

Matrix Reduction in Persistent Homology: Theory, Examples, and Applications

Sahil Kumar
PhD in Mathematics and Artificial Intelligence
Yeshiva University, New York, USA
skumar4@mail.yu.edu

ABSTRACT

Matrix reduction is a fundamental computational process in persistent homology, a key method in topological data analysis (TDA) for studying the shape of data. It simplifies boundary matrices to reveal topological features such as connected components, loops, and voids, allowing the computation of Betti numbers and the Euler characteristic, which describe the structure of a space. This report elaborates on the theoretical foundations of matrix reduction, including the Euler-Poincaré formula, and demonstrates its application through worked examples and Python implementation. By reducing the complexity of boundary computations, matrix reduction enables efficient analysis of real-world data in fields such as biology, image processing, and network science. References include foundational texts like Edelsbrunner and Harer’s *Computational Topology: An Introduction* and recent computational advancements.

1 Introduction

Topological Data Analysis (TDA) is an innovative field that applies principles from algebraic topology to analyze and interpret the structure of data. Among its cornerstone techniques is persistent homology, which identifies and quantifies topological features such as connected components, loops, and voids across multiple scales. Central to this process is *matrix reduction*, a computational method that simplifies boundary matrices, enabling efficient extraction of topological information from high-dimensional data.

A simplicial complex, composed of vertices, edges, triangles, and higher-dimensional simplices, serves as the fundamental structure for representing relationships within a dataset. These relationships are encoded in boundary matrices, which capture how

higher-dimensional simplices connect to lower-dimensional ones. By reducing these matrices, key topological invariants such as Betti numbers—representing independent topological features—can be calculated efficiently. This reduction process forms the foundation of computational topology³.

The mathematical theory underpinning matrix reduction involves homology groups, chain complexes, and the Euler-Poincaré formula. Homology groups characterize independent cycles and voids within a space, while the Euler-Poincaré formula provides a succinct relationship between these features and the simplicial structure. It is expressed as:

$$\chi = \sum_{p \geq 0} (-1)^p (z_p - b_p),$$

where z_p represents the rank of the cycle group, and b_p denotes the rank of the boundary group. This formula elegantly connects geometric intuition with algebraic constructs, making it a cornerstone of topological analysis².

A notable advancement in matrix reduction is the use of Smith normal form, which facilitates the computation of Betti numbers by transforming boundary matrices into simplified forms. This approach not only reduces computational complexity but also reveals the intrinsic relationships between simplicial elements, enhancing the understanding of complex datasets¹. The Smith normal form has also been extended to handle integer matrices with applications in diverse areas such as cryptography and network analysis⁵.

This report explores matrix reduction through:

- An explanation of the mathematical theory and algorithms behind boundary matrix reduction.
- Demonstrative examples of matrix reduction applied to simplicial complexes.
- A practical Python implementation of the reduction process, highlighting its utility in real-world applications.

Applications of matrix reduction span various domains, including point cloud analysis for clustering and void detection, time-series analysis for trend identification, and image processing for extracting structural features. Foundational texts, such as Edelsbrunner and Harer’s *Computational Topology: An Introduction*, and computational tools like Ripser and GUDHI³ have significantly contributed to advancing persistent homology as an accessible technique for analyzing large datasets^{3,4}.

2 Matrix Reduction

This section delves into the theoretical and computational aspects of matrix reduction, elaborating on its mathematical foundations, algorithmic steps, and practical applications. The theory is connected to key concepts in algebraic topology, while the computational methodology is demonstrated using illustrative examples.

2.1 Euler-Poincaré Formula

The Euler-Poincaré formula is a cornerstone of algebraic topology, providing a bridge between the combinatorial structure of a simplicial complex and its topological invariants. The Euler characteristic, denoted by χ , is defined as the alternating sum of the number of simplices in each dimension:

$$\chi = \sum_{p \geq 0} (-1)^p n_p,$$

where n_p represents the number of p -dimensional simplices in the complex.

2.1.1 Connection to Homology Groups

Homology groups provide a way to measure topological features, such as connected components, loops, and voids, at different dimensions. The Euler characteristic can also be expressed in terms of the ranks of the cycle groups and boundary groups:

$$\chi = \sum_{p \geq 0} (-1)^p (z_p + b_{p-1}) = \sum_{p \geq 0} (-1)^p (z_p - b_p),$$

where:

- z_p : Rank of the p -dimensional cycle group Z_p , representing features that do not bound.
- b_p : Rank of the p -dimensional boundary group B_p , representing features that bound higher-dimensional simplices.

This representation highlights the relationship between the combinatorial structure of the simplicial complex and its algebraic properties.

2.1.2 Interpretation with Betti Numbers

The Euler characteristic can also be written as the alternating sum of the Betti numbers, β_p , which count the independent p -dimensional features (e.g., connected components, loops, and voids):

$$\chi = \sum_{p \geq 0} (-1)^p \beta_p.$$

This equivalence shows that the Euler characteristic encapsulates the topological essence of the space by summarizing its Betti numbers.

2.1.3 Significance and Invariance

The Euler characteristic is an invariant under homotopy equivalence, meaning that two spaces which can be continuously deformed into each other share the same Euler characteristic. This property is particularly useful in applications, as it ensures that χ does not depend on the specific triangulation of the space.

For example, consider a continuous map $f : X \rightarrow Y$ that defines a homotopy equivalence. If f has a homotopy inverse, it implies that the homology groups of X and Y are isomorphic. Consequently, the Euler characteristic of X is equal to that of Y . This invariance under homotopy equivalence underscores the robustness of χ as a topological invariant.

2.1.4 Implications of the Formula

- The Euler characteristic provides a concise summary of the topology of a space, independent of the triangulation used.
- It serves as a bridge between geometry and topology, connecting the number of simplices to intrinsic algebraic invariants.
- Applications include analyzing the shape of data in fields like image processing, point cloud analysis, and network science.

Theorem 1 (Euler-Poincaré). *The Euler characteristic of a topological space is the alternating sum of its Betti numbers:*

$$\chi = \sum_{p \geq 0} (-1)^p \beta_p.$$

This theorem highlights the dual roles of the Euler characteristic as both a geometric and algebraic descriptor of topological spaces, making it an essential tool in topological data analysis.

2.2 Boundary Matrices and Their Construction

To compute homology, we combine information from two sources: one representing the cycles and the other the boundaries. This information is represented using boundary matrices, which play a key role in homological computations.

2.2.1 Definition and Construction

Let K be a simplicial complex. The p -th boundary matrix represents the $(p-1)$ -simplices as rows and the p -simplices as columns. Assume a fixed ordering of the simplices. For each dimension, the boundary matrix is denoted by $\partial_p = [a_{ij}]$, where:

$$a_{ij} = \begin{cases} 1 & \text{if the } i\text{-th } (p-1)\text{-simplex is a face of the } j\text{-th } p\text{-simplex, oriented consistently,} \\ 0 & \text{otherwise.} \end{cases}$$

In words, the matrix entry indicates the relationship between p -simplices and their $(p-1)$ -dimensional faces.

The boundary matrix ∂_p can be computed as:

$$\partial_p c = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n_p} \\ a_{21} & a_{22} & \cdots & a_{2n_p} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n_{p-1}1} & a_{n_{p-1}2} & \cdots & a_{n_{p-1}n_p} \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \\ \vdots \\ c_{n_p} \end{bmatrix}.$$

Here, a collection of columns represents a p -chain, and the sum of these columns gives its boundary.

2.2.2 Row and Column Operations

The rows of the boundary matrix ∂_p form a basis of the $(p-1)$ -st chain group C_{p-1} , and the columns form a basis of the p -th chain group C_p . Two column operations are used for modifying the matrix while preserving its rank:

- **Exchanging columns k and l .**
- **Adding column k to column l (mod 2).**

To exchange two columns, we multiply the matrix by a matrix V from the right:

$$V = [v_{ij}], \text{ where } v_{ii} = 1, v_{ij} = 1 \text{ if } i = k, j = l.$$

Figure 1 and Figure 2 illustrate the effects of these operations:

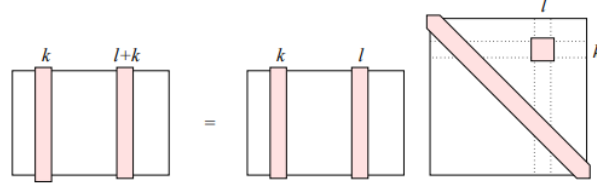


Figure 1: The effect of exchanging columns in the boundary matrix.

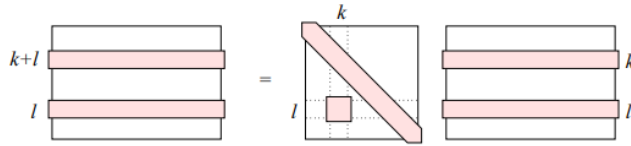


Figure 2: Adding one row in the boundary matrix. The sum of the rows gives the new representation of p -chains.

2.3 Smith Normal Form

2.3.1 Definition

Using row and column operations, the boundary matrix ∂_p can be reduced to Smith normal form. This is achieved as follows:

- The diagonal entries are 1, and everything else is 0.
- The matrix decomposes into components that represent boundary and cycle groups.

2.3.2 Key Insights and Applications

Once the matrix is in Smith normal form:

- The Betti numbers can be extracted from the rank:

$$\beta_p = \text{rank } Z_p - \text{rank } B_p.$$

- The structure of homology groups is revealed.

Figure 3 shows the structure of the matrix after transformation into Smith normal form.

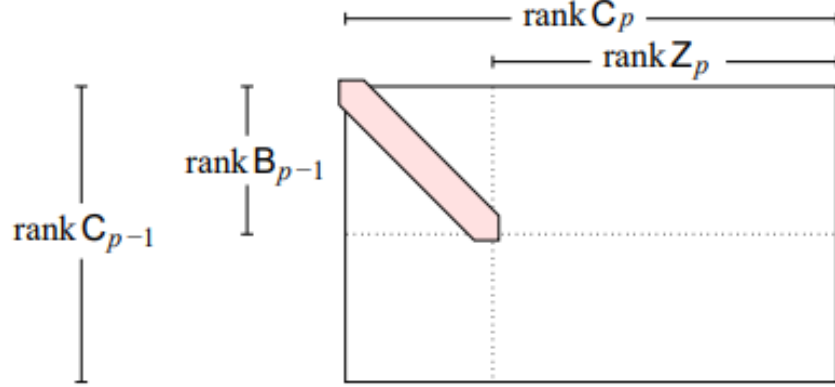


Figure 3: Smith Normal Form: Decomposition of boundary matrix into ranks of cycle and boundary groups.

2.3.3 Worked Example

Let us consider a simplicial complex with:

- Vertices: $\{A, B, C, D\}$.
- Edges: $\{AB, BC, CD, DA\}$.
- A square region bounded by the edges.

The boundary matrix ∂_1 for edges to vertices is:

$$\partial_1 = \begin{bmatrix} 1 & 0 & 0 & 1 \\ -1 & 1 & 0 & 0 \\ 0 & -1 & 1 & 0 \\ 0 & 0 & -1 & 1 \end{bmatrix}.$$

After performing column operations:

$$\text{Smith Normal Form: } \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}.$$

From this, we observe:

- $\beta_0 = 1$: One connected component.
- $\beta_1 = 1$: One loop (the square boundary).

2.4 Applications of Matrix Reduction

Matrix reduction has broad applications in various domains:

- **Point Cloud Analysis:** Identifying clusters and voids in high-dimensional data.
- **Time Series Analysis:** Detecting periodicity and trends.
- **Image Processing:** Extracting topological features from pixel-based data.

This combination of mathematical rigor and computational practicality highlights the importance of matrix reduction in topological data analysis.

2.5 Reduction of Boundary Matrices

The reduction of the boundary matrix ∂_p is analogous to Gaussian elimination, aiming to simplify the matrix while preserving its homological information. This process involves systematically performing row and column operations to reduce the matrix to a simplified form.

The reduction algorithm starts with a matrix $N_p[i, j] = a_j^i$ for all i and j . The function 'Reduce(x)' works on the position of the considered diagonal element and performs the following steps:

1. If there exists a nonzero element in column x , exchange rows and columns to bring it to the diagonal.
2. Use row additions to eliminate all other entries in the current column.
3. Use column additions to eliminate all other entries in the current row.
4. Recurse to the next submatrix by incrementing x .

The pseudocode for this reduction is provided below:

```
void Reduce(x):
    if there exists k >= x, l >= x with Np[k, l] = 1 then:
        exchange rows x and k, exchange columns x and l;
        for i = x + 1 to np-1:
```



```

    if  $N_p[i, x] = 1$  then add row  $x$  to row  $i$ ;
endfor;
for  $j = x + 1$  to  $n_p$ :
    if  $N_p[x, j] = 1$  then add column  $x$  to column  $j$ ;
endfor;
Reduce( $x + 1$ );
endif;

```

The complexity of this algorithm is at most $2n_{p-1}n_p \min(n_{p-1}, n_p)$, where n_{p-1} and n_p are the dimensions of the boundary matrix. The resulting diagonal form provides the ranks of the cycle and boundary groups directly.

3 Results

3.1 Illustrative Example: Homology of a Tetrahedron

To demonstrate the reduction algorithm, we compute the reduced homology of a tetrahedron. The process involves reducing the boundary matrices for the 0-skeleton, 1-skeleton, 2-skeleton, and 3-skeleton sequentially.

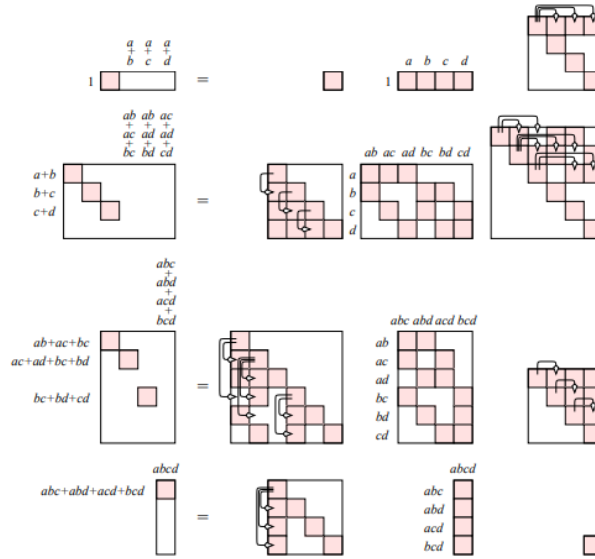


Figure 4: Matrix reduction for the tetrahedron. Rows and columns are shaded to indicate operations, and the diagonal form reveals the Betti numbers for each skeleton.

Figure 4 illustrates the reduction steps. The matrices for each skeleton are re-

duced, and the Betti numbers β_p are computed as follows:

- $\beta_0 = 0$: The 0-skeleton consists of four vertices, all connected, resulting in no independent components.
- $\beta_1 = 0$: The 1-skeleton consists of edges with no independent loops, as all loops are filled by the triangular faces of the 2-skeleton.
- $\beta_2 = 0$: The 2-skeleton consists of triangular faces enclosing no independent hollow surfaces, as the volume is filled by the 3-skeleton.
- $\beta_3 = 0$: The 3-skeleton adds a tetrahedron with no independent voids.

These results align with the theoretical Betti numbers expected for a 3-ball, confirming its simple, connected topology.

3.2 Python Code Implementation

Below is the Python implementation of the matrix reduction algorithm, applied to the tetrahedron example. The code outputs the reduced boundary matrices and the Betti numbers for each skeleton of the tetrahedron. The results match the theoretical values derived earlier. The complete Python code can be found on the following [GitHub repository](#).

4 Conclusion

Matrix reduction bridges algebraic theory with practical applications in TDA. By simplifying boundary matrices, it enables efficient computation of topological invariants, offering insights across diverse domains. Future work could explore applications in dynamic datasets and higher-dimensional simplicial complexes.

References

- [1] Kannan, R. & Bachem, A. (1979). Polynomial algorithms for computing the Smith and Hermite normal forms of an integer matrix. *SIAM J. Comput.*, 8, 499–507.
- [2] Lakatos, I. (1984). *Proofs and Refutations: The Logic of Mathematical Discovery*. Cambridge Univ. Press, Cambridge, England.

- [3] Munkres, J. R. (1984). *Elements of Algebraic Topology*. Addison-Wesley, Redwood City, California.
- [4] Smith, H. J. (1861). On systems of indeterminate equations and congruences. *Philos. Trans.*, 151, 293–326.
- [5] Storjohann, A. (1997). New optimal algorithm for computing Smith normal forms of integer matrices. *Proc. Internat. Sympos. Symbol. Algebraic Comput.*, 267–274.