

Public-Key Encryption: Definition

- Three parts:

- $\text{KeyGen}() \rightarrow PK, SK$: Generate a public/private keypair, where PK is the public key, and SK is the private (secret) key
- $\text{Enc}(PK, M) \rightarrow C$: Encrypt a plaintext M using public key PK to produce ciphertext C
- $\text{Dec}(SK, C) \rightarrow M$: Decrypt a ciphertext C using secret key SK

- Properties

- **Correctness**: Decrypting a ciphertext should result in the message that was originally encrypted
unique M
 - $\text{Dec}(SK, \text{Enc}(PK, M)) = M$ for all $PK, SK \leftarrow \text{KeyGen}()$ and M
- **Efficiency**: Encryption/decryption should be fast
- **Security**:
 - 1. Alice (the challenger) just gives Eve (the adversary) the public key, and Eve doesn't request encryptions. Eve cannot guess out anything;
 - 2. computationally infeasible to recover M with PK and ciphertext (PK, C)

$$\begin{array}{ccc} \checkmark & y = x^2 ? & \times \\ \nearrow & \uparrow & \nearrow \\ \text{ciphertext} & M & \text{plaintext} \\ x=1 & & x=-1 \end{array}$$

$$\begin{array}{ccc} & \cancel{\xrightarrow{SK}} & \\ PK & \cancel{\xrightarrow{\text{decrypt}}} & \end{array}$$

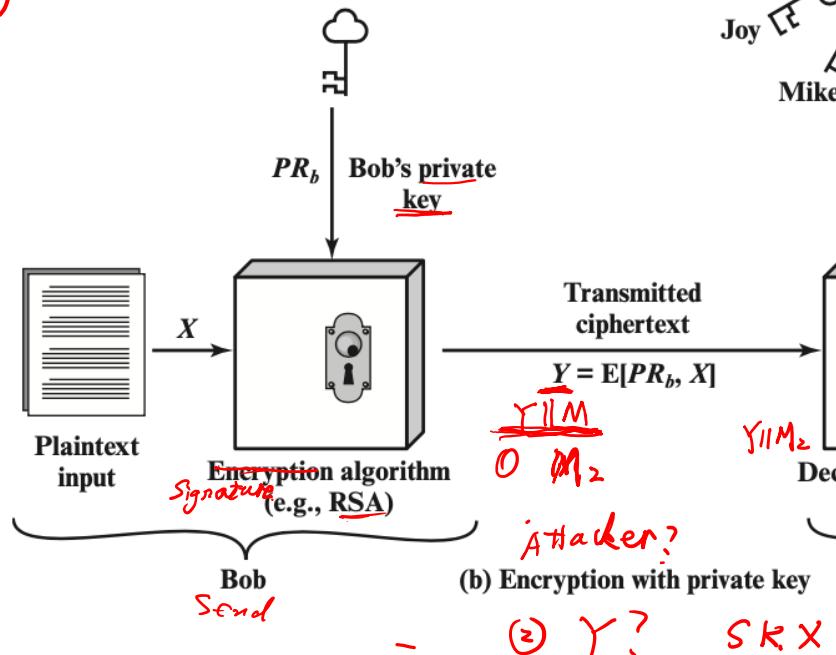
$$\cancel{(PK, C) \rightarrow M}$$

Public-Key Cryptography - Signature

Motivation:

- ✓ 1. prove source or id of doc
- ✓ 2. detect modification → integrity

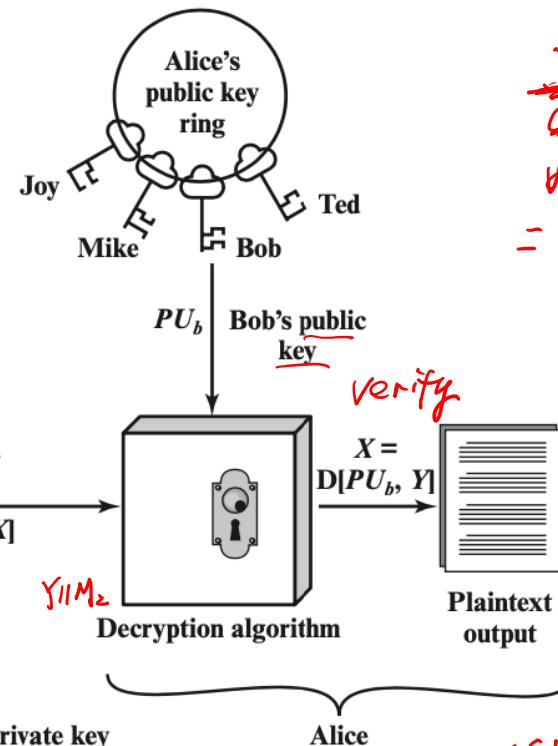
~~Confidentiality~~



$$Y = E[PR_b, X]$$

$Y \parallel M_2$

Attacker?
② Y? SK X



$$X = D[PU_b, Y]$$

$Y \parallel M_2$

③ $Y_2 \parallel M_2 \leftarrow^{SK_{\text{attacker}}}$
 $\leftarrow_{\text{man-in-middle attack}}$

$$\begin{array}{c} Y \parallel M \\ \downarrow Q \\ \text{Verify } [PK, Y] \end{array}$$

$= M'$

$M' \neq M$
if Yes, No modification,
else, attack.

Source?
(SK, PK) ✓
only owned by
Bob

$M_2 \neq M'$

At least one signature has problems. Please fill out the following form.

cd0d6521-c54d-4827-a20e-c455 09/25/2025

Signature validation status

Signature validity is UNKNOWN. *integrity*

- The revision of the document that was covered by this signature has not been altered; however, there have been subsequent changes to the document.
- The signer's identity is unknown because it has not been included in your list of trusted certificates and none of its parent certificates are trusted certificates.
- Click Signature Properties and then click View Signed Version to see what is covered by this signature.

Signature Properties... Close

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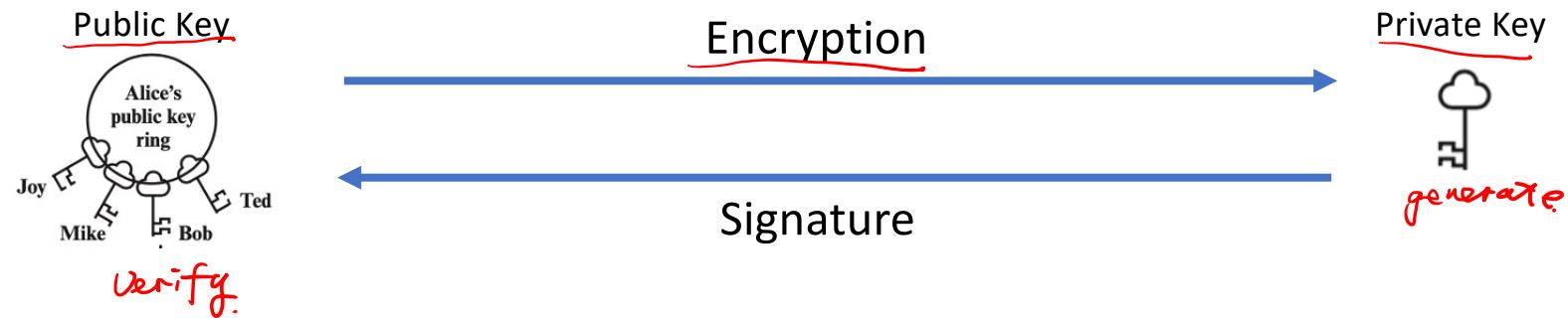
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Signature Panel Highlight Existing Fields

7 9 8 5 6 4 3 2 1

Review



Public-Key application

- can classify uses into 3 categories:
 - encryption/decryption (provide secrecy)
 - digital signatures (provide authentication)
 - key exchange (of session keys)
- some algorithms are suitable for all uses; others are specific to one
- Either of the two related keys can be used for encryption, with the other used for decryption

Algorithm	Encryption/Decryption	Digital Signature	Key Exchange
RSA	Yes	Yes	Yes
Diffie–Hellman	No	No	Yes
DSS	No	Yes	No
Elliptic curve	Yes	Yes	Yes

TLS 1.2 – Use Public Key for Session Key Exchange

~~TLS 1.2~~
TCP/IP stack

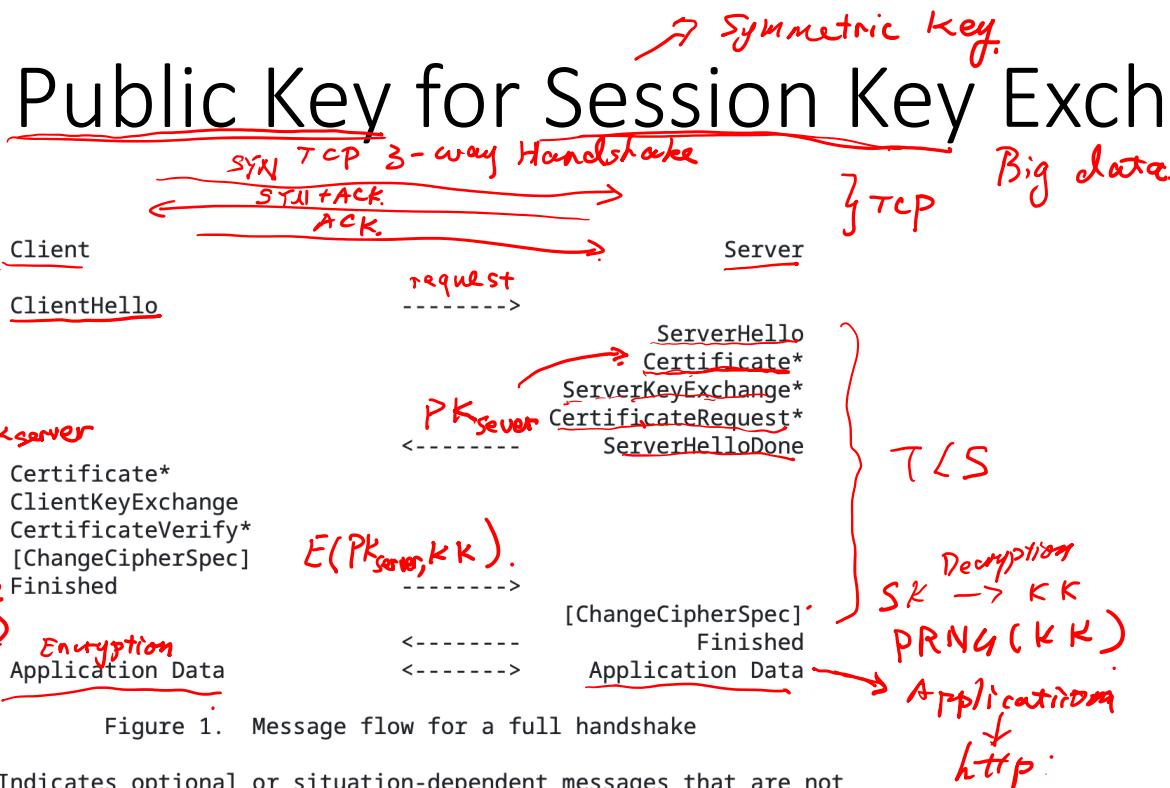
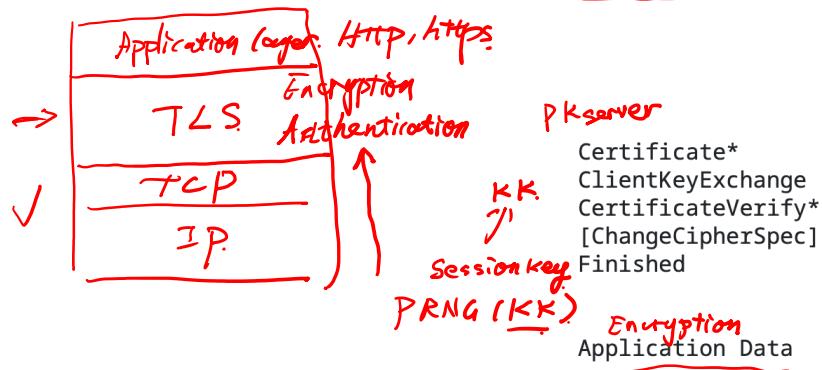


Figure 1. Message flow for a full handshake

* Indicates optional or situation-dependent messages that are not always sent.

Note: To help avoid pipeline stalls, ChangeCipherSpec is an independent TLS protocol content type, and is not actually a TLS handshake message.

RFC 5246: The Transport Layer Security (TLS) Protocol - Version 1.2

Security of Public Key Schemes

- Keys used are very large (>512 bits) 4096 bits
 - like private key schemes brute force **exhaustive search** attack is always theoretically possible
- Security relies on a large enough difference in **difficulty** between easy (en/decrypt) and hard (cryptanalyze) problems
 - more generally the hard problem is known, it's just made too hard to do in practice
 $p \cdot q \rightarrow \text{prime}$
- Requires the use of very large numbers, hence is slow compared to private/symmetric key schemes

