Data Structures and Algorithms CS 213- Lecture 1

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Broad Topics

- Introduction to data structures,
 - Abstract data types, analysis of algorithms.
- Creation and manipulation of data structures: arrays, matrices, sparse matrices, lists, stacks, queues, trees, heaps, hash tables, balanced trees, tries, graphs.
- Algorithms for sorting and searching,
- Depth-first and breadth-first search,
- Paradigms:
 - The greedy method.
 - Divide-and-conquer.
 - Dynamic programming.
 - Backtracking.
 - Branch-and-bound.

Text Books

- S. Sahni, Data Structures,
 Algorithms and Applications in C++,2nd edition, Universities
 Press,2005
- T. Cormen, C. Leiserson, R. Rivest, C. Stein, Introduction to Algorithms, 2nd edition, Prentice-Hall India, 2001.

Organization

- Lectures Tuesday, Friday: 330-500PM
- Quizes 3-- 30%
- Mid Term 30%
- Final Exam 40%
- Attendance Necessary
- Academic Honesty Policy:

http://www1.iitb.ac.in/newacadhome/ rules.jsp

TAs

BS Radhika Amit Goel Sundaram Gupta

- MTP-1 --6 TAs??
- BTP 3 TAs??

Tentative Agenda

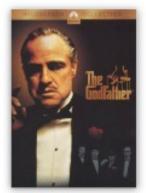
- Introduction
 - Role of Data structures
 - Concerns of program Construction
 - Analysis of Algorithms: Notations
 - Complexity Analysis forms
 - Recurrence Relations/Recursive Programs
- Merge Sort, asymptotic analysis and performance measurement of programs.
- Data representation methods and linear lists
- Arrays, matrices, sparse matrices
- Stacks.
- Queues.
- Hashing
- LZW compression.
- Binary trees, AVL Trees

- Priority queues.
- Tournament trees.
- Search trees.
- Graphs.
- The greedy method.
- Divide-and-conquer.
- Dynamic programming.
- Backtracking.
- Branch-and-bound.

What is it about?

- Data structures: Representation and manipulation of data.
- Programs manipulate data.
- Programs as data
- How do we manipulate Data?

Digital Data



Movies



Music





Photos

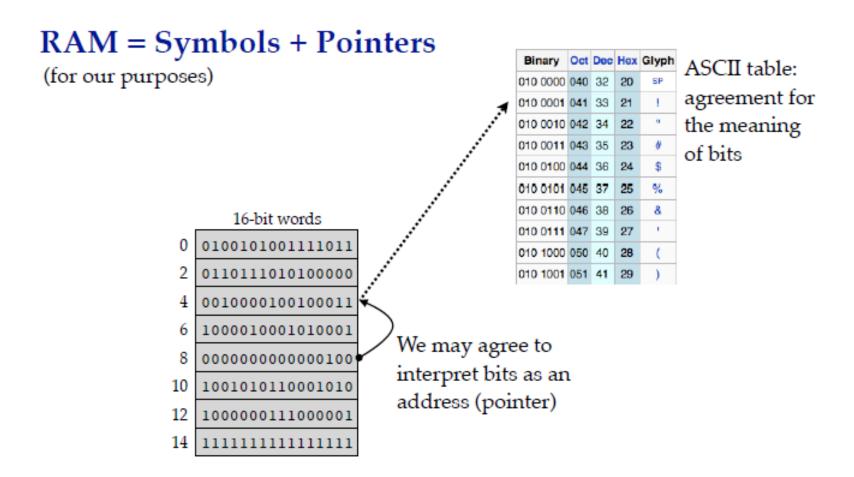


Maps

DNA

gatottttta tttaaaogat ototttatta gatotottat taggatoatg atoototgtg gataagtgat tattoaoatg goagatoata taattaagga ggatogtttg ttgtgagtga ooggtgatog tattgogtat aagotgggat otaaatggoa tgttatgoao agtoaotogg oagaatoaag gttgttatgt ggatatotao tggttttaoo otgotttaa goatagttat aoaoattogt togogogato tttgagotaa ttagagtaaa ttaatooaat otttgaoooa

00101010010101010101001001001010100000100100100100100....



Physically, RAM is a random accessible array of bits

=> We can store and manipulate arbitrary symbols (like letters) and associations between them.

Digital Data Must Be ...

• Encoded (e.g. 01001001 <->

Arranged

- Stored in an orderly way in memory / disk

Accessed

- Insert new data
- Remove old data
- Find data matching some condition

The focus of this class

Processed

Algorithms: shortest path, minimum cut, FFT, ...

Data Structure Example Applications

- 1. How does Google quickly find web pages that contain a search term?
- 2. What's the fastest way to broadcast a message to a network of computers?
- 3. How can a subsequence of DNA be quickly found within the genome?
- 4. How does your operating system track which memory (disk or RAM) is free?
- 5. In the game Half-Life, how can the computer determine which parts of the scene are visible?

What is a Data Structure Anyway?

- It's an agreement about:
 - how to store a collection of objects in memory,
 - what operations we can perform on that data,
 - the algorithms for those operations, and
 - how time and space efficient those algorithms are.
- Ex. vector in C++:
 - Stores objects sequentially in memory
 - Can access, change, insert or delete objects
 - Algorithms for insert & delete will shift items as needed
 - Space: O(n), Access/change = O(1), Insert/delete = O(n)

Abstract Data Types (ADT)

```
class Dictionary {
   Dictionary();
   void insert(int x, int y);
   void delete(int x);
   ...
}
```

```
insert()
delete()

find_min()

find()

int main() {
    D = new Dictionary()
    D.insert(3,10);
    cout << D.find(3);
</pre>
```

- Data storage & operations encapsulated by an ADT.
- ADT specifies permitted operations as well as time and space guarantees.
- User unconcerned with how it's implemented (but we are concerned with implementation in this class).
- ADT is a concept or convention:
 - not something that directly appears in your code
 - programming language may provide support for communicating ADT to users (e.g. classes in Java & C++)

Dictionary ADT

- Most basic and most useful ADT:
 - insert(key, value)
 - delete(key, value)
 - value = find(key)
- Many languages have it built in:

 Insert, delete, find each either O(log n) [C++] or expected constant [perl, python]

C++ STL

 STL (Standard Template Library) is a powerful set of C++ template classes to provides general-purpose templatized classes and functions that implement many popular and commonly used algorithms and data structures like vectors, lists, queues, and stacks.

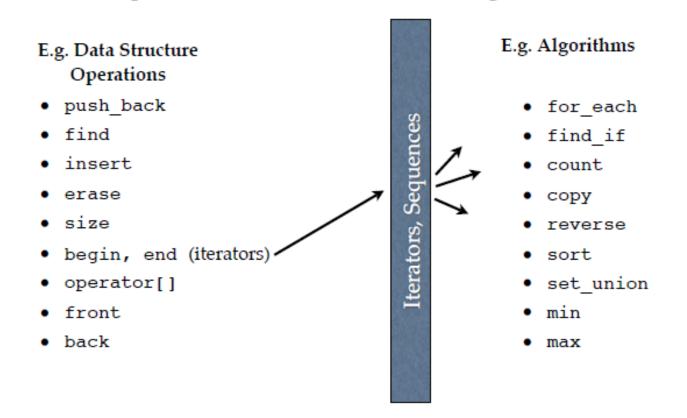
C++ STL

- Data structures = "containers"
- Interface specifies both operations & time guarantees

Container	Element Access	Insert / Delete	Iterator Patterns
vector	const	O(n)	Random
list	O(n)	const	Bidirectional
stack	const (limited)	O(n)	Front
queue	const (limited)	O(n)	Front, Back
deque	const	O(n), const @ ends	Random
map	O(log n)	O(log n)	Bidirectional
set	O(log n)	O(log n)	Bidirectional
string	const	O(n)	Bidirectional
array	const	O(n)	Random
valarray	const	O(n)	Random
bitset	const	O(n)	Random

Some STL Operations

- Select operations to be orthogonal: they don't significantly duplicate each other's functionality.
- Choose operations to be useful building blocks.



Consider Google Maps

You want to store data about cities (location, elevation, population)...



What kind of operations should your data structure(s) support?

Operations to support these Operations

- Finding addresses on map?
 - Lookup city by name...
- Mobile iPhone user?
 - Find nearest point to me...
- Car GPS system?
 - Calculate shortest-path between cities...
 - Show cities within a given window...
- Political revolution?
 - Insert, delete, rename cities



Data Organizing Principles

Ordering:

- Put keys into some order so that we know something about where each key is are relative to the other keys.
- Phone books are easier to search because they are alphabetized.

Linking:

- Add pointers to each record so that we can find related records quickly.
- E.g. The index in the back of book provides links from words to the pages on which they appear.

Partitioning:

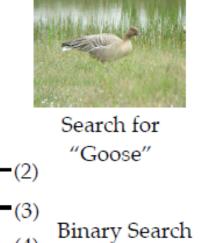
- Divide the records into 2 or more groups, each group sharing a particular property.
- E.g. Multi-volume encyclopedias (Aa-Be, W-Z)
- E.g. Folders on your hard drive

Ordering

Duck, Woodpecker, Robin, Cardinal, Eagle,	10 89 55 3 32 18 50 89 102 43
Duck, Woodpecker,	50
Cardinal,	102
Chicken, Pigeon,	7 201 57
Swan, Loon, Turkey,	213 99
Albatross, Ptarmigan, Finch,	0 22 38
Bluejay, Heron,	24 70
Egret, Goose,	88 67

Sequential Search – O(n)

Albatross, Bluejay,	0 24
Cardinal,	102
Chicken,	7
Duck,	18 ←
Eagle,	43
Egret,	88◀
Finch,	38
Goose,	67 ←
Grouse,	89
Heron,	70 ←
Loon,	213
Partridge,	32
Pelican,	3
Pheasant,	10
Pigeon,	201
Ptarmigan,	22
Quail,	55
Robin,	89
Swan,	57
Turkey,	99
Woodpecker,	50



O(log n)

Every step discards half the remaining entries:

-(1)

$$n/2^k = 1$$
$$2^k = n$$
$$k = \log n$$

Paradigms

- Algorithms + Data Structures = Programs
 - Imperative Programs

- Logics + Control = Programs
 - Logic Programs Declarative
 Specifications
 - Example Finding a transitive closure of a

Broad Coverage

 Algorithm design methods needed to develop programs that do the data manipulation.

 Study of data structures and algorithms: Crux of Computer Science.

Programs

Two Concerns

- Correctness
- Efficiency

Writing an convincing program

- Three arrays
- F[i], G[j], H[k] monotonically decreasing
- Given that there exists a i,j, k such that F[i] = G[j] = H[k], write program computing the points of common intersection

A Simple Example

if $x \ge y \rightarrow m := x$ $[] y \ge x \rightarrow m := y$ fi.

Permutation

```
q\hat{1}, q2, q3, q4 := Q1, Q2, Q3, Q4;

do q1 > q2 \rightarrow q1, q2 := q2, q1

\boxed{\qquad q2 > q3 \rightarrow q2, q3 := q3, q2}

\boxed{\qquad q3 > q4 \rightarrow q3, q4 := q4, q3}

od.
```

Another Example

```
k := 0; j := 1;

\operatorname{do} j \neq n \to \operatorname{if} f(j) \leq f(k) \to j := j + 1

[]f(j) \geq f(k) \to k := j; j := j + 1

\operatorname{fi}

od.
```

GCD

$$x := X; y := Y;$$

 $do x > y \rightarrow x := x - y$
 $y > x \rightarrow y := y - x$
 $od.$

GCD (Classical versions)

```
x := X; y := Y;
                                           (version A)
while x \neq y do if x > y then x := x - y
                        else y := y - x fi od
and
x := X; y := Y;
                                           (version B)
while x \neq y do while x > y do x := x - y od;
               while y > x do y := y - x od
            od.
```

What about computation of LCM?

Ackerman Function

```
    A(m,n) = n+1 if m=0
    = A (m-1,1), m > 0 and n=0
    = A (m-1, A(m, n-1)), m > 0 n > 0
```

Tower of Hanoi