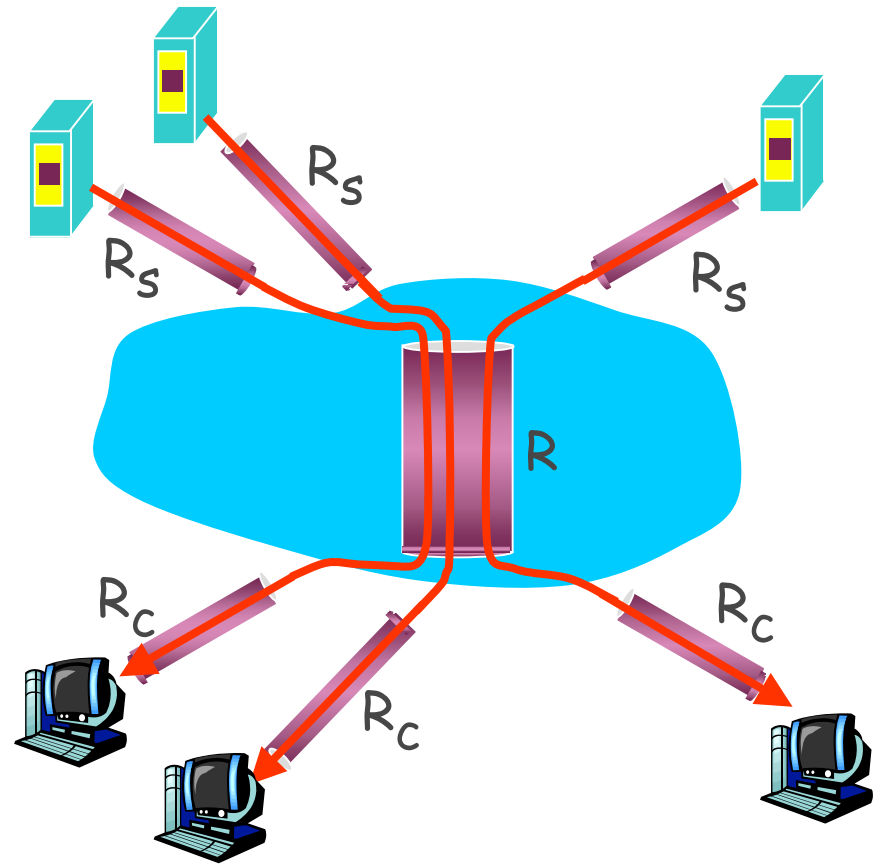


Bottleneck Link

link on end-end path that constrains end-end throughput

# Example on Throughput

- per-connection end-end throughput:  
 $\min(R_c, R_s, R/3)$
- In practice:  $R_c$  or  $R_s$  is often bottleneck
- Trunks have huge BW (i.e.  $R$  is v. Large)



3 connections (fairly) share backbone bottleneck link  $R$  bits/sec

# Bandwidth/Capacity

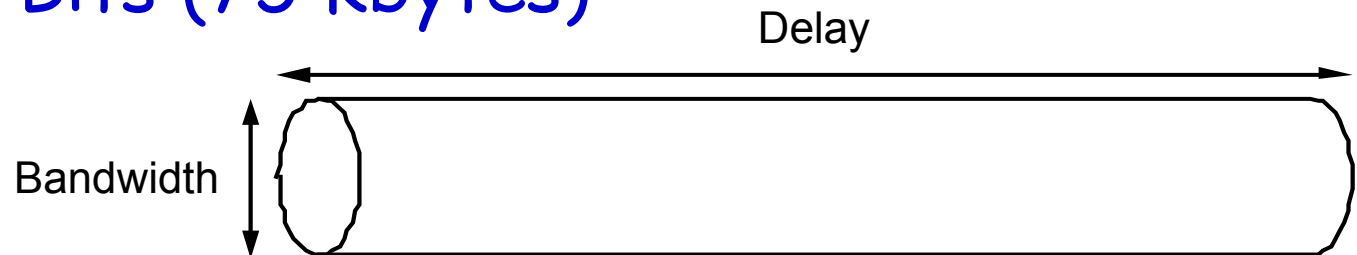
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- The bandwidth or the data rate is the number of bits that can be transmitted over a certain period of time.
  - For example, 10 Mbps means that 10 million bits are transmitted every seconds.
- Link Capacity is the maximum data rate possible on the link with negligible error rate (Shannon Theorem, to be discussed later)



# Bandwidth X Delay Product

- Pipe Size: The maximum amount of data present on the line, usually in an interval of RTT
- Example: If the line bandwidth (data rate) is 10 Mbps and the end-to-end delay is 30 msec, the amount of data found on the line is 600K Bits (75 Kbytes)



# Networking Perspective

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- Application Programmer / End User
  - Guaranteed timely, reliable and recognizable delivery of message/information
- Network Designer
  - Cost-effective design. Resources (Bandwidth, Memory and CPUs) must be used efficiently and are fairly allocated
- Network Provider
  - Administration & management effort, fault detection/fault isolation, easy to account for usage

# The Internet Today

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- ~ 500 million hosts
- Voice, Video over IP
- P2P applications: Napster, BitTorrent (file sharing) Skype (VoIP), PPLive (video)
- More applications: YouTube, gaming, social networking
- Wireless, mobility, networked embedded sensors,...