

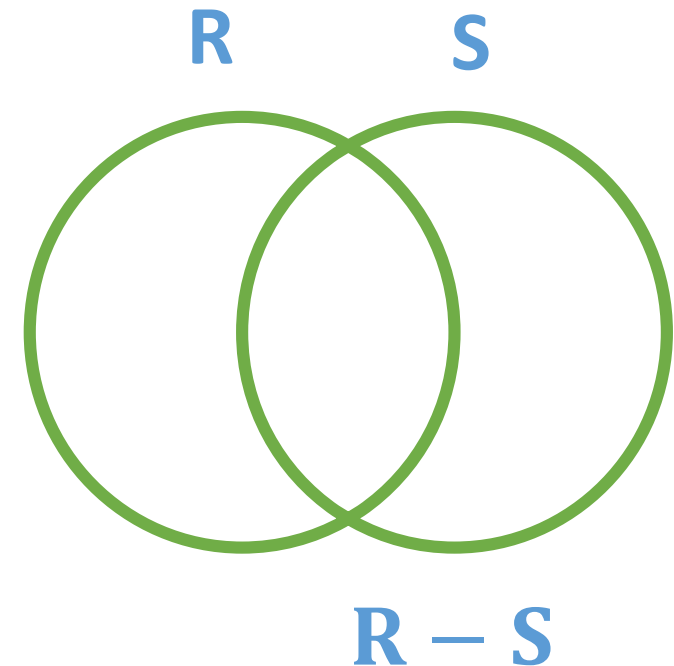
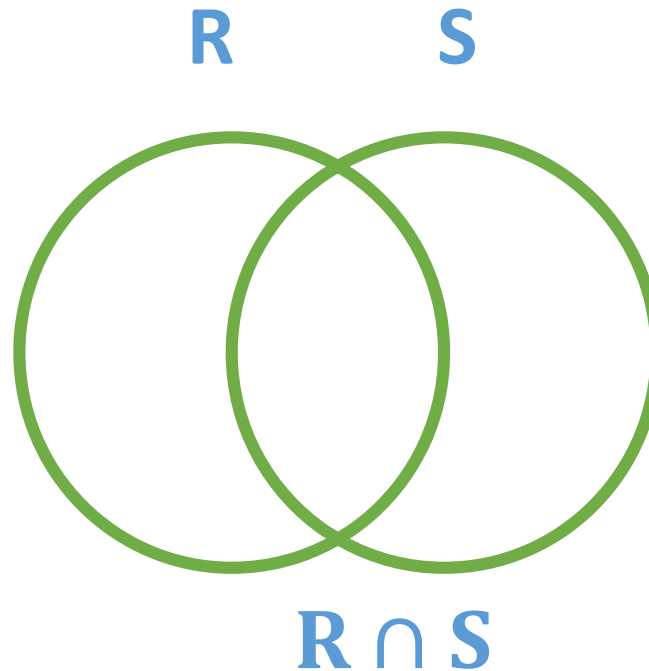
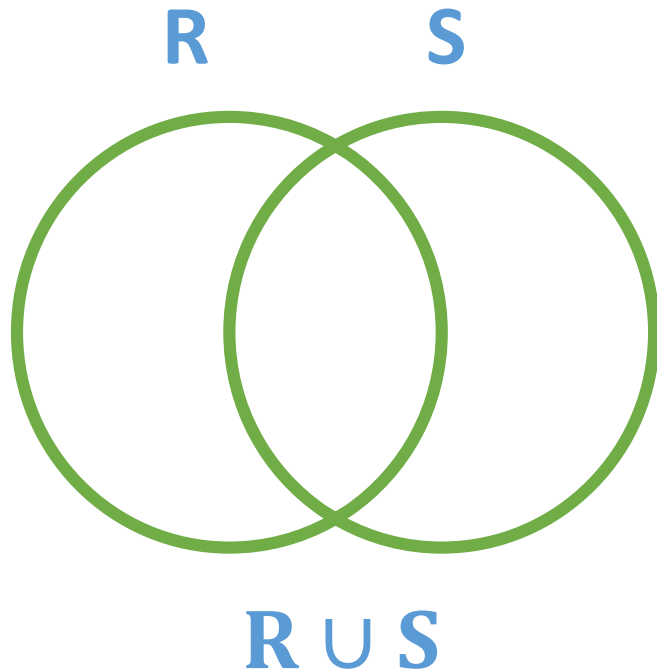
# UNIT 2

## Lecture 16

### Relational Algebra

# Set Operations

- Union ( $R \cup S$ )
- Intersection ( $R \cap S$ )
- Set difference ( $R - S$ )



# Set Operations

R		S		$R \cup S$		$R \cap S$		$R - S$		$S - R$	
A	B	A	B	A	B	A	B	A	B	A	B
1	2	1	2	1	2	1	2	3	4	5	6
3	4	5	6	3	4						
				5	6						

Relations must have union compatible

1. They have same number of columns.
2. Their domains should be same.

$$R \cup S \equiv S \cup R$$

$$R \cap S \equiv S \cap R$$

$$R - S \not\equiv S - R$$

# Union

R

A	B
1	2
3	4

S

A	B
1	2
5	6

$R \cup S$

A	B
1	2
3	4
5	6

$R \cap S$

A	B
1	2

$R - S$

A	B
3	4

$S - R$

A	B
5	6

Equivalent SQL Queries

Select \* from R

union

Select \* from S;

*OUTPUT*

A	B
1	2
3	4
5	6

# Union all in SQL

R

A	B
1	2
3	4

S

A	B
1	2
5	6

$R \cup S$

A	B
1	2
3	4
5	6

$R \cap S$

A	B
1	2

$R - S$

A	B
3	4

$S - R$

A	B
5	6

Equivalent SQL Queries

Select \* from R

union all

Select \* from S;

*OUTPUT*

A	B
1	2
3	4
1	2
5	6

# Intersect

R

A	B
1	2
3	4

S

A	B
1	2
5	6

$R \cup S$

A	B
1	2
3	4
5	6

$R \cap S$

A	B
1	2

$R - S$

A	B
3	4

$S - R$

A	B
5	6

*OUTPUT*

A	B
1	2

Equivalent SQL Queries

Select \* from R

intersect

Select \* from S;

# Set Difference (minus in oracle SQL)

R

A	B
1	2
3	4

S

A	B
1	2
5	6

$R \cup S$

A	B
1	2
3	4
5	6

$R \cap S$

A	B
1	2

$R - S$

A	B
3	4

$S - R$

A	B
5	6

Equivalent SQL Queries

Select \* from R

minus / except

Select \* from S;

*OUTPUT*

A	B
3	4

# Set Difference (minus in oracle SQL)

R

A	B
1	2
3	4

S

A	B
1	2
5	6

$R \cup S$

A	B
1	2
3	4
5	6

$R \cap S$

A	B
1	2

$R - S$

A	B
3	4

$S - R$

A	B
5	6

Equivalent SQL Queries

Select \* from S

minus / except

Select \* from R;

*OUTPUT*

A	B
5	6



# Example : 1

Display the name of branches in which project 121 or 122 or both are running.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{pno} = 121}(\text{STUDENT})) \cup \Pi_{\text{branch}}(\sigma_{\text{pno} = 122}(\text{STUDENT}))$

SQL > select distinct branch from student where pno = 121

Union

select distinct branch from student where pno = 122;

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**OUTPUT**

Branch	Branch	Branch
CSE	CSE	CSE
IT	MECH	IT
		MECH

# Example : 1

Display the name of branches in which project 121 or 122 or both are running.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{pno} = 121 \vee \text{pno} = 122}(\text{STUDENT}))$

SQL > select distinct branch from student where pno = 121 or pno = 122;

or

SQL > select distinct branch from student where pno in (121, 122);

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**OUTPUT**

Branch
CSE
IT
MECH

## Example : 2

Display the name of branches in which project 121 and 122 are running.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{pno} = 121}(\text{STUDENT})) \cap \Pi_{\text{branch}}(\sigma_{\text{pno} = 122}(\text{STUDENT}))$

SQL > select distinct branch from student where pno = 121

intersect

select distinct branch from student where pno = 122;

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**OUTPUT**

Branch	Branch	Branch
CSE	CSE	CSE
IT	MECH	

## Example : 2 (Wrong query)

Display the name of branches in which project 121 and 122 are running.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{pno} = 121 \wedge \text{pno} = 122}(\text{STUDENT}))$

SQL > select distinct branch from student where pno = 121 and pno = 122;

### STUDENT

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

### OUTPUT

Invalid Query

# Example : 3

Display the name of branches in which student got 50 and project 122 is running.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{marks} = 50}(\text{STUDENT})) \cap \Pi_{\text{branch}}(\sigma_{\text{pno} = 122}(\text{STUDENT}))$

SQL > select distinct branch from student where marks = 50

intersect

select distinct branch from student where pno = 122;

## STUDENT

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

## OUTPUT

Branch	Branch	Branch
CSE	CSE	CSE
	MECH	

## Example : 3

Display the name of branches in which student got 50 and project 122 is running.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{marks} = 50 \wedge \text{pno} = 122}(\text{STUDENT}))$

SQL > select distinct branch from student where marks = 50 and pno = 122;

### STUDENT

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

### OUTPUT

Branch
CSE

# Example : 4

Display the name of branches in which project 121 is running but 122 is not.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{pno} = 121}(\text{STUDENT})) - \Pi_{\text{branch}}(\sigma_{\text{pno} = 122}(\text{STUDENT}))$

SQL > select distinct branch from student where pno = 121

minus / except

select distinct branch from student where pno = 122;

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**OUTPUT**

Branch	Branch	Branch
CSE	CSE	IT
IT	MECH	

# Example : 5

Display the name of branches in which project 122 is running but 121 is not.

RA :  $\Pi_{\text{branch}}(\sigma_{\text{pno} = 122}(\text{STUDENT})) - \Pi_{\text{branch}}(\sigma_{\text{pno} = 121}(\text{STUDENT}))$

SQL > select distinct branch from student where pno = 122

minus / except

select distinct branch from student where pno = 121;

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**OUTPUT**

Branch	Branch	Branch
CSE	CSE	MECH
MECH	IT	



# Example : 6

RA :  $\Pi_{\text{branch}}(\sigma_{\text{pno} = 121}(\text{STUDENT}))$

$(\Pi_{\text{branch}}(\sigma_{\text{pno} = 121}(\text{STUDENT})))$

$\Pi_{\text{branch}}(\sigma_{\text{pno} = 122}(\text{STUDENT}))$

## STUDENT

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

## OUTPUT

Branch	Branch	Branch	Branch	Branch
CSE	CSE	CSE	IT	CSE
IT	IT	MECH		

# Intersection is derived from set difference

$$R - (R - S) \equiv S \cap R$$

That is why, intersection is not a basic (fundamental) operation of Relational Algebra

# Example : 7

$\Pi_{pno} (STUDENT) \cup \Pi_{pno} (PROJECT)$

## STUDENT

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

## PROJECT

Pno	Pname	Duration
121	P1	10
122	P2	20
123	P3	30
124	P4	40

## OUTPUT

Pno
121
122
123
124

# Example : 8

$\Pi_{pno} (STUDENT) \cap \Pi_{pno} (PROJECT)$

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**PROJECT**

Pno	Pname	Duration
121	P1	10
122	P2	20
123	P3	30
124	P4	40

**OUTPUT**

Pno
121
122
123

# Example : 9

$\Pi_{pno} (STUDENT) - \Pi_{pno} (PROJECT)$

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**PROJECT**

Pno	Pname	Duration
121	P1	10
122	P2	20
123	P3	30
124	P4	40

**OUTPUT**

Pno

No rows selected

# Example : 9

$$\Pi_{pno} (PROJECT) - \Pi_{pno} (STUDENT)$$

**STUDENT**

Rollno	Sname	Sem	Branch	Marks	Pno
1	RAM	3	CSE	40	121
2	SHYAM	5	CSE	50	122
3	MOHAN	7	CSE	55	123
4	GOPAL	5	IT	65	121
5	RINKI	3	MECH	40	122
6	PINKI	3	ETC	90	123

**PROJECT**

Pno	Pname	Duration
121	P1	10
122	P2	20
123	P3	30
124	P4	40

**OUTPUT**

Pno
124

# GATE question

Suppose R1 (A, B) and R2 (C, D) are two relation schemas. Let r1 and r2 be the corresponding relation instances. B is a foreign key that refers to C in R2. If data in r1 and r2 satisfy referential integrity constraints, which of the following is ALWAYS TRUE?

- (A)  $\Pi_B(r1) - \Pi_C(r2) = \phi$
- (B)  $\Pi_C(r2) - \Pi_B(r1) = \phi$
- (C)  $\Pi_B(r1) = \Pi_C(r2)$
- (D)  $\Pi_B(r1) - \Pi_C(r2) \neq \phi$

**[GATE 2012]**

# GATE question

Which of the following query transformations (i.e. replacing the l.h.s. expression by the r.h.s. expression) is incorrect? R1 and R2 are relations, C1, C2 are selection conditions and A1, A2 are attributes of R1?

(A)  $\sigma_{c_1}(\sigma_{c_2}(R1)) \rightarrow \sigma_{c_1}(\sigma_{c_2}(R1))$

(B)  $\sigma_{c_1}(\Pi_{A_1}(R1)) \rightarrow \Pi_{A_1}(\sigma_{c_1}(R1))$

(C)  $\sigma_{c_1}(R1 \cup R2) \rightarrow \sigma_{c_1}(R1) \cup \sigma_{c_2}(R2)$

(D)  $\Pi_{A_2}(\sigma_{c_1}(R1)) \rightarrow \sigma_{c_1}(\Pi_{A_2}(R1))$

**[GATE 1998]**



For Video lecture on this topic please subscribe to my youtube channel.

The link for my youtube channel is

[https://www.youtube.com/channel/UCRWGtE76JITp1iim6aOTRuW?sub\\_confirmation=1](https://www.youtube.com/channel/UCRWGtE76JITp1iim6aOTRuW?sub_confirmation=1)