

Substitution Techniques



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Outline

- ❖ **Substitution Techniques**
- ❖ **Types of Substitution Cipher**
- ❖ **Mono-alphabetic Cipher**
 - ✓ **Caesar Cipher**
 - ✓ **Analysis of Caesar Cipher**
 - ✓ **Monoalphabetic Substitution Ciphers**

Operations on Encryption Algorithm

- ❖ All **Encryption Algorithms** are based on **two general principles**:
 - ✓ **Substitution**
 - ✓ **Transposition**
- ❖ **Substitution** : Each element in the **Plaintext (Bit, Letter, Group of Bits Or Letters)** is mapped into **Another Element**,
- ❖ **Transposition** : Each elements in the **plaintext** are **rearranged**.

Substitution Techniques

Introduction

- ❖ Before Computers, **Cryptography** consisted of **Character-based Algorithms**.
- ❖ The **better algorithms** always use both **Substitution** and **Transposition**

Substitution Cipher

- ❖ A **Substitution Cipher** is one in which each **character in the plaintext** is **SUBSTITUTED** for **another character in the ciphertext**.
- ❖ The **receiver inverts** the **substitution** on the **ciphertext** to **recover the plaintext**.

Types of Substitution Cipher

Types of Substitution Cipher

❖ In **Classical Cryptography**, there are **Four Types of Substitution Ciphers**:

1. **A Simple Substitution Cipher or Mono-alphabetic Cipher.**
2. **A Homophonic Substitution Cipher.**
3. **A Polygram Substitution Cipher.**
4. **A Polyalphabetic Substitution Cipher.**

Mono-alphabetic Cipher

Mono-alphabetic Cipher

- ❖ A **Mono-alphabetic Cipher**,
 - ✓ Each **character of the plaintext** is **replaced uniquely** with a **corresponding character** of **cipher text**.
- ❖ The **famous Caesar Cipher** falls under **simple substitution cipher**

Types of Mono-alphabetic Cipher

Types of Mono-alphabetic Cipher

- ❖ A **Mono-alphabetic Cipher** falls under **three categories**
 - ✓ **Caesar Cipher**
 - ✓ **Substitution Cipher**
 - ✓ **The Affine Cipher**

Caesar Cipher

Caesar Cipher

- ❖ **Caesar Cipher** is also called as **Shift Cipher** which is Primitive(basic) Cipher
- ❖ **Caesar substitution Cipher** was introduced by **Julius Caesar**.

Working Model of Caesar Cipher

- ❖ The **Caesar cipher** involves **replacing each letter of the alphabet** with the **Unique letter** standing **THREE** places further down the alphabet.

Working Model of Caesar Cipher

❖ Let us Consider the set

$\mathbf{Z_{26} = \{0,1,2,3,\dots\}}$ usually represented with Alphabets **A-Z**

❖ Let us also consider the tuples

$P=C=K= Z_{26}$ Here K is the Piece of Information called **KEY**

which also contains 26 Character($0 < K < 25$)

Simplified the Caesar Cipher

- ❖ We can define the **transformation by listing all possibilities**, as follows:

Plain Text																									
a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Cipher Text																									
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C

- ❖ Let us assign a **numerical equivalent** to each letter:

Plain Text																									
a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

Working Model of Caesar Cipher

- ❖ We define **Encryption Function** with help of **KEY**

$$E_K(x) = x + K \bmod 26$$

- ❖ We define **Decryption Function** with help of **KEY**

$$D_K(y) = y - K \bmod 26$$

- ❖ Finally We define **Caesar Cipher** with help of **KEY**

$$P = D_K(E_K(x))$$

Example of Working Model

- ❖ **P** = {**A=0,B=1,C=2,D=3.....**}
- ❖ **K** = {**A=0,B=1,C=2,D=3.....**}
- ❖ We define **Encryption Function** with help of **KEY**

$$E_K(x) = x + K \bmod 26$$

$$E_K(x) = 0(A) + 3(D) \bmod 26 = 3(D) \text{ is a Cipher text}$$

Example of Working Model

❖ We define Encryption Function with help of KEY

$$D_K(y) = y - K \bmod 26$$

$$E_K(x) = 3(D) - 3(D) \bmod 26 = 0(A) \text{ is a Plain text}$$

Analysis of Caesar Cipher

Analysis of Caesar Cipher

❖ Plain Text : meet me after the toga party

❖ Cipher Text : PHHW PH DIWHU WKH WRJD SDUWB

Analysis of Caesar Cipher

If it is **Cryptanalyst** knows that a given **Cipher text** is a **Caesar Cipher**,

- ✓ **Brute-force Cryptanalysis** is easily performed
- ✓ Simply try all the **25 possible keys**

Characteristics of Caesar Cipher

❖ **Three important characteristics** of this **Caesar Cipher** enabled to use a **brute force cryptanalysis**:

1. The **encryption and decryption algorithms** are **known**.
2. There are **only 25 keys** to try.
3. The language of the **plaintext is known** and **easily recognizable**

What makes brute-force cryptanalyst impractical

- ❖ **Caesar Cipher Algorithm** should employs a **large number of keys**.
- ❖ For example, the **triple DES algorithm**, makes use of a **168-bit key**,
- ❖ The key space of **2^{168} or $3.7 * 10^{50}$ possible keys**.

Brute-Force Cryptanalysis of Caesar Cipher

KEY		PHHW	PH	DIWHU	WKH	WRJD	SDUWB
1		oggv	og	chvgt	vjg	vqic	rectva
2		nffu	nf	bgufs	uif	uphb	qbsuz
3		meet	me	after	the	toga	party
4		ldds	ld	zesdq	sgd	snfz	ozqsx
5		kccr	kc	ydrp	rfe	rmey	nyprw
6		jbbq	jb	xcqbo	qeb	qldx	mxoqv
7		iaap	ia	wbpan	pda	pkcw	lwnpu
8		hzzo	hz	vaozm	ocz	ojbv	kvmot
9		gyyn	gy	uznyl	nby	niau	julns
10		fxxm	fx	tymxk	max	mhzt	itkmr
11		ewwl	ew	sxlwj	lzw	lgys	hsjlg
12		dvvk	dv	rwkvi	kyv	kfxr	grikp
13		cuuj	cu	qvjuh	jxu	jewq	fqhjo
14		btti	bt	puirg	iwt	idvp	epgin
15		assh	as	othsf	hvs	hcuo	dofhm
16		zrrg	zr	nsgre	gur	gbtn	cnegl
17		yqqf	yq	mrfqd	ftq	fasm	bmdfk
18		xppe	xp	lqepc	esp	ezrl	alcej
19		wood	wo	kpdob	dro	dyqk	zkbdl
20		vnnc	vn	jocna	cqn	cxpj	yjach
21		ummb	um	inbmz	bpm	bwoi	xizbg
22		tlla	tl	hmaly	aol	avnh	whyaf
23		skkz	sk	glzlx	znk	zumg	vgxze
24		rjjy	rj	fkyjw	ymj	ytlf	ufwyd
25		qiix	qi	ejxiv	xli	xske	tevxc

Observation

- ❖ The third characteristic is also **significant**. If the language of the **plaintext is unknown**, then plaintext **output may not be recognizable**

```
~+Wµ"- Ω-0)≤4{∞‡, ë~Ω%ràu.-í Ø-Z-  
Ú#2Ò#Åæð æ«q7,Ωn·@3NÔÚ Ez'Y-f∞í[±Ŧ_ èΩ,<NO¬±«~xã Åäfèü3Å  
x)ö$ksÂ  
_yÍ ^ΔÉ] .¤ J/'iTê&1 'c<uΩ-  
ÄD(G WÄC~y_iöÄW PÔ1«îÜ†ç],¤;~î^uÑπ~≈~L~9OgfiO~&E≤ ¬≤ ØÔ$~:  
~E!SGqèvo^ ú\,S>h<-*6ø‡%x'~|fiÓ#≈~my%~≥fiP<,fi Áj ÅÔ¿~Zù-  
Ω"Ö-6EY{§ „ΩÊó .Y π+Áî'ú02çSY'O-  
2Äfißi /@^"ΠK≈PCEπ,úé^'3Σ~ø~ÔZî"Y¬YΩæY> Ω+eô/'<Kf¿*+~"≤û~  
B ZøK~Qßyüf,!ÒfiîzsS/]>ÈQ ü
```

Sample of Compressed Text

Limitations of Caesar Ciphers

- ❖ **Caesar cipher** is far from **secure**, With **only 25 possible keys** are used, as **Cryptanalyst** can easily **deduce** the key by **Bruce-force attack**.

Monoalphabetic Substitution Ciphers

Substitution Ciphers

- ❖ One way to **increase the key space** and **improve the security of the cipher** is to allow **arbitrary substitution**.
- ❖ In this case, the “**cipher**” can be any **PERMUTATION** of the **26 alphabetic characters**.

PERMUTATION

❖ A **Permutation** of a **finite set of elements** is an **ordered sequence** of all the **elements of S**, with each element **appearing exactly once**.

❖ **For Example**

✓ if $S = \{a, b, c\}$ there are six permutations of abc, acb, bac, bca, cab, cba

Monoalphabetic Substitution Cipher

- ❖ In a **monoalphabetic cipher**, our substitution characters are a **random permutation** of the **26 letters of the alphabet**

Plain Text																									
a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Substitution Text																									
J	I	C	A	X	X	E	Y	V	D	K	W	B	Q	T	Z	R	H	F	M	P	N	U	L	G	O

- ❖ The **key** now is the **sequence of substitution letters**. In other words, the **key** in this case is the **actual random permutation** of the alphabet used.

Monoalphabetic Substitution Cipher

- ❖ Arbitrary **substitution of letters**
- ❖ Number of keys **$26 \times 25 \times \dots \times 1 = 26!$** (Over 4×10^{26})
- ❖ Note that there are **$26!$ permutations** of the alphabet. That is a **number larger than 4×10^{26} .**

Advantages of Substitution Cipher

- ❖ **Substitution Cipher** will **eliminate brute-force techniques** for cryptanalysis.
- ✓ Requires $>2^{88}$ **Possible Keys** for **brute-force attacks** which take **zillions of years** to try out even half the keys

Limitation of Substitution Cipher

- ❖ Any **Substitution Cipher**, regardless of the **size of the key space**, can be **broken easily** with a **Statistical Attack**.

Frequency Analysis of Substitution Ciphers

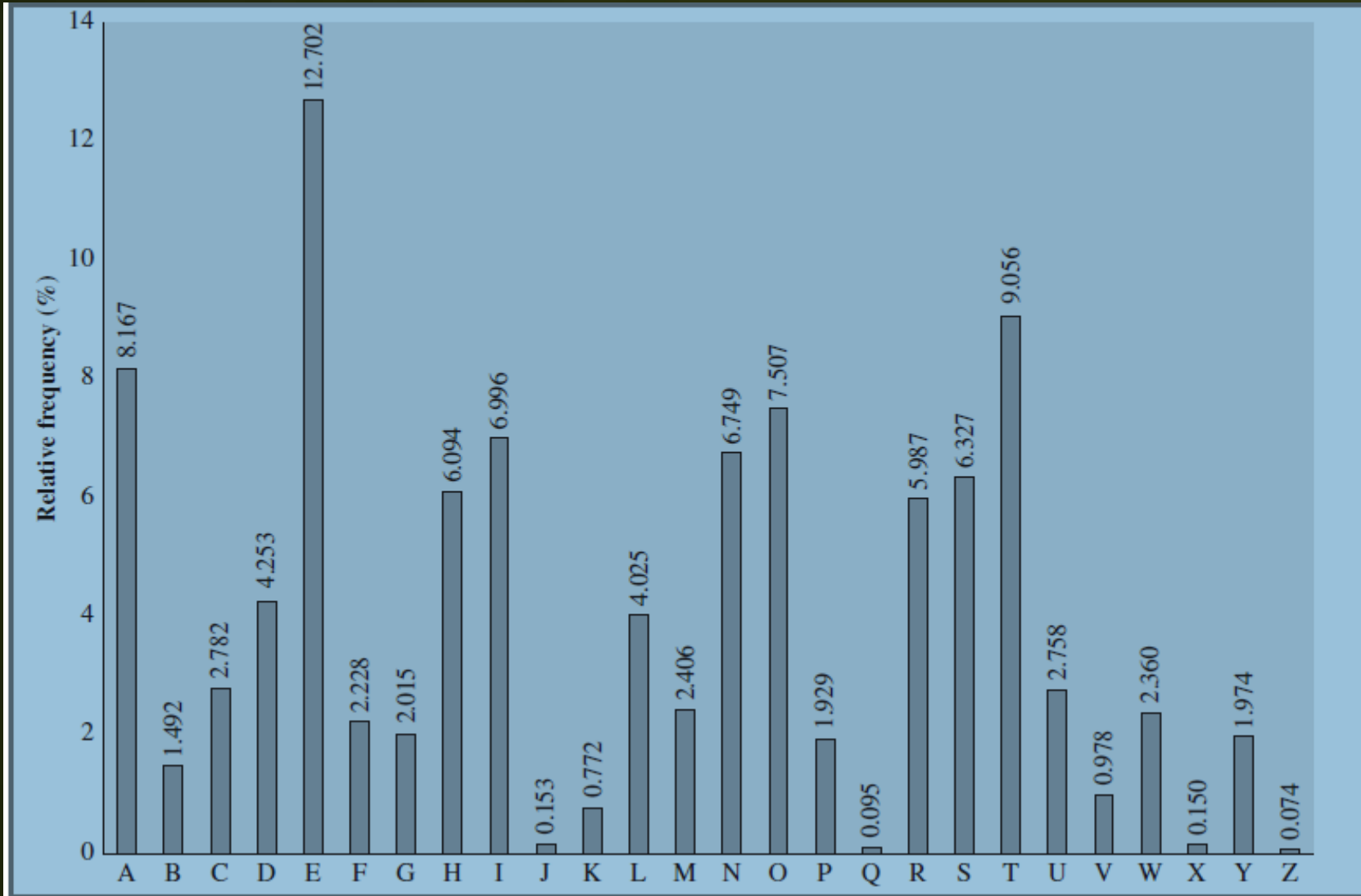
Frequency Analysis

- ❖ **Substitution Cipher** suffers **Statistics attack**,
 - ✓ If the **cryptanalyst** knows the **nature of the plaintext**, then the **analyst can exploit** the **regularity of the language**, called **FREQUENCY ANALYSIS**.
- ❖ **Frequency Analysis** studies the **frequency of letters or groups of letters** in a **ciphertext**.

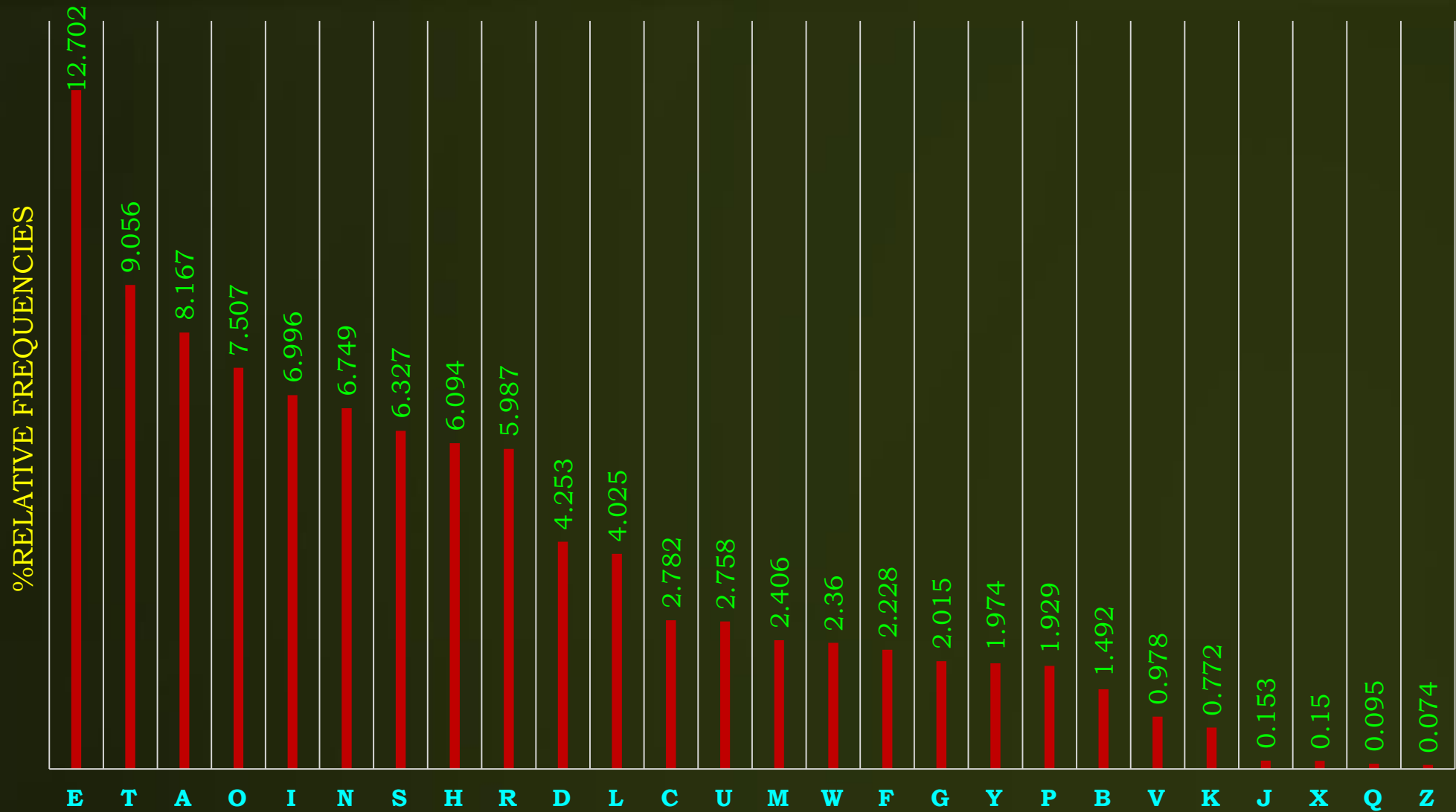
Frequency Analysis

- ❖ When the **plaintext is plain English**, a simple form of **statistical attack** consists **measuring the frequency distribution** for
 - ✓ **Single Characters,**
 - ✓ **Pairs of Characters,**
 - ✓ **Triples of Characters,** etc.,
- ❖ Comparing **all character** with **similar statistics** for English

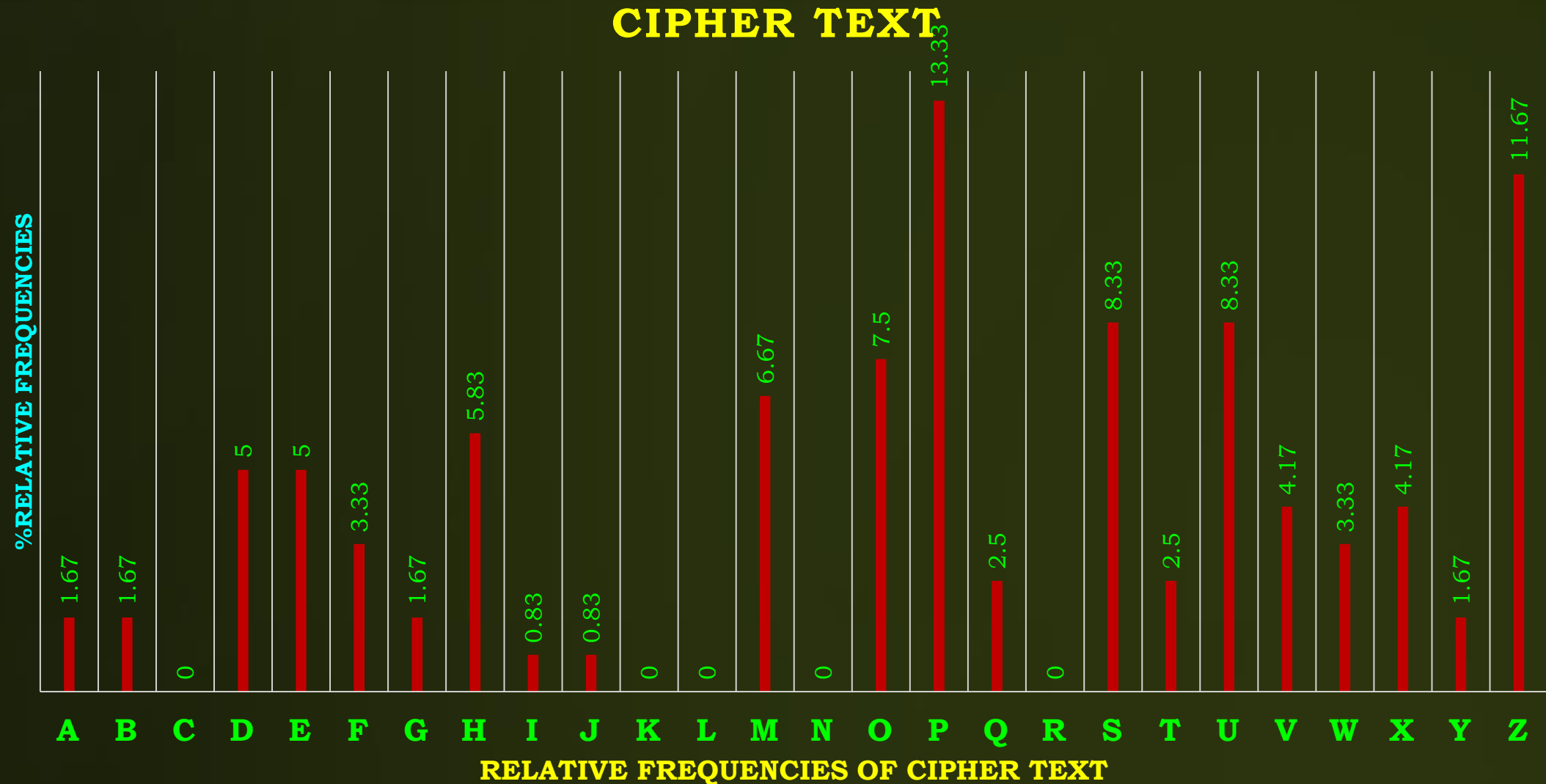
Relative Frequency of the letters of English text



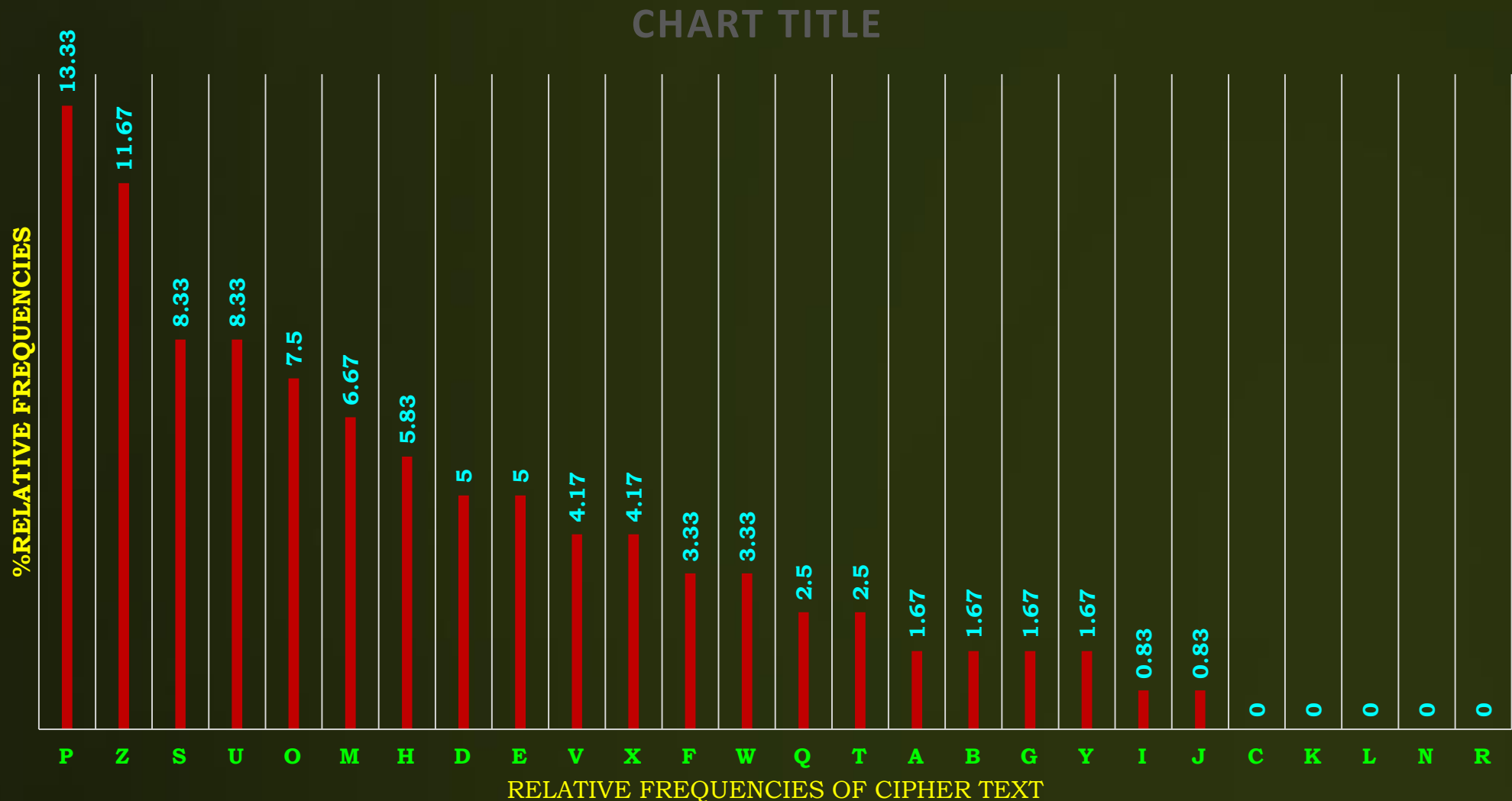
Relative Frequency of the letters of English text



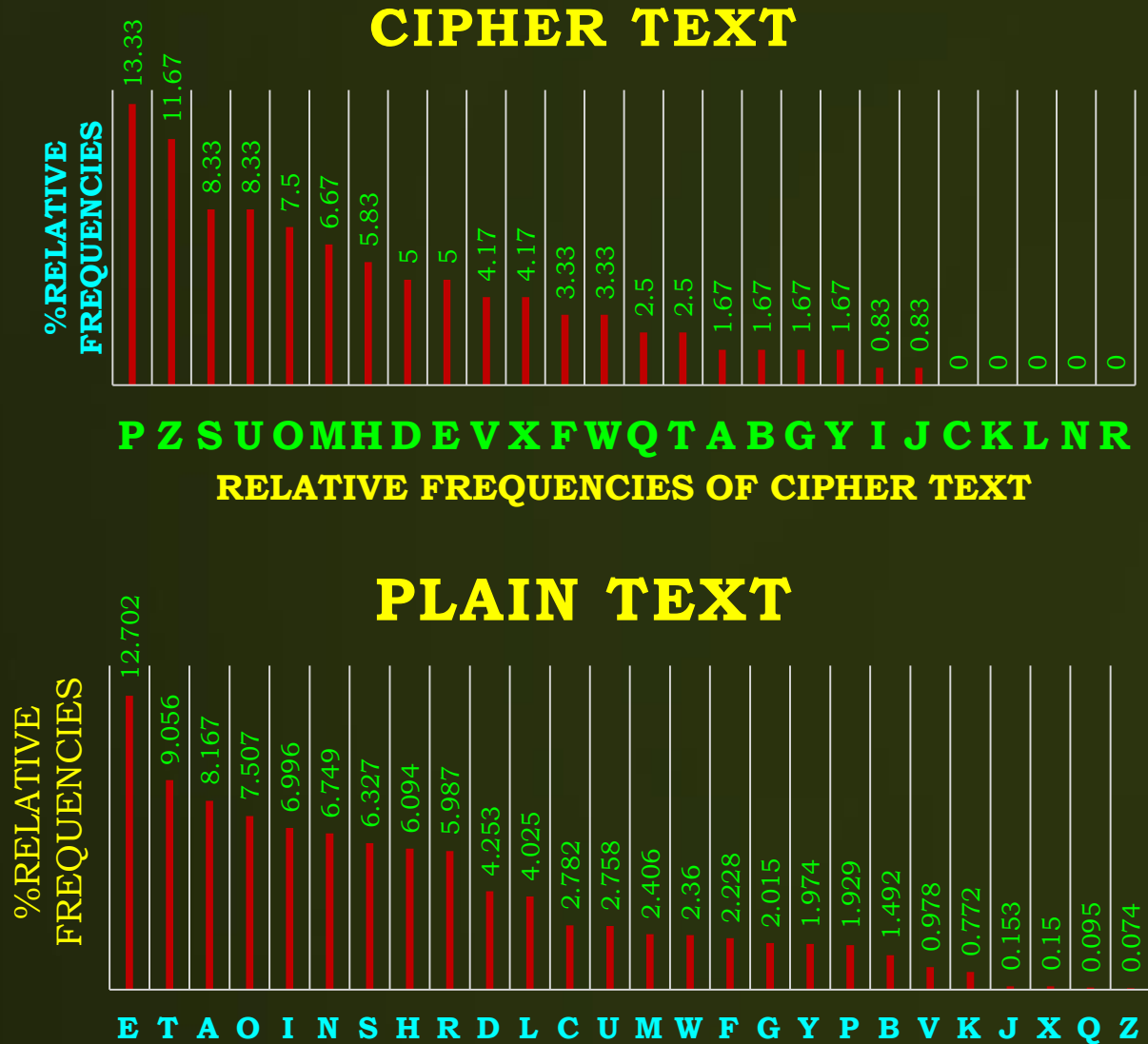
The Relative Frequencies of the Letters in the Ciphertext (In Percentages)



The Relative Frequencies of the Letters in the Ciphertext (In Percentages) in Decreasing Order



Comparing The Relative Frequencies of the Plain Text & Ciphertext



Digram Frequencies

<i>digram</i>	<i>frequency</i>	<i>digram</i>	<i>frequency</i>	<i>digram</i>	<i>frequency</i>	<i>digram</i>	<i>frequency</i>
th	3.15	to	1.11	sa	0.75	ma	0.56
he	2.51	nt	1.10	hi	0.72	ta	0.56
an	1.72	ed	1.07	le	0.72	ce	0.55
in	1.69	is	1.06	so	0.71	ic	0.55
er	1.54	ar	1.01	as	0.67	ll	0.55
re	1.48	ou	0.96	no	0.65	na	0.54
es	1.45	te	0.94	ne	0.64	ro	0.54
on	1.45	of	0.94	ec	0.64	ot	0.53
ea	1.31	it	0.88	io	0.63	tt	0.53
ti	1.28	ha	0.84	rt	0.63	ve	0.53
at	1.24	se	0.84	co	0.59	ns	0.51
st	1.21	et	0.80	be	0.58	ur	0.49
en	1.20	al	0.77	di	0.57	me	0.48
nd	1.18	ri	0.77	li	0.57	wh	0.48
or	1.13	ng	0.75	ra	0.57	ly	0.47

Trigrams Frequencies

- ❖ A powerful tool is to look at the **frequency of THREE-letter combinations**, known as **Trigrams**.
- ❖ The most **frequently occurring trigrams** ordered by decreasing frequency are:
***the, and, ent, ion, tio, for, nde** are some of the Trigram*

For Example

The Ciphertext to be solved is

**UZQSOVUOHXMOPVGPOZPEVSGZWSZOPFPESXUDBMETSXAIZVU
EPHZHMDZSHZOWSFPAPPDTSVPQUZWYMXUZUHSXEPYEPOPDZS
ZUFPOMBZWPFUPZHMDJUDTMOHMQ**

For Example

- ❖ The most common such **digram is th**. In our ciphertext, the most common **digram is ZW** which appears three times. So we make the correspondence of **Z with t and W with h**. we can equate **P with e**.

UZQSOVUOHXMOPVGPOZPEVSG**ZWS**SZOPFPESXUDBMETSXAIZVUEPHZHMDZSHZO
WSFPAPPDTSVPQU**ZW**YMXUZUHSXEPYEPOPDZSZUFPOMB**ZW**PFUPZHMDJUDTMO
HMQ

UZQSOVUOHXMOPVGPOZPEVSG**th**SZOPFPESXUDBETSXAIZVUEPHZ
HMDZSHZOWSFPAPPDTSVPQU**th**YMXUZUHSXEPYEPOPDZSZUFPOM
BthPFUPZHMDJUDTMOHMQ

For Example

- ❖ Now notice that the **sequence ZWP** appears in the **ciphertext**, and we can translate that **sequence as “the.”**

UZQSOVUOHXMOPVGPOZPEVSG**ZW**SZOPFPESXUDBMETSXAIZVUEPHZHMDZSHZ
OWSFPAPPDTSVPQU**ZW**YMXUZUHSXEPYEPOPDZSZUFPOMB**ZWP**FUPZHMDJUDTMOHMQ

UZQSOVUOHXMOPVGPOZPEVSG**th**SZOPFPESXUDBETSXAIZVUEPHZ
HMDZSHZOWSFPAPPDTSVPQU**th**YMXUZUHSXEPYEPOPDZSZUFPOM
B**the**FUPZHMDJUDTMOHMQ

For Example

- ❖ Next, notice the **sequence ZWSZ** in the first line. We do not know **that these four letters form a complete word**, but if they do, it is of the form **th_t**. If so, **S equates with 'a'**.

UZQSOVUOHXMOPVGPOZPEVSG**ZWSZ**OPFPESXUDBMETSXAIZVUEPHZHMDZSHZ
OWSFPAPPDTSVPQUZWYMXUZUHSXEPYEPOPDZSZUFPOMBZWPFUPZHMDJUDTMO
OHMQ

UZQSOVUOHXMOPVGPOZPEVSG**that**OPFPESXUDBETSXAIZVUEPHZ
HMDZSHZOWSFPAPPDTSVPQU**th**YMXUZUHSXEPYEPOPDZSZUFPOM
B**the**FUPZHMDJUDTMOHMQ

For Example

❖ we have

```
UZQSOVUOHXMOPVGPOZPEVSGZWSZOPFPESXUDBMETSXAIZ
t a           e e te a that e e a           a
VUEPHZHMDZSHZOWSFPAPPDTSVPQUZWYMXUZUHSX
e t   ta t ha e ee a e th   t a
EPYEPOPDZSZUFPOMBZWPFPUPZHMDJUDTMOHMQ
e e e tat e   the t
```

For Example

- ❖ Finally, The **complete plaintext, with spaces** added between words as follows:

```
it was disclosed yesterday that several informal but  
direct contacts have been made with political  
representatives of the viet cong in moscow
```

Limitations of Monoalphabetic Substitution ciphers

- ❖ **Monoalphabetic Substitution ciphers** are easy to **break** because they **reflect the frequency data** of the **original alphabet**

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Thank U