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Saishnu Ramesh Kumar (300758706)
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CSCI 117 – Lab 11
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Part 1a:

```
local SumListS SumList Out1 Out2 SumList2 in
  fun{SumList L} //Declarative iterative
    case L of nil then 0
    []'|'(1:H\ 2:T) then (H+\{SumList\ T\})
    end
  end
  fun{SumListS L} //Stateful iterative
  C = newCell 0
  {SumList2 C L}
  end
  fun{SumList2 C L}
    case L of nil then @C
    []'|(1:H 2:T) \text{ then } C := (H + @C) \{SumList2 C T\}
    end
  end
    Out1 = {SumList[1 2 3 4]}
    Out2 = {SumListS [1 2 3 4]}
```

```
skip Browse Out1
     skip Browse Out2
     skip Full
end
local foldL2 foldL Out1 Out2 Z in
  fun\{foldL\ F\ Z\ L\}
     case L of nil then Z
     [] '|' (1:H 2:T) then
       \{foldL\ F\ \{F\ Z\ H\}\ T\}
     end
  end
  fun\{foldL2\ F\ Z\ L\}\{
     FL2 = newCell 0
     foldL3 in
     proc {foldL3 F Z L}
     case L of nil then FL2 := @FL2
     [] '|'(1:H 2:T) then
     FLS := 0
     FLS := \{F Z H\}
     {foldL3 F @FLS T}
     end
```

```
end
```

```
{foldL3 F Z L}
@FLS
end

Out1 = {foldL fun P {$ X Y} (X + Y) end 3 [1 2 3 4]}
Out2 = {foldL2 fun {$X Y} (X+Y) end 3 [1 2 3 4]}
skip Browse Out1
skip Browse Out2
skip Full
}
```

Part 1b:

With skip Full in the SumList and FoldL program, it outputs the values for Out1 and Out2, as well as the kernel syntax. This will also output the stores and the environment for the two functions as well.

Part 2a:

local Generate Num GenF Out1 Out2 Out3 in

```
fun{Generate}
```

Num = newCell - 1

```
fun{$}
       Num := (@Num + 1)
       @Num
       end
end
GenF = \{Generate\}
Out1 = \{GenF\} // returns 0
Out2 = \{GenF\} // returns 1
Out3 = \{GenF\} // returns 2
skip Browse Out1
skip Browse Out2
skip Browse Out3
```

end

Part 3a:

```
local NewQueue S Pu PPo IsE Av A1 A2 B1 B2 V1 V2 V3 Out Appened Out 1 in
```

```
Append = fun \{ Ls Ms \}
case Ls
  of nil then (Ms | nil)
  []'|'(1:X 2:Lr) then Y in
  Y = \{Append Lr Ms\}
    (X|Y)
    end
  end
  fun {NewQueue L}
    C = newCell nil
    S = newCell 0
    Push Pop IsEmpty SlotsAvailable in
    proc {Push X}
    if(@S == L) then
       B = @C in
       case B of '|' (1:Y 2:S1) then C:= S1 end
       C := \{Append @C X\}
```

```
S := (@S + 1)
  else
    C:=\{Append @C X\}
    S := (@S + 1)
  end
end
fun \{Pop\}\ B = @C \text{ in }
  case B of '|' (1:X \ 2:S1) then C:=S1 \ X end
end
fun{IsEmpty} (@C == nil) end
fun{SlotsAvailable} B in
  B = (L - @S)
  В
end
ops(push:Push pop:Pop isEmpty:IsEmpty avail:SlotsAvailable)
end
S = \{NewQueue 2\}
```

S = ops(push:Pu pop:Po isEmpty:IsE avail:Av)

```
B1 = \{IsE\}
  A1 = \{Av\}
  {Pu 1}
  {Pu 2}
  A2 = \{Av\}
  {Pu 3}
  B2 = \{IsE\}
  V1 = \{Po\}
  V2 = \{Po\}
  V3 = \{Po\}
  Out = [V1 \ V2 \ V3 \ B1 \ B2 \ A1 \ A2]
  skip Browse Out // Out : [ 2 3 Unbound true() false() 2 0 ]
end
```

Part 3b:

This is a secure ADT because it aids in hiding information. The data that is located inside the NewQueue function cannot be changed unless you have some form of token which would grant you access into the data, the tokens in this case are: Push, Pop, isEmpty, SlotsAvailable. Those tokens were not given access to the user. In this case, we are only using the push and pop for the elements that are in the queue, but we are not changing any of the values of the elements that are in the queue currently.

Part 3c:

In comparison to the secure declarative ADT on page 431, the program given on the page is related to the memory usage that is needed only for unbundled ADTs as wrapping is needed for those types. Moreover, the function StackOps takes the list S and it returns the values of the procedure ops (pop:Pop push:Push isEmpty:IsEmpty). This can also be considered to be a declarative type of object-oriented programming where the stack is a declarative object.