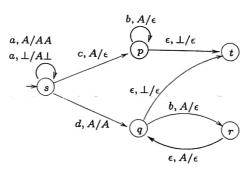
## Automata Theory and Computability

## Final Examination 2024

Max Marks: 70, Weightage: 40%, Time: 3 hours. Write your answers neatly, briefly, and to the point. If required write your answers on a rough sheet first. There is no need to write proofs, except when explicitly asked to justify your answer.

- A run in a string is a maximal contiguous subsequence of length 1 or more of the same letter. For example the string "aabaaaabb" has runs (of a's) of length 2 and 4, and runs (of b's) of length 1 and 2; but no run of length 3. Give a DFA for the language of all strings over the alphabet {a, b} in which every run is of length 3 or more.
- 2. Give a context-free grammar for the language of strings over {0,1} which have strictly more 1's than 0's.

  (5)
- 3. Consider the PDA M below, which runs on the input alphabet  $A = \{a, b, c, d\}$ , has stack alphabet  $\Gamma = \{\bot, A\}$ , and accepts by empty stack. Describe the language accepted by the PDA. (5)



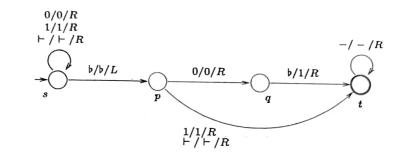
4. Consider the language L of all strings of the form

$$w_1 \# w_2 \# \cdots \# w_n \quad (n \ge 1)$$

where each  $w_i \in \{a, b\}^+$  and all  $w_i$ 's are of the same length. Is L a context-free language? Justify your answer. (10)

- 5. Recall that two languages L and L' over an ordered alphabet A are called "letter equivalent" iff  $\psi(L) = \psi(L')$ , where  $\psi$  is the Parikh map corresponding to A. In the questions below assume that A is the alphabet  $\{a,b\}$ .
  - (a) Prove every CFL L over A is letter-equivalent to a regular language R over A.
  - (b) Is it true that every language over A is letter-equivalent to a regular language over A? Justify your answer.

6. Describe the function on binary strings computed by the Turing machine M below. (5)



- 7. Let L and M be arbitrary languages over the alphabet  $\{a,b\}$  such that L is recursive and M is recursively enumerable. Then which of the following statements must be true?
  - (a)  $L \cap M$  is recursive.
  - (b)  $L \cap M$  is recursively enumerable but not necessarily recursive.
  - (c)  $L \cap M$  is not recursively enumerable.

- 8. Consider the languages containing encodings of Turing machines below. One of them is recursively enumerable and the other is not. Which is which? Justify your answer. (10)
  - TOTAL =  $\{enc(M) \mid M \text{ halts on all inputs}\}$
  - PARTIAL =  $\{enc(M) \mid M \text{ halts on some input}\}$
- 9. Is it decidable whether a given CFG accepts a *regular* language? Justify your answer. (10)
- 10. In the proof of Gödel's Incompleteness Theorem done in class, an important step is to be able to pick out the digit at a particular position in a number. Give a formula, 3digit(v, y, a, b, c), which says that in the base-7 representation of the number v, the 7-ary digits at position y, y + 1, and y + 2 are a, b and c respectively. Assume that the position y is represented as a power of 7, and position 0 corresponds to the least significant 7-ary digit.