

Pseudocode:

Array[j] refers to the light side whereas array[j+1] refers to dark side

Alternating algorithm:

Make a temp disk state	1 tu
Make a local variable that will function as swap count: swap	1 tu
For i from 0 thru n in array	n/2+1 tu
For j from 0 thru n-1 in array	n tu
If array[j] is greater than array[j+1] then	2 tu
Increment swap	2 tu
Swap array[j] and array[j+1]	3 tu
Return array, swap	

Lawnmower algorithm:

Make a temp disk state	1 tu
Make a local variable that will function as swap count: swap	1 tu
For i from 0 to n in array	- n/2 + 1 tu
For s from total count - 1 downto 1 in array	1-n tu
If array[s] is less than array[s-1]	- 2 tu
Increment swap	- 2 tu
Swap array[s] and array [s-1]	- 3 tu
For j from 0 to n-1 in array	- n tu
If array[j] is greater than array[j+1]	- 2 tu
Increment swap	-2 tu
Swap array[j] and array [j+1]	- 3 tu

Analysis:

Alternate algorithm

$$\begin{aligned}
 T(n) &= 1 + 1 + (n/2+1)(n)(2+\max(5, 0)) \\
 &= 2 + (n/2+1)(n)(7) \\
 &= 2 + (n/2+1)(7n) \\
 &= 2 + (7n^2)/2 + 7n
 \end{aligned}$$

$$\begin{aligned}
 &= (7n^2)/2 + 7n + 2 \\
 &O((7n^2)/2 + 7n + 2) \\
 &O((7n^2)/2) \\
 &O(n^2)
 \end{aligned}$$

Lawnmower algorithm

$$\begin{aligned}
 T(n) &= 1 + 1 + ((n/2) + 1) ([(1-n)(2 + \max(6, 0))] + [n(2 + \max(5, 0))]) \\
 &= 2 + ((n/2) + 1) ([8 - 8n] + [7n]) \\
 &= 2 + ((n/2) + 1) (8 - 1n) \\
 &= 2 + ((8n/2) + 1)(8 - 1n) \\
 &= 2 + (8n/2) - (n^2/2) + 8 - n \\
 &= 10 + (8n/2) - (n^2/2) + 8 - n \\
 &= 10 + 3n - (n^2/2) \\
 &= (n^2/2) - 3n - 10 \\
 &O((n^2/2) - 3n - 10) \\
 &O((n^2/2)) \\
 &O(n^2)
 \end{aligned}$$