

# CSE 2233 Theory of Computing

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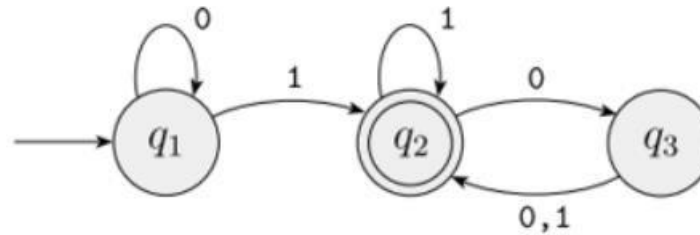
UIU



# Deterministic Finite Automaton (DFA)

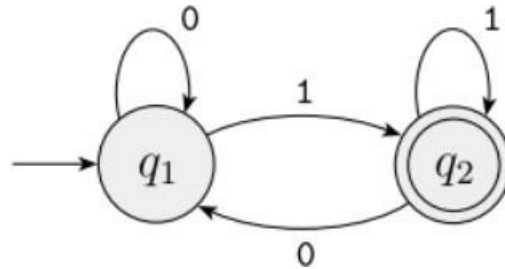
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The term “deterministic finite automata” refers to the fact that on each input there is one and only one state to which the automaton can transition from its current state.



# Two state Finite Automaton $M1$

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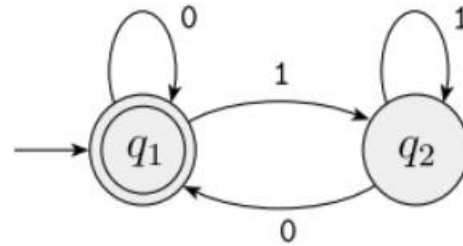


Defined by 5-tuple (  $Q$ ,  $\Sigma$ ,  $\delta$ , Start state,  $F$  )

- $Q = ?$
- $\Sigma = ?$
- $\delta = ?$
- Start state = ?
- Accept state,  $F = ?$
- $L(M1) = ?$
- Acceptability Check : 01011

# Two state Finite Automaton $M2$

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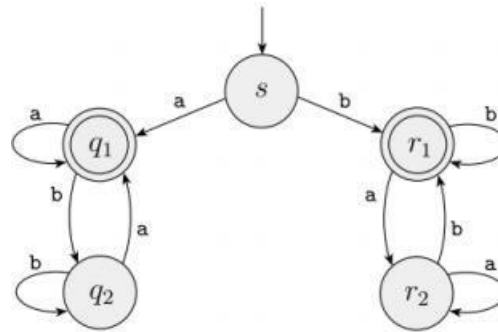


Defined by 5-tuple (  $Q$ ,  $\Sigma$ ,  $\delta$ , Start state,  $F$  )

- $Q = ?$
- $\Sigma = ?$
- $\delta = ?$
- Start state = ?
- Accept state,  $F = ?$
- $L(M2) = ?$
- Acceptability Check :

# Finite Automaton $M3$

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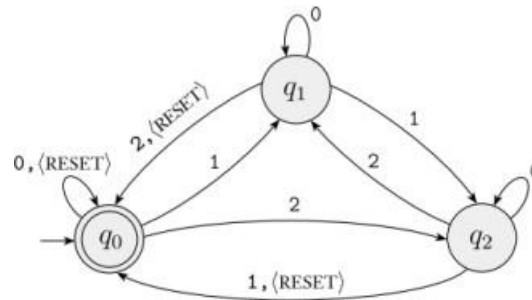


Defined by 5-tuple (  $Q$ ,  $\Sigma$ ,  $\delta$ , Start state,  $F$  )

- $Q = ?$
- $\Sigma = ?$
- $\delta = ?$
- Start state = ?
- Accept state,  $F = ?$
- $L(M3) = ?$
- Acceptability Check :

# Finite Automaton $M_4$

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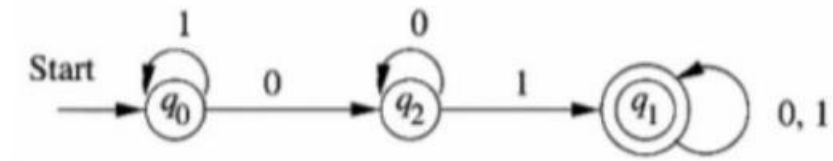


Defined by 5-tuple (  $Q$ ,  $\Sigma$ ,  $\delta$ , Start state,  $F$  )

- $Q = ?$
- $\Sigma = ?$
- $\delta = ?$
- Start state = ?
- Accept state,  $F = ?$
- $L(M_4) = ?$
- Acceptability Check :

# Finite Automaton $M5$

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Defined by 5-tuple (  $Q$ ,  $\Sigma$ ,  $\delta$ , Start state,  $F$  )

- $Q = ?$
- $\Sigma = ?$
- $\delta = ?$
- Start state = ?
- Accept state,  $F = ?$
- $L(M5) = ?$
- Acceptability Check :

# Designing Finite Automata

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Machine, *E1*

Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts any language consists of all strings with an odd number of 1s.



# Do we need to remember the entire string seen so far?

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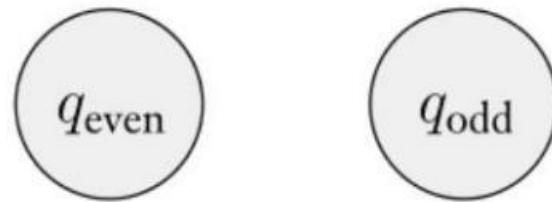
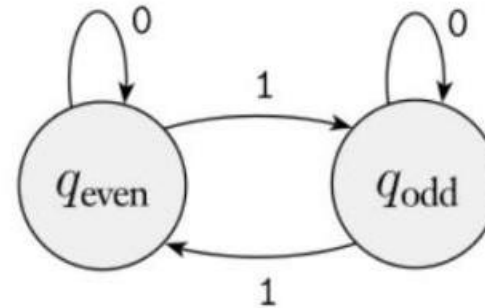


Fig: The two states  $q_{\text{even}}$  and  $q_{\text{odd}}$

Simply remember whether the number of 1s seen so far is even or odd and keep track of this information as you read new symbols. If you read a 1, flip the answer; but if you read a 0, leave the answer as is.

# Apply Transitions

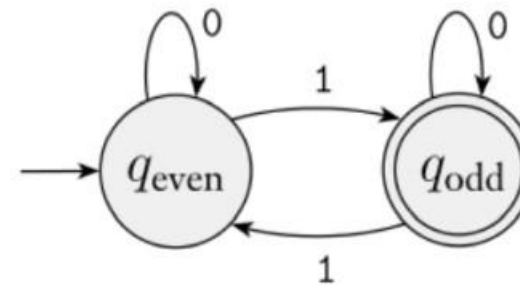
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**FIGURE 1.19**  
Transitions telling how the possibilities rearrange

# Add accept and reject states

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**FIGURE 1.20**  
Adding the start and accept states

# Machine, *E2*

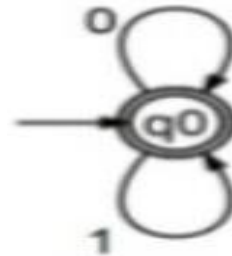
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Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts any string containing any number of 0's, 1's or empty string.

# Machine, *E2*

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Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts any string containing any number of 0's, 1's or empty string.



# Machine, *E3*

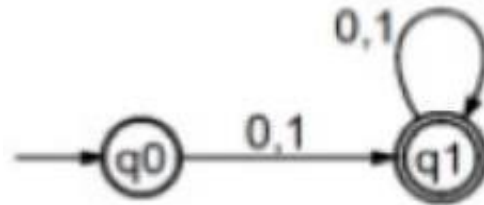
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Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts any string containing any number of 0's, 1's except empty string.

# Machine, $E3$

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Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts any string containing any number of 0's, 1's except empty string.



# Machine, *E4*

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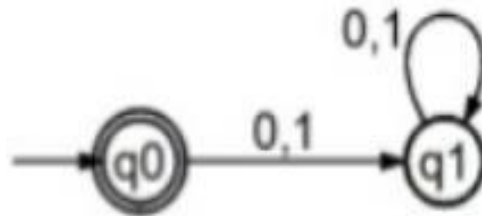
Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts only the empty string.



# Machine, *E4*

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Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts only the empty string.



# Machine, *E5*

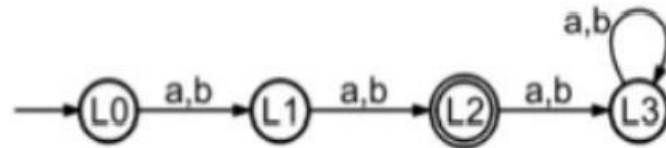
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Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have exact length 2.

# Machine, *E5*

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Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have exact length 2.



# Machine, *E6*

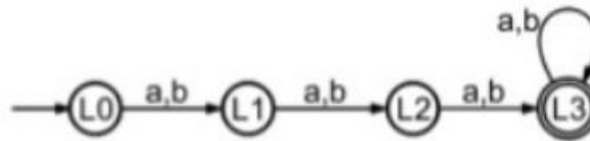
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Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have length at least 3.

# Machine, *E6*

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Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have length at least 3.



# Machine, *E7*

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Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have length at most 2.

# Machine, $E7$

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Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have length at most 2.



Machine  $E_{8.1}$  -  
What about length at least 7

Machine  $E_{8.2}$  -  
What about length at most 5

# Machine, *E8*

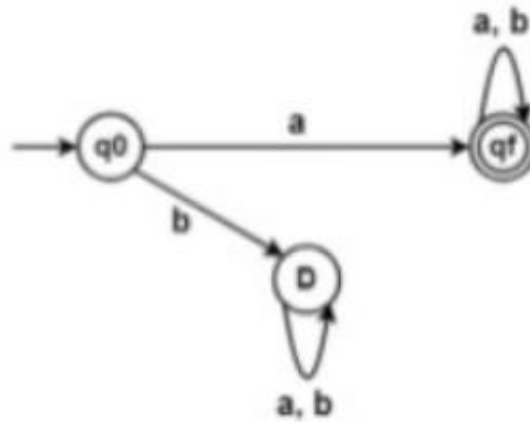
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Draw a DFA for the language accepting strings starting with 'a' over input alphabets  $\Sigma = \{ a, b \}$ .



# Machine, *E8*

Draw a DFA for the language accepting strings starting with 'a' over input alphabets  $\Sigma = \{ a, b \}$ .



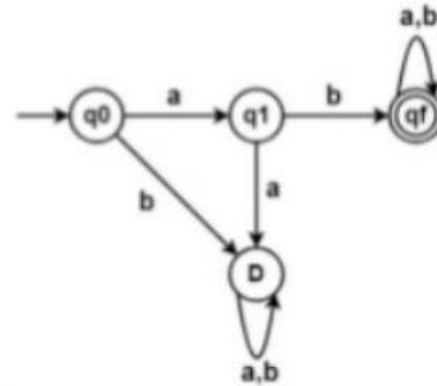
# Machine, *E9*

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Draw a DFA for the language accepting strings starting with 'ab' over input alphabets  $\Sigma = \{ a, b \}$ .

# Machine, *E9*

Draw a DFA for the language accepting strings starting with 'ab' over input alphabets  $\Sigma = \{ a, b \}$ .



# Machine, *E10*

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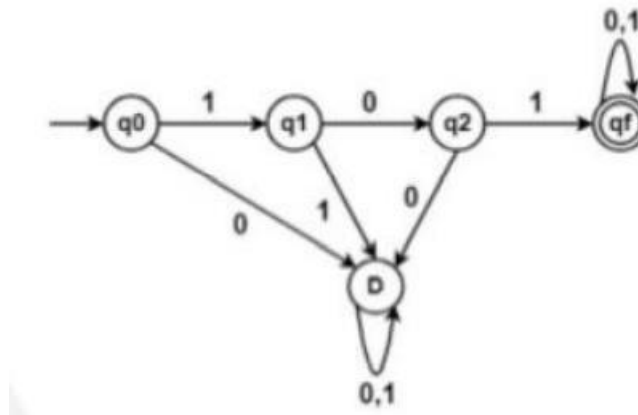
Draw a DFA for the language accepting strings starting with '101' over input alphabets

$\Sigma = \{ 0, 1 \}$ .

# Machine, *E10*

Draw a DFA for the language accepting strings starting with '101' over input alphabets

$\Sigma = \{ 0, 1 \}$ .



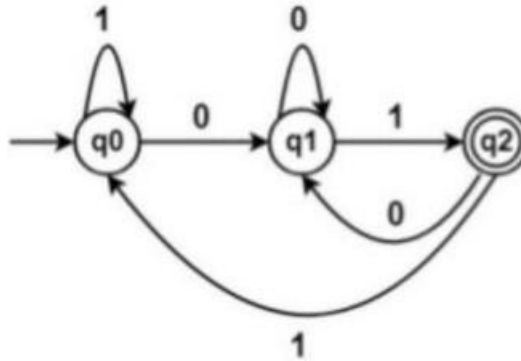
# Machine, *E11*

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Draw a DFA for the language accepting strings ending with '01' over input alphabets  $\Sigma = \{ 0, 1 \}$

# Machine, *E11*

Draw a DFA for the language accepting strings ending with '01' over input alphabets  $\Sigma = \{ 0, 1 \}$



# Machine, *E12*

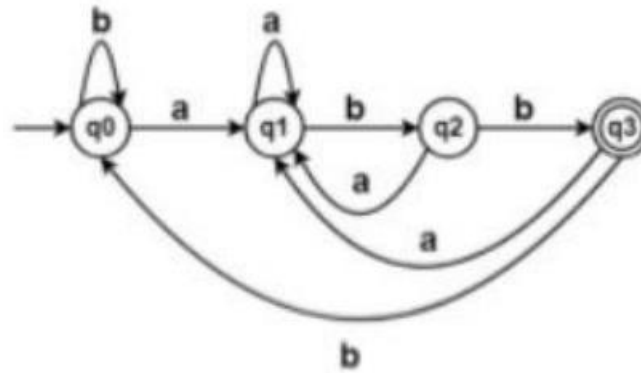
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Draw a DFA for the language accepting strings ending with 'abb' over input alphabets  $\Sigma = \{ a, b \}$



# Machine, *E12*

Draw a DFA for the language accepting strings ending with 'abb' over input alphabets  $\Sigma = \{ a, b \}$



# Machine, *E13*

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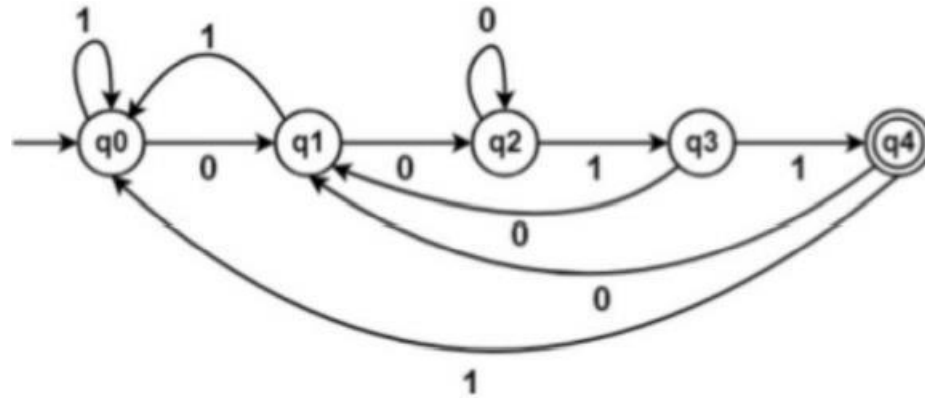
Draw a DFA for the language accepting strings ending with '0011' over input alphabets

$\Sigma = \{ 0, 1 \}$ .

# Machine, *E13*

Draw a DFA for the language accepting strings ending with '0011' over input alphabets

$\Sigma = \{ 0, 1 \}$ .



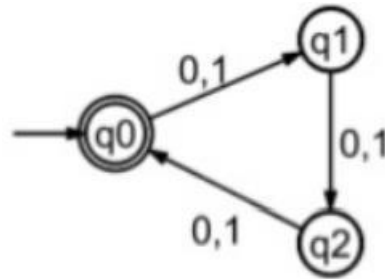
# Machine, *E14*

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Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts only those strings (including empty string) whose length is a multiple of 3.

# Machine, *E14*

Give a DFA for  $\Sigma = \{ 0, 1 \}$  that accepts only those strings (including empty string) whose length is a multiple of 3.



Machine  $E_{5.1}$  -

$L(M) = \{ w \mid \text{the length of } w \bmod 3 = 1 \}$

Machine  $E_{5.2}$  -

$L(M) = \{ w \mid \text{the length of } w \bmod 4 = 0 \}$

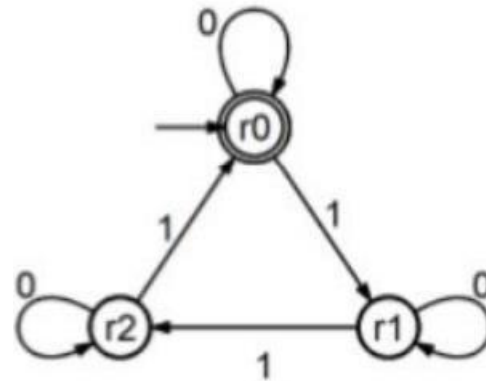
# Machine, *E15*

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Give a DFA for  $\Sigma = \{ 0, 1 \}$  and strings that contain any number of 0's and the total number of 1's is a multiple of 3.

# Machine, *E15*

Give a DFA for  $\Sigma = \{ 0, 1 \}$  and strings that contain any number of 0's and the total number of 1's is a multiple of 3.



Machine  $E_{9.1}$  -

What about strings with number of 0's is a multiple of 4 and any number of 1's

# Machine, $E_{16}$ and $E_{17}$

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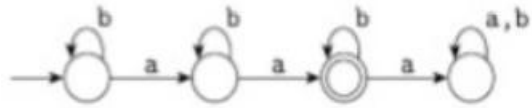
- Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have exactly 2 a's.
- Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have at least 2 b's.



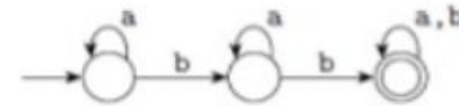
# Machine, *E16* and *E17*

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- Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have exactly 2 a's.

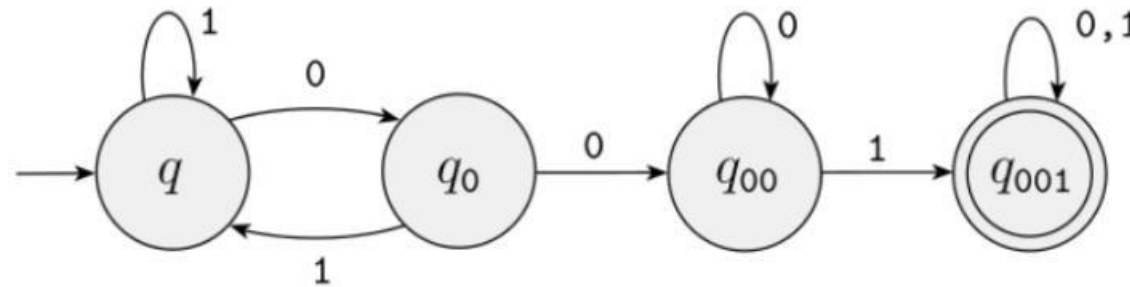


- Give a DFA for  $\Sigma = \{ a, b \}$  and strings that have at least 2 b's.



# Machine, *E18*

- Design a finite automaton to recognize the regular language of all strings that contain the string 001 as a substring.
- For example, 0010, 1001, 001, and 11111110011111 are all in the language, but 11 and 0000 are not.



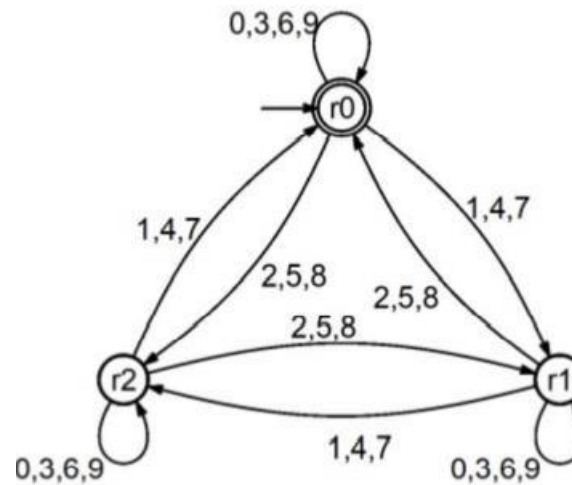
# Machine, *E19*

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Give a DFA for  $\Sigma = \{ 0, 1 \}$  and only accepts binary strings those are a multiple of 3.

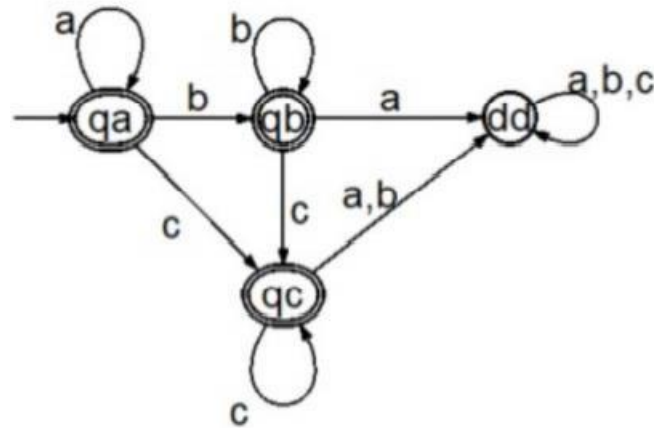
# Machine, *E20*

Give a DFA for  $\Sigma = \{ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9 \}$  that accepts strings whose decimal equivalent is a multiple of 3



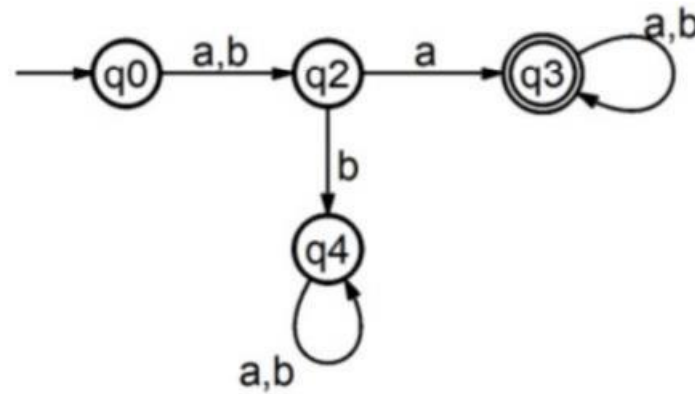
# Machine, *E21*

Give a DFA for  $\Sigma = \{ a, b, c \}$  and recognizes strings having patterns  $= \{ a^n b^m c^l \mid n, m, l \geq 0 \}$



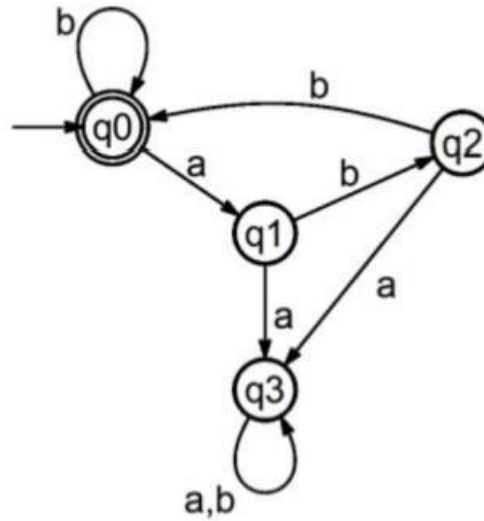
# Machine, *E22*

Give a DFA for  $\Sigma = \{ a, b \}$  and the second symbol from the left side is 'a'.



# Machine, *E23*

Give a DFA for  $\Sigma = \{ a, b \}$  and contains strings where each 'a' is followed by 'bb'



# Try yourself =>

## Machine, *E24.1* and Machine, *E24.2*

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- Give a DFA for  $\Sigma = \{ a, b \}$  and recognizes strings having patterns =  $\{a^n b^m \mid n+m = \text{even} \}$ .
- Give a DFA for  $\Sigma = \{ a, b \}$  and recognizes strings having patterns =  $\{a^n b^m \mid n+m = \text{odd} \}$ .



Thanks