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Year: - E2-SEM1

Semester: - 1

Online Batch: - Batch 1

Offline Section: - Batch 1

Subject Name: - OLD

MIDTEST-1

Date: - 17-07-2021

message: - 1110

Parity: - even

Hamming code formula: -

$$2^k = n + k + 1$$

$$\text{here } n=4 \Rightarrow 2^k = 4 + k + 1$$

$$k=3$$

Here are 3 parity bits are added for this message

| | | | | | | |
|----------------|----------------|----------------|----------------|----------------|----------------|----------------|
| 111 | 110 | 101 | 100 | 011 | 010 | 001 |
| D ₇ | D ₆ | D ₅ | P ₄ | D ₃ | P ₂ | P ₁ |
| 1 | 1 | 1 | ? | 0 | ? | ? |

⇒ For finding P₁ collect the bits which has first

bit as 1 → (i.e., 1, 3, 5, 7)

P₁ D₃ P₅ D₇ even parity

$$P_1 \ 0 \ 1 \ 1 \Rightarrow P_1 = 0$$

⇒ For finding P₂ collect the bits which has second bit as 1.

P₂ D₃ D₆ D₇

$$P_2 \ 0 \ 1 \ 1$$

Even no. of ones are there so, P₂ = 0

for finding P_4 , collect the bits which has third bit as 1

$P_4 \ D_5 \ D_6 \ D_7$

$P_4 \ 1 \ 1 \ 1$

odd no. of ones are there so, $P_4 = 1$

original Message 1110

Transmitted message

$D_7 \ D_6 \ D_5 \ P_4 \ D_3 \ P_2 \ P_1$

1 1 1 1 0 0 0

\therefore Transmitted Message is 1111000

1)
b

$$(101101110001.00101)_2 = (\quad)_{10}$$

$$\begin{array}{cccccccccccccccc} 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 1 & . & 0 & 0 & 1 & 0 & 1 \\ 2^{11} & 2^{10} & 2^9 & 2^8 & 2^7 & 2^6 & 2^5 & 2^4 & 2^3 & 2^2 & 2^1 & 2^0 & 2^{-1} & 2^{-2} & 2^{-3} & 2^{-4} & 2^{-5} \end{array}$$

$$\Rightarrow (1 \times 2^{11} + 0 \times 2^{10} + 1 \times 2^9 + 1 \times 2^8 + 0 \times 2^7 + 1 \times 2^6 + 1 \times 2^5 + 1 \times 2^4 + 0 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0) \cdot (0 \times 2^{-1} + 0 \times 2^{-2} + 1 \times 2^{-3} + 0 \times 2^{-4} + 1 \times 2^{-5})$$

$$\Rightarrow (2929.15625)_{10}$$

i) $(292)_{16} = ()_2$

$$\begin{array}{ccc} 2 & 9 & 2 \\ 0010 & 1001 & 0010 \end{array}$$

$$\Rightarrow (292)_{16} = (001010010010)_2$$

iii) $(163.789)_{10} = ()_8$

$$\begin{array}{r} 8 \overline{) 163} \\ 8 \overline{) 20} - 3 \\ \quad \overline{) 2} - 4 \end{array}$$

$$\therefore (163)_{10} = 243$$

| | <u>Integer part</u> | <u>Decimal part</u> |
|------------------|---------------------|---------------------|
| 0.789×8 | 6 | 0.312 |
| 0.312×8 | 2 | 0.496 |
| 0.496×8 | 3 | 0.968 |
| 0.968×8 | 7 | 0.744 |
| 0.744×8 | 5 | 0.952 |
| 0.952×8 | 7 | 0.616 |
| 0.616×8 | 4 | 0.928 |
| 0.928×8 | 7 | 0.424 |
| 0.424×8 | 3 | 0.392 |
| 0.392×8 | 3 | 0.136 |

$$\therefore (163.789)_{10} = (243.6237574733)_8$$

2)

a)

Ans

$$F(x, y, z) = x'y'z' + x'y'z + xy'z = y'(x' + z)$$

$$\text{L.H.S} \Rightarrow x'y'z' + x'y'z + xy'z$$

$$\Rightarrow x'y'(z' + z) + xy'z$$

$$\Rightarrow x'y'(1) + xy'z \quad (\because A + \bar{A} = 1)$$

$$\Rightarrow x'y' + xy'z$$

$$\Rightarrow y'(x' + xz)$$

$$\Rightarrow y'((\bar{x} + x) \cdot (x' + z)) \quad (\because A + BC = (A + B) \cdot (A + C))$$

$$\Rightarrow y'(1(x' + z))$$

$$\Rightarrow y'(x' + z)$$

$$\Rightarrow \text{R.H.S}$$

$$\therefore \text{L.H.S} = \text{R.H.S}$$

2b)

Ans

$$F(a, b, c) = (a' + b'), (a' + b), (a + b') = a' \cdot b'$$

$$\text{L.H.S} \Rightarrow (a' + b'), (a' + b), (a + b')$$

$$\Rightarrow [(a' + b'), (a' + b)] \cdot (a + b') \quad (\because (A + B) \cdot (C + D) = AC + BC + AD + BD)$$

$$\Rightarrow (a'a' + a'b + b'a' + b'b) \cdot (a + b') \quad (\because \bar{a} \cdot a = 0 \text{ domination law})$$

$$\Rightarrow (a'a'a + a'a'b' + a'ba + a'bb' + b'a'a +$$

$$b'a'b' + b'ba + b'bb)$$

$$\Rightarrow 0 + a'b' + 0 + 0 + 0 + a'b' + 0 + 0$$

$$\Rightarrow a'b' + a'b = a'b' \quad (\text{as } a+a = a \text{ idempotent law})$$

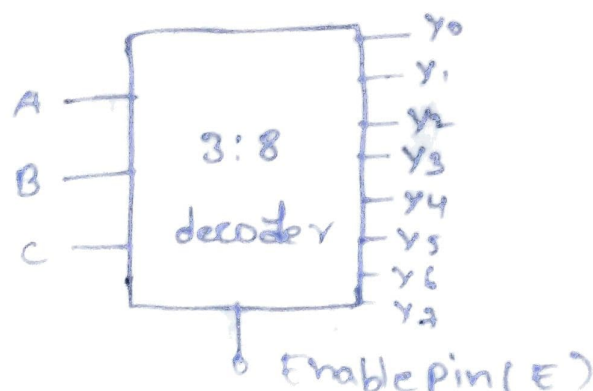
$$\therefore (a' + b')(a + b) = a'b'$$

3) Ans \Rightarrow 3 to 8 decoder

Decoder - It is combinational logic circuit. It is used to decode the data, which is already encoded

\Rightarrow It is exact opposite to Encoder

Block-Diagram



\Rightarrow To Represent 8 inputs 3 bits are enough so

encoder encodes these up to 3 bits

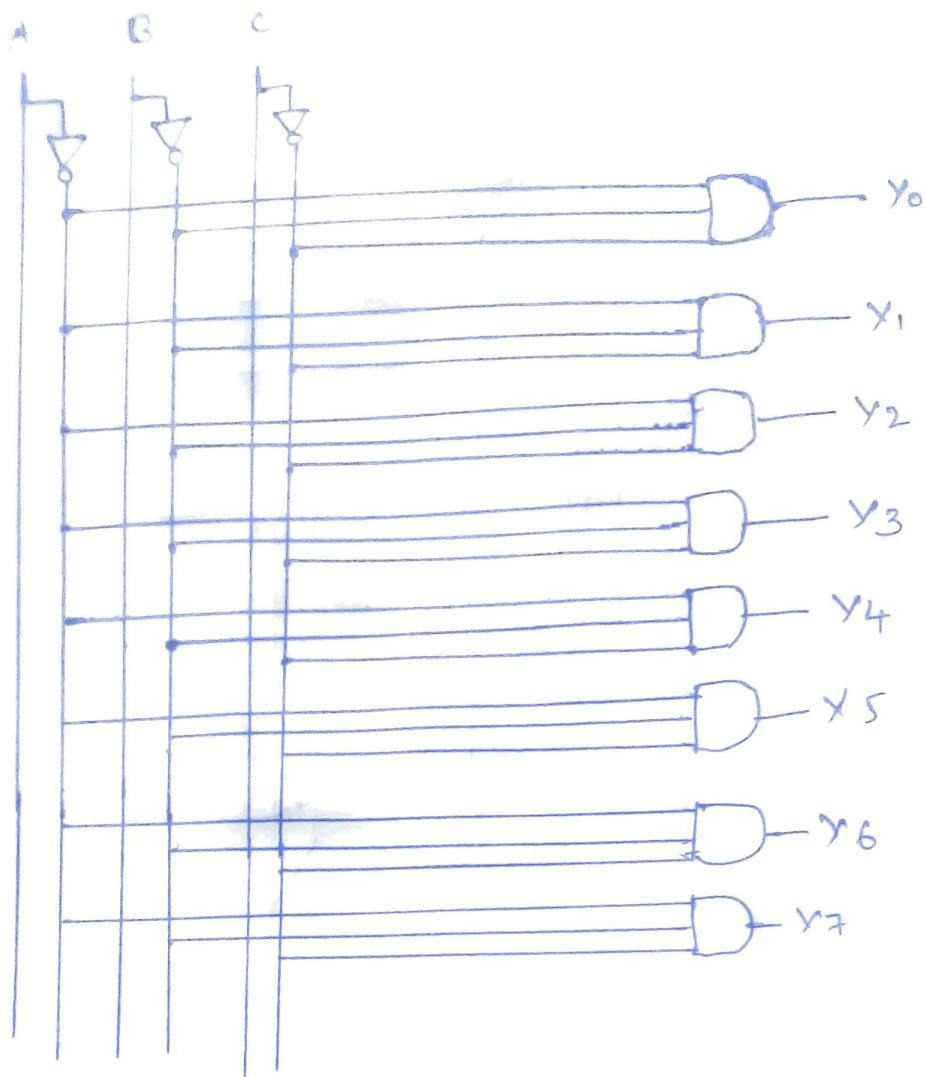
Decoder decodes the encoded data into 8 inputs

Truth table

| E | A | B | C | Y_0 | Y_1 | Y_2 | Y_3 | Y_4 | Y_5 | Y_6 | Y_7 |
|---|---|---|---|-------|-------|-------|-------|-------|-------|-------|-------|
| 0 | x | x | x | x | x | x | x | x | x | x | x |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |

⇒ If Enable (E) is 0 then decoder won't be working

⇒ To make decoder active we have to give Enable as 1.



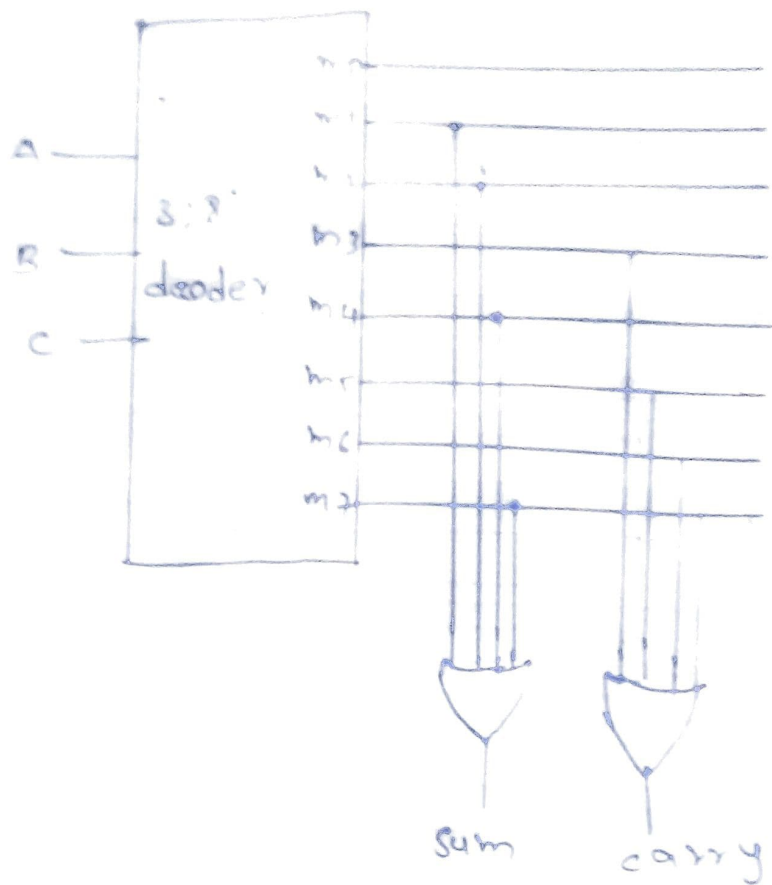
Full Adder -

| A | B | C | Sum | Carry |
|---|---|---|-----|-------|
| 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 |

$$\text{Sum} = \sum m(1, 2, 4, 7)$$

$$\text{Carry} = \sum m(3, 5, 6, 7)$$

Implementation of Full Adder using 3:8 decoder

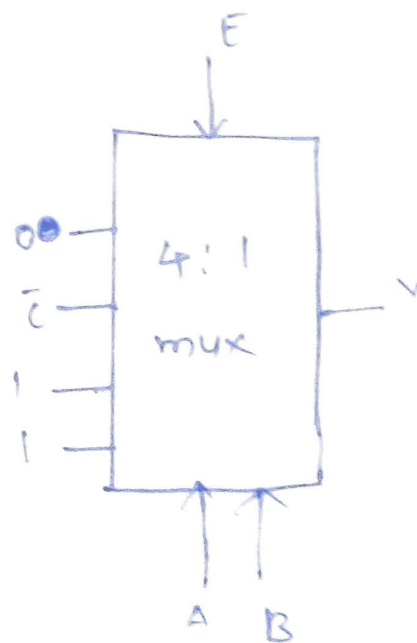


8)

Ans)

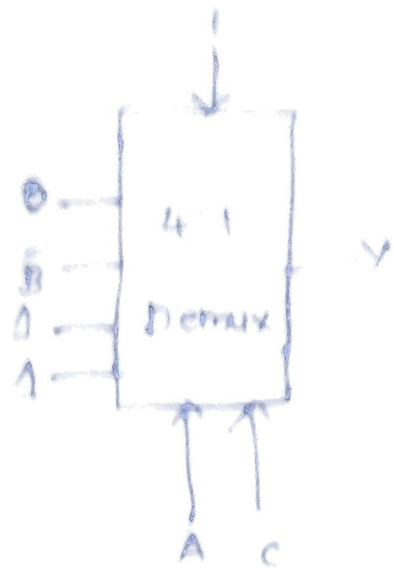
i) AB as select line

| | D_0 | D_1 | D_2 | D_3 |
|-----------|-------|-----------|-------|-------|
| \bar{C} | 0 | ① | ② | ③ |
| C | 4 | 5 | ⑥ | ⑦ |
| | 0 | \bar{C} | 1 | 1 |



ii) AC as select line

| | D_0 | D_1 | D_2 | D_3 |
|-----------|-------|-----------|-------|-------|
| \bar{A} | 0 | 1 | 5 | 3 |
| A | 4 | 2 | 6 | 7 |
| | 0 | \bar{A} | 1 | 1 |



iii) BC as select line

| | D_0 | D_1 | D_2 | D_3 |
|-----------|-------|-----------|-------|-------|
| \bar{A} | 0 | 1 | 2 | 3 |
| A | 4 | 5 | 6 | 7 |
| | 0 | \bar{A} | 1 | 1 |

