

## Recursive Minimax

Minimax is used to dynamically evaluate a position in the game. The procedure is described here as a depth first search. The nodes in the search tree are positions. A node is a “**max**” node if it describes a position where the “maximizer” is to make a move. It is a “**min**” node if it describes a position where the “minimizer” is to make a move. The children of a node are all the possible positions that can be reached after one move. With each position  $x$  we associate the value  $V_x$ . The function **static**( $x$ ) gives a static evaluation of the position. The following two recursive procedures evaluate “dynamically” the value  $V_x$  of the node  $x$ :

**if  $x$  is a max node:**  $V_x = \text{MaxMin}(x)$ , where:

$$\text{MaxMin}(x) = \begin{cases} \text{static}(x) & \text{if } x \text{ is a leaf} \\ \max_{\text{children } y \text{ of } x} \text{MinMax}(y) & \text{otherwise} \end{cases}$$

**if  $x$  is a min node:**  $V_x = \text{MinMax}(x)$ , where:

$$\text{MinMax}(x) = \begin{cases} \text{static}(x) & \text{if } x \text{ is a leaf} \\ \min_{\text{children } y \text{ of } x} \text{MaxMin}(y) & \text{otherwise} \end{cases}$$

These routines can also be written as:

<pre> MaxMin(x): if x is a leaf return static(x). else {     set v = -∞     for each child y of x:         v = max(v, MinMax(y))     return v }</pre>	<pre> MinMax(x): if x is a leaf return static(x). else {     set v = +∞     for each child y of x:         v = min(v, MaxMin(y))     return v }</pre>
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## Recursive Alpha-Beta pruning

For each node  $x$  retain the range of values of its parents so that:  $\alpha \leq V_{\text{parents of } x} \leq \beta$ . If the root of the tree is a MAX node it is evaluated by:  $\text{MaxMin}(\text{root}, -\infty, +\infty)$ .

<p><b>MaxMin(<math>x, \alpha, \beta</math>):</b></p> <ol style="list-style-type: none"> <li>1. if <math>x</math> is a leaf return <math>\text{static}(x)</math>.</li> <li>2. else do steps 2.1, 2.2, 2.3 <ol style="list-style-type: none"> <li>2.1. set <math>v = -\infty</math></li> <li>2.2 for each child <math>y</math> of <math>x</math>: <ol style="list-style-type: none"> <li>2.2.1 <math>v = \max(v, \text{MinMax}(y, \alpha, \beta))</math></li> <li>2.2.2 if <math>(v \geq \beta)</math> return <math>v</math></li> <li>2.2.3 else <math>\alpha = \max(v, \alpha)</math></li> </ol> </li> <li>2.3. return <math>v</math></li> </ol> </li> </ol>	<p><b>MinMax(<math>x, \alpha, \beta</math>):</b></p> <ol style="list-style-type: none"> <li>3. if <math>x</math> is a leaf return <math>\text{static}(x)</math>.</li> <li>4. else do steps 2.1, 2.2, 2.3 <ol style="list-style-type: none"> <li>4.1 set <math>v = +\infty</math></li> <li>4.2 for each child <math>y</math> of <math>x</math>: <ol style="list-style-type: none"> <li>4.2.1 <math>v = \min(v, \text{MaxMin}(y, \alpha, \beta))</math></li> <li>4.2.2 if <math>(v \leq \alpha)</math> return <math>v</math></li> <li>4.2.3 else <math>\beta = \min(v, \beta)</math></li> </ol> </li> <li>4.3 return <math>v</math></li> </ol> </li> </ol>
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**Notes:**

- . Step 2.2.2 is a  $\beta$  cut. The value returned at 2.2.2 will not affect the outcome since it will never be the minimum at 4.2.1.
- . Same observation holds for the  $\alpha$  cut at Step 4.2.2.