

# Topics

- 80x86 Stack
- 32-bit Procedures with Value Parameters
- Additional 32-bit Procedure Options
- 64-bit Procedures
- Macro Definition and Expansion

32-bit Procedures with Value  
Parameters

Courtesy: UMBC and JB Learning

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# Terminology

- *Procedure*
  - a subprogram
  - essentially a self-contained unit
  - Well-defined interface
- Called by
  - Main program
  - another subprogram
- Objective
  - perform a set of tasks
  - may or may not return a value

# Procedure Concepts

- Transfer of control from calling program to procedure and back
- Passing parameter values to procedure and results back from the procedure
- Having procedure code that is independent of the calling program

The *cdecl* protocol provides one standard implementation scheme in the 32-bit environment

EAX-->RETURN

Call--> by default we push  
ret addr

ret---->by default pop EIP

.code

main proc

call addproc

main endp

Addproc Proc

ret

Addproc endp

# Procedure Definition

- In a code segment with body statements bracketed by PROC and ENDP directives giving procedure name

```
.CODE  
procName PROC  
; procedure body  
    ...  
procName ENDP
```

## Transferring Control to a Procedure

- In the “main”(or any other) proc , use  
`call procName`
- The next instruction executed will be the first  
instruction in the procedure.

## Returning from a Procedure

- In the procedure, use  
ret
- The next instruction executed will be the one following the call in the “main”(or any other) proc.

1) EXP X

call → pushed → ret → pop EP

Code

Add Proc

====

ret

Add Endt

Main Proc

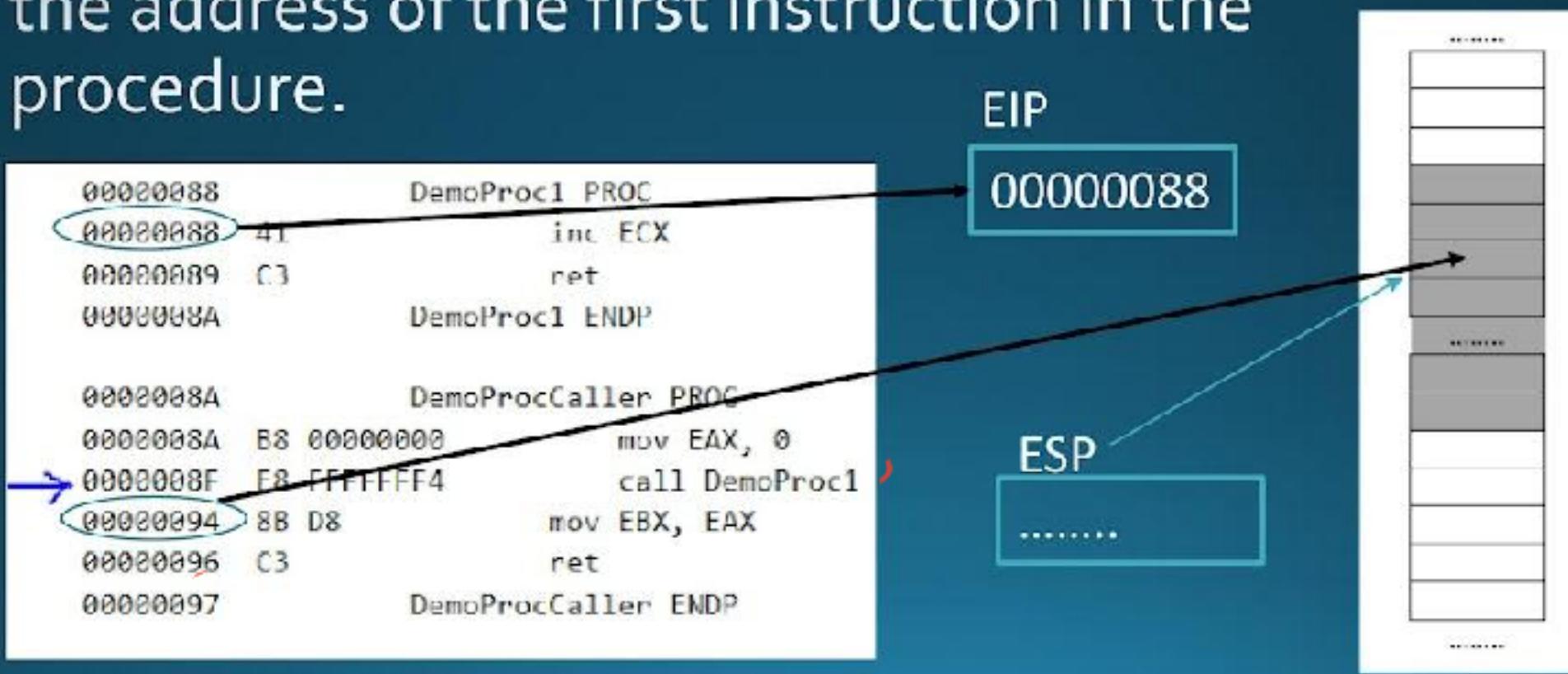
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→ call Add

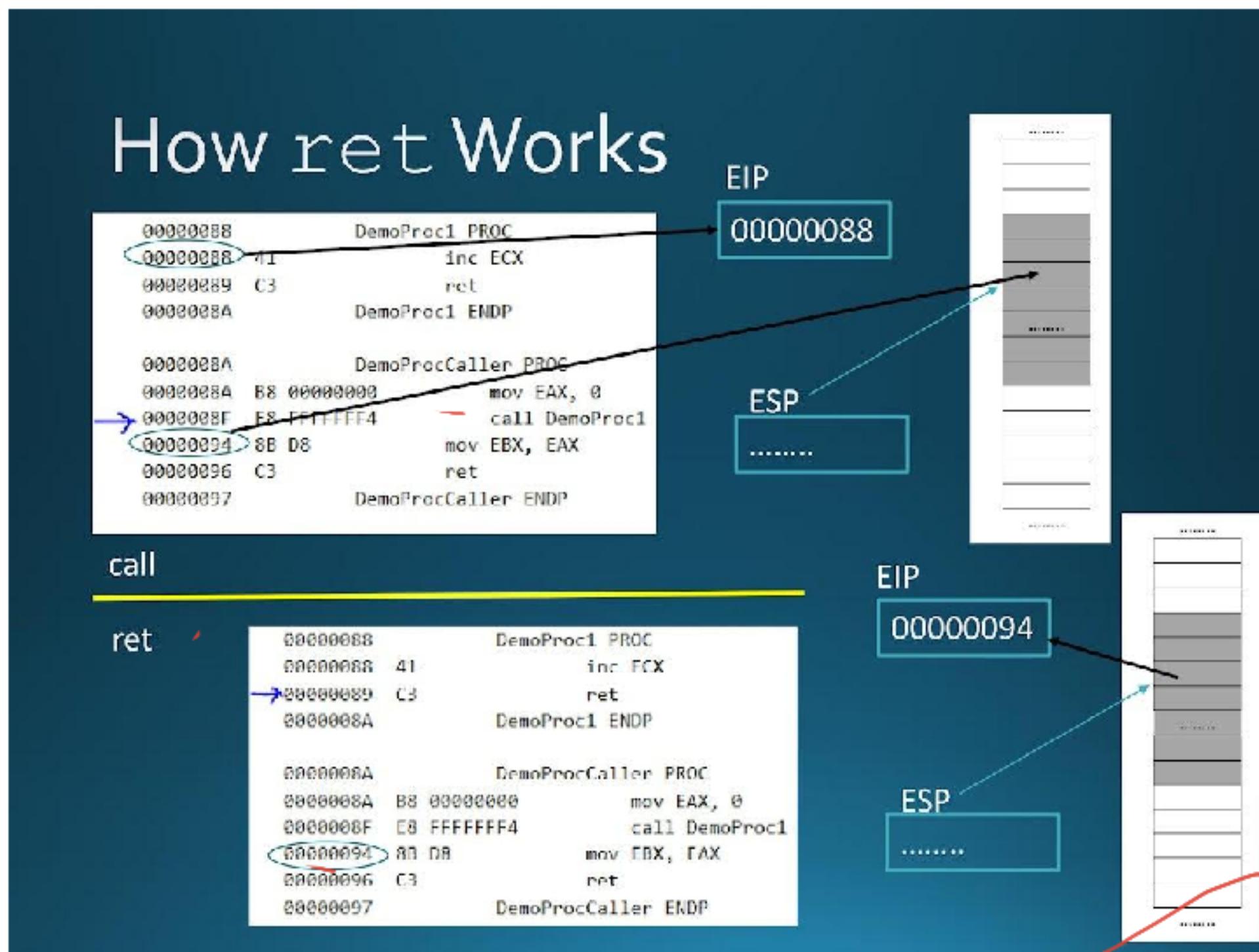
Main Endt

# How call Works

- The address of the instruction following the call is pushed on the stack.
- The instruction pointer register EIP is loaded with the address of the first instruction in the procedure.



EIP-- has the address of next instruction to be executed



## ret – instruction encoding

Type	Operand	Opcode	Byte length
Near	None	C3	1
Near	Immediate	C2	3
Far	None	CB	1
Far	Immediate	CA	3

```
00000097 C3          ret
00000098 C2 0014      ret 20
```

# Parameter Terminology

- A procedure definition often includes *parameters* (also called *formal parameters*).
- These are associated with *arguments* (also called *actual parameters*) when the procedure is called.
- For a procedure's *in* (*pass-by-value*) parameters, values of the arguments are copied to the parameters when the procedure is called.
  - These values are referenced in the procedure using their local names (the identifiers used to define the parameters).

Add (Value<sup>1</sup>, Value<sup>2</sup>)

Pass by value

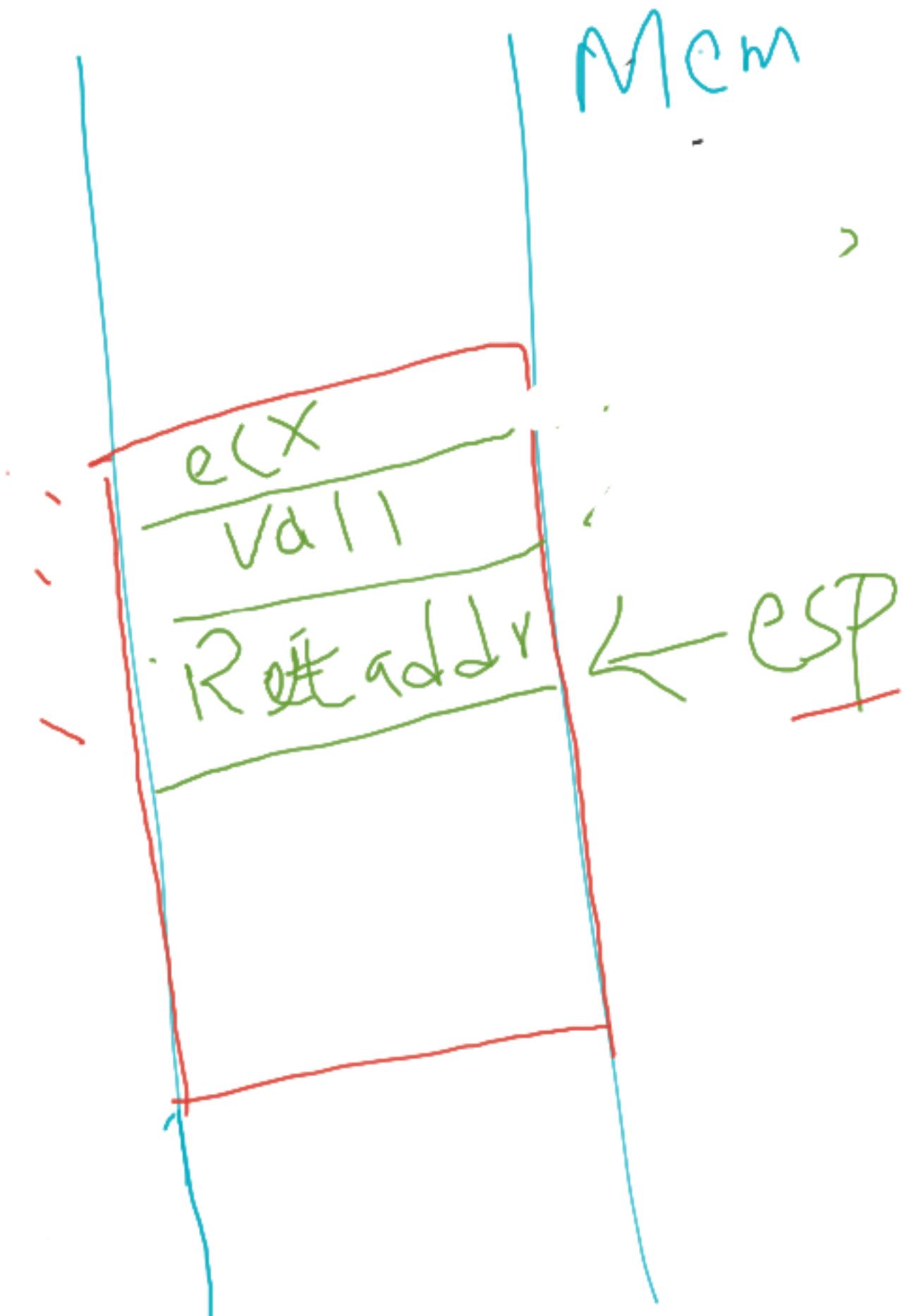
Pass by reference

# Implementing Value Parameters

- Parameter values normally passed on the stack
  - Pushed in reverse order from argument list
- Example
  - Design or high-level code:  $sum := add2(value1, value2)$
  - 80x86 implementation:

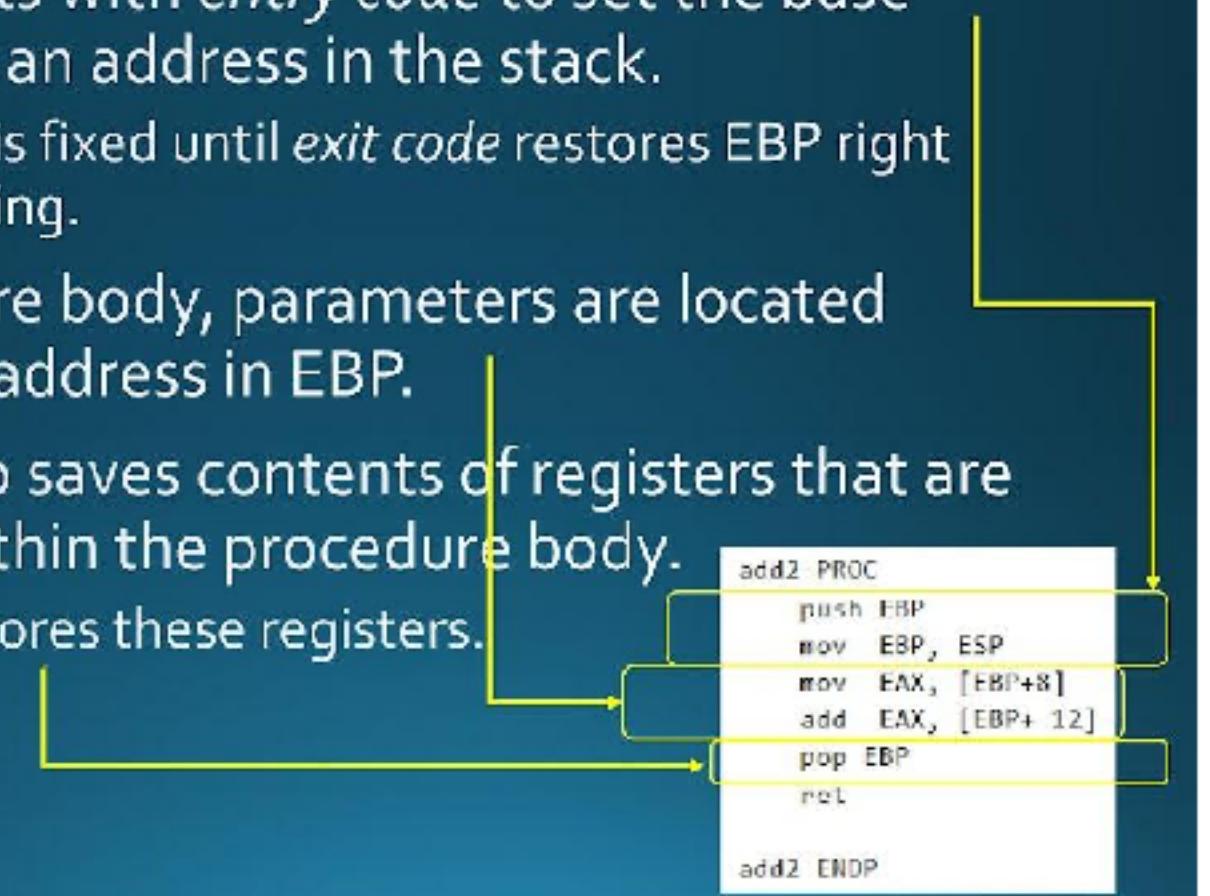
```
push    ecx      ; assuming value2 in ECX
push    value1    ; assuming value1 in memory
call    add2      ; call procedure to find sum
add    esp, 8     ; remove parameters from stack
mov    sum, eax   ; sum in memory
```
- With the *cdecl* protocol, the calling program must remove parameters from the stack.
- A single integer value is returned in EAX.

procedure result need to be stored  
in EAX



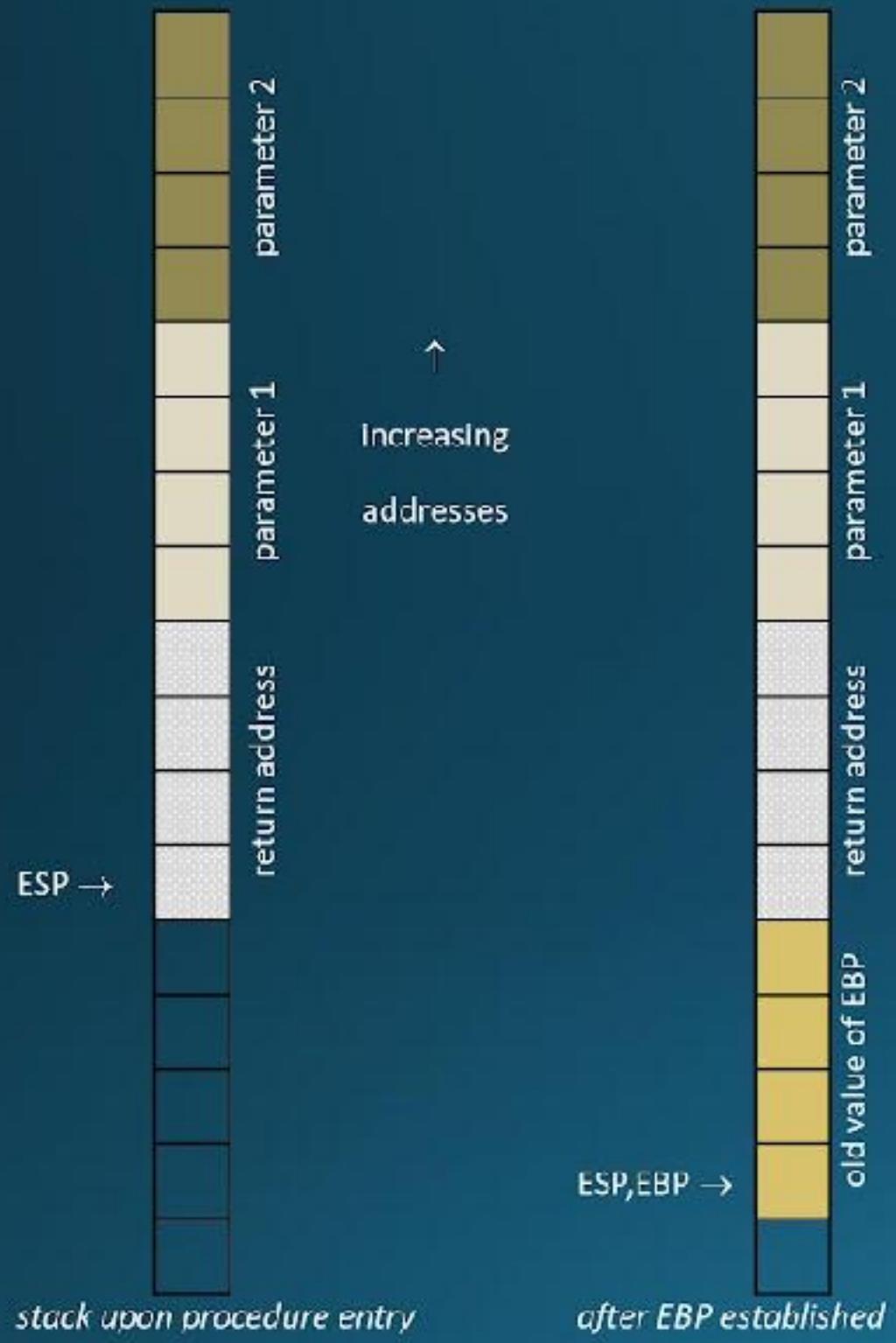
# Procedure Entry Code

- Because the stack pointer ESP may change, a procedure starts with *entry code* to set the base pointer EBP to an address in the stack.
  - This location is fixed until *exit code* restores EBP right before returning.
- In the procedure body, parameters are located relative to the address in EBP.
- Entry code also saves contents of registers that are used locally within the procedure body.
  - Exit code restores these registers.



# Example Procedure

```
add2    PRCC          ; add two words passed on the stack
              ; return the sum in the EAX register
    push   ebp          ; save EBP
    mov    ebp,esp       ; establish stack frame
    mov    eax,[ebp+8]   ; copy first parameter value
    add    eax,[ebp+12]  ; add second parameter value
    pop    ebp          ; restore EBP
    ret               ; return
add2    ENDP
```



Stack frame upon  
entry to  
procedure *add2*  
and after EBP  
established in  
entry code

## Accessing Parameters in a Procedure

- Use *based addressing*
- Because the value is actually on the stack, [EBP+*n*] references the value
- Example:  
`mov eax, [ebp+8]`  
copies the last parameter pushed to EAX

# Saving/Restoring Registers

- Push registers in entry code after EBP is established.
- Exit code pops registers in the opposite order.
- Calling program's flag values from can be saved with pushfd/popfd.

# Entry and Exit Code Summary

- *Entry code:*

```
push    ebp      ; establish stack frame
mov     ebp, esp
push    ...      ; save registers
...
push    ...
pushfd  ; save flags
```

- *Exit code:*

```
popfd  ; restore flags
pop    ...      ; restore registers
...
pop    ...
pop    ebp      ; restore EBP
ret    ; return
```

# Procedure Call Summary

## *calling program code*

- push arguments on stack in right-to-left order
- call procedure

## *procedure code*

- save EBP and establish stack frame
  - save registers used by procedure
  - access parameter values using based addressing in procedure body
  - return value, if any, goes in EAX
  - restore saved registers and EBP
  - return
- 
- add number of bytes of parameters to ESP
  - for value-returning procedure, use value in EAX

.4096 stack

.data

value1 dword ?

value2 dword ?

.code

add proc

push ebp

mov ebp,esp

mov eax,[ebp+8]

add eax,[ebp+12]

pop ebp

ret

add endp

main proc

push value2

push value1

call add

add esp,8

mov sum,eax

main endp

pass by value  
example

Ebp is a global register it may have valuable  
information needed for another procedure

add(value1,value2)