

# How can we speak math?

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## Abstract

It is likely that most people can communicate mathematics to a computer more effectively (rapidly and accurately) by speaking than they can by using a stylus on a computer tablet. This may seem surprising, but is our speculation based on trying various alternative input methods. An even better setup may be to speak and simultaneously use pointing or handwriting. Unfortunately, building a properly functioning prototype using this concept is difficult. Yet a successful implementation of such a “multimodal” combination should allow the computer to reinforce correct recognition while identifying and perhaps repairing “unimodal” errors. In some cases speaking *may be more convenient than typing, even for rapid typists*: many mathematical symbols are missing from the keyboard but can be easily spoken and recognized. Even without venturing into Greek, or alternative fonts, just handwriting or even typing a number, say “fifty million” may be slower and more error-prone than speaking.

Pursuing the goal of effectively speaking and recognizing small pieces of mathematics, led to a study of how hard it would be to speak arbitrarily long sections of mathematics, including nested complex expressions.

We first describe programs for the inverse problem: computer generation of mathematical speech. This requires that we address some speaking conventions to overcome the unfortunately ambiguous and inconsistent common usages of mathematics.

Then we consider tools and guidelines to make it more plausible for humans to speak full mathematical formulas unambiguously so they can be recognized by a computer using a speech recognizer program.

We describe our prototype programs which do somewhat less than we propose, but are effective in that speech can either be used alone, or used to fill in boxes (superscripts, etc.) or larger pieces. Speech can also be used for choosing alternatives from plausible symbols resulting from uncertain recognition from handwriting (or speech). We believe the principal barriers to engineering a more complete program can be overcome, though a driving application may be essential for refining prototypes into useful programs. This paper is not intended to be the last word on the subject, but simply exposes problems and approaches relevant to the task. Demonstrations of partial implementations are available as Window (XP) programs.

## 1 Introduction

Handwriting mathematics seems natural because it is what we have been taught in school. We find it natural to view mathematics in typeset form because that too is commonplace and familiar. If asked, most professional users of mathematics will opine that speaking mathematics is difficult, since the “hard parts” come to mind. In fact users of math routinely speak *small pieces* quite comfortably. Often a paper introducing new written notation specifies how it should be pronounced! These small bits can often easily be combined to medium-sized sections. We do not hesitate to vocalize “the quadratic formula”<sup>1</sup>. Given that speech input to computers is becoming more common as it is better supported by technical advances, the question arises: when is it useful to speak mathematics into the computer? One argument is that if we could do so, persons with disabilities in writing or typing should be able to more easily communicate mathematics to a computer, just as they might dictate business correspondence. Yet even for non-disabled, there may

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<sup>1</sup>Even though most people who nominally know it are likely to speak it in a manner that is arguably wrong or ambiguous, given inadequate “brackets”.

be advantages for speech in some circumstances. We contend that speech can be used in three ways: as a primary method for conveying mathematics, a supportive auxiliary method in a “multimodal” context, or an error-correction command language.

The reverse operation, namely a computer speaking mathematics and the human listening, has more of a successful history. So-called Text-to-Speech (TTS) but adapted for math, is, so far as we can tell, not widely adopted except as an assistive technology for sight-disabled. The two notable successes are AsTeR [16] and Design Sciences’ MathPlayer [4]. We first discuss this material as background and then proceed to our main results where humans speak aloud and the computer listens to mathematical discourse.

## 2 Computers speaking math

The program AsTeR [16] is an excellent prototype for speaking mathematics; indeed it seems quite worthy of use for the reading of  $\text{\TeX}$  mathematics to visually disabled persons<sup>2</sup> Nevertheless, there is a problem with this approach:  $\text{\TeX}$  does not provide an encoding of the semantics for the mathematical material, since  $\text{\TeX}$  is only a presentation view of mathematics supported by  $\text{\TeX}$ . Semantics must be derived from some (external) context or encoded in extra data attached to the encoding. Thus  $f^{-1}$  might be  $f$  to the power  $-1$  or it might be  $f$  inverse, or even, in the case  $\sin^{-1}$ , the function named “arcsine”. There may even be homonyms (“sign” and “sin”). If the speech is generated from a computer algebra system, or encoded in a semantic description (even MathML, a computer algebra system form), there is a better chance of getting it right. In fact, Design Science, [www.dessci.com](http://www.dessci.com) has a “speak expression” option that allows Internet Explorer to read math aloud from a MathML expression if the (free) MathPlayer plug-in is available. Its effectiveness depends on a browser/operating system capability for text-to-speech. Given the underlying support, it then feeds locutions like “begin fraction a+b over c+d end fraction.” It seems to us plausible that one might do somewhat better by directly speaking from a computer algebra system (CAS) rather than through a browser. In the CAS case, the system could contain more context including line labeling schemes, aliasing of symbols to names, or abbreviations (e.g. let  $r = \sqrt{x^2 + y^2}$  in an expression). It could also make reasonable and consistent choices as for  $x^{-1}$  vs.  $1/x$ . It might even describe expressions in a preliminary “outline” to prepare the listener. For example “a fraction with a long numerator of 25 summands and a denominator which is the product of 5 terms.” Instructing the computer to provide more details could be done by keyboard, handwriting, or speaking. For example, the computer might advise, “To hear the terms in the numerator one at a time, say *next*. ...” This segmented approach has been explored in the Universal Speech Interface project<sup>3</sup>.

A back-and-forth interaction between a remote CAS and a local browser speaking MathML via MathPlayer could probably simulate this situation fairly well, so a browser cannot be discounted entirely.

An application other than the sight-disabled motivation, and one that strikes us as more compelling for advanced mathematics is proofreading (perhaps of  $\text{\TeX}$ ). A (sighted, hearing) human need not glance between two written versions to see if they are the same. Certainly for the unreliable handwriting input method, a math-to-speech program could be useful as a proofreading or interactive-feedback assistant for input methods.

Just as a side note; humans are fairly sensitive to oddities in speech. Typical computer-generated speech is generally easily identified as unnatural. This does not mean it is necessarily difficult to understand or distressing to listen to, at least for technical material. We are not reading poetry.

### 2.1 Speaking on the Internet

Stepping back from math specifically, how hard is speech production? Given the state of the art today, it is possible, even easy, to have a web browser speak (in one of various available voices of your choosing) the XML encoding of a speech utterance. It is possible to encode speed, pitch, volume, and other voice characteristics. How adaptable is this to mathematics? We have experimented with this, and have written a program suite providing the translation of algebraic expressions given as Lisp prefix data into words. For example `(* r s t)` would be spoken as “r times s times t.” More specifically our Lisp-to-speech-XML program would

<sup>2</sup>The author is blind; Aster was the name of his seeing-eye dog.

<sup>3</sup><http://www.cs.cmu.edu/usi>

produce this underlying encoding for  $r \cdot s \cdot t$ .

"<p><spell>r</spell> times <spell>s</spell> times <spell>t</spell> </p>" Similarly,  $f(x, y)$  would be

"<p><spell>f</spell> of <spell>x</spell> and <spell>y</spell> </p>".

Not all the nuances of AsTeR may be available, but the XML encoding in fact provides considerable opportunities for speech variation: changes in volume, speed and pitch. We have not seen an additional feature which might be cute: using stereo, proceeding from the left speaker to the right as the expression is read aloud.

A curiosity that we did not anticipate in our initial design is the extent to which most listeners and speakers leave out critical information, even when they think they are speaking unambiguously, and how overbearing a complete and unambiguous rendering sounds when we produce it from our own program. This become apparent when the program, naturally set up to be unambiguous in its utterances, is given common middlingly-complex expressions.

The well-known quadratic formula can be written as a Lisp prefix expression as  $(/ (\pm (- b) (^ (- (^ b 2) (* 4 a c)) 1/2)) (* 2 a))$  where  $\pm$  means  $\pm$ . This can be read in a variety of ways. Here we remove the `<spell>` pieces, as well as a change in pitch for the denominator and other minor items in order to make the text more perspicuous. In school you might get full credit if you recite it as `minus b plus or minus square root of b squared minus 4 a c divided by 2 a`.

Without prior knowledge of this formula how could you know if the  $4ac$  or even the  $2a$  belongs within the square root? You don't from this reading. Is the  $-b$  in the numerator or outside the fraction? Again you don't know. In fact, is the  $a$  in the denominator, or is it a multiplier for the whole previous expression?

Our punctilious program insists on bracketing, by inserting "the quantity" and "end" around components so it can provide non-ambiguous renderings<sup>4</sup> But for this formula our program needs to put in three sets of brackets, making it *seem* excessively pedantic. Judicious omission of bracketing on output seems advantageous, and so our original default speaking program does not always insert brackets. Instead there is a explicit insertion of tags required for enunciating brackets. As an example  $(* (+ a b) c)$  could be spoken identically with  $(+ a (* b c))$ , which is clearly unsatisfactory. Our fix is to use a bracket constructor which is spoken by the computer (and to keep the listener on guard). The example would be  $(* (\text{bracket } (+ a b)) c)$ , and would be pronounced "The quantity a plus b end times c." The commercial product MathPlayer speaks the quadratic formula by talking about fractions, end-square-roots, and yet leaves out operators like "times". Here is an ML version of the quadratic, taken from a Design Science demonstration page:

```
<p id="MPEqn">
<m:math>
<m:mstyle displaystyle="true">
<m:semantics>
<m:mrow>
<m:mi>x</m:mi><m:mo>=</m:mo><m:mfrac>
<m:mrow>
<m:mo>±</m:mo><m:mi>b</m:mi><m:mo></m:mo><m:msqrt>
<m:mrow>
<m:msup>
<m:mi>b</m:mi>
<m:mn>2</m:mn> </m:msup>
<m:mo>±</m:mo><m:mn>4</m:mn><m:mi>a</m:mi><m:mi>c</m:mi> </m:mrow>
</m:msqrt>
</m:mrow>
<m:mrow>
<m:mn>2</m:mn><m:mi>a</m:mi> </m:mrow>
</m:mfrac>
```

<sup>4</sup>The exact phrasing is under constant reappraisal: e.g. inserting "begin square-root" and "end square-root" may be better.

```

</m:mrow>
<m:annotation encoding="MathType-MTEF">
MathType@MTEF@5@5@+ .... truncated...
</m:semantics>
</m:mstyle>
</m:math></p>
<p id="MPEqnAlt" style="display:none">[MathML Equation -- requires MathPlayer]</p>

```

We have truncated some material above: it is a compact encoding of the speech version.

It may be feasible to disambiguate expressions by the use of prosody – intonation, timing, volume, etc. We can speak “French bread and cheese” in different ways to distinguish the case that both the bread and the cheese are French, and the case that the bread is French but the cheese is of unknown origin. We could propose to pronounce “three x plus y” by analogy, distinguishing  $3(x + y)$  or  $3x + y$ , depending on whether there is a detectable pause after the “x”.

## 2.2 Non-speech approaches to natural math

This is necessarily a brief review. On the *output side*, in recent years computers have essentially replaced older typesetting technology for mathematical printing. Software can now support the whole workflow from the original creation and composition, perhaps with the aid of a computer algebra system, through interpretation by some typesetting program, to the point of printing on paper or display on a browser. Most readers of this paper will be aware of such editors (using keyboard and mouse) and printers or screen displays (using raster graphics).

On the *input side*, most mathematics programs are heavily keyboard-dependent, with perhaps mouse/menu assists. Among current computer algebra systems, Maple version 10 (2006) allows limited handwriting input of single symbols.

Yet looking back at research programs, since at least 1965 programs [1] there have been demonstrations of software which serve as intermediaries for the conversion of (hand)written material into typeset material. More recently it has become plausible to actually make use of such programs on the much-more powerful computers of today.

Today’s demonstration programs [20, 14, 3, 13] show that while it is fairly easy to recognize a *subset* of simple math symbols and expressions as *usually* written by hand, there remain substantial barriers to usefulness. While a short demonstration may show remarkable effectiveness, these program work best when used by their authors on pre-tested examples. It is expected that novices attempting more complex tasks will suffer from a higher error rate. This is a consequence of understandable difficulties. Trouble distinguishing many pairs: (p vs P, 0 vs O, 5 vs S, 1 vs l vs i vs — vs [ vs ] etc), means that some demonstration programs may work only by requiring special gestures, or taking steps such as simply excluding the letters S, l, and O from the vocabulary. Other confusions are possible with positioning or stroke identification. Thus  $1 < 2$  could easily be written so as to be confused with K2.

We suggest the following experiment might illustrate some of the difficulties, easily performed by a college student or teacher.

Walk in to a mathematics or physics classroom at the end of the lecture and see if you can read all the mathematics on the blackboard.

You probably can’t understand it all. Expecting a computer to understand it, devoid of mathematical or physical context is unrealistic. Additionally, a computer post-processing of a blackboard has another handicap compared to the student in the classroom. The computer does not have the benefit of the lecturer’s simultaneous speech while writing on the board, nor the sequence of writings and erasures. Of course it also does not have the opportunity to ask clarifying questions.

Other methods for input and editing of math via templates, menus, keyboarding, and other non-handwriting forms are surveyed by Kajler and Soiffer [9]. A skilled user can generally do quite well with such systems, but systems can be frustrating to the novice. In some cases they can also be frustrating to the expert who requires close adherence to some format unanticipated by the system designers. Our proposals

generally would complement existing systems like  $\text{\TeX}$  and  $\text{\TeX}$ macs which is an interactive system inspired by  $\text{\TeX}$  ; some later editing might be needed for small corrections if there is a need to precisely control the typesetting.

## 2.3 Speaking mathematics

One reviewer of an earlier version of this paper claimed that “the use of dual input (speech and pen) is not much different than pen and keyboard or pen and palets[sic]”. This reviewer missed two points:

- One cannot type on a keyboard simultaneously with using a mouse. One must lose time moving a hand to the mouse and then later re-establishing a position over the keyboard. Picking up a stylus is harder in this respect than grabbing a mouse, requiring more complicated motions. Once picked up, a stylus has quite a different feel from a mouse, a feel which is much superior for writing. We could solve this problem by learning to type with our feet, point with our nose or eyes, growing a third hand, or using a keyboard with some mouse-substitute as on laptop computers. *Or speaking*. Fortunately many people can use a keyboard or pen and *simultaneously speak* without any genetic engineering or special training.
- Speaking “bold italic capital gamma” is probably faster than typing or writing. Most writers no longer know the markup conventions historically used by (human) typesetters (Years ago, authors were told to write a wavy line under symbols intended to be typeset in bold, and a single underline for italic).

Speaking mathematics could be used for a spectrum of uses from educational testing to search in digital systems containing mathematics. Suitably instrumented, it could be used as a testbed for evaluation of interaction via speech with web-based services.

While this paper emphasizes speech, research has, more generally, been looking at “multimodal” techniques using speech in combination with handwriting [7, 6].

Before addressing, in the next section, an apparently simple question: “How do we (intuitively) speak math?” we briefly review in the section some useful speech processing ideas for the uninitiated.

### 2.3.1 Brief digression on speech processing

There are substantial research efforts on speech and computing, a number of competing commercial products for speech recognition, and a WWW standard for speech markup (text to speech or TTS). From a relatively naive standpoint, but one which we think is adequate here, the speech issues seem to be separable into

- Output aids for the visually impaired. The audience may be computer users (programmers, too) who are unable to see text as routinely displayed by a computer. Text-to-Speech (TTS) makes it possible for a computer to “read aloud” to a blind person, or to speak to a person who has no other display, which includes a sighted person using a telephone. A truly useful audio interface for a structured domain like mathematics or a graphical display will require rather more elaborate design [16, 4] than just reading a text basically because there is no standard translation of math to text suitable for speaking.
- Input aids to the keyboard-typing impaired. The user may suffer from some temporary or permanent disability. Automatic Speech Recognition (ASR) makes it possible for a user to “speak” words and phrases, constituting dictation of content (perhaps intermixed with commands such as “new paragraph” or “file save”) to the computer. Generally the user is able to see a display for feedback, but not always. A user of such a system might be at a telephone speaking commands to a computer. (If a handset is separate from a keypad, simple numeric input from a sighted person might best be provided through the keypad. Alphabetic input is trickier, as is input from a one-piece cellular phone. Not too tricky for the millions of people who use text messaging via phone, though.)
- “Multimodal” assistance, for example for the task of correction (proofreading) of material that may have been entered into the computer by some error-prone method. The first method might be document image analysis, handwriting, or speech. Both TTS and ASR may be used. Proofreading data entry of tables of numbers by having them read back by the computer seems quite straightforward with today’s

technology. Even reading math formulas out loud to see if they have been typed (or typeset) could be an application.

There are notable simplifications possible. Consider a system trained on a single voice (easier) or one which must work with all speakers (harder). Consider a system to recognize a small vocabulary and grammar (say digits, or telephone numbers, or dates) versus a larger language such as “business letter English” (harder). The least accurate recognition would be expected of a system for arbitrary users on unconstrained vocabulary.

### 2.3.2 The trivial non-solutions

One solution for “speaking mathematics” that immediately presents itself as unambiguous is to merely spell expressions as though you were typing them—character by character—on a single line. All the disambiguation must be done prior to spelling. In this way the problem has been reduced to that of the previously “solved” problem, namely the parsing of a programming language that is typed into a computer, and all that is needed is a mapping of sounds to keyboard elements. If the encoding language is  $\text{\TeX}$ , then the appearance of almost any mathematical notation can be provided, on almost any computer system, thanks to the continuing work on maintaining  $\text{\TeX}$ . If the programming language is the painfully-verbose MathML, simulating a keyboard by voice would be very time-consuming. Even with the much more concise  $\text{\TeX}$ , entering  $\beta$  would require saying something like “dollar backslash b e t a dollar” or once you realize how close certain sounds are (a, eight) or (b, d, p) or (s, f), you might use a “military alphabet” for spelling. (In practice a military<sup>5</sup> spelling option uses more phonemes but is nearly error-free. It is not too difficult to learn.) Thus for a higher accuracy, you might learn to say “dollar backslash bravo echo tango able dollar”. Of course it would be easier to say “beta”!

(We note in passing that the usual programming language notations, such as Fortran, while adequate for specifying “arithmetic” are grossly inadequate notationally for serious math, and we cannot seriously consider “speaking Fortran” as a substitute for math<sup>6</sup>. We also note once again that the interpretation of  $\text{\TeX}$  as math can be ambiguous, but at least it is as good as mathematicians usually see; a spoken version will not necessarily be semantically unambiguous either!)

## 3 Developing an intuitive speech model

First we discuss speaking numbers, which is surprisingly tricky. Then non-numeric symbolism follows.

### 3.1 Reading numbers aloud

If we wish to enter content consisting of applied mathematics we need to be able to read numbers. It may surprise you that the reading (and hence the speaking) of numbers is rife with special cases and ambiguity. At the risk of belaboring the trivial yet non-obvious, we include the following observations.

The TTS (Text To Speech) program from Microsoft which we use has some interesting features for reading numbers aloud. We review its behavior not only for amusement, but for illustrating these issues. After all, if we hope to have the computer listen to us speak numbers, perhaps we should attempt to understand the rules that TTS uses for pronouncing numbers (starting from text) as guidelines.

The following examples (from Microsoft speech SDK 5.1) suggest that sometimes this provides a plausible guideline. Microsoft does not provide access to the *complete rule-set* for TTS, and so we cannot be definite about how TTS speak every number given to it as `ascii` text.

Here are some examples. We’ve marked with a (\*) those that seem open to debate.

- 123 is one hundred twenty-three.
- 123.123 is one hundred twenty-three point one two three.

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<sup>5</sup>NATO uses Alpha Bravo Charlie Delta Echo Foxtrot Golf Hotel India Juliet Kilo Lima Mike November Oscar Papa Quebec Romeo Sierra Tango Uniform Victor Whiskey Xray Yankee Zulu.

<sup>6</sup>Of course, speaking Fortran qua Fortran, or using speech as source input in any programming language is a possibility, with many of its own difficulties not necessarily related to math.

- 1,000.00 is one thousand.(\*)
- 1,000.000 is one thousand point zero zero zero.
- 3.1415929 is three point one four one five nine two six.
- 3.14.15929 is three point fourteen point fifteen thousand nine hundred twenty-six. (\*)
- 3.14.1592 is March fourteenth, fifteen ninety-two. (Note the use of ordinal 14th).(\*) The program knows that the nearby “number” 3.32.1592 is an invalid date, and thus spells it out. It does not know that September has only 30 days, much less the rules about leap years. In fact it is not possible to speak this into the standard dictation grammar, which will produce a sequence of two numbers, 3.14 and 0.1592. But see the related date *fractions* below.
- 1/10 is one tenth.
- 9/10 is nine tenths.
- 10/11 is ten over eleven.
- 14/100 is fourteen hundredths.
- 14/10000 is fourteen over ten thousand.
- 14/100000 is fourteen slash ten oh oh oh oh. (\*)
- 14/1000000 is fourteen slash one oh oh oh oh oh oh. (\*)
- 14/1000000000000000 is fourteen slash one zero zero ... zero.
- 14/ 1000000000000000 is fourteen slash ten trillion.
- 3/100 and 300 sound almost the same: “three hundredths” versus “three hundred.”
- 2-2 as well as 2-2-2 is two to/two two.
- 1-3, as well as 1-2-3, is one to/two three.
- 1-2-9 is one two nine, but 1-2-10 is January second, ten.
- 40/500 and 45/100 are indistinguishable. (The second can only be spoken as 45 slash 100 or 45 over 100. forty-five hundredths yields 40/500.)
- 3/14/1592 which might appear to be (3/14) divided by 1592, is not. It is March 14, 1592.
- 0.0 is zero point zero.
- 0.00 is just zero.
- 1,500,000 is 1 point 5 million.

Integers up to “9999999999999999” (999 trillion and change) are spoken, but above that are spelled out digit by digit. There are different rules for integers appearing in denominators.

Numbers that do not have commas set out “correctly” are spelled out. Thus 5,10.0 is five comma ten point zero.

Floating point numbers such as “5.00d0” are handled as separate components, namely “5.00” or five, and “d0” (dee zero). -1/2 is dash one slash two.

Who would have thought it was so complicated? Of course just reading off the digits and punctuation would be unambiguous, but who wants to speak like a cheap robot<sup>7</sup>.

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<sup>7</sup>Mr. Data on Star Trek isn’t programmed to speak contractions!

### 3.2 How humans *should* speak numbers to computers

The TTS rules are too complicated. Would a subset of the rules be adequate? Which utterances are acceptable? Do you want to use numbers like “three and a quarter” or “one point five million.” Our advice is to use easily-parsed “full” natural numbers including properly indicated steps like “one hundred twenty three thousand”. An alternative is a string of single digits. Full numbers may be combined with decimal points (“.” pronounced “point”) or for fractions, the virgule (“/” pronounced “slash” or “over”). We also permit “oh” for zero. How important is it to recognize words like “million”? The purely digit-list prescription is easy to program but saying a number like 3 million, saying all digits, is painful: it has an excessive number of zeros to pronounce and recognize accurately.

There are other problems if numbers occur adjacent without intervening punctuation. This can happen with single digits perhaps more often: “The single-digit primes are 2, 3, 5, and 7” does not mean “The single-digit primes are 235 and 7.” Thus the commas must be enunciated, or the speaker must force the recognizer to accept the phrase in pieces. “US paper currency includes fifty, one-hundred and five-hundred dollar denominations” could be read as “5100 and 500 dollar.”

We tried several approaches.

- A pattern-matching heuristic program we have written is perfectly happy with numbers constructed like “one hundred twenty-three thousand four hundred fifty-six point seven eight” for 123,456.78. We recommend “one slash two” for  $1/2$ , since generalizations of fractions are tricky. Being written in Common Lisp, our program has essentially no limits on the number of digits in a number, though it tends to reduce  $3/6$  to  $1/2$ .

- For most uses, we expect that the Microsoft published `cmnrules` grammar<sup>8</sup> for various kinds of numbers including natural numbers, fractions, floating-point, could be used. Much to our relief this can be included rather painlessly in a speech recognition program by specifying (in an SASDK/ SALT application that can, for example, be run with a browser plug-in), a `listen` tag.

```
<item> <ruleref uri="cmnrules.grxml#number"/>
<tag> $. _value = $$._value </tag> </item>
```

It would be even better for our use if the SASDK allowed for multiple return values for a speech recognition task (that is, with ranked alternates); at the moment this is only possible for the default Microsoft grammar, a default suitable for typical business applications, but which is unsuitable for mathematics. We understand that this limitation may be lifted in the VISTA version of Windows, which we have avoided for reasons not directly related to speech.

- The principal defect in `cmnrules` from our exact mathematics perspective is that it is limited to numbers less than  $10^{15}$  and fractions are converted to decimal numbers of limited precision. This is an artifact of using the arithmetic in the underlying J++ scripting language which is the default (and at the time of writing of this paper, sole) programming technology in the Microsoft grammar implementation of the W3C recommendations for XML speech grammar. We have constructed a modification of the grammar to maintain exact ratios for numbers like  $1/3$ , where numerator and denominator can only be represented exactly by strings. This is passed on to Lisp for further evaluation. Thus the string “six quintillion plus one” is parsed to “(+ (\* 6 (expt 10 18)) 1)” which is exactly evaluable in Lisp. (There is a disappointment at a different level in the grammar XML processing, in that true context-free grammars are not acceptable.)
- A third possibility, also easily implemented by reference to `cmnrules` is to use lists of digits for numbers. As illustrated in examples above, this is occasionally in conflict with the other common usage rules, but could easily be used instead of, or in preference to, the more general usage. In fact the digit-list convention *is* used in conjunction with other parts of the grammar for decimal fractions. Consider “seventeen hundred point oh four five”. To the right of the point we speak in digit lists.

Who would have anticipated such complications for numbers? It is much easier to write a demonstration program that works only for single digits, or integers, but would that be sufficiently useful?

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<sup>8</sup>We found, reported and corrected two bugs in this. June, 2004.



### 3.3 Non-numeric tokens

In our experiments to date, starting with a short list, dissimilar words can be recognized very accurately. Given a larger word list, especially if context (e.g. grammar) does not play a role, the recognition can be more error-prone. Given that our list of mathematical notation includes the presence of easily-confused short words, we have a choice.

- Satisfaction with relative poor initial accuracy, relying on rapid correction.
- Resolution of ambiguity based on context. Given our formula context, we prefer “eight equals two times four” to the identical phonemes in “ate equals to times for”. Unfortunately “Pick a number from one to ten” and “Pick a number from 1, 2, 10.” are rather close. Sometimes the context may be quite small “Capital a” is a plausible sequence, while “Capital 8” is less. If the recognizer is supplied with a grammar for complete formula utterances, or a grammar for phrases, this can be helpful context.
- Removing some ambiguity at the source: rename or provide synonyms for all letters via a military alphabet, as suggested earlier. We choose names one that do not conflict with other math tokens such as Greek letters. Thus (adam or able, ..., dog or david, ...) rather than (alpha, ..., delta, ...).

Other token considerations:

The well-used spoken tokens include not only letters of the Roman alphabet (optionally modified with “bold,” “Roman,” “Italic,” “capital”, “upper-case”, etc), but other alphabets as well. Symbols taken from sources include the T<sub>E</sub>X typesetting repertoire, computer algebra systems such as Mathematica, and selected parts of Unicode. Even among the common names, there are ambiguities. Consider the homonyms “sign” and “sin” which are equally plausible in many contexts.

Words for spaces are handy as well, such as “quadspace”.

Typically these tokens can be separated into operators and operands, but we *cannot depend on such classifications* for rigid parsing.

It is also quite likely that macro-expressions defined verbally will be useful for the serious speaker. Thus “let big Adam equal script capital bold adam sub Greek nu” allows an abbreviation<sup>9</sup>. Clearly this could be made as elaborate as any macro language, although here we propose simple constant non-parametric substitutions.

### 3.4 Caution on complete forms

Imagine how annoying it would be if, as you were typing at a computer keyboard, every one of your pauses were treated as an end-of-sentence marker and the computer immediately made an observation that your sentence was incomplete, or if it appeared to be complete, it immediately whisked it off and processed it. We must refrain from insisting that math be spoken all in one breath, or else  $x + y + z$  would be impossible:  $x + y$ , being complete, would be gobbled up first. We can signal explicitly by a mouse click<sup>10</sup> or alternatively, the computer will just wait, and proceed after a short pause when you are presumed to be finished speaking *for the moment*. In such circumstances it cannot be too authoritarian about preventing what you say next to be appended to, or somehow modify, the previous utterance<sup>11</sup>.

### 3.5 Expressions

In this section we describe variations for speaking a prototypical expression that would seem to be at first glance non-linear in appearance. We omit the “OK” needed at the end of each expression:

$$\frac{a + b}{c + d}.$$

This can be linearized in various ways. In T<sub>E</sub>X it is spelled out as `$\frac{a+b}{c+d}$...`

<sup>9</sup>Using arbitrary words, e.g. “let *doodah* equal...” requires that “doodah” be in our speech grammar’s wordlist.

<sup>10</sup>We can signal the end of a phrase by a word marker such as “OK”, but the program will wait for a pause following the “OK”.

<sup>11</sup>(What’s your favorite color? Blue. No, yellow; <http://www.sacred-texts.com/neu/mphg/mphg.htm>)

Or spelling it out we could say, “dollar, backslash eff arr ay see open brace, ay plus bee close ...”. In a military alphabet ... foxtrot romeo adam charlie ....

We assume here that “close” is adequate to match the previous still-open bracket, and we can save quite a few syllables if we do not have to say “close parenthesis” or “right parenthesis”.

In future examples in this paper we won’t use spelling, even though it may be inevitable for peculiar words.

Instead of spelling  $\TeX$  we can spell a linearized form  $(a+b)/(c+d)$ , which is shorter, unambiguous, but still uncomfortable. Instead of a dollar sign we use “begin math” and “end math”. Instead of targeting  $\TeX$  we are targeting a typical programming language (perhaps a computer algebra system, or a “natural” math input system [17, 15]. )

```
begin math ( a + b ) / ( c + d ) end math.
```

This requires saying open/close four times. To preview our proposal in this regard, in this paper we suggest that the expression above be spoken this way:

```
begin math
a+b quantity
over
quantity c+d
end math.
```

or perhaps

```
begin math
adam + bravo quantity
over
quantity charlie + david
OK
```

(We will refrain from using the military alphabet subsequently because it is a distraction; however, in our limited experiments, an otherwise irksome level of erroneous recognition of some letters can be effectively remedied this way.)

Grouping based on the embedded key words *quantity*, *over* and *end* can be done by some simple transformations on the stream of tokens. We start by implicitly enclosing every begin/end math expression with a default  $(\dots ( \text{ and } \dots )$ . The word “quantity” immediately *after* an *operator* (defined below), can be changed to the insertion of a “(”. “Quantity” *before* an operator, is equivalent to “)”. If the speaker says “quantity” between two operands (which are presumably going to be multiplied together by a “silent times”) then we propose the same result as “quantity times quantity”. This may not be the speaker’s intention, so some extra feedback or warning may be advisable.

The extra prefix “(” and suffix “)” are appended only as needed to balance the brackets. Operators are not necessarily unique. That is, “over” and “divided-by” are synonyms. We include

- infix such as plus, times, over, slash, divided-by, raised-to, to-the-power, space, quadspace
- prefix such as sum, product, *function* of (e.g. sine of), bold, italic, roman, upper, lower, big, capital, script, Greek
- suffix such as factorial, squared, cubed, prime, double prime,
- overhead, which in  $\TeX$  constitute prefix such as hat, bar. In common math speech, these would generally be voiced as suffix operations.  $\hat{x}$  in TeX is  `$\hat{x}$`  but probably pronounced **x hat**.
- matchfix such as left/right square brackets, left/right angle brackets, open, close (paren, bracket, square bracket) These matchfix operators can come in many sizes like **big** or **big big**, and presumably must be matched in size.

There are large tables of additional operators in The  $\text{\TeX}$ book, and similar references, each attempting to be encyclopedic; see also the menus in Mathematica.

Typical *operands* are essentially *everything else*, including syntactic components like symbols, numbers, and (recursively) subexpressions.

Given these rules, our spoken expression is transformed to text as

(a+b)  
/  
(c+d)

### 3.6 Math on a line

It seems at first that any math expression that fits on a single line without up/down excursions would not be problematical, since it has an “obvious” order in which to read characters<sup>12</sup>. It seems that difficulties could only occur if the speaker leaves out characters necessary for grouping, or declines to pronounce the brackets. Unfortunately, *leaving out such characters is entirely conventional*, even when the result is ambiguous, as shown by later examples.

Simple Examples:

Display	Spoken
$ab \sin x$	a b sine of x
$a + \frac{b}{c} + d$	a + b over c + d
$\frac{a+b}{c} + d$	a + b quantity over c + d
$a + \frac{b}{c+d}$	a + b over quantity c + d

This next set of examples is insufficient to tell us how to deal with extra cases that require groupings “in the middle”.

Most of what we have said up to this point does not get much of a rise out of most readers who may have been only mildly surprised by some of the difficulties encountered. Not having tried to program speech recognizers for math, this is reasonable all around.

This next proposal is more controversial: We believe we may have to add only *one* additional linguistic marker, *all*, or alternatively, *end* or *close*. In fact, all three terms, *all*, *end*, and *close* are synonymous [to the computer]. This would work with the term “quantity” previously used. Let us argue in favor of this.

The term “all” or its alternatives essentially *jumps out a level*.

Display	Spoken
$a + \frac{b}{c+d} + e$	a + b over quantity c + d all + e
$a + \frac{b}{c+d} \times e$	a + b over quantity c + d all times e

We can also use “all” without “quantity”

Display	Spoken
$(a + b)/(c + d)$	a + b all over quantity c + d
$(a + b)/c + d$	a + b all over c + d

Consider this:  $a + \frac{b}{c+d} \times \frac{e}{f} + g$ . We could try grouping this using prosody, inserting pauses: a + pause b over quantity c + d pause times e over f pause + g. Raman’s AsTeR program [16] can use prosody, changing pitch upward for superscripts for *output*, but human speakers, and the programs listening to them may not be so capable of such small distinctions. And sometimes one would need several pauses at the same place. Nevertheless, in combination with a geometric handwriting interface and feedback, perhaps this could work.

<sup>12</sup>Actually a linear sequence is possibly ambiguous in a larger sense of conveying mathematics.  $1/2\pi$  sometimes means  $\pi/2$  and sometimes  $1/(2 \times \pi)$ . But this is not a speech problem.

Display	Spoken
$a + \frac{b}{c+d} \times \frac{e}{f} + g$	a + quantity b over quantity c + d all times e over f + g
$(a + \frac{b}{c+d}) \times \frac{e}{f} + g$	a + b over quantity c + d all all times e over f + g

This last expression is peculiar in requiring “all all”, but we see no especially intuitive shorthand around this occasional need. No one said that reading mathematics, especially deeply-nested mathematics, was going to be simple!

Let us return to the quadratic formula. We can say it “The quantity minus b plus or minus the square-root of the quantity b squared minus 4 a c end all over the quantity 2 a end.”

## 4 Speaking Integrals and Sums

The integral [from x=a to b] of f(x)+g(x) d x has the advantage of the closing “d x”, and so in most (not all) traditional notations we can try to read or listen, anticipating that somewhere ahead we will find the “d”.

The  $\sum f(i)$  construction doesn’t have any close, so  $\sum fg + h$  is ambiguous. We could just leave it that way and say that our job is over when the speech is changed to text, but can we fix it with a modest effort? It is unlikely to mean  $(\sum f) \times g + h$  but could be  $(\sum f \times g) + h$  or  $\sum(f \times g + h)$ . These could be spoken, respectively as

sum of f all times g + h;

sum of f times g all + h;

sum of f times g+h all.

Of these, the last seems a strain, but only because there is no operator after the h. If this is truly the end of the expression, the “all” could be left out! The others seem fairly natural, and perhaps more natural if we allow “f times g” to be simply “f g”. It may also be preferable, as mentioned earlier to say “sum of f times g+h end sum.

The advantage of a multimodal input model is that the computer system can display what has been recognized so far. The longest delay is likely to be the pause while the computer waits to determine the end of the utterance. The translation and typesetting should be quite rapid by comparison. Thus the feedback of the choice made by the system may provide a valuable learning experience in these less common forms.

## 5 Additional Examples

We promised some additional examples to fill out the description.

Display	Spoken
$a_0 + x(a_1 + x(a_2 + \dots))$	a sub 0 + x times quantity a sub 1 + x times quantity a sub 2 + dot dot dot
	or — x quantity a sub 2 + dot dot dot
$((a_3x + a_2)x + a_1)x + a_0$	a sub 3 x + a sub 2 quantity times x + a sub 1 quantity times x s+ a sub 0
$a_{(n-1)^2} + 1$	a sub quantity n minus 1 all squared all + 1
$a_n^2 - a_{n-1}^2$	“a sub n squared minus quantity a sub quantity n minus 1 all all squared
$a_n - a_{n-1}^2$	quantity a sub n minus a sub quantity n minus 1 all squared
$(a_n - a_{n-1})^2$	a sub n minus a sub quantity n minus 1 all all squared
$x^3$	x to the third (note ordinal 3rd)
	x to the third power
	x to the power 3
	x raised to the power 3
	x cubed

For convenience in the next few examples we will just say “ $x \uparrow 3$ ” for  $x^3$ .

Display	Spoken
$x^{34}e^z$	$x \uparrow 34 \text{ e } \uparrow z$
	$x \uparrow 34 \text{ times e } \uparrow z$
$x^{n+1}e^z$	$x \uparrow \text{ quantity } n + 1 \text{ all times e } \uparrow z$
$x^{1+n^{2+k}+r}e^z$	$x \uparrow \text{ quantity } 1 + n \uparrow \text{ quantity } 2 + k \text{ all } + r \text{ all times e } \uparrow z$
$\frac{x^n}{y^m} + 4$	$x \uparrow n \text{ over } y \uparrow m + 4$
$f(x) + g(y, z) + r_{i,j} + h(s_i, j)$	$f \text{ of } x + g \text{ of } y \text{ comma } z \text{ all } + r \text{ sub } i \text{ comma } j + h \text{ of quantity } s \text{ sub } i \text{ all comma } j$

The next section presents a small controversy as to how to bracket argument lists of functions. Consider either form  $\sin x$  or  $\sin(x)$  can be spoken “sine of  $x$ ” or “sine  $x$ ”. The comparable “anonymous” functional application form, assuming there is a function named  $p$ , could be written  $px$  or  $p(x)$ . Either of these can be “ $p$  of  $x$ ”, but only the first of these,  $px$  can plausibly be spoken as “ $p$   $x$ ”. And in that case it might be the product of two items. What gives?

Here is one proposal:

There is *no special meaning* for “*of*” in “sine of ... ” or any function known to be a univalent function. If the single argument is compound it must be introduced by “open” or “quantity” Thus  $\sin x$  is simply “sine  $x$ ” but  $\sin(2x)$  is “sine of quantity  $2x$  [close]”.

If the function is not well known, then there is a significance to the “*of*”.  $px$  should be pronounced “ $p$  of  $x$ ” and probably should be written  $p(x)$ . Saying “ $p$   $x$ ” is hazardous. It looks like a multiplication. This is not a prohibition; spoken math as well as written math can be ambiguous! “ $p$  of  $x + 1$ ” means  $p(x) + 1$ . “ $p$  of quantity  $x + 1$  close  $+ 2$ ” means  $p(x + 1) + 2$ .

Display	Spoken
$a \sin x$	a sine of $x$
	a sine $x$
$a \sin x + y$	a sine of $x + y$
	a sine $x + y$
$\frac{a \sin(x+y)}{2} + 1$	a sine of quantity $x + y$ all quantity over $2+1$
$a \sin \frac{x+y}{2} + 1$	a sine of quantity $x + y$ quantity over $2$ all $+ 1$
$\sin x + \cos y$	sine of $x + \cos$ of $y$
$\sin x + \cos y$	sine $x + \cos y$
$\sin(x \cos y)$	sine quantity $x + \cos y$ [unusual]
$f(x, y) + 1$	$f$ of quantity $x$ comma $y$ all $+ 1$ [bad]

Here is an alternative proposal:: *The word “of” has special significance and always carries with it an implicit “open”.* Thus “ $p$  of  $x + 1$ ” means  $p(x + 1)$ , and “ $p$  of  $x + 1$  close  $+ 2$ ” means  $p(x + 1) + 2$ .

The alternative here allows a distinction between “sine of  $x$ ” which is  $\sin(x)$  and “sin  $x$ ” which is  $\sin x$ , but that may be too subtle for users/ speakers.

Display	Spoken
$a \sin x$	a sine of $x$ [implicit close]
	a sine $x$
$a \sin x + y$	a sine of $x$ all $+ y$
$\frac{a \sin(x+y)}{2} + 1$	a sine of $x + y$ all quantity over $2$ all $+ 1$
$a \sin \frac{x+y}{2} + 1$	a sine of $x + y$ quantity over $2$ close $+ 1$
$\sin x + \cos y$	sine of $x$ all $+ \cos$ of $y$ all
$\sin x + \cos y$	sine $x + \cos y$
$\sin(x \cos y)$	sine of $x + \cos y$ close [unusual]
$f(x, y) + 1$	$f$ of $x$ comma $y$ all $+ 1$ [better]

## 6 Extensions

You may not be entirely comfortable with the limited vocabulary, and prefer other words and phrases. These should also be allowed, certainly to the extent that they do not interfere with the existing mechanisms.

Other locutions such as “the fraction  $a + b$  divided by  $c$ ” are easily accomodated. The word “fraction” has exactly the same meaning as “quantity”, and the phrase “divided by” means the same as “over”.

We might plausibly say, and parse,  $(a + bc)(d + ef) + 1$  as “the product of  $a + b$  times  $c$  and  $d + e$  times  $f$  all  $+1$ ” as an alternative to “ $a + b$   $c$  quantity times quantity  $d + e$   $f$  all  $+1$ .”

There are other common locutions such as “ $d$  squared by  $d$  x squared of  $f$  of  $x$ ” for  $\frac{d^2}{dx^2}f(x)$ . A phrase such as “the second derivative of  $f$  of  $x$  with respect to  $x$ ” could be worked into a parser.

Here are some example additional phrases. Consider  $y'' + a^2y = 0$  spoken as “ $y$  double prime  $+ a$  squared  $y$  equals  $0$ ”. Thus “prime”, “double prime”, and “triple prime” seem like “prime candidate” for additions.

To show the flexibility or perversity of notation, here is an expression approximating original notation from Hoare logic [8] describing the semantics of **while** loops. It looks like this:

$$\frac{P \wedge b \{S\} P}{P \{\mathbf{while} \ b \ \mathbf{do} \ S\} P \wedge \neg b}$$

Note that adjacent symbols here are not multiplied. In the form  $A\{B\}C$  indicates a kind of temporal ordering, proceeding from left to right. The curly braces have specific meanings (separating predicates from program sections), and the horizontal line means something like “implies” and does not have any relationship with division. We can nevertheless speak it as “ $\text{cap } p \text{ wedge } b \{ \text{cap } s \} \text{ cap } p$  quantity over quantity  $\text{cap } p \{ \text{roman while } b \text{ roman do cap } s \} \text{ cap } p \text{ wedge neg } b$ ”. We will have to say open/close curly brace; we would have a higher accuracy if we used “Bravo” and “Papa” for the letters  $b$  and  $p$  respectively.

Any programming structure for the recognition of math must be extensible in two directions:

- Allowing new spoken tokens to be introduced to the speech grammar, and
- Allowing new phrases to be parsed, at least in a rudimentary fashion.

Thus one might need to add the words “Poisson bracket” and appropriate pronunciation to the speech engine and then also introduce a grammar rule to allow the “Poisson bracket of  $x$  and  $y$ ” to be typeset in its usual notation  $(x, y)$ , or perhaps come up with some (almost any!) more perspicuous form.

The spoken form as we have specified it does not have any precedence rules of its own, and perhaps surprisingly tends to just pass along ambiguities: That is, you can often speak or typeset an expression without complete concern for its meaning. For example,  $a = b \vee c$  could denote one of the Boolean expressions  $(a = b) \vee c$ , or  $a = (b \vee c)$ , or could be a programming language assignment expression, or something else. Even as we use the expressions to convey different meanings to the reader in the previous sentence, we feel compelled to point out that it requires some kind of external interpretation to distinguish those expressions. When computer typesetting was more of a novelty, it provided the unfortunate illusion that a formula undergoing typesetting gains authority; as is obvious to the modern reader, typesetting does not mean correct or even defined.

### 6.1 A disappointment

It would appear that the speech grammar XML formalism provided by the W3C organization and its conforming implementation by Microsoft would naturally provide a framework for context-free grammar (CFG) parsing. In fact, Microsoft uses a file-extension “.cfg” for compiled grammars. Sadly, the grammars are not permitted to be context-free, but only finite-state. This means that a grammar that allows for nested subexpressions, (using “quantity” and “end” or equivalently open and close parentheses) *cannot* be described completely in the given speech grammar. This throws us back to a more primitive stage in which we can use the Microsoft grammar for recognizing tokens (numbers, symbols), but cannot actually depend on it for parsing. It is true that with some effort one can write a grammar that looks like it will work for expressions, but only if they are of finite depth. Thus one could come up with a grammar that allowed no parentheses, or some fixed depth. This requires essentially duplicating the grammar for each nested depth. W3C *permits*

but does not require a context free grammar. We do not know if the Microsoft VISTA implementation will exceed its current capabilities.

Because of perceived limitations, we began in Sept, 2006 experimentation with the free open-source Java implementation of the speech recognition system Sphinx-4 [18]. This has simplified some issues, but has not been adequately incorporated into the current prototypes.

## 7 Comparisons

We have found only one existing commercial computer program with related goals, Mathtalk [11]. This program is an attempt to provide a facility for humans to speak math to a computer. The on-line demonstrations suggest that the method used is rather sluggish, requiring pauses before and after each operation to wait for recognition. It is not clear how to correct any mistakes. The engineering is quite limiting: Mathtalk *requires* referring to letters by military names (e.g.  $f$  becomes “function foxtrot” and  $\lim_{n \rightarrow \infty}$  is “limit November goes to infinity”.) It is unclear how much of the limitations are inherent in the design or in the lower-level support from their perhaps primitive speech tool.

## 8 Status of spoken math recognition

We have written a rudimentary translator from strings of some spoken symbols into strings again, but of more conventional symbols, essentially replacing words like “quantity” and “all” by parentheses, and inserting some parentheses in other places as appropriate. We can also parse numbers from spoken words (twenty one hundred becomes 2100). Since we anticipate that words and symbols will be misrecognized or missed entirely, we cannot just walk away from the task when the speaker halts. There is a feedback step in which the computer attempts to display—to the extent possible—what has been recognized or not. This feedback is described in a separate paper on the display of incomplete expressions [6]. This feedback is based on transforming a string of tokens into T<sub>E</sub>X and typesetting an approximation to what has been spoken, with placeholders for parts that were not recognized. It may seem odd to a programming-language trained reader that one can truthfully declare complete speech recognition success upon nicely typesetting something “symbolic” as a mere string of words. Yet it cannot be the recognizer’s fault if the human speaker has uttered partial or complete nonsense posing as mathematics. Given some partial display, it may be plausible for the speaker to abandon the original utterance and instead patch what the computer heard. If an unmatched open parenthesis is spoken, the matching one will be displayed, and need not be spoken. However, the insides may need to be filled in.

As of June, 2004, student Kevin Lin and R. Fateman connected a speech recognizer with a simple grammar to a prototype mathematics editing system (SKEME), which allows the user to speak simple tokens or expressions instead of typing. The tokens may be compounded as in “script capital A” and they may also be concatenated with simple operators as in “a plus b”. The locations for the insertion of spoken or written symbols is governed by cursor or attention point which is positioned with a mouse. The mechanism used (Microsoft Speech SDK5.1) does not provide alternate (less confident) speech recognition results, and so is overly fragile. Another group of students using a similar design produced a more robust yet still “demoware” system called Math Speak and Write, which can be accessed at <http://www.cs.berkeley.edu/~fateman/msw/msw.html>.

We continued, in 2006, to explore a more effective speech grammar definition permitting alternative recognitions, and variations on user-interface issues, especially for correcting errors in conjunction with handwriting input so that the such methods have adequate appeal for users to try these novel interfaces. The Sphinx-4 re-engineering of the project was halted primarily by the graduation of students. Also in 2006, undergraduate students (principally Sherman Lee), have shown how to link the math recognizer to virtually any Windows program via the clipboard. That is, a spoken (or handwritten + spoken + typed) formula from MSW can be inserted into Powerpoint, Word, Excel, etc. There is an additional barrier to overcome, in that a mere textual version of a formula does not have interoperability with the objects produced by using (say) the expression templates available in Microsoft’s Equation 3.0. Lee’s program produces text similar to the

input for T<sub>E</sub>X. Further students (in 2006) Albert Shau and Eric Chang were studying alternatives for TTS, and David Poll was helping simplify the programming by using a Lisp to .NET tool, thereby eliminating some other layers of languages, allowing improvement to be more easily incorporated into prototypes. The project has lain dormant now for about 3 years.

## 9 Summary and Conclusion

This paper reports on incomplete work primarily on design, but describes partial implementations (which are available from the author). It does not include human-factors experiments. While such exercises may be useful in the future, we are not convinced that the substantial effort to mount such experiments would be worthwhile just yet: most of our observations, based on our implementations, have led us unambiguously in particular design directions for improvement. It is yet time to seek the reactions of a class of naive users who might very well concentrate on trivial implementation issues rather than design; we believe that the effort, at the outset, of training of speech recognition (and associated handwriting) on a person-by-person basis would turn away all but the most highly motivated persons. In a future system which is largely “pre-trained” such problems may be reduced. At this point we would rather not delay the presentation of this material for others to consider.

We can still experiment with the prospects of multimodal input of mathematics, but would probably bite the bullet and move the project to Vista, which has superior speech facilities. We see clear benefits of speech when saying “bold script capital R” or for distinguishing among the large number of symbols and collections of strokes that look essentially similar when written (recall 1<2 and K2 as one example [14]). If speech can also be used for accurately conveying larger multi-token arrangements as indicated here, that represents a possible benefit. Speaking is “hands free” so that it can complement either a stylus or keyboard-form of input<sup>13</sup>.

While we do not expect speaking to be a “unimodal” mode of choice for very long expressions in a single gulp, we believe that in combination with pointing, speech can “fill in the boxes” which would be pointed-to and otherwise constructed or corrected via templates in an interactive input system. We have also written program modules to implement handwriting correction in which a token is displayed along with alternatives. By saying “no. alternate 2” the token is replaced by another, in this case, the second-ranked one.

We hope this paper and the open-source availability of our earlier programs will encourage others to join us to pursue this approach as well. We are quite aware that it requires substantial resources to raise “demonstration” programs or prototypes to the level of general usefulness; identification of a critical (and funded) application would be key.

## 10 Acknowledgments

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<sup>13</sup>A head-mounted microphone/earphone set is nearly imperative on most computers for good accuracy, and is rather expensive. At least one Toshiba Tablet PC overcomes the need for a headset by having 3 microphones arrayed around its screen.



improving speech interfaces to intelligent speech processing capability in the future. Such processing need not be restricted to computing on the smart phone itself.)

Several of the papers referenced below are unpublished but accessible from the author's home page:  
<http://www.cs.berkeley.edu/~fateman/papers/>

Thanks to Neil Soiffer for useful suggestions. Chu Yaru, a graduate student at Lanzhou University made some suggestions as well (2011). We were also made aware of the TalkMaths project at Kingston University, London (UK) <http://talkmaths.sourceforge.net> which is listed as a current project (2011) of the CLASP group at Kingston. We have not done a comprehensive update of references to date for this paper; an internet web search may reveal yet other projects.

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## 11 Appendix: Ambiguity, Syntax and Semantics

Humans tend to write ambiguous mathematics, expecting that the context, imposed by the human reader, will disambiguate.

For example if  $e_1$  and  $e_2$  are expressions, how do we parse arguments to functions like  $\sin$  and  $\cos$ ? Does  $\sin e_1 e_2$  mean

1.  $(\sin e_1) \times e_2$  or
2.  $\sin(e_1 \times e_2)$ ?

Please choose one and then examine the well-known equation displayed below.

$$\sin 2z = 2 \cos z \sin z$$

This is a formula *typeset exactly as given in a standard reference* (4.3.24 Abramowitz and Stegun [2]). You might read it out loud. Now note that on the left,  $e_1 = 2$ ,  $e_2 = z$  uses convention 2. On the right,  $e_1 = z$ ,  $e_2 = \sin z$  uses convention 1. Two different parsing conventions are used in the same equation. This is not unusual. Our point here is that linearizing the token stream is actually not sufficient to guarantee unambiguous syntax. (Additional rules about the precedence of invisible multiplication between numbers and symbols can solve this particular problem, but there are others in which the writer and the typesetter conspire to confuse even the skilled reader.) For an amusing account of ambiguity in classical mechanics see the introduction to an on-line mechanics book by Sussman and Wisdom [19] <http://mitpress.mit.edu/SICM/book-Z-H-5.html>.

On a semantic note, there are at least two proposals for a semantic encoding of mathematics, the most prominent of which is (the semantic component of) MathML [12]. On the face of it, writing a new paper in which one is using some speech or handwriting or keyboarding to produce a totally new notation makes it logically impossible to properly encode it with respect to its (up to this moment undefined) semantics. A kind of meta-language relating both new notation and the new semantics to that of existing notions is required. This meta-language too may have similar limitations, rather like explaining colors to a sightless person.

Where does this lead us? We are personally more inclined to first look for mappings of mathematical notations to some operational semantics such as computations in a given computer algebra system. Such a system generally imposes limits but within its realm of discourse is at least definite. Beyond that level, we may be forced to deal with a largely syntactic or geometric appearance of mathematics or aural representation!

For many purposes, including the obvious application of typesetting for consumption by mathematicians at a future time, this operational semantics is an additional boon, and can be combined with older material or material currently produced in a conventional manner. This traditional material has been limited to mathematical syntax encoded as typeset material, plus natural language commentary. Most current computer algebra systems provide some kind of documentary framework or notebook which can handle such conventional material. The issue is how much further we can push in the semantics direction.