

2407. Longest Increasing Subsequence II

Problem Statement: [Longest Increasing Subsequence II - LeetCode](#)

Disclaimer: Please try your best before jumping into the solution.

We store the Longest Increasing Subsequence ended by each number in array LIS (1 – indexed). Let's say the input **nums = [4, 2, 4, 5, 9]**. In other words, indices of LIS represent actual number in input array.

Initially, **LIS[i] = 0** as we haven't added any number yet.

N	1	2	3	4	5	6	7	8	9
LIS	0	0	0	0	0	0	0	0	0

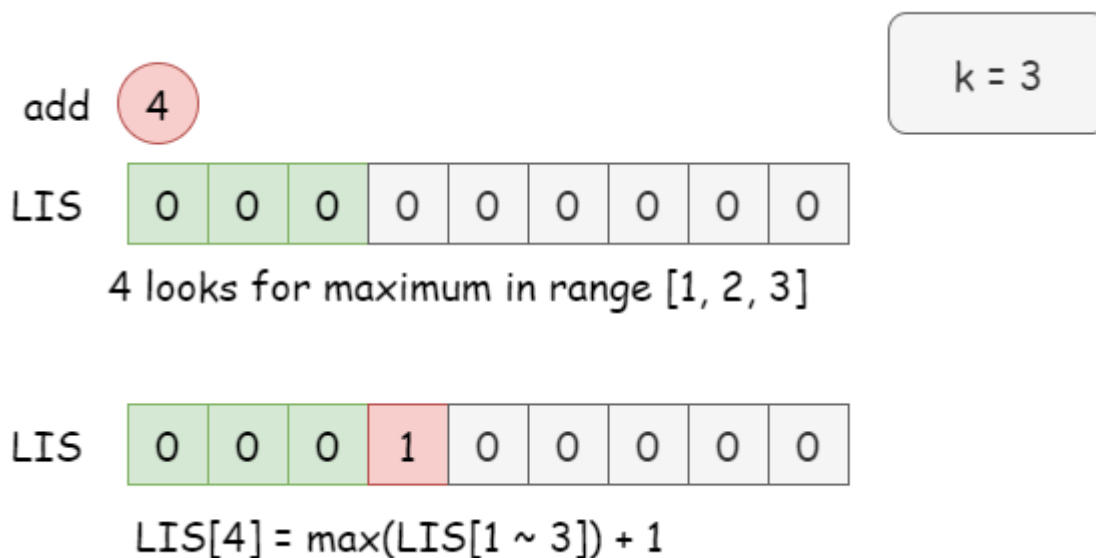
k = 3

nums = [4, 2, 4, 5, 9]

Now, the key is for a given value x , we should find the maximum value from **LIS[a - k.....x-1]**, then **LIS[x] = 1 + max(LIS[a-k.....x-1])**. In other words, we will get maximum LIS from previous subsequence that ends at a number y which lies in the range **[x-k.....x-1]** since it is given that $x - y \leq k$

Let's take the above example and see how it works:

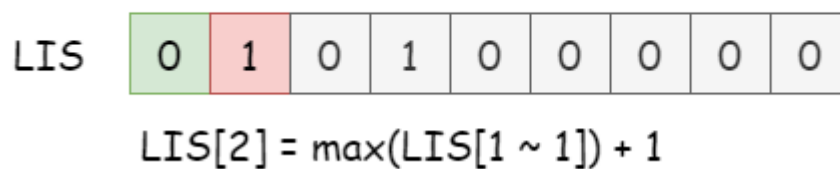
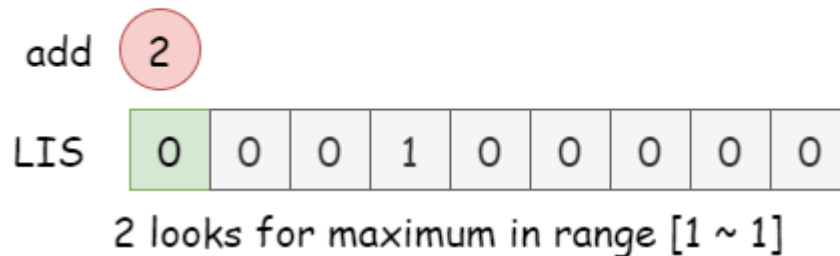
- For the first number **4**, the maximum length is the maximum of **LIS[1], LIS[2], LIS[3]** plus **1** (4 itself). Thus we shall look for the **max(LIS[1...3])**. Apparently, **LIS[4] = 1** which stands for **4** itself.



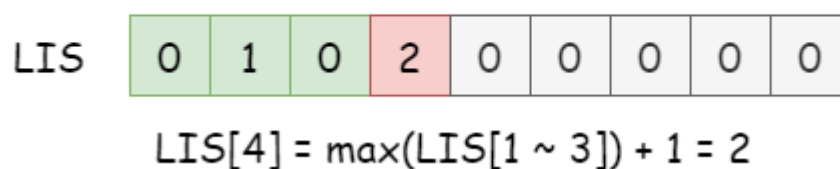
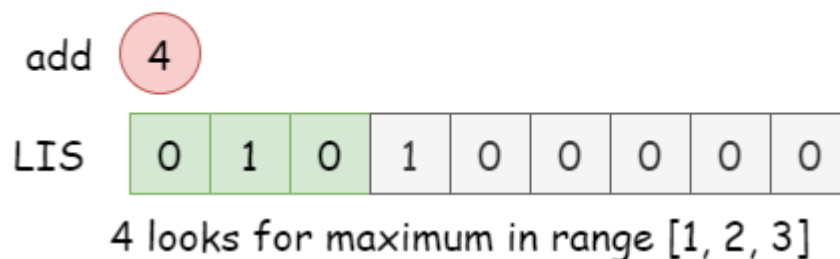
2407. Longest Increasing Subsequence II

Problem Statement: [Longest Increasing Subsequence II - LeetCode](#)

- Then update LIS for 2, we shall look for $\max(\text{LIS}[1\dots 1])$, $\rightarrow \text{LIS}[2] = 1 + \max(\text{LIS}[1\dots 1])$



- Then update LIS for 4, we look for $\max(\text{LIS}[1\dots 3])$ $\rightarrow \text{LIS}[4] = 1 + \max(\text{LIS}[1\dots 3])$. Since there is an 2 updated in LIS, thus the maximum value from the same range $\text{LIS}[1\dots 3]$ gives us 1. Then we can update $\text{LIS}[4] = 2$, implying a subsequence of 2, 4.



2407. Longest Increasing Subsequence II

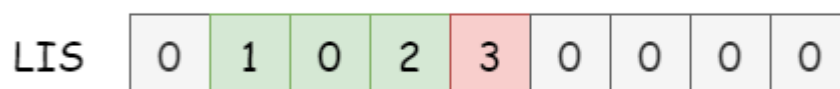
Problem Statement: [Longest Increasing Subsequence II - LeetCode](#)

- Update LIS for 5, $\rightarrow \text{LIS}[5] = 1 + \max(\text{LIS}[2 \dots 4])$ as $k = 3$, implying a subsequence of 2, 4, 5.



$k = 3$

5 looks for maximum in range $[2 \sim 4]$



$$\text{LIS}[5] = \max(\text{LIS}[2 \sim 4]) + 1 = 3$$

- Then finally, update LIS for 9, $\rightarrow \text{LIS}[9] = 1 + \max(\text{LIS}[6 \dots 8])$



$k = 3$

9 looks for maximum in range $[6 \sim 8]$



$$\text{LIS}[9] = \max(\text{LIS}[6 \sim 8]) + 1 = 1$$

Note: Since we are querying the max value in a range $\text{LIS}[x-k \dots x-1]$, we can use Segment Tree which gives us better complexity.

2407. Longest Increasing Subsequence II

Problem Statement: [Longest Increasing Subsequence II - LeetCode](#)

Implementation in C++

```
1  const int mxN = 1e5;
2  int n;
3
4  struct SegTree {
5      int a[4*mxN];
6      SegTree() {
7          memset(a, 0, sizeof(a));
8      }
9
10     void upd(int idx, int x, int i=1, int l2=0, int r2=n) {
11         if(l2 == r2) {
12             a[i] = x;
13             return;
14         }
15
16         int m2 = (l2 + r2)/2;
17         if(idx <= m2)
18             upd(idx, x, 2*i, l2, m2);
19         else
20             upd(idx, x, 2*i+1, m2+1, r2);
21         a[i] = max(a[2*i], a[2*i+1]);
22     }
23
24     int qry(int l1, int r1, int i=1, int l2=0, int r2=n) {
25         if(l2>r1 || r2<l1) return 0;
26         if(l2>=l1 && r2<=r1) return a[i];
27
28         int m2 = (l2 + r2)/2;
29         return max(qry(l1, r1, 2*i, l2, m2), qry(l1, r1, 2*i+1, m2+1, r2));
30     }
31 };
32
33 class Solution {
34 public:
35     int lengthOfLIS(vector<int> &nums, int k) {
36         n = *max_element(nums.begin(), nums.end());
37
38         SegTree st;
39         int answer = 0;
40         for(int x: nums) {
41             // get maximum lis which ends with a number in a range [x-k.....x)
42             int l1 = max(0, x-k), r1 = max(0, x-1);
43             int cnt = st.qry(l1, r1);
44             st.upd(x, cnt+1);
45
46             answer = max(answer, cnt+1);
47         }
48         return answer;
49     }
50 };
```

Credit: LeetCode