

# Diffie-Hellman: issues

- Diffie-Hellman is not secure against a MITM adversary
- Diffie-Hellman does not provide *authentication*
  - You exchanged keys with someone, but Diffie-Hellman makes no guarantees about who you exchanged keys with; it could be David!
- DHE is an active protocol: Alice and Bob need to be online at the same time to exchange keys
  - What if Bob wants to encrypt something and send it to Alice for her to read later?

TLS 1.3

# Diffie-Hellman Key Exchange: Summary

- Algorithm:

P2P

- Alice chooses  $X_A$  and sends  $\alpha^{X_A} \bmod q$  to Bob
- Bob chooses  $X_B$  and sends  $\alpha^{X_B} \bmod q$  to Alice
- Their shared secret is  $(\alpha^{X_A})^{X_B} = (\alpha^{X_B})^{X_A} = \alpha^{X_A X_B} \bmod q$
- Diffie-Hellman provides forwards secrecy: Nothing is saved or can be recorded that can ever recover the key
- Diffie-Hellman can be performed over other mathematical groups, such as elliptic-curve Diffie-Hellman (ECDH)
- Issues
  - **Not** secure against MITM
  - Does not provide authenticity
  - Both parties must be online

# DHKE in Python Cryptography Library

- <https://cryptography.io/en/latest/hazmat/primitives/asymmetric/>

# Take home exercises

- SW, “Network Security Essentials”, 6<sup>th</sup> Edition, 2017

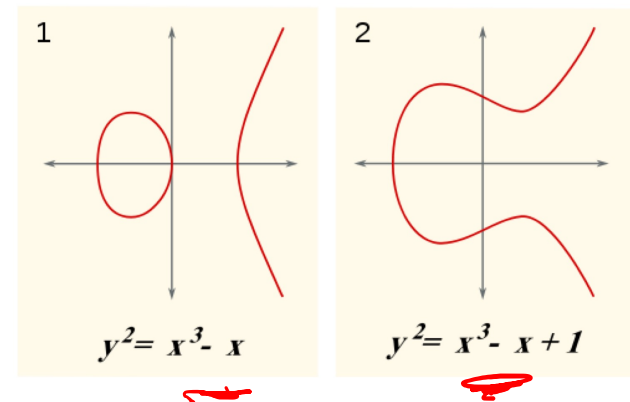
- Problems – 3.21

Consider a Diffie-Hellman scheme with a common prime  $q = 11$  and a primitive root  $\alpha = 2$ .

- a. if user A has public key  $Y_A = 9$ , what is A's private key  $X_A$ ?
- b. If user B has public key  $Y_B = 3$ , what is the shared secret key  $K$ ?

# Elliptic Curve Cryptography (ECC)

- Originally independently proposed by Neal Koblitz (University of Washington) and Victor Miller (IBM) in 1985.
- ECC was proposed as an alternative to other public key encryption algorithms, for example RSA.
- All ECC schemes are public key



CPSC 467: Cryptography and Computer Security , Michael J. Fischer, 2017,  
<https://zoo.cs.yale.edu/classes/cs467/2017f/lectures/ln13.pdf>

# The Elliptic Curve Equation

An elliptic curve over real numbers:

where:




$$y^2 = x^3 + ax + b$$

- $4a^3 + 27b^2 \neq 0$  (to avoid singularities)

Example curve:

$$y^2 = x^3 - 4x + 1$$

## Common ECC Curves

Curve	Field	Bit Strength	Usage
 secp256k1	$\mathbb{F}_p$	128-bit	<u>Bitcoin</u>
 P-256	$\mathbb{F}_p$	128-bit	<u>TLS, NIST</u>
 Curve25519	$\mathbb{F}_p$	128-bit	Signal, SSH

# Why ECC?

- In case of ECC, we are able to use smaller primes, or smaller finite fields, and achieve a level of security comparable to that of RSA
- ECC has lower computational requirements. For this reason, ECC algorithms can be easily implemented on smart cards, paggers, or mobile devices. Some smart cards can only work with ECC.

<u>Symmetric Key Size</u> (bits)	<u>RSA and Diffie-Hellman Key Size</u> (bits)	<u>Elliptic Curve Key Size</u> (bits)
80	1024	160
112	2048	224
128	3072	256
192	7680	384
256	15360	521



# ECC Key Generation

$\rightarrow$  private key

$\rightarrow$  generator

$$PK = k \cdot G$$

- Let  $k$  be an integer and  $G$  a point on  $E$ .  $k \times G$  is defined as adding  $G$  to itself  $k$  times.
- Once we calculate  $Q = k \times G$ , it is extremely difficult to recover  $k$  from  $Q$ . The only way to recover  $k$  from  $Q$  is to try every possible repeated addition of  $G$ .

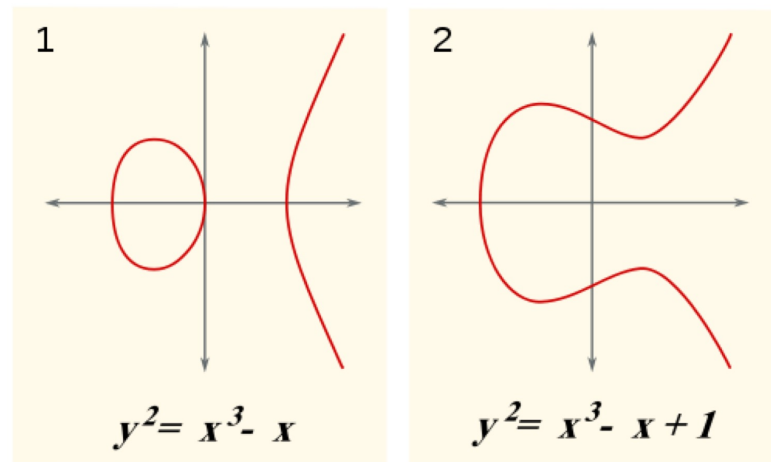
$G$

- Q: Does it sound familiar?

$$e \cdot d \equiv 1 \pmod{\phi(n)}$$

# Elliptic Curve Discrete Logarithm Problem (ECDLP)

- Let  $G$  be a point on  $E$ . Compute  $Q = k \times G$ . Then, ECDLP: given  $G$  and  $Q$  compute  $k$ . ← private key PK
- This allows us to translate crypto schemes based on DLP to EC-based schemes.
- $Q$  is a public key.  $k$  is a private key.  $G$  is a generator point on  $E$ .

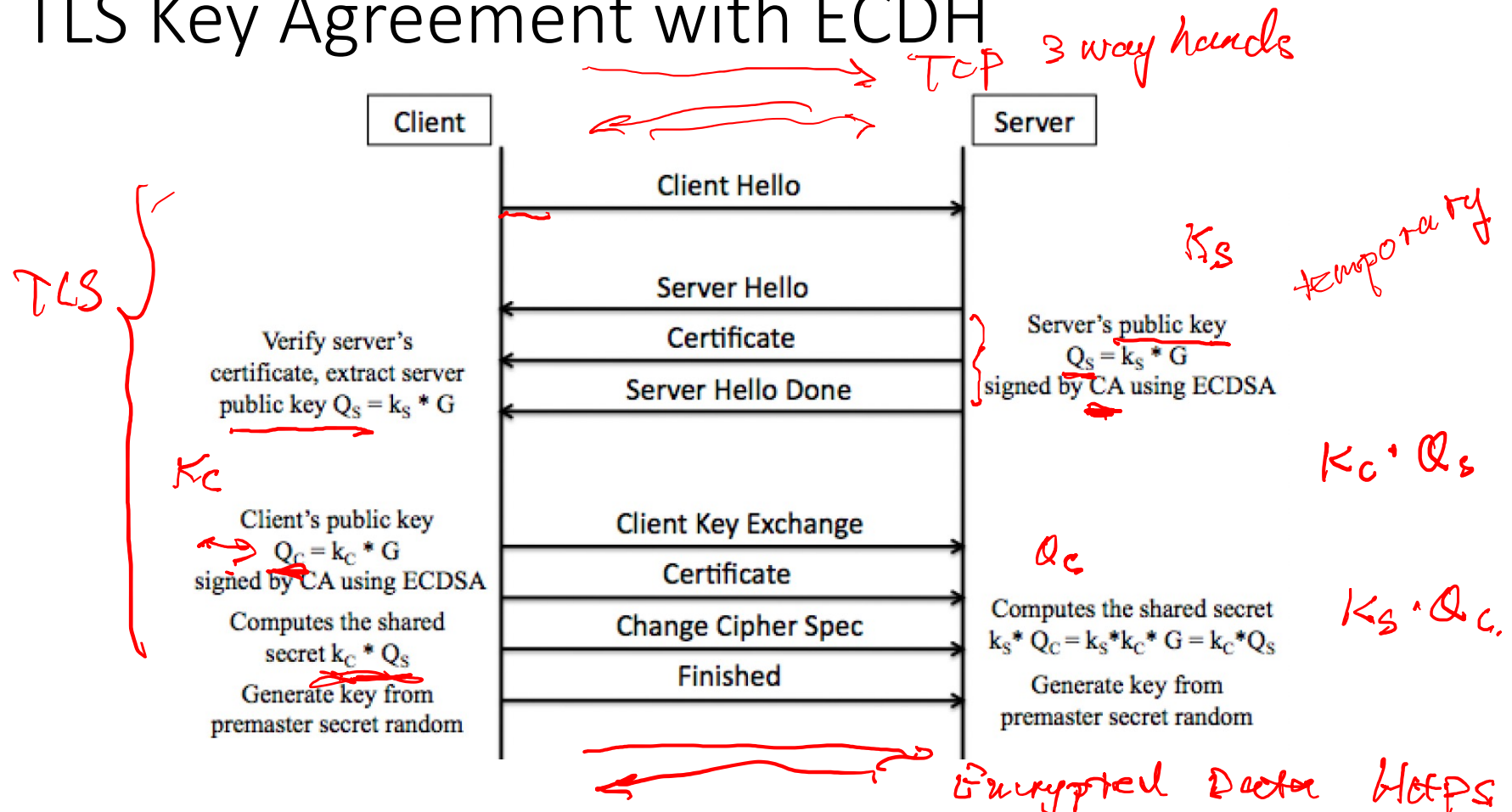


$K \cdot G$   
 $Q = K \cdot G$

## ECC in Practice

- TLS 1.3: Uses ECDHE for forward secrecy.
- **Cryptocurrencies:** Bitcoin, Ethereum (secp256k1).
- **Messaging apps:** Signal, WhatsApp (Curve25519).
- **Government/NIST:** P-256, P-384, P-521.

# TLS Key Agreement with ECDH



# Summary

- ECC achieves **strong security with small keys**.
- Based on the **hardness of ECDLP**.
- Powers many modern systems (TLS, blockchain, mobile apps).