

Answer to the question no. 2

Implementation 1

```
def fibonacci(n):
```

```
    if (n <= 0):
```

```
        print("Invalid input!")
```

```
    elif (n <= 2):
```

```
        return n-1
```

```
    else:
```

```
        return fibonacci(n-1) + fibonacci(n-2)
```

```
n = int(input("Enter a number: "))
```

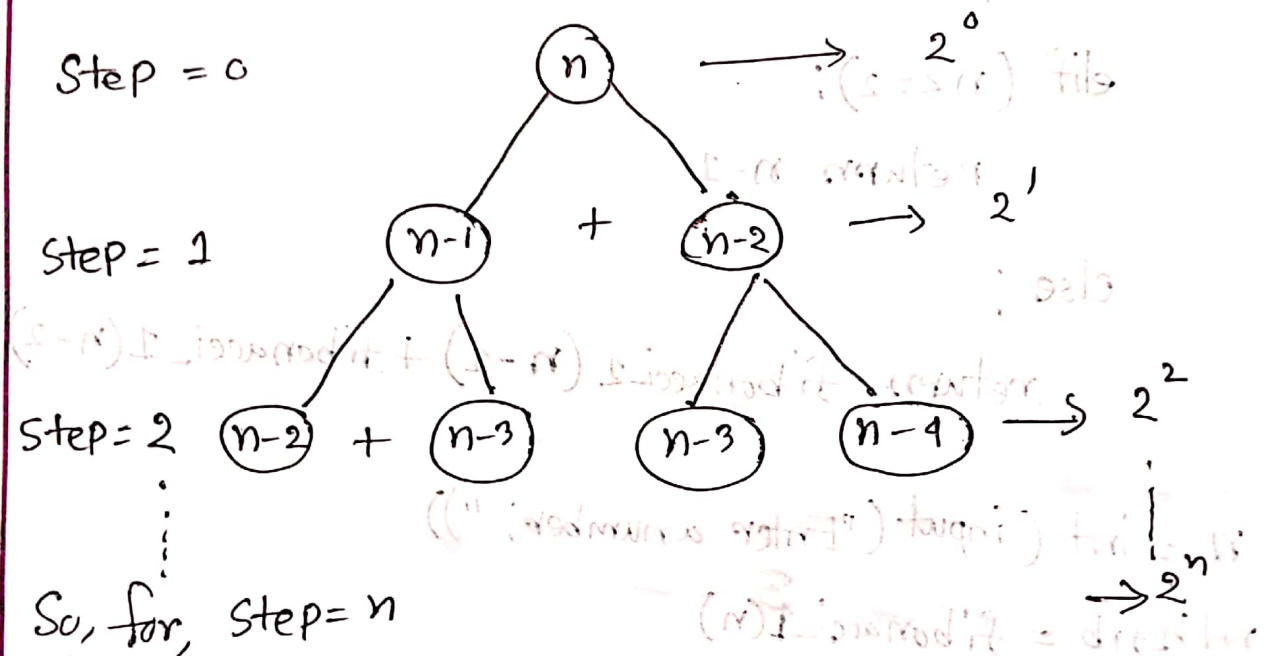
```
nth_fib = fibonacci(n)
```

```
print("The %d-th fibonacci number is %d" % (n, nth_fib))
```

Here,

$$T(n) = \begin{cases} 1 & n \leq 0 \\ 1 & n \leq 2 \\ T(n-1) + T(n-2) & n > 2 \end{cases}$$

Using recursive tree:



So, $2^0 + 2^1 + 2^2 + \dots + 2^n = 2^{n+1} - 1$

$\therefore O(2^{n+1} - 1) = O(2^n)$

Implementation - 2

```
def fibonacci_2(n):  
    fibonacci_array = [0, 1]  
    if n < 0:  
        print("Invalid input!")  
  
    elif n <= 2:  
        return fibonacci_array[n-1]  
  
    else:  
        for i in range(2, n):  
            fibonacci_array.append(fibonacci_array[i-1] +  
                                    fibonacci_array[i-2])  
        return fibonacci_array[-1]  
  
n = int(input("Enter a number: "))  
nth_fib = fibonacci_2(n).  
print("The {}th fibonacci number is {}".format(n, nth_fib))
```

P.T.O

$$\begin{aligned}
 \therefore O(1) + O(1) + O(1) + O(n) \times O(1) \\
 = O(1) + O(1) + O(1) + O(n) \\
 = O(n)
 \end{aligned}$$

So,

Time complexity for 1st implementation is $O(2^n)$.

And, for the 2nd implementation is $O(n)$.

So, by comparing these two, 2nd implementation $O(n)$ is better than the 1st one $O(2^n)$.

Ans no : 9

<u>i</u>	<u>j</u>	<u>k</u>
0	0	n
0	1	n
0	2	n
⋮	⋮	⋮
⋮	⋮	⋮
0	n	n
⋮	⋮	⋮
⋮	⋮	⋮
n	n	n

Here in the Multiply-matrix method, when $i=0$, then $j=0$, then the k loop will run for n times. Similarly, when $i=0$, j loop will run for n times. Also, Outer loop will run for n times. So, there are 3 loops in the method which will run for n times which makes the complexity of program $O(n^3)$. (Ans)