

# CS 252 - Lab 03 Tele-Communication System Design

## Report

### GROUP DETAILS

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### TITLE

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#### Tele-communication System Design

The principle we used to transmit messages is transferring the signal in form of high and low voltages between two raspberrypis.

### DESIGN

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The User gives input to the Rasp-berry pi through a computer connected to its command line through secure shell.To indicate which bits to be flipped the user adds letter 'e' just afer that bit.We used hamming encoding/decoding. The reciever checks the message with help of redundant bits and corrects the 1 bit error,and detects the 2bit error asking for retransmission(NACK).we also added start flag and end flag.

### IMPLEMENTATION

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#### Sender end

The user gives input to the raspberry pi command line. The code computes redundant bits using **hamming encoding**.The final message consists of redundant bits and payload and start/end flags.These are sent in form of high and low voltages through GPIO pins.After it sends message it starts listening for ACK,If it recieves a NACK(0) it retransmits the message and process is repeated else sending one message is complete

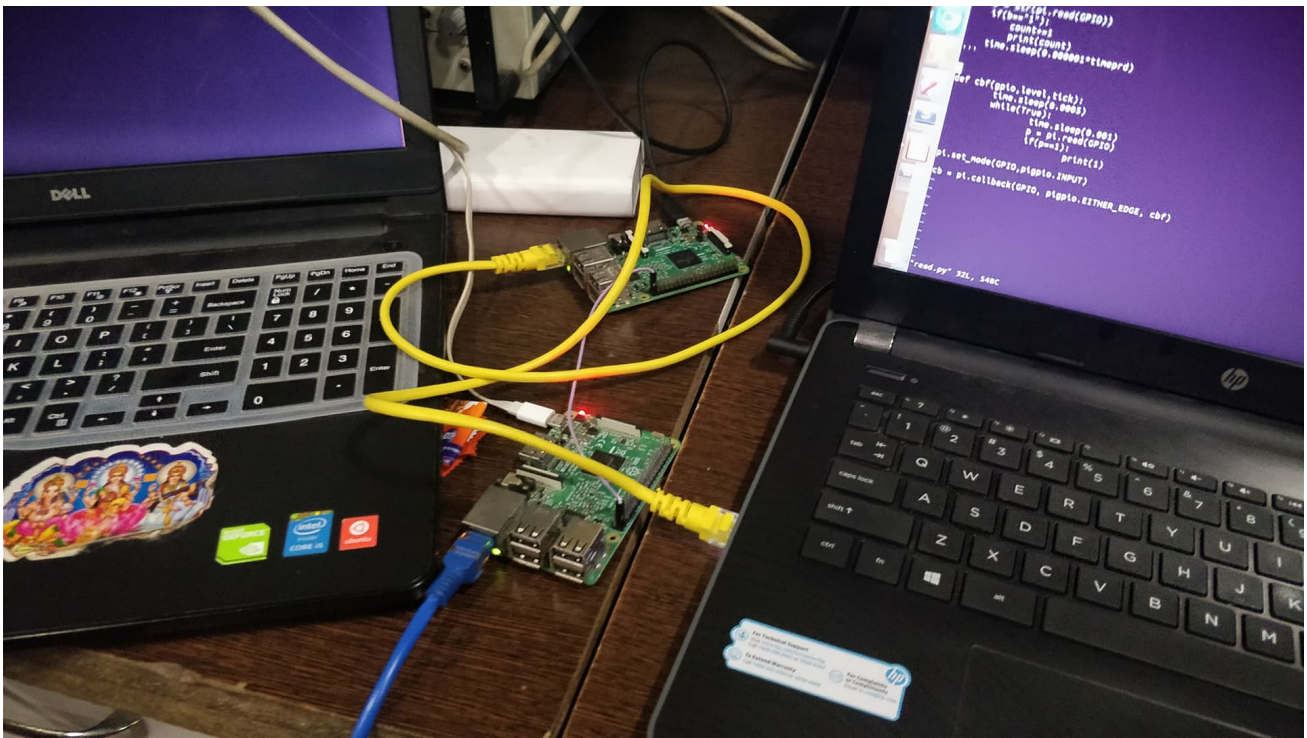


Figure 1: Connection of two raspberry pis

## Receiver end

After the message is received, it reads the bits and error check occurs through hamming decoding. If there is an error of 2 bits NACK is sent otherwise errors of 1 bit are automatically corrected. All the signals of gpio pins are handled using pigpio library

## Hamming code

Hamming codes are a class of binary codes. If the length of the code is  $n = 2^r - 1$  then the length of the message will be  $2^r - r - 1$  and the number of parity bits will be  $r$ . The position of the parity bits are  $2^0, 2^1, 2^3, \dots, 2^{r-1}$ . The method of finding the value of parity bit at the position  $2^k$  is as follows :

- 1) Considers all the bits in the respective positions which contain '1' in their  $k^{\text{th}}$  least significant bit (the positions are in the binary form).
- 2) Add all the bits in those positions modulo 2 (i.e. add as normal decimal numbers and divide by 2 at the end).
- 3) The remainder is the required value of the parity bit.

The method of recovering the bits from the parity bits is as follows:

- 1) Collect all the bits with 1 as their position's 1st least significant bit and add them modulo 2.
2. Continue this for 2nd least significant, 3rd, ...
- 2) Now put all the bytes next to each other from left to right in the decreasing order of their significance.
- 3) Convert this value into decimal number and this is the required position of the error bit.

If the code length is between  $2^r - 1$  and  $2^{r+1} - 1$ , then there will be  $r$  parity bits at  $1, 2, 4, \dots, 2^{r-1}$ . 1 bit can be detected and corrected using these many parity bits. But if an additional bit is added at the end, then 2 bits can be detected and 1 bit can be corrected. The additional parity bit is the even parity of the previously encoded string of bits.

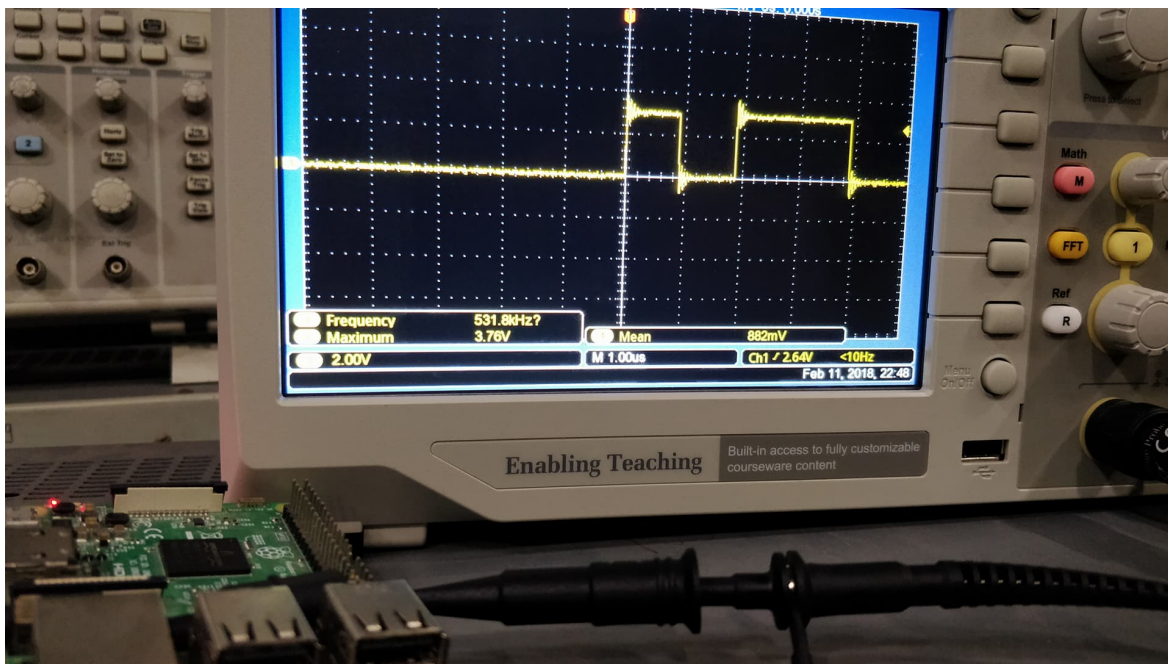


Figure 2: Checking input generation through oscilloscope

## REFERENCES

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Hamming code - [GeeksforGeeks](#)  
RaspberryPi functioning - [RaspberryPi official site](#)  
pigpio package - [Pigpio library](#)