1. Convert the following grammar to a PDA that accepts the same language.

$$S \rightarrow 0S1 \mid A$$
  
 $A \rightarrow 1A0 \mid S \mid \epsilon$ 

## **Solution:**

The CFG can be first simplified by eliminating unit productions:

$$S \rightarrow 0S1 \mid 1S0 \mid \ \epsilon$$

Now we will convert this CFG to GNF:

$$S \rightarrow 0SX \mid 1SY \mid \epsilon$$
  
 $X \rightarrow 1$   
 $Y \rightarrow 0$ 

The PDA can be:

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R1: \delta(q, \epsilon, S) = \{(q, 0SX) \mid (q, 1SY) \mid (q, \epsilon)\}

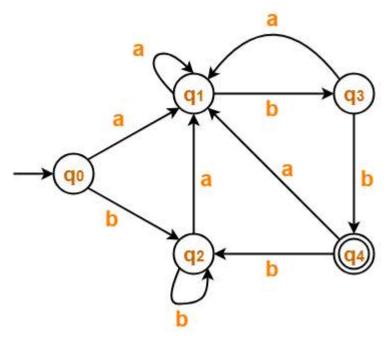
R2: \delta(q, \epsilon, X) = \{(q, 1)\}

R3: \delta(q, \epsilon, Y) = \{(q, 0)\}

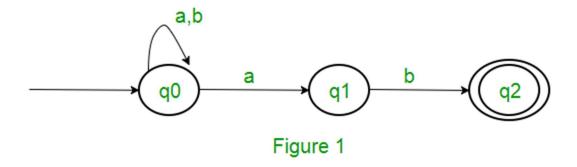
R4: \delta(q, 0, 0) = \{(q, \epsilon)\}

R5: \delta(q, 1, 1) = \{(q, \epsilon)\}
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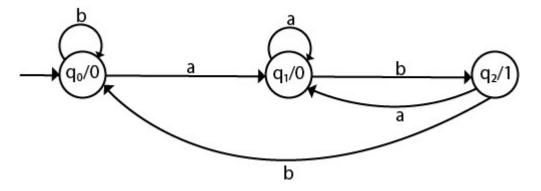
- 2. Design the Following Finite Automata.
- A) Design a FA with  $\{0, 1\}$  accepts the only input 101.
- b) Design an NFA with  $\{0, 1\}$  accepts all string ending with 01.
- c) Design an NFA with  $\{0, 1\}$  in which double 11 is followed by double 00.
- d) Design an NFA in which all the string contain a substring 1110.
- 3. Minimize the given DFA. Using Table Filling Method OR Equivalence method.



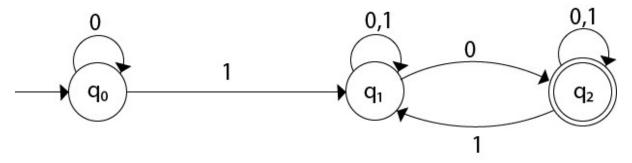
 ${\tt Q4}\,.$  Write Procedure for converting NFA to DFA. Convert the Following NFA to DFA shown in Figure-1



- Q5. Describe Type-3 ,Type-2 ,Type-1 and Type-0. Grammar with Example.
- Q6. Design DFA with {0, 1} accepts
  - a) even number of 0s and even number of 1s.
  - b) even number of 0s and odd number of 1s.
  - c) odd number of 0s and even number of 1s.
  - d) odd number of 0s and odd number of 1s.
- Q7. Convert the given Moore machine into its equivalent Mealy Machine.

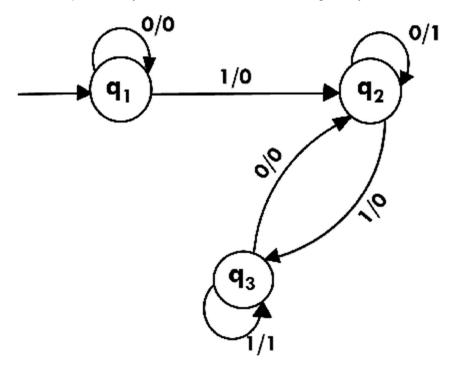


Q8. Define DFA and covert the given NFA into DFA.



- Q9. Prove Ardens's Theorem. Find the regular expression corresponding to figure.
- Q.10 Prove that every finite language is regular with example. State Kleens Theorem.
- Q11. Let M1 and M2 be the FA that accept languages L1 and L2 then find L1^L2.
- Q12. Construct FA for a regular expressions:
  - 1. (a+b)\*ab(a\*+b)\*
  - 2. aab(a+b)\* baaa\*
- Q13. Write Regular Expressions for the following languages of all strings in  $\{x,y\}^*$ 
  - (i) Strings that contains odd number of x.
  - (ii) Strings that begin or end with xx or yy.
- Q14. Use the pumping lemma to show that following language is not regular:  $L = \{xyx, x \text{ \&epsilon}; \{0,1\}^* \text{ and } y \text{ is } R\}$

Q.15 Explain Mealy machine. Convert the following mealy machine in its equivalent Moore machine



- Q 16. State Moore Machine in terms of tuples. Explain the Procedure to Convert Mealy Machine into Moore Machine with the help of an example.
- Q.17. Explain the procedure of Minimizing Automata with the help of example of your own choice.
- Q18. Minimize the following FA.

	a	b
->A	В	A
B C D E *F *G	A	C
С	D	В
D	D	A
E	D	F
*F	G	E G
*G	F G	G
*H	G	D

- Q.19. Define NFA. Explain Conversion of NFA -ε into NFA without ε with example.
- Q.20. Discuss the following terms with Example.
  - a) Alphabets
  - b) String
  - c) Symbols
  - d) Language
  - e) Grammar
- Q 21. Construct a PDA for language L = {0^n1^n2^m3^m; n=1, m=1}

- Q22. Explain Non deterministic PDA with example.
- Q. 23 Explain two stack Pushdown Automata with example.
- Q24. Differentiate between Deterministic PDA and Non Deterministic PDA. (CO4)
- Q25. Construct a PDA that accepts the language L over {0, 1} by empty stack which accepts all the string of 0's and 1's in which a number of 0's are twice of number of 1's.
- Q26. Design a PDA for accepting a language {a^n b^2n n=1}
- Q 27. Compare FM and PDA. Construct a PDA accepting all palindromes over {a, b}.
- Q. 28. Explain about the graphical notation of PDA. Construct an equivalent PDA for the following CFG.

S→aAB bBA

A→bS a

B→aS b

Q.29 Design a PDA for the following CFG:

S ----> E

S ----> SS

S ----> (S)

Show the transition relation on string ( ) ( ) for constructed PDA.

Q30. Construct a CFG equivalent to the following PDA.

PDA= $\{(p, q), (0, 1), \delta, p, q, (Z, X)\}$ , where p is initial state, q is final state.

 $\delta$  is defined as

 $\delta(p,0,Z)=(p,XZ),$ 

 $\delta(p,0,X)=(p,XX),$ 

 $\delta(p,1,X){=}(q,\epsilon),$ 

 $\delta(p,1,X)=(p,\epsilon),$ 

 $\delta(p,\epsilon,Z)=(p,\epsilon).$ 

Q31. Construct PDA from the following Grammar

S→0BB

B→0S/1S/0.

Show Instantaneous description for the string 010000 is generated for PDA?

Q32. Construct the PDA accepting the language L={(ab)^n n=1} by empty stack.

- Q33. Compare Deterministic and Non deterministic PDA. Is it true that non deterministic PDA is more powerful than deterministic PDA? Justify your answer.
- Q34. Convert grammar  $S \rightarrow aSA$ ,  $A \rightarrow bSa$ ,  $S \rightarrow \Lambda$  to pda that accepts the same language by empty stack.
- Q35. Write down the rules for Construction of PDA equivalent of CFG.

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S ---> aSa
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S ---> bSb

S ----> c

Q36. Construct a deterministic pda accepting  $L=\{w \in \{a,b\}^* \mid \text{the number of a's in } w \text{ equal number of b's in } w\}$  by final state.

- Q37. Design PDA for Language W c W<sup>R</sup>, W  $\in$  (a, b) \*
- Q38. Construct PDA from the following Grammar.

S →aB

 $B \rightarrow bA/b$ 

A →aB

- Q39. A PDA is more powerful than a finite automaton. Justify this statement.
- Q40. Construct a PDA which recognizes all strings that contain equal number of 0's and 1's.
- Q41. Construct Pushdown automata for L = {  $a^{(2*m)}c^{(4*n)}d^n b^m m, n \ge 0$ }
- Q42. Construct a PDA for language  $L = \{wcw' w = \{0, 1\}^*\}$  where w' is the reverse of w.
- Q43. Design a non deterministic PDA for accepting the language  $L = \{a^n b^n n=1\}$
- Q44. Construct a PDA for language L = {0^n 1^m 2^m 3^n n=1, m=1}
- Q45. Construct PDA for the language of all strings of 0's and 1's where number of 1's is less than number of 0's accepting by final state.
- Q46. Explain push down automata with an example.
- Q47. Convert to PDA, CFG with productions

S -> aAS/bAB/aB,

A -> bBB/aS/a,

B -> bA/a

- Q48. How will you convert PDA into CFG? Give Examples
- Q49. Describe the following in detail: (CO1)
- (a) Language of NFA and DFA
- (b) ε- Closure of a State
- (c) NFA with  $\varepsilon$  transition

- (d) Accepters and Transducers
- Q50. Draw an NFA that accepts a language L over an input alphabet  $\Sigma = \{a, b\}$  such that L is the set of all strings where  $S^{rd}$  symbol from the right end is 'b'. Also convert the same to DFA.
- Q51. Design DFA for Following Languages:
- (a) Design a FA with  $\Sigma = \{1, 2\}$  which accepts those string which starts with 0 and ends in 1.
- (b) Draw a DFA that accepts a language L over input alphabets  $\Sigma = \{a, b\}$  such that L is the set of all strings that does not contain 'ba' as substring.
- (c) Design a FA with  $\Sigma = \{0, 1\}$  that accepts those string where number of 1's are even.
- (d) Design a FA with  $\Sigma = \{0, 1\}$  accepts the only input 101.
- Q52. Describe the procedure to convert NFA with epsilon to without epsilon?
- Q53. a) Describe Multi tape Turing and Universal Machines.

Explain informally how they can simulate the moves of a Turing Machine(with proper examples.).

- Q54. Discuss halting problem in Turing Machine? Prove that halting problem of turing machine is undecidable.
- Q55. Explain PCP problem. Show that the PCP with two lists  $x = (b, bab^3, ba)$  and  $y = (b^3, ba, a)$  has a solution. Give the solution sequence.
- Q56. Design a Turing Machine to accept the strings having equal number of 0's and 1's.Describe all instantaneous descriptions (ID) from initial ID:q001011 with respect to constructed Turing Machine (assume q as initial state).
- Q57. Explain various types of Turing Machines with example.
- Q58. Convert following Left linear grammar into FA. Also write right linear grammar

A -> Ba/Ab/b

B -> Ca/Bb

C -> Aa/Cb

Q59. Is the following grammar is ambiguous? Justify your answer.

 $S \rightarrow S+S/S - S/id$ 

Q60. Describe the following: (CO3)

- (i) Eliminating the Use Less Symbols in CFG
- (ii) Removal of Unit Production in CFG
- (iii) Removal of Null Production in CFG

Find the Reduced Grammar that is equivalent to the CFG given below :

S ---> AB

A ---> a

B ---> C / b

C ---> D

D ----> E

E ----> a