

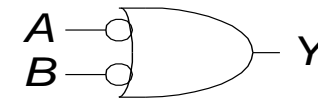
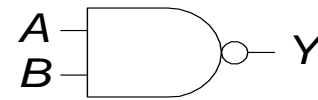
ETE 212/EEE 211

Introduction to Digital Electronics

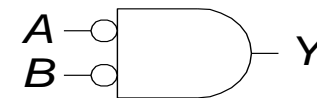
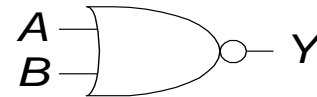
***Boolean Expressions and
Canonical Forms***

DeMorgan's Theorem

- $Y = \overline{AB} = \overline{A} + \overline{B}$



- $Y = \overline{A + B} = \overline{A} \cdot \overline{B}$



Complementing Functions

- Use DeMorgan's Theorem:
 1. Interchange AND and OR operators
 2. Complement each constant and literal
- Example: Complement $F = \bar{x}y\bar{z} + x\bar{y}\bar{z}$
$$\bar{F} = (x + \bar{y} + z)(\bar{x} + y + z)$$
- Example: Complement $G = (\bar{a} + bc)\bar{d} + e$
$$\bar{G} = (a(\bar{b} + \bar{c}) + d)\bar{e}$$

Expression Simplification

- An application of Boolean algebra
- Simplify to contain the smallest number of literals (variables that may or may not be complemented)

$$\begin{aligned} & \mathbf{AB + \bar{A}CD + \bar{A}BD + \bar{A}C\bar{D} + ABCD} \\ &= AB + ABCD + \bar{A}CD + \bar{A}C\bar{D} + \bar{A}BD \\ &= AB + AB(CD) + AC(D + \bar{D}) + \bar{A}BD \\ &= AB + \bar{A}C + \bar{A}BD = B(A + \bar{A}D) + \bar{A}C \\ &= B(A + D) + \bar{A}C \quad (\text{has only 5 literals}) \end{aligned}$$

Canonical Forms

- Minterms and Maxterms
- Sum-of-Minterm (SOM) Canonical Form
- Product-of-Maxterm (POM) Canonical Form
- Representation of Complements of Functions
- Conversions between Representations

Minterms

- Minterms are AND terms with every variable present in either true or complemented form.
- Given that each binary variable may appear normal (e.g., x) or complemented (e.g., \overline{x}), there are 2^n minterms for n variables.
- Example: Two variables (X and Y) produce $2 \times 2 = 4$ combinations:

\mathbf{XY} (both normal)

$\mathbf{X\overline{Y}}$ (X normal, Y complemented)

$\mathbf{\overline{X}Y}$ (X complemented, Y normal)

$\mathbf{\overline{X}\overline{Y}}$ (both complemented)

- Thus there are four minterms of two variables.

Maxterms

- Maxterms are OR terms with every variable in true or complemented form.
- Given that each binary variable may appear normal (e.g., x) or complemented (e.g., \overline{x}), there are 2^n maxterms for n variables.
- Example: Two variables (X and Y) produce $2 \times 2 = 4$ combinations:

$X + Y$ (both normal)

$X + \overline{Y}$ (x normal, y complemented)

$\overline{X} + Y$ (x complemented, y normal)

$\overline{X} + \overline{Y}$ (both complemented)

Minterms & Maxterms for 2 variables

- Two variable minterms and maxterms.

x	y	Index	Minterm	Maxterm
0	0	0	$m_0 = \bar{x} \bar{y}$	$M_0 = x + y$
0	1	1	$m_1 = \bar{x} y$	$M_1 = x + \bar{y}$
1	0	2	$m_2 = x \bar{y}$	$M_2 = \bar{x} + y$
1	1	3	$m_3 = x y$	$M_3 = \bar{x} + \bar{y}$

- The minterm m_i should evaluate to 1 for each combination of x and y.
- The maxterm is the complement of the minterm

Minterms & Maxterms for 3 variables

x	y	z	Index	Minterm	Maxterm
0	0	0	0	$m_0 = \bar{x} \bar{y} \bar{z}$	$M_0 = x + y + z$
0	0	1	1	$m_1 = \bar{x} \bar{y} z$	$M_1 = x + y + \bar{z}$
0	1	0	2	$m_2 = \bar{x} y \bar{z}$	$M_2 = x + \bar{y} + z$
0	1	1	3	$m_3 = \bar{x} y z$	$M_3 = x + \bar{y} + \bar{z}$
1	0	0	4	$m_4 = x \bar{y} \bar{z}$	$M_4 = \bar{x} + y + z$
1	0	1	5	$m_5 = x \bar{y} z$	$M_5 = \bar{x} + y + \bar{z}$
1	1	0	6	$m_6 = x y \bar{z}$	$M_6 = \bar{x} + \bar{y} + z$
1	1	1	7	$m_7 = x y z$	$M_7 = \bar{x} + \bar{y} + \bar{z}$

Maxterm M_i is the complement of minterm m_i

$$M_i = \overline{m_i} \text{ and } m_i = \overline{M_i}$$

Purpose of the Index

- Minterms and Maxterms are designated with an index
- The index number corresponds to a binary pattern
- The index for the minterm or maxterm, expressed as a binary number, is used to determine whether the variable is shown in the true or complemented form
- For Minterms:
 - ‘1’ means the variable is “Not Complemented” and
 - ‘0’ means the variable is “Complemented”.
- For Maxterms:
 - ‘0’ means the variable is “Not Complemented” and
 - ‘1’ means the variable is “Complemented”.

Standard Order

- All variables should be present in a minterm or maxterm and should be listed in the same order (usually alphabetically)
- Example: For variables a, b, c:
 - Maxterms $(a + b + \bar{c})$, $(\bar{a} + b + \bar{c})$ are in standard order
 - However, $(b + \bar{a} + c)$ is NOT in standard order
 $(\bar{a} + c)$ does NOT contain all variables
 - Minterms $(a b \bar{c})$ and $(\bar{a} b \bar{c})$ are in standard order
 - However, $(b a \bar{c})$ is not in standard order
 $(\bar{a} c)$ does not contain all variables

Sum-Of-Minterm (SOM)

- Sum-Of-Minterm (SOM) canonical form:
Sum of minterms of entries that evaluate to '**1**'

x	y	z	F	Minterm
0	0	0	0	
0	0	1	1	$m_1 = \bar{x} \bar{y} z$
0	1	0	0	
0	1	1	0	
1	0	0	0	
1	0	1	0	
1	1	0	1	$m_6 = x y \bar{z}$
1	1	1	1	$m_7 = x y z$

Focus on the
'**1**' entries

$$F = m_1 + m_6 + m_7 = \sum (1, 6, 7) = \bar{x} \bar{y} z + x y \bar{z} + x y z$$

Sum-Of-Minterm Examples

- $F(a, b, c, d) = \sum(2, 3, 6, 10, 11)$
- $F(a, b, c, d) = m_2 + m_3 + m_6 + m_{10} + m_{11}$
 $\bar{a} \bar{b} c \bar{d} + \bar{a} \bar{b} c d + \bar{a} b c \bar{d} + a \bar{b} c \bar{d} + a \bar{b} c d$
- $G(a, b, c, d) = \sum(0, 1, 12, 15)$
- $G(a, b, c, d) = m_0 + m_1 + m_{12} + m_{15}$
 $\bar{a} \bar{b} \bar{c} \bar{d} + \bar{a} \bar{b} \bar{c} d + a b \bar{c} \bar{d} + a b c d$

Product-Of-Maxterm (POM)

- Product-Of-Maxterm (POM) canonical form:
Product of maxterms of entries that evaluate to ‘0’

x	y	z	F	Maxterm
0	0	0	1	
0	0	1	1	
0	1	0	0	$M_2 = (x + \bar{y} + z)$
0	1	1	1	
1	0	0	0	$M_4 = (\bar{x} + y + z)$
1	0	1	1	
1	1	0	0	$M_6 = (\bar{x} + \bar{y} + z)$
1	1	1	1	

Focus on the
‘0’ entries

$$F = M_2 \cdot M_4 \cdot M_6 = \prod (2, 4, 6) = (x + \bar{y} + z) (\bar{x} + y + z) (\bar{x} + \bar{y} + z)$$

Product-Of-Maxterm Examples

- $F(a, b, c, d) = \prod(1, 3, 6, 11)$
- $F(a, b, c, d) = M_1 \cdot M_3 \cdot M_6 \cdot M_{11}$
 $(a+b+c+\bar{d})(a+b+\bar{c}+\bar{d})(a+\bar{b}+\bar{c}+d)(\bar{a}+b+\bar{c}+\bar{d})$
- $G(a, b, c, d) = \prod(0, 4, 12, 15)$
- $G(a, b, c, d) = M_0 \cdot M_4 \cdot M_{12} \cdot M_{15}$
 $(a+b+c+d)(a+\bar{b}+c+d)(\bar{a}+\bar{b}+c+d)(\bar{a}+\bar{b}+\bar{c}+\bar{d})$

Observations

- We can implement any function by "ORing" the minterms corresponding to the '**1**' entries in the function table. A minterm evaluates to '**1**' for its corresponding entry.
- We can implement any function by "ANDing" the maxterms corresponding to '**0**' entries in the function table. A maxterm evaluates to '**0**' for its corresponding entry.
- The same Boolean function can be expressed in two canonical ways: Sum-of-Minterms (SOM) and Product-of-Maxterms (POM).
- If a Boolean function has fewer '**1**' entries then the SOM canonical form will contain fewer literals than POM. However, if it has fewer '**0**' entries then the POM form will have fewer literals than SOM.

Converting to Sum-of-Minterms Form

- A function that is not in the Sum-of-Minterms form can be converted to that form by means of a truth table
- Consider $F = \bar{y} + \bar{x} \bar{z}$

x	y	z	F	Minterm
0	0	0	1	$m_0 = x \bar{y} \bar{z}$
0	0	1	1	$m_1 = x \bar{y} z$
0	1	0	1	$m_2 = x y \bar{z}$
0	1	1	0	
1	0	0	1	$m_4 = x y \bar{z}$
1	0	1	1	$m_5 = x y z$
1	1	0	0	
1	1	1	0	

$$F = \sum(0, 1, 2, 4, 5) =$$

$$m_0 + m_1 + m_2 + m_4 + m_5 =$$

$$\bar{x} \bar{y} \bar{z} + \bar{x} \bar{y} z + \bar{x} y \bar{z} +$$

$$x \bar{y} \bar{z} + x \bar{y} z$$

Converting to Product-of-Maxterms Form

- A function that is not in the Product-of-Minterms form can be converted to that form by means of a truth table
- Consider again: $F = \bar{y} + \bar{x} \bar{z}$

x	y	z	F	Maxterm
0	0	0	1	
0	0	1	1	
0	1	0	1	
0	1	1	0	$M_3 = (x + \bar{y} + \bar{z})$
1	0	0	1	
1	0	1	1	
1	1	0	0	$M_6 = (\bar{x} + \bar{y} + z)$
1	1	1	0	$M_7 = (\bar{x} + \bar{y} + \bar{z})$

$$F = \prod(3, 6, 7) =$$

$$M_3 \cdot M_6 \cdot M_7 =$$

$$(x + \bar{y} + \bar{z}) (\bar{x} + \bar{y} + z) (\bar{x} + \bar{y} + \bar{z})$$

Conversions Between Canonical Forms

x	y	z	F	Minterm	Maxterm
0	0	0	0		$M_0 = (x + y + z)$
0	0	1	1	$m_1 = \bar{x} \bar{y} z$	
0	1	0	1	$m_2 = \bar{x} y \bar{z}$	
0	1	1	1	$m_3 = \bar{x} y z$	
1	0	0	0		$M_4 = (\bar{x} + y + z)$
1	0	1	1	$m_5 = x \bar{y} z$	
1	1	0	0		$M_6 = (\bar{x} + \bar{y} + z)$
1	1	1	1	$m_7 = x y z$	

$$F = m_1 + m_2 + m_3 + m_5 + m_7 = \sum(1, 2, 3, 5, 7) =$$

$$\bar{x} \bar{y} z + \bar{x} y \bar{z} + \bar{x} y z + x \bar{y} z + x y z$$

$$F = M_0 \cdot M_4 \cdot M_6 = \prod(0, 4, 6) = (x + y + z)(\bar{x} + y + z)(\bar{x} + \bar{y} + z)$$

Algebraic Conversion to Sum-of-Minterms

- Expand all terms first to explicitly list all minterms
- AND any term missing a variable v with $(v + \bar{v})$

- Example 1: $f = x + \bar{x} \bar{y}$ (2 variables)

$$f = x (y + \bar{y}) + \bar{x} \bar{y}$$

$$f = x y + x \bar{y} + \bar{x} \bar{y}$$

$$f = m_3 + m_2 + m_0 = \sum(0, 2, 3)$$

- Example 2: $g = a + \bar{b} c$ (3 variables)

$$g = a (b + \bar{b})(c + \bar{c}) + (a + \bar{a}) \bar{b} c$$

$$g = a b c + a b \bar{c} + a \bar{b} c + a \bar{b} \bar{c} + a \bar{b} c + \bar{a} \bar{b} c$$

$$g = \bar{a} \bar{b} c + a \bar{b} \bar{c} + a \bar{b} c + a b \bar{c} + a b c$$

$$g = m_1 + m_4 + m_5 + m_6 + m_7 = \sum(1, 4, 5, 6, 7)$$

Algebraic Conversion to Product-of-Maxterms

- Expand all terms first to explicitly list all maxterms
- OR any term missing a variable v with $v \cdot \bar{v}$
- Example 1: $f = x + \bar{x} \bar{y}$ (2 variables)

Apply 2nd distributive law:

$$f = (x + \bar{x}) (x + \bar{y}) = 1 \cdot (x + \bar{y}) = (x + \bar{y}) = M_1$$

- Example 2: $g = a \bar{c} + b c + \bar{a} \bar{b}$ (3 variables)

$$g = (a \bar{c} + b c + \bar{a}) (a \bar{c} + b c + \bar{b}) \quad (\text{distributive})$$

$$g = (\bar{c} + b c + \bar{a}) (a \bar{c} + c + \bar{b}) \quad (x + \bar{x} y = x + y)$$

$$g = (\bar{c} + b + \bar{a}) (a + c + \bar{b}) \quad (x + \bar{x} y = x + y)$$

$$g = (\bar{a} + b + \bar{c}) (a + \bar{b} + c) = M_5 \cdot M_2 = \prod (2, 5)$$

Function Complements

- The complement of a function expressed as a sum of minterms is constructed by selecting the minterms missing in the sum-of-minterms canonical form
- Alternatively, the complement of a function expressed by a Sum of Minterms form is simply the Product of Maxterms with the same indices
- Example: Given $F(x, y, z) = \sum (1, 3, 5, 7)$
 $\overline{F}(x, y, z) = \sum (0, 2, 4, 6)$
 $\overline{F}(x, y, z) = \prod (1, 3, 5, 7)$

Summary of Minterms and Maxterms

- There are 2^n minterms and maxterms for Boolean functions with n variables.
- Minterms and maxterms are indexed from 0 to $2^n - 1$
- Any Boolean function can be expressed as a logical sum of minterms and as a logical product of maxterms
- The complement of a function contains those minterms not included in the original function
- The complement of a sum-of-minterms is a product-of-maxterms with the same indices

Standard Forms

- Standard Sum-of-Products (SOP) form: equations are written as an OR of AND terms
- Standard Product-of-Sums (POS) form: equations are written as an AND of OR terms
- Examples:
 - SOP: $\mathbf{A B C + \bar{A} \bar{B} C + B}$
 - POS: $\mathbf{(A + B) \cdot (A + \bar{B} + \bar{C}) \cdot C}$
- These “mixed” forms are neither SOP nor POS
 - $\mathbf{(A B + C) (A + C)}$
 - $\mathbf{A B \bar{C} + A C (A + B)}$

Standard Sum-of-Products (SOP)

- A sum of minterms form for n variables can be written down directly from a truth table.
 - Implementation of this form is a two-level network of gates such that:
 - The first level consists of n -input AND gates
 - The second level is a single OR gate
- This form often can be simplified so that the corresponding circuit is simpler.

Standard Sum-of-Products (SOP)

- A Simplification Example:

$$\mathbf{F(A,B,C) = \sum (1,4,5,6,7)}$$

- Writing the minterm expression:

$$F = \overline{A} \overline{B} C + A \overline{B} \overline{C} + A \overline{B} C + A B \overline{C} + A B C$$

- Simplifying:

$$F = \overline{A} \overline{B} C + A (\overline{B} \overline{C} + \overline{B} C + B \overline{C} + B C)$$

$$F = \overline{A} \overline{B} C + A (\overline{B} (\overline{C} + C) + B (\overline{C} + C))$$

$$F = \overline{A} \overline{B} C + A (\overline{B} + B)$$

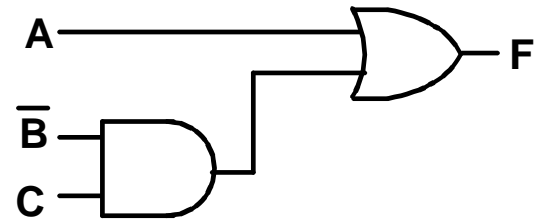
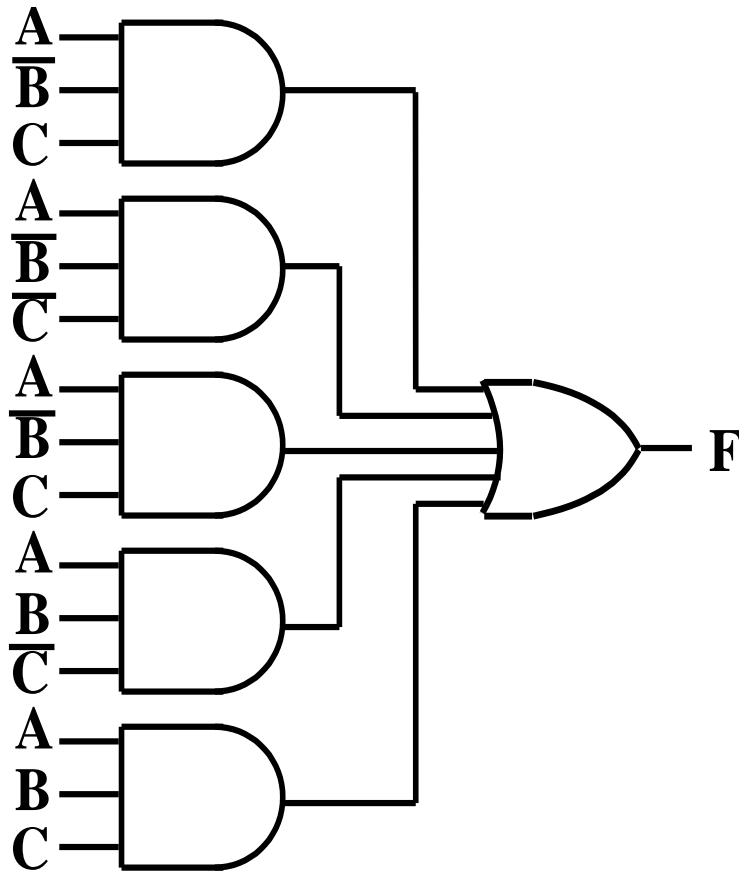
$$F = \overline{A} \overline{B} C + A$$

$$F = \overline{B} C + A$$

- Simplified F contains 3 literals compared to 15

AND/OR Two-Level Implementation

- The two implementations for F are shown below



It is quite
apparent which
is simpler!

SOP and POS Observations

- **The previous examples show that:**
 - Canonical Forms (Sum-of-minterms, Product-of-Maxterms), or other standard forms (SOP, POS) differ in complexity
 - Boolean algebra can be used to manipulate equations into simpler forms
 - Simpler equations lead to simpler implementations
- **Questions:**
 - How can we attain a “simplest” expression?
 - Is there only one minimum cost circuit?
 - The next part will deal with these issues