

2001. Number of Pairs of Interchangeable Rectangles

Medium

Topics

Companies

Hint

You are given n rectangles represented by a 0-indexed 2D integer array `rectangles`, where `rectangles[i] = [widthi, heighti]` denotes the width and height of the i^{th} rectangle.

Two rectangles i and j ($i < j$) are considered **interchangeable** if they have the **same** width-to-height ratio. More formally, two rectangles are **interchangeable** if `widthi/heighti == widthj/heightj` (using decimal division, not integer division).

Return the **number** of pairs of **interchangeable** rectangles in `rectangles`.

Example 1:

Input: `rectangles = [[4,8],[3,6],[10,20],[15,30]]`

Output: 6

Explanation: The following are the interchangeable pairs of rectangles by index (0-indexed):

- Rectangle 0 with rectangle 1: $4/8 == 3/6$.
- Rectangle 0 with rectangle 2: $4/8 == 10/20$.
- Rectangle 0 with rectangle 3: $4/8 == 15/30$.
- Rectangle 1 with rectangle 2: $3/6 == 10/20$.
- Rectangle 1 with rectangle 3: $3/6 == 15/30$.
- Rectangle 2 with rectangle 3: $10/20 == 15/30$.

Example 2:

Input: `rectangles = [[4,5],[7,8]]`

Output: 0

Explanation: There are no interchangeable pairs of rectangles.

Constraints:

- `n == rectangles.length`
- `1 <= n <= 105`
- `rectangles[i].length == 2`
- `1 <= widthi, heighti <= 105`

Approach 1: Using `sort()`

→ store width/height ratio of each

→ store width/height ratio of each rectangle in a ratios vector.

→ now sort ratios vector.

→ now use two pointers approach to count.
[0.5 0.5 0.5 0.6 0.6 0.7]

```
l=0, r=0
while(r < ratios.size())
{
    if(ratios[l] == ratios[r])
    {
        ans = ans + r - l
        r++
    }
    else
    {
        l = r
        r++
    }
}
```

$T(n): O(n) + O(n \log n) + O(n)$
 $S(n): O(n)$

Approach 2: Using hashmap

→ Store width/height of each rectangle in a hashmap with width/height as key and frequency of width/height as value.

→ now traverse map and count ans using frequencies.

```
for (auto i : map)
{
    x = i.second
    ans = ans + x * C2
}
```

$$T(n) : O(n) + O(n)$$

$$S(n) : O(n)$$