ASSIGNMENT 2: ADDING SYSTEM CALLS

Aim: To add new system calls to a given kernel version and observe the results

Requirement: All work should be done on linux kernel version 4.19.210.

ADDING KERNEL 4.19.210 TO SYSTEM

- 1. Created an azure student account using iiit mail id.
- 2. Created vm with **UBUNTU 16**, 4 core (linux-4.19.210).
- 3. Deployed vm on azure and used the public ip address to run the vm on the local machine.(using SSH client)
- 4. Created new folder in sudo mode and installed the pre-requisites required to run the kernel.
- 5. Go to above directory, created a new folder which contains all the required .c files, and then built and ran `sudo make` and `sudo make modules_install install`.
- 6. Made some configurations and rebooted the vm.

ADDING A SYSTEM CALL TO KERNEL

- 1. Created a new directory in /linux-4.19.210 for a system call. In this text, let it be 'ass2'.
- Created a .c files (eg q1.c) and add the necessary kernel space c code to it. We can use
 either asmlinkage long functionname(void) or the macro
 SYSCALL_DEFINE#(syscall_name) where # is the number of arguments expected and
 syscall_name is the name given to it.
- 3. Create a file named 'Makefile' and add the line 'obj-y := q1.o' to it, if the name of the c file is q1.c.
- 4. Navigate to / linux-4.19.210/include/linux and open syscalls.h
- 5. Add the definition of the function created in q1.c
- 6. Navigate to / linux-4.19.210/ and open Makefile
- 7. Add the name of the folder as 'ass2/' to core-y assignments as shown in subsequent headers.
- 8. Navigate to /linux-4.19.210/arch/x86/entry/syscalls and open syscall_64.tbl (or syscall_32.tbl if 32 bit system) and add appropriate entries to it. If function was made directly, we can make entry as shown for number 548, but if SYSCALL_DEFINE# macro was used, the following 3 lines should be used as references (x64(or32)__sys_ prefix should be added to the last column).
- 9. Finally, navigate to /linux-4.19.210 and run the following commands:
 - `sudo make modules_install`
 - `sudo make install`
- 10. Updated the GRUB menu and restarted the vm and, then enter the same kernel again
- 11. Test the system call via a c code.

SYSCALLS TO CREATE:

Four system calls are to be created, given below:

- 1. Printing a welcome message to kernel logs
- 2. Printing a given string to kernel logs
- 3. Printing current and parent process id to kernel logs
- 4. Recreating an existing syscall and implementing that (in this case, getpid() has been recreated)

SYSCALL_64.TBL, MAKEFILE AND SYSCALLS.H

This list contains syscall entries and reference numbers for them. For each new syscall that we create, we need to add the corresponding entry. For the above 4 tasks, the last 4 entries in the given image were added (548, 549, 550, 551)

```
cs. root@myVM: ~/linux-4.19.210/arch/x86/entry/syscalls
  GNU nano 2.5.3
                                                                         File: syscall_64.tbl
                                                    x32_compat_sys_preadv64
         x32
                   pready
                                                   x32_compat_sys_pwritev64
535
         x32
                   pwritev
                                                   x32_compat_sys_rt_tgsigqueueinfo
         ×32
                   rt_tgsigqueueinfo
536
         x32
                   recvmmsg
                                                   x32_compat_sys_recvmmsg
                                                   538
         x32
                   sendmmsg
                  process_vm_readv
process_vm_writev
539
         x32
540
         x32
                                                   _x32_compat_sys_process_vm_writev
                  setsockopt
541
         x32
                                                    x32_compat_sys_setsockopt
                   getsockopt
542
         x32
                                                   x32_compat_sys_getsockopt
                                                   x32_compat_sys_io_setup
x32_compat_sys_io_submit
543
         x32
                   io_setup
io_submit
544
         x32
545
                                                   x32_compat_sys_execveat/ptregs
x32_compat_sys_preadv64v2
         x32
                   execveat
546
                   preadv2
         x32
547
                   pwritev2
                                                   ___x32_compat_sys_pwritev64v2
_x64_sys_saihello
         x32
548
                   saihello
                                                   _x64_sys_saiprint
_x64_sys_saiprocess
_x64_sys_saigetpid
                   saiprint
                   saiprocess
                   saigetpid
```

Each system call was added in a different folder under the linux-4.19.210 directory, so each of the folders have to be added in the Makefile for the kernel itself (cust1 through 4 for each question 1 to 4)

Finally, the required definitions are also added into syscalls.h

WELCOME MESSAGE TO KERNEL LOGS

1. The following code was added to saihello.c in /linux-4.19.210/ass2, along with the following Makefile containing 'obj-y := saihello.o`.

```
GNU nano 2.5.3
#include<linux/syscalls.h>
#include<linux/kernel.h>

SYSCALL_DEFINEO(saihello)
{
    printk("Hello, I am Sai Teja\n");
    return 0;
}
```

2. A new syscall entry was added in syscall_64.tbl

PRINTING SOME STRING FROM USER SPACE TO KERNEL SPACE

1. The following code was added to saiprint.c in /linux-4.19.210/ass2, along with the following Makefile containing 'obj-y := saiprint.o`.

Here, SYSCALL_DEFINE2 is a macro that is creating a function asmlinkage long saiprint(char __user * str,int len), with one argument str and one int. The '__user' marks the pointer as one to the user space. 'Copy_from_user' is an API call that allows copying data from a pointer in userspace to kernelspace buffer. If there is some sort of error while fetching information, EFAULT error flag will be returned.

A new syscall entry is added in syscall_64.tbl. It is different from the first entry in that it
has a new sys_ prefix and __x64_sys_ prefix for alias and syscall name respectively.
This is a result of using the SYSCALL_DEFINE# macro.

PRINTING CURRENT AND PARENT PROCESS ID

1. The following code was added to saiprocess.c in /linux-4.19.210/ass2, along with the following Makefile containing 'obj-y := saiprocess.o`.

```
Select root@myVM: ~/linux-4.19.210/ass2

GNU nano 2.5.3

#include<linux/syscalls.h>
#include<linux/kernel.h>
#include<linux/cred.h>
#include<linux/sched.h>

SYSCALL_DEFINEO(saiprocess)
{
    struct task_struct *parent=current->parent;
    printk("parent_process_pid: %d \n", parent->pid);
    printk("current_process_pid: %d \n", current->pid);
    return 0;
}
```

'Current' refers to a structure of type 'task' that stores information about the current process (that includes the entirety of the Process Control Block). The current->pid element will have the current process id. The parent->pid call will return the pid of the parent process given a task structure for current process.

2. Appropriate entry in syscall_64.tbl is added

RECREATING A SYSCALL: 'getpid()'

1. The following code was added to saigetpid.c in /linux-4.19.210/ass2, along with the following Makefile containing 'obj-y := saigetpid.o`.

```
GNU nano 2.5.3 File
#include<linux/syscalls.h>
#include<linux/kernel.h>
#include<linux/cred.h>
#include<linux/sched.h>

SYSCALL_DEFINE@(saigetpid)
{
    printk("%u\n",task_tgid_vnr(current));
    return task_tgid_vnr(current);
}
```

The uidgid.h header contains the k_uid structure that will have the user id stored in it. Syscall getuid() itself fetches this information from this place. Current_uid() will return the k_uid structure and val will return the uid unsigned long from it.

Appropriate entry in syscall_64.tbl is added

OUTPUTS

The above code has been used to test the newly created syscalls 548 to 551.

The kernel logs are as follows

```
root@myVM:~# ./a.out
2062root@myVM:~# dmesg
[ 422.979801] Hello, I am Sai Teja
[ 422.979806] Hello World
[ 422.979807] parent_process_pid: 1998
[ 422.979808] current_process_pid: 2062
[ 422.979809] 2062
root@myVM:~#
```

Parent and current process ids are different. Why?

The process id refers to the process that is handling both the syscall and the function that called the syscall. That is, in this case saiprocess.c and test.c are part of the same process. Calling another function does not necessarily create a new process, hence 2062 in this case refers to our currently running functions.

The parent process id 1998 refers to the parent process of test.c, the code within which the system call is being tested.