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INTRODUCTION:

This activity is to convert the fifteen-instruction 32-bit single-cycle MIPS processor into a five-stage, pipelined design. Apart from that, we interface the processor with a factorial accelerator and a general-purpose I/O unit using memory-mapped interface registers.

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GIVEN TASK:

The task is to complete the design and functional verification of a pipelined MIPS processor. Also, we are further required to complete the design and functional verification of a small system-on-chip (SoC) system which includes the pipelined MIPS processor, a factorial accelerator (that we have completed in Assignment 1), and a simple general-purpose IO (GPIO) unit, all connected via direct interface to the MIPS local bus, using memory-mapped interface registers. This report must contain the following:

- a) Draft of pipelined MIPS microarchitecture schematic.
- b) Draft of SoC interface schematic.
- c) Tables for MIPS control unit and memory maps for SoC interface.
- d) Performance analysis of hardware accelerated $n!$
- e) The unit-level including interface wrappers simulation.
- f) Completed interface design for the SoC with the pipelined MIPS processor, the factorial unit (from Assignment 1) and the simple GPIO.

Note that apart from the above given tasks, we must do the following two activities.

- a) Based on the enhanced single-cycle MIPS processor (which supports the additional instructions multu, mfhi, mflo, jr, jal, sll, slr) you successfully accomplished in Assignment 6, use the test program (factorial calculation via nested procedure calls) included in the assignment to test the pipelined MIPS processor.
- b) For the SoC design, we must be able to use the factorial accelerator to handle 4-bit input data (n). In addition to that, we should use the simple GPIO for validation i.e., the MIPS should read input (n) via GPIO and return the result ($n!$) also via GPIO.

STEPS TAKEN TO COMPLETE THE TASK:

- a. Spent time to understand the question carefully.
- b. Discussed on how to proceed with the pipelined MIPS microarchitecture. This took us more than 1 week.
- c. After coming up with a rough design, we started designing or drawing the architecture in Microsoft Visio tool. This too took us a week or so to complete the design in Visio.
- d. Since, the control unit of the pipelined MIPS processor is pretty much like the single-cycle processor, we directly started preparing the tables for both auxiliary and main decoders by differentiating them with respect to their values in each pipeline stage.
- e. Gone through the IO lecture notes to revisit some of the topics that are required to complete this assignment such as interface wrappers for both factorial accelerator and GPIO,

memory-mapped IO design (using *lw*, *sw* instructions), memory map tables, and the verilog code to understand the core logic behind the interface wrappers.

- f. Using the above knowledge, we started preparing the memory map tables for MIPS processor, Factorial Accelerator and the GPIO.
- g. The above lecture notes also provide insight into the system-on-chip (SoC) design with a schematic already given in the notes. Spend some time to understand the logic behind the SoC schematic diagram.
- h. Performance analysis is done for both single-cycle and MIPS pipelined processor with respect to the number of cycles factorial accelerator took to get output of the factorial computation.
- i. After the above analysis, we started designing the verilog modules and writing corresponding logic for the modules.
- j. To test the design, we created a new testbench for all major components such as Factorial Accelerator, GPIO unit, Pipelined MIPS and the SoC address decoder.
- k. After that, we simulated the design and verified the output waveforms for the design correctness and concluded that each component is working fine on its own (had to make additional changes to the datapath design of the pipelined MIPS).
- l. Now, it was time to test the whole SoC design, for which we had written a new testbench which accomplishes the above task of testing the whole system.
- m. Finally, we simulated the final testbench and after getting the desired output, the whole pipelined datapath is updated using the Visio tool to match the logic written in verilog.

DISCUSSION SECTION:

1) System-on-chip (SoC) Design:

The design explanation of the system-on-chip (SoC) will facilitate in explaining how the factorial accelerator and GPIO unit were designed and how they were interfaced with the CPU. The SoC contains the following elements: instruction memory, MIPS Processor (pipelined CPU), data memory, factorial accelerator, general purpose I/O (GPIO), an address decoder, and a 4-input multiplexer to control writeback data to the MIPS processor. The following design of this SoC is shown in *Figure 1.1*.

The MIPS processor fetches an instruction from instruction memory and processes the instruction accordingly. If the instruction is a save word (*sw*) or load word (*lw*) instruction, the processor will issue an address to the address decoder along with a memory write enable signal. The address bus is also connected to each I/O device. A write data bus is connected to the data memory, factorial accelerator, and GPIO unit. However, these I/O devices essentially do nothing until given a write enable signal from the address decoder. The address of the target I/O device is given in the *lw* or *sw* instruction. In this way, the MIPS processor uses memory-mapped I/O to communicate with connected I/O devices like the factorial accelerator, data memory, or GPIO unit. A

table showing what addresses are mapped to each device is shown in *Table 1.2*. For example, a *lw \$t1, 0x800(\$0)* instruction addresses the factorial accelerator.

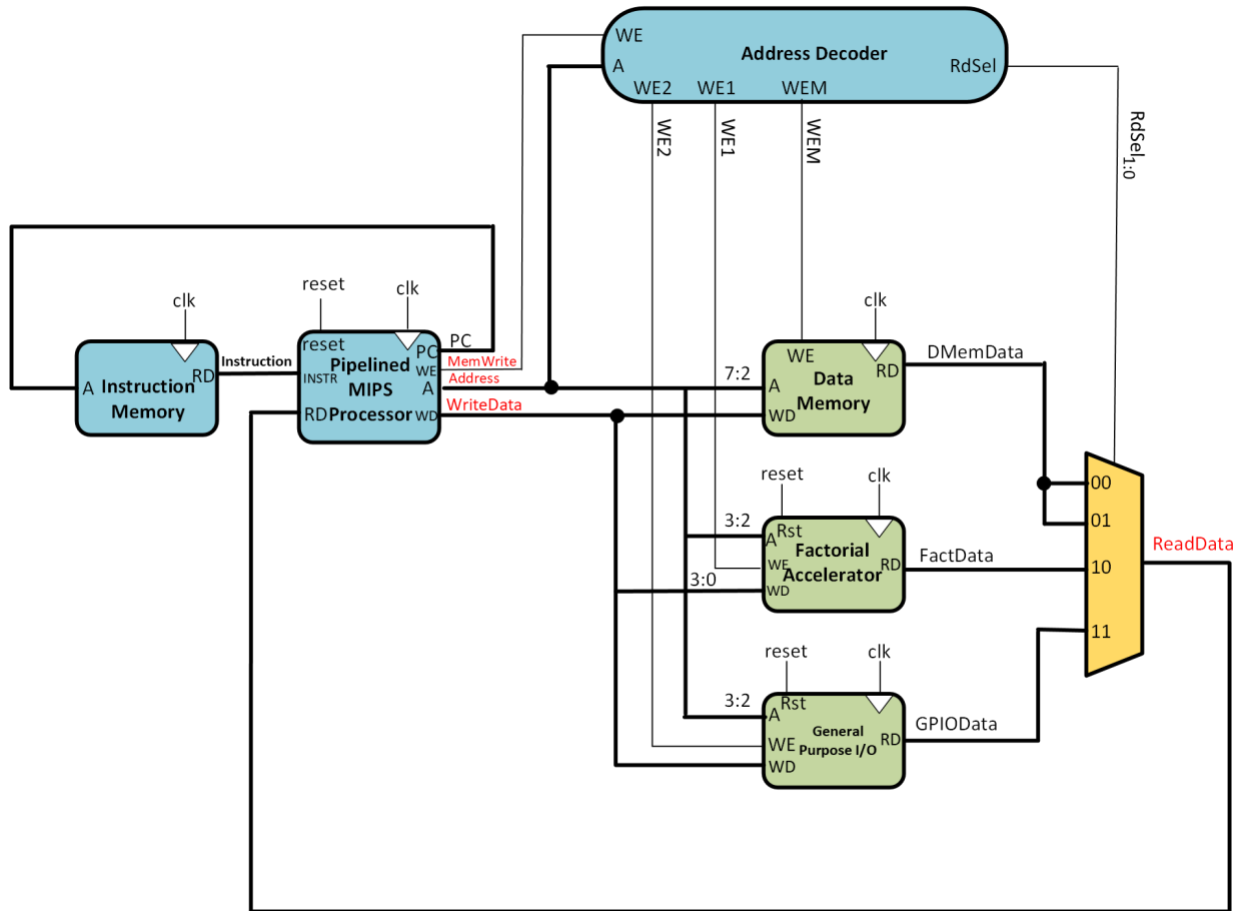


Figure1.1 – System-on-chip (SoC) schematic diagram

MIPS Control			
Base Address	Address Range	R/W	Description
0x00000000	0x00 - 0xFC	R/W	Data Memory
0x00000800	0x00 - 0x0C	R/W	Factorial Accelerator
0x00000900	0x00 - 0x0C	R/W	General Purpose I/O

Table1.2 – Memory-mapped I/O design table

The following *Table 1.3* explains each module used in the SoC design in detail.

Modules	Description
Instruction Memory	Holds machine code from assembled source program. The MIPS processor retrieves an instruction from this

	memory module one at a time. This is a read-only module.
MIPS Processor	Either the single-cycle MIPS processor or pipelined MIPS processor (pipelined processor is used here in this assignment). This module processes all instructions and communicates with the connected I/O devices.
Data Memory	An I/O device that stores data when the MIPS processor issues a <i>lw</i> or <i>sw</i> instruction with the address range 0x000 - 0x7FF. This module is read or write module.
Factorial Accelerator	An I/O device that calculates $n!$ given an input n and a Go signal. The device provides the result and status bits Done and Error via a write back data bus. This I/O device is only active when the MIPS processor issues a <i>lw</i> or <i>sw</i> instruction with the address range 0x800 - 0x80C. This module is read or write module.
General Purpose I/O	An I/O device that interfaces with connected input and output devices on external FPGA devices. The device takes two 32-bit inputs and outputs two 32 bit values and can deliver data to the MIPS processor via a writeback data data bus. This I/O device is only active when the MIPS processor issues a <i>lw</i> or <i>sw</i> instruction with the address range 0x900 - 0x90C. This module is read or write module.
Address Decoder	A module that contains simple logic to ensure that only a specific I/O device is active and writing data back to the MIPS processor. It receives an address from the MIPS processor and issues write enable signals to the respective I/O devices and controls the output of a connected 4-input multiplexer that is connected to the input read data port of the processor.
4-Input Mux	A module that controls what data is written back to the MIPS processor. It takes data from the data memory, factorial accelerator, and GPIO unit as inputs.

Table 1.3 – Explanation of all SoC modules

2) Factorial Accelerator with Interface Wrapper:

The factorial accelerator used in this assignment was based on the design of the factorial accelerator from Assignment 1. That is, the factorial module is designed from the bottom-up based on the following algorithm:

INPUT n


```

product = 1
WHILE (n > 1) {
    product = product * n
    n = n - 1
}
OUTPUT product

```

As shown in *Figure 1.4*, our Factorial module takes an input n and input Go signal to begin the factorial calculation. The control unit drives the calculation of $n!$ inside the datapath. The *product* or result of the calculation is outputted along with a *Done* bit. If the input n is greater than 12, the factorial calculation does not take place. Instead, the datapath outputs an *Error* bit to the control unit which then outputs an *Error* bit. This was essential to Assignment 8 as our single-cycle CPU and pipelined CPU would rely on these status bits to find out when the factorial accelerator was done processing with the calculation. Therefore, as our factorial accelerator provided everything we needed, we opted to not change the design. *Figure 1.5* and *1.6* shows the design of the datapath and the ASM chart of the control unit for the factorial accelerator.

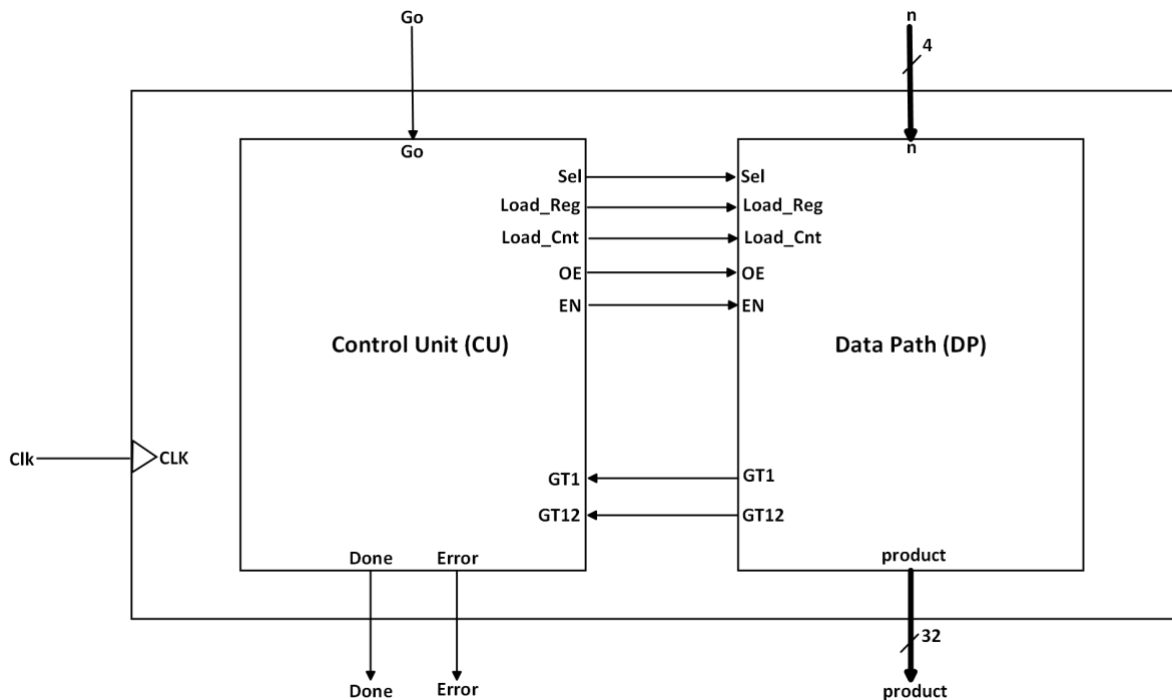


Figure 1.4 – Factorial Accelerator CU-DP diagram

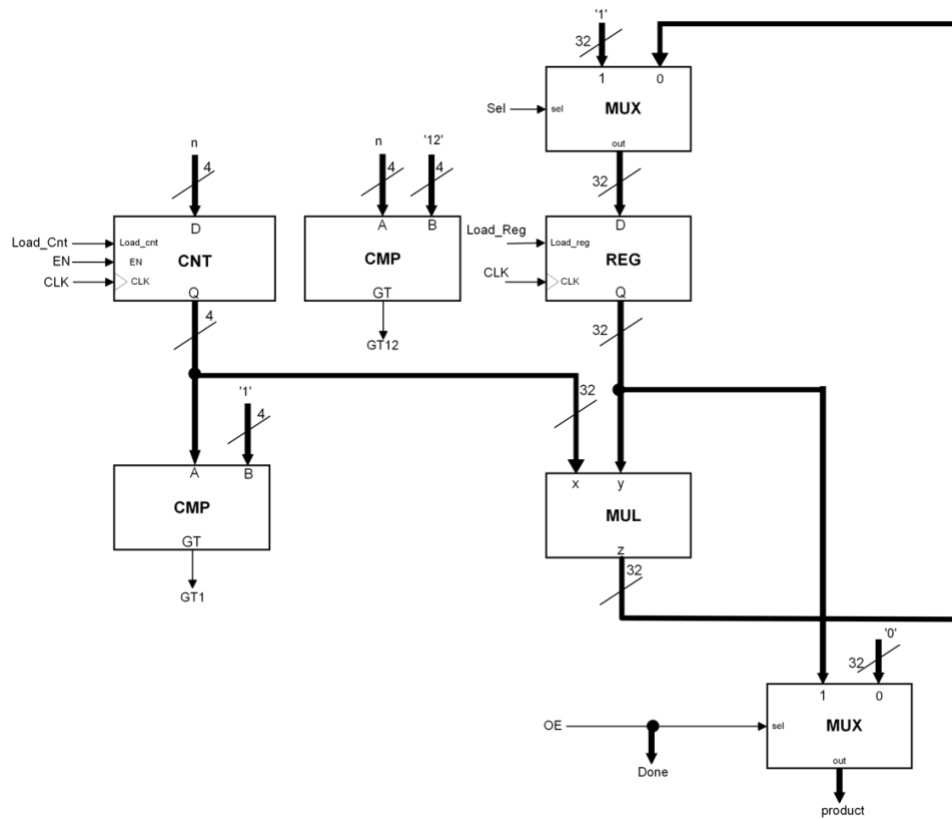


Figure1.5 – Factorial Accelerator Datapath diagram

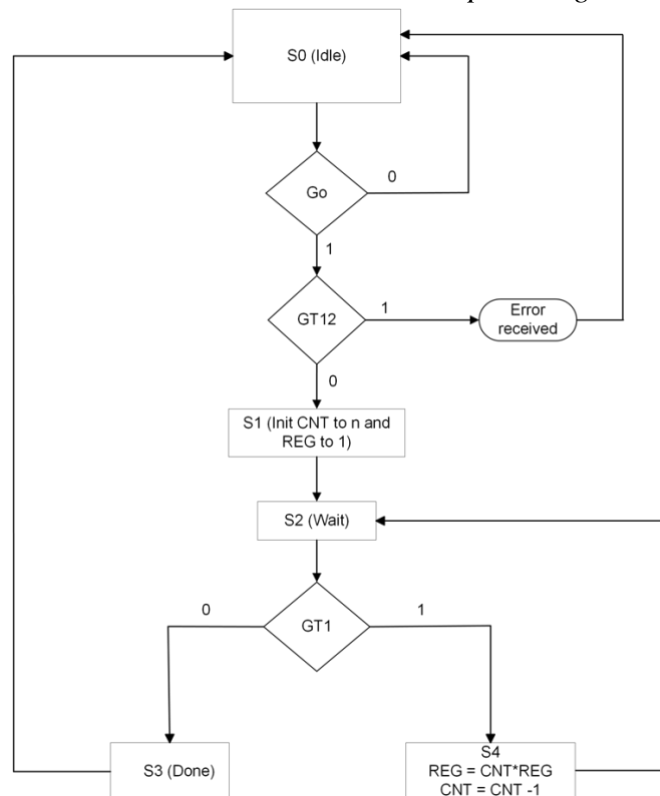


Figure1.6 – Factorial Accelerator ASM chart

We could not simply connect our factorial module to the pipelined CPU. The factorial accelerator would need a way to know when it was being addressed by a CPU. It would also need something to convert address values from the CPUs to something the factorial module could understand, i.e., an input 'n' or a Go signal to kick off the calculation. Therefore, we created an interface wrapper for the factorial unit. *Figure 1.7* shows the design of this interface wrapper.

The design of this interface wrapper was based on a design given to the class during the IO interface lecture. It works by accepting bits 3 and 2 from the connected address bus from the MIPS processor. Bits 3 and 2 of the input address were selected as these are the most important bits when the CPU addresses the factorial accelerator. This is because the program counter inside the MIPS processor always changes by a factor of 4. The binary value of 4 is 100, the binary value of 8 is 1000, and the binary value of 32 is 100000. As shown, the rightmost bits of all numbers of factor 4 end in 00 in binary. Therefore, it's safe to disregard these bits and only worry about bits 3 and 2 as they are the only bits that change in numbers of factor 4. This design principle is applied in all the other I/O devices attached to the MIPS processor.

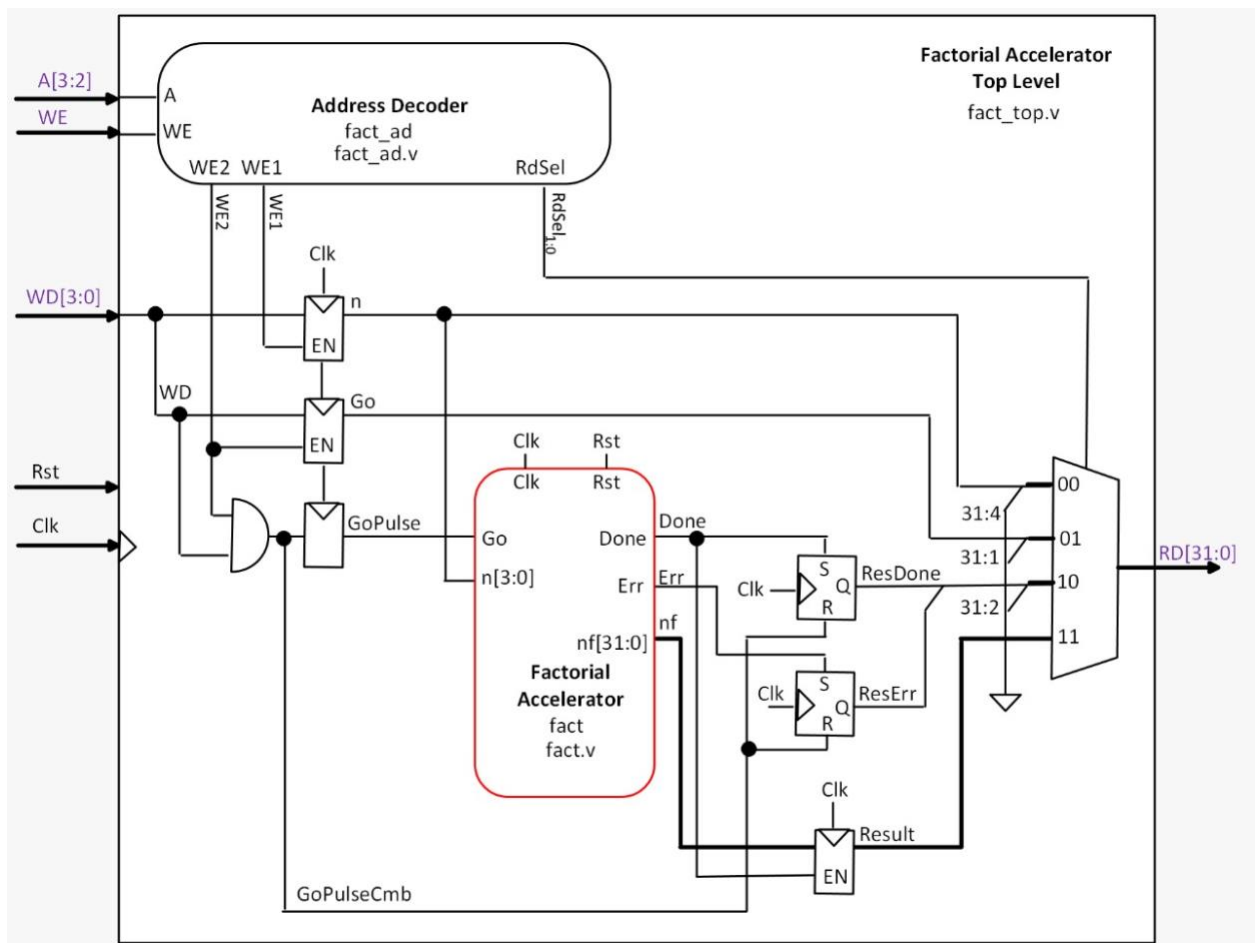


Figure 1.7 – Interface Wrapper with Factorial Accelerator

The interface wrapper also accepts a write enable (WE) signal from the address decoder inside the SoC. Additionally, the factorial accelerator wrapper also takes an input ‘n’ given by the first four bits of the attached write data (WD) bus from the MIPS processor. In this way, whenever the MIPS processor processes a *lw* or *sw* instruction with an address between 0x800 and 0x80C, it will address the factorial accelerator interface. *Table 1.8* below shows the relevant address bits (A[3:2]) that determine what happens inside the interface. For example, if the MIPS processor processes a *sw \$t1, 0x800(0)*, the processor will address the factorial accelerator and give it the data in \$t1 as the input ‘n’. The address decoder in the SoC will also issue WE to enable the interface.

Inside the interface exists another address decoder, the factorial accelerator from Assignment 1, a few registers to hold input data and output data from the factorial accelerator, and a 4-input mux to control output data from the interface. The basic idea is that registers must hold input data from the MIPS processor (such as the Go signal or ‘n’) so that the MIPS processor can do other things while the factorial accelerator performs a calculation. Output registers must hold the result data from the factorial accelerator until the MIPS processor is ready to read the data with a *lw* instruction. The address decoder has simple logic to issue write enable signals to the relevant registers depending on the function the MIPS processor is telling the interface to perform. For example, in a *sw \$t1, 0x804(0)*, the address decoder must enable a Go register and a GoPulse register to hold the input Go signal from the MIPS processor. Registers hold input data *n* and *Go* and output data *Done*, *Error*, and *Result*. The 4-input mux has a select input controlled by the address decoder. Output data from the interface wrapper is shown in *Table 1.8*.

Factorial Accelerator Memory Map							
Address	R/W	Register Name	Bits	Bit Definition	Decoder Output		
					WE1	WE2	RdSel
00	R/W	Data Input (n)	31:4	Unused			
			3:0	n[3:0]	1	0	00
01	R/W	Control Input (n)	31:1	Unused			
			0	Go Bit	0	1	01
10	R	Control Output (Done, Err)	31:2	Unused			
			1	Err Bit	0	0	10
			0	Done Bit	0	0	10
11	R	Data Output (Result)	31:0	nf[31:0]	0	0	11

Table 1.8 – Memory map table for the Factorial Accelerator Wrapper

The following *Table 1.9* explains each module used in the Factorial Accelerator Interface Wrapper design in detail.

Modules	Description
Address Decoder	This module controls which register is addressed by the MIPS processor. It issues write enable signals to the respective registers inside the interface wrapper. It also controls which register outputs its data to the read data bus.
Factorial Accelerator	This module performs the calculation of $n!$ given an input 'n' and a Go signal. The result and status bits are stored in status registers until the MIPS processor is ready to read the data.
n Register	This module holds input 'n' data until the MIPS processor instructs the factorial accelerator to begin with the Go signal.
Go Register	This module holds the input Go bit if the MIPS processor ever wants to read the stored Go bit.
GoPulse Register	This module provides the Go pulse necessary to begin the calculation of $n!$. $\text{GoPulse} = 1$ when both the address decoder issues a write enable signal for this register and the MIPS processor issues the Go bit with a <code>sw \$s, 0x804(\$0)</code> instruction.
ResDone Register	This module stores the Done bit output of the factorial accelerator until ready to be read by the MIPS processor.
ResError Register	This module stores the Error bit output of the factorial accelerator until ready to be read by the MIPS processor.
4-Input Mux	This module controls which data is output from the interface back to the MIPS processor. Options include input 'n', the Go bit, the status bits Done and Error, and the result of $n!$

Table 1.9 – Explanation of all modules in the Factorial Accelerator Wrapper

3) General Purpose I/O (GPIO) with Interface Wrapper:

The design of the GPIO unit mirrors that of the design of the factorial accelerator interface wrapper. The GPIO unit is responsible for interfacing with external input and output

device. When addressed by the MIPS processor, the GPIO unit will either read an input from its input gpI1 or gpI2 input ports or write to its output ports gpO1 or gpO2. The design of this interface is shown in *Figure 1.10*.

Inside the interface wrapper exists another address decoder like the one in the SoC and the factorial accelerator interface wrapper. The address decoder takes bits 3 and 2 from the connected address bus from the MIPS processor. Like the factorial accelerator address decoder, only bits 3 and 2 are relevant to the operation of the GPIO unit. A write enable signal also ensures the GPIO unit is only active when enabled by the correct address range 0x900 - 0x90C. Additionally, two registers hold input data from the write data to data memory bus from the MIPS processor. These two registers only store new data when given a write enable signal by the address decoder within the GPIO interface wrapper. In this way, one register stores data input from the MIPS processor to be output to gpO1 port while the other register holds input data from the MIPS processor to output to the gpO2 port. Furthermore, this design principle enables each register to be accessed by a specific address from the MIPS processor. These two registers were designed to hold the output of the factorial accelerator's calculation and the status bits from the factorial accelerator. For example, gpO2 holds the computation result, while gpO1 holds the status bits to output such as error bit. Finally, a 4-input mux like as in the SoC and factorial accelerator controls the output of the GPIO unit. It's select port is controlled by the address decoder. *Table 1.11* shows what happens inside the GPIO unit when it is addressed by the MIPS processor. It also shows what data is output on the read data bus from the GPIO unit.

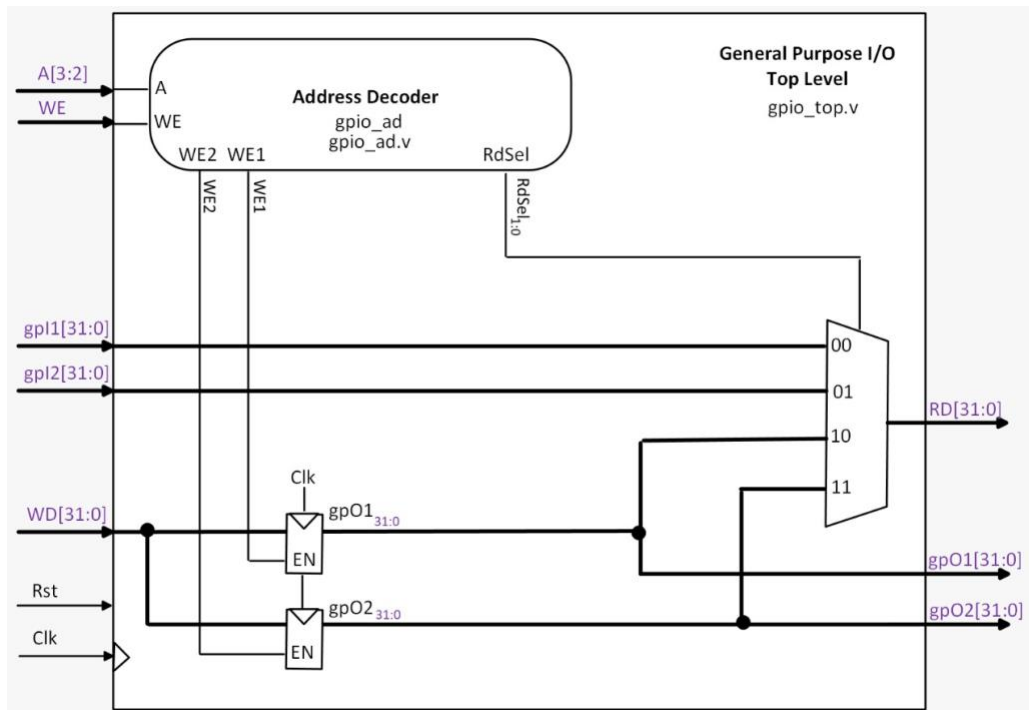


Figure 1.10 – Schematic of the Interface Wrapper for the GPIO Unit

GPIO Memory Map						
Address	R/W	Register Name	Bits	Decoder Output		
				Bit Definition	WE2	Rdsel
00	R	Input1	31:0	General Input	0	00
01	R	Input2	31:0	General Input	0	01
10	R/W	Output1	31:0	General Output	0	10
11	R/W	Output2	31:0	General Output	1	11

Table1.11 – Memory map table for the GPIO unit Interface Wrapper

The following *Table 1.12* explains each module used in the GPIO Interface Wrapper design in detail.

Modules	Description
Address Decoder	This module controls which register is addressed by the MIPS processor. It issues write enable signals to the respective registers inside the interface wrapper. It also controls which register outputs its data to the read data bus.
gpO1 Reg	This register holds input data from the write data bus from the MIPS processor. Data from this register is output via the gpO1 port to an external device outside of the SoC.
gpO2 Reg	This register holds input data from the write data bus from the MIPS processor. Data from this register is output via the gpO2 port to an external device outside of the SoC.
4-Input Mux	This input controls what data is written back to the MIPS processor. Options include input data on gpI1, gpI2 or output data from the gpO1 or gpO2 register. The select signal for this mux is given by the address decoder inside this interface.

Table1.12 – Explanation of all modules in the GPIO Interface Wrapper

4) Pipelined MIPS Processor Design:

Designing the pipelined MIPS processor proved to be the most challenging task. We spent a lot of time grouping the respective modules by the following stages: fetch, decode, execute, memory, and writeback. Each stage processes part of one instruction per clock cycle, forwarding relevant data necessary for the next step of execution via a state register to the next stage. For example, in an *add* instruction, the instruction must be fetched in the fetch stage, the operands must be fetched in the decode stage, the addition must occur

in the execute stage, the data must pass through the memory stage, and finally the result must be written back to the destination register in the writeback stage.

In the fetch stage, the MIPS processor tries to fetch an instruction from the instruction memory. Because instruction memory is external to the CPU, it must calculate the program counter value and output this address to the instruction memory to retrieve the next instruction. Therefore, the fetch stage contains several multiplexers and an adder to calculate program counter + 4. This adder is responsible for calculating the next instruction to retrieve. A series of multiplexers chooses which address is output to the instruction memory. They choose between the program counter + 4 address, branch address, or jump address, or jump return address. For example, a *j target* instruction will send the address of the instruction to jump to the instruction memory. Once the instruction memory releases an instruction back to the MIPS processor, the instruction is stored in a state register. Other data stored in the state register includes the program counter + 4 address. This data is stored in this fetch-to-decode state register until the next clock cycle to forward information to the decode stage.

In the decode stage, the MIPS processor retrieves the relevant operand data from the register addresses encoded in the machine code. The machine code comes from the instruction now stored in the fetch-to-decode state register. In this stage, the control unit also generates relevant control signals that are forwarded through the pipeline until they are necessary. Adhering to the principles of pipelining, data retrieved from the specified registers is clocked into the decode-to-execute state register to be processed on the next clock cycle. The control signals determined by the machine code inside the control unit are also stored in the same state register for the execute stage as well as the writeback address of the destination register. This address is also encoded in the machine code and must be forwarded along with the operands of this instruction (and eventually the result of the execution). The actual address given in the machine code changes depending on the type of instruction being executed. Therefore, a multiplexer chooses between bits 20:16 for R-Type instructions and bits 15:11 for all other instructions. A final multiplexer chooses between the previous multiplexer's out register address or register 31 for *jal* instructions. This writeback address is stored along with all other data in the decode-to-execute state register. Finally, to compute branch control value and to reduce the branch overhead cycles, the logic for branch instruction is moved to the decode stage using a comparator between two read data outputs of the register file. This comparator value is then anded with the branch control signal to determine the next program counter value. Note that, during all this, branch target address will be calculated and ready.

In the execute stage, the actual processing of an execution takes place. For example, an *add* instruction's operands are added inside the ALU and the operands of a *multu* instruction are multiplied here. Operand data for each instruction is output from the decode-to-execute state register to the ALU, and Multiplier modules. The ALU processes *and*, *or*, *add*, *sub*, and *slt* instructions. The Multiplier processes *multu* instructions. The *slr* and *sll* instructions will be taken care by the ALU itself. A multiplexer chooses between a second operand from the register file and a sign-extended immediate value to forward as the second operand for the ALU. The sign-extended immediate value comes from a SEXT module in

the decode stage that extends a 16-bit immediate value given in the machine code (*addi* instructions) into a 32-bit value. This value also comes through the same state register. The ALU takes a “shamt” input given in the machine code. Therefore, the machine code is also forwarded through the pipeline for use in the execute stage. The output of all these modules is stored in another state register, execute-to-memory. Control signals needed for later in the pipeline are also forwarded from the previous decode-to-execute state register to the execute-to-memory state register. Refer to Appendix to see all control signals that are forwarded through each state register until needed. Data from previous stages that were not operated upon in this stage but are needed later are also forwarded through the execute stage and the state registers. This data includes the destination register address, data from source register 1 (for *jr* instructions), data from source register 2 to write to data memory or an I/O device, and program counter + 4.

In the memory stage, data is either output from the MIPS processor to memory or an I/O device or read from memory or an I/O device. This only occurs if a *sw* or *lw* instruction is being processed. Although the diagram in Appendix shows the data memory within the datapath, it’s important to note that memory is actually external to the CPU. The following signals are outputs of the processor: *alu_outM* (Address) , *wd_dmM* (WriteData), and *we_dmM* (MemWrite). *rd_dmM* (ReadData) is an input to the processor. The signals in parentheses correspond with the signals in *Figure 1.1* of the System-on-Chip Design section. Input data from the SoC (ReadData) is pipelined in the memory-to-writeback state register. Aside from memory, data from the *multu* instruction is also stored to a high and low register (*HiLo_reg*) in this stage. In an earlier iteration of this design, we placed these registers in the decode stage. Recall from Assignment 6 that the *hi* register holds the upper 32 bits of the result of the multiplication while *low* register holds the lower 32 bits. Data from these registers are also stored in the memory-to-writeback state register. Control signals not needed in this stage are also stored in the memory-to-writeback state register for use later. Data from previous stages that were not operated upon in this stage but are needed later are also forwarded through the memory stage and the state registers. This data includes the output of the ALU to be stored in the register file, data from source register 1 (for *jr* instructions), program counter + 4, and the machine code from the fetch stage.

Finally, in the writeback stage, data is written back to the register file. R-Type instructions primarily use this stage. This stage is made up of several buses connected to previous portions in the pipeline and several multiplexers that control what data is written back to the register file. A 4-input multiplexer is chosen to transfer the data from the high register, low register, ALU output, and input from data memory or I/O devices to writeback depending on the instruction being processed. For example, *mflo \$t1*, will forward data from the low register through this multiplexer. Additionally, several buses are directly connected from the memory-to-writeback state register back to various ports in the datapath. These include the register destination address that is given to the register file with the data to be written back in an R-Type or *jal* instruction and the jump address given by bits 25:0 of the machine code.

In each stage, certain modules require control signals. As mentioned before, control signals are created in the decode stage based on the input machine code. All control signals are pipelined and stored in state registers until needed. The control signals are unchanged from Assignment 6 and are only grouped by their respective stage in Appendix. These design choices were motivated by the following thought process: control signals should travel through the pipeline with the instruction until a module needs that control signal. In this way, processing of an instruction flows smoothly through the pipeline and does not have to wait for control signals as they are readily available.

As architecture diagram and control unit truth tables for the design of this pipelined CPU are too large to insert here, they have been added to Appendix. The explanation of the pipeline makes more sense as you follow along with the diagram of the datapath. Signals are named according to the state they originate in. For example, “instr” is a 32 bit bus that travels through all state registers. In each stage, a letter for the name of that stage is appended to the name of the signal: “instrF” for fetch, “instrD” for decode, “instrE” for execute, “instrM” for memory, and “instrW” for writeback. This naming convention applies to nearly all signals and buses in the datapath.

Because of the design of our datapath, few NOPs (no operation MIPS instruction) were added to the driver assembly code. NOPs were added whenever a data dependency was detected. That is, if an instruction saved data to a register that the following instruction required for its operation, we added NOPs to allow time for the data to flow through the entire pipeline and writeback to the register file. Because it takes a total of 3 clock cycles to travel from the decode stage to the writeback stage, we added 3 NOPs in cases of data dependencies. NOPs were also added for jump and branch instructions for the reasons described earlier in this paragraph. Time was needed for the address to jump or branch to be given back to the fetch stage. The modified code for this driver assembly code can be found in the Appendix. A Gantt chart showing the result of our pipelining efforts is present below in *Figure 1.13*. The diagram shows all instructions from the modified MIPS driver code with few NOPs inserted for instructions with data dependencies.



Figure 1.13 – Gantt Chart of the Pipelined MIPS processor

5) Pipelined MIPS Processor with Data Forwarding:

To improve the performance of our pipelined MIPS processor, we attempted to implement a form of data forwarding. This would solve situations where instructions required data from registers that were not yet updated from the previous instruction. Essentially, we were trying to solve a data dependency hazard. Unfortunately, we could not get a working design complete before the demonstration deadline. The basic idea is that if an instruction requires data that has not been written back to the register file, but the data is present somewhere in the pipeline, then that data should be forwarded back to the execution stage to be used in execution. Therefore, there are two paths to retrieve data from the pipeline: a path from the memory stage and a path from the writeback stage. The logic equations are as follows:

Forwarding from memory stage:

if (we_regM and (instrE[25:21] != 0) and (rf_waM == instrE[25:21])) then ForwardAE = 10
if (we_regM and (instrE[20:16] != 0) and rf_waM == instrE[20:16])) then ForwardBE = 10

Forwarding from writeback stage

if (we_regW and (instrE[25:21] != 0) and (rf_waW == instrE[25:21])) then ForwardAE = 01
if (we_regW and (instrE[20:16] != 0) and (rf_waW == instrE[20:16])) then ForwardBE = 01

ForwardAE and ForwardBE are control signals from a data forwarding logic unit. These control signals are fed to two 3-input multiplexer that controls which data is given to the ALU. One mux controls data for input port A of the ALU while the other controls data for port B of the ALU. The ForwardAE mux chooses between data from a register in the register file, the output of the ALU from the memory stage, or the writeback data from the writeback data. The ForwardBE mux chooses between the same input data as ForwardAE except with the second register output port of the register file. It's important to note that no ports from the datapath modules are directly connected to these mux, only ports from the state registers. As data forwarding was not a core requirement of the assignment, it is not shown in the Waveforms section. To reiterate the above explanation, a figure from the lecture slides has been referenced below in *Figure 1.14*.

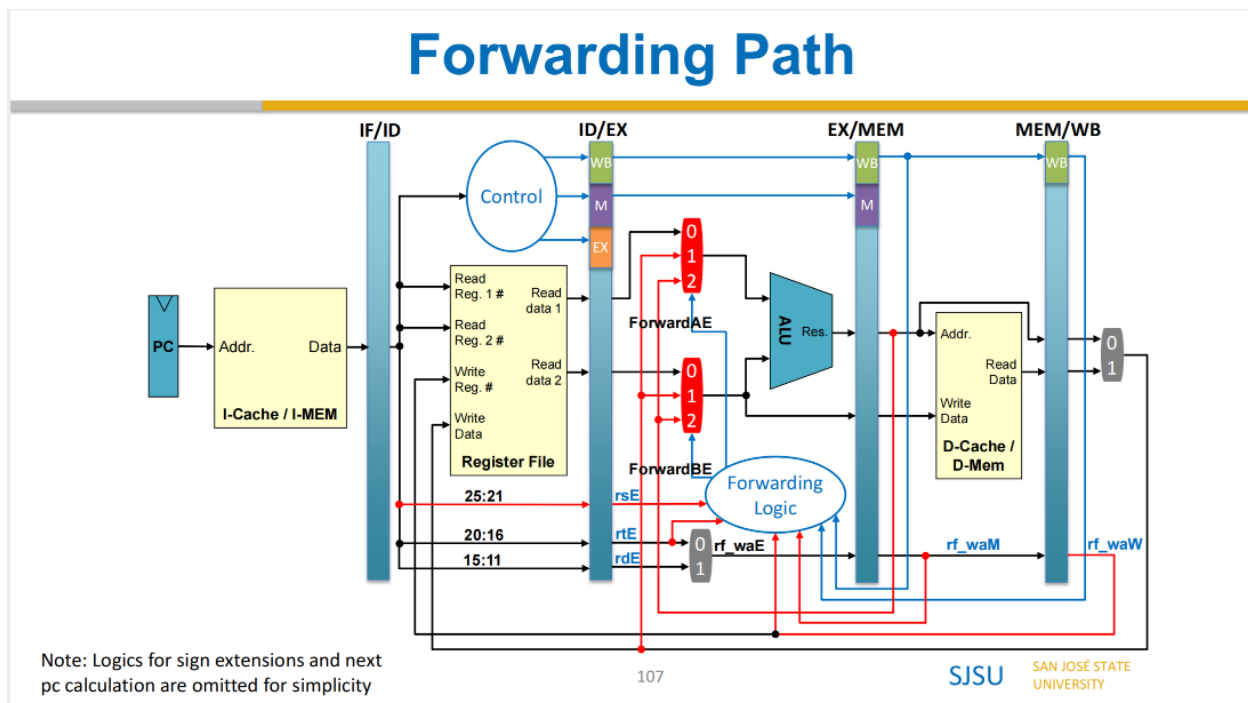


Figure 1.14 – Data Forwarding using the Pipelined MIPS processor

6) Performance Analysis:

The performance analysis involved the analysis of the pipelined MIPS processor design against the single-cycle (recursive) MIPS processor design for the factorial input of n . The input values of n included a range from 1 to 12 for computing the predicted results of $n!$ of each design as depicted in Figure 1.15. With each transition to a new state, an additional cycle occurred. The runtime was calculated for the factorial accelerator by totaling the number of cycles for values of 1 to 12, equation (1) was used respectively. For the single-cycle MIPS design, the recursive version was assembled in the MIPS IDE. The instruction counter tool was relied upon to count the number of instructions executed for each input value. The total number of executed instructions totaled to the runtime of the single-cycle design. Overall, the total runtime for the single-cycle MIPS design was greater than the total runtime of the factorial accelerator for each factorial input of N . The total runtime for the factorial accelerator and the single-cycle MIPS design was computed for each factorial input of N and depicted in Figure 1.15 using an excel sheet. The below equations are calculated by first setting up a few data points and then using interpolations techniques to get the approximated equations.

Formula for Factorial Accelerator

$$y = 2n+2$$

Formula for Single-Cycle MIPS

$$y = -3.5 + 14 \cdot x$$

The Figure 1.15 below shows the graphical representation of the runtime cycles taken by the single-cycle processor to perform the factorial computation and runtime cycles taken by the factorial accelerator to do the same.

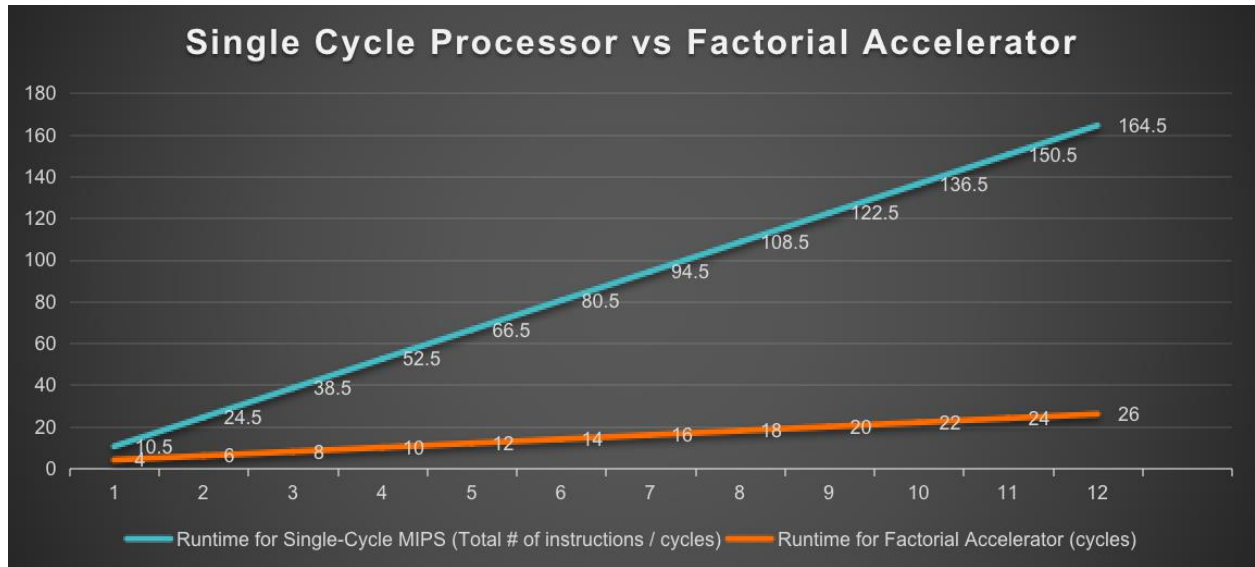


Figure 1.15 – Runtime cycles taken by each for the factorial computation

Similar procedure has been done to get the runtime cycles taken by the pipelined MIPS processor. The *Figure 1.16* below shows the graphical representation of the runtime cycles taken by the single-cycle processor to perform the factorial computation and runtime cycles taken by the factorial accelerator to do the same. (Note the scale change in the below graph)

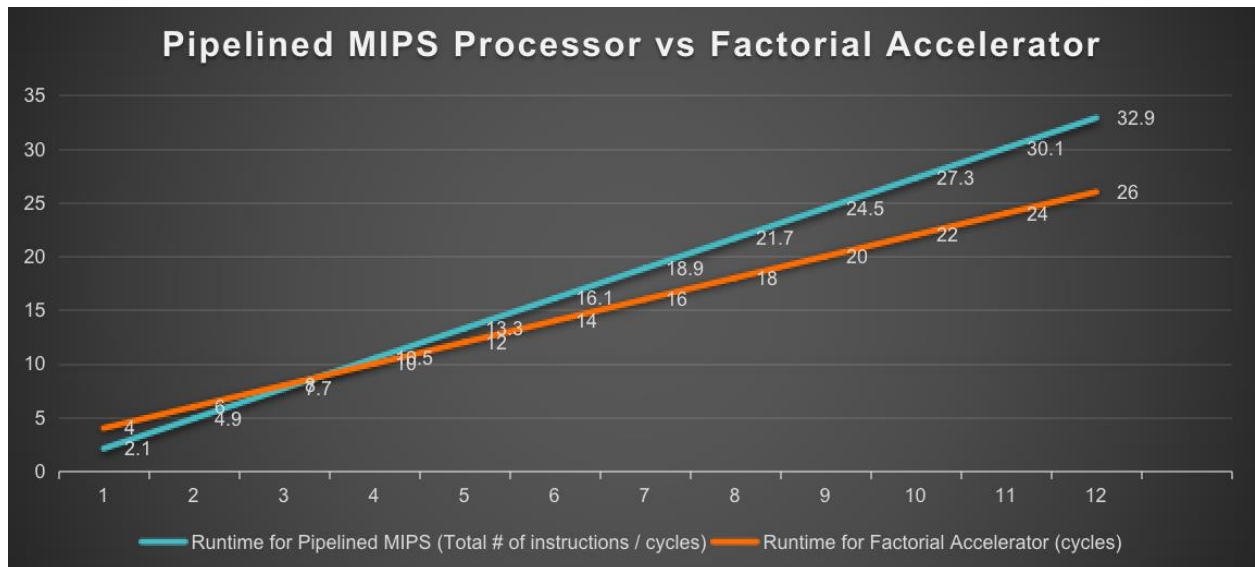


Figure 1.16 – Runtime cycles taken by each for the factorial computation

7) Simulation Waveforms (with Unit-Level simulations):

To functionally verify the designs of our factorial accelerator with interface wrapper, GPIO unit with interface wrapper, and pipelined SoC, me and my teammate wrote testbenches that stimulate the input of the respective top module and verified the outputs

were as expected. In cases where the design did not produce the expected results, we debugged our design by probing more signals to find errors in our logic. This section is divided by the testbenches we created to test each major component of our SoC.

a) Testbench for the Factorial Accelerator Interface Wrapper:

Testing the factorial accelerator with an interface wrapper involved writing a testbench that tested values for n from 0 to 13. This would ensure that our design was tested with edge cases where $n = 0, 12$, with normal values, and where $n > 12$. Our testbench instantiation of the factorial accelerator with interface wrapper as the design under test can be seen below. It then iterates from 0 to 13. In each loop iteration, the write data port (WD) is given an input based on the iteration number (i.e. $n = i = 0, 1, 2, \dots$). Then the address input to the module is changed through all possible values as given in *Table 1.8* with a clock tick between each change. In this way, we examine all values of n . The following waveforms show verification of our design. This testbench is an eyeballing testbench where we must examine the waveforms to verify the design.



Figure1.17 – Waveform of $n=4$ for the Factorial Accelerator Wrapper

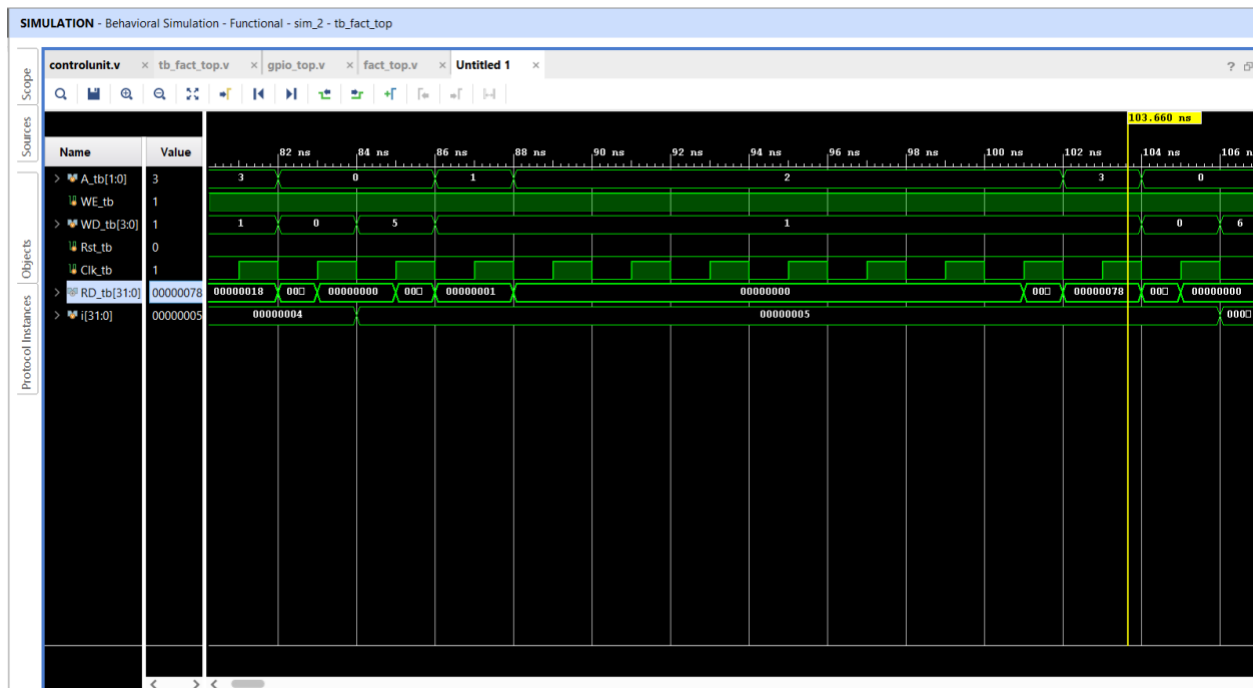


Figure 1.18 – Waveform of $n=5$ for the Factorial Accelerator Wrapper

As shown in the figures 1.17 and 1.18, RD_tb shows the output of the factorial accelerator interface wrapper. When $A = 0$, WD_tb gives $n = 4$ to the wrapper, which is stored in a register. When $A = 1$, the Go bit is given to the wrapper, which is stored in a register. When $A = 2$, the calculation takes place until RD_tb shows 1 indicating the Done bit is set. Finally, when $A = 3$, the result is shown on RD_tb as $0x18_H$ indicating the correct output of $4! = 24_d$.

b) Testbench for the GPIO Unit Interface Wrapper:

Testing of the GPIO unit with interface wrapper involved writing a testbench that tested the inputs of the GPIO interface module and viewing the outputs with an eyeball test. As the design of the GPIO interface wrapper was much simpler than the design of the factorial accelerator with interface wrapper, it follows that the testing of GPIO unit with interface wrapper was much simpler. To test, we gave the gpI1 and gpI2 inputs a value and examined the outputs of RD as we changed the input address (A) between all possible values as shown in Table 1.11. We also examined the gpO1 and gpO2 outputs. The waveform of this functional verification is shown in Figure 1.19.

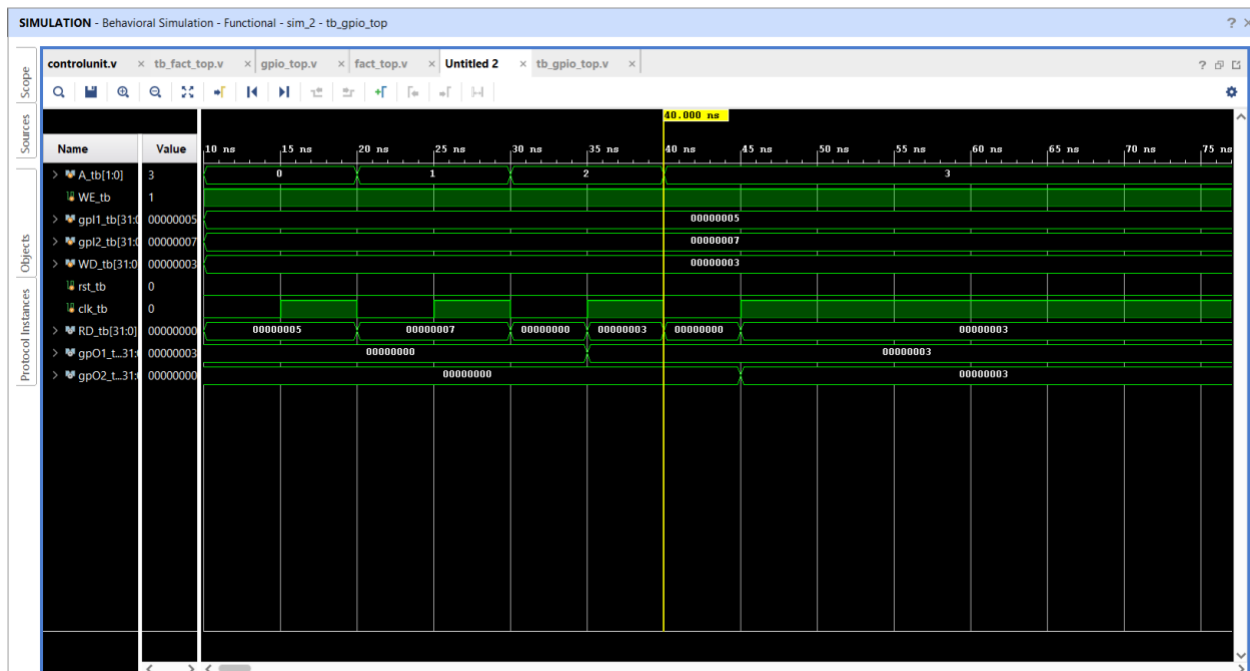


Figure 1.19 – Waveform of GPIO Unit Interface Wrapper after testing

To safely confirm that our design works as intended, RD must output the following: the value of gpI1 when A = 0, the value of gpI2 when A = 1, the value of WD when A = 2, the value of WD when A = 3. Furthermore, gpO1 must output the value of WD after A = 2 and gpO2 must output the value of WD after A = 3. As seen in the waveform output, this holds true. Note that output data doesn't appear right after A changes. This is because the registers internal to the interface wrapper do not store input values until given a write enable signal determined by the value of A and a high clock pulse.

c) Testbench for the System-on-Chip (SoC) Address Decoder:

Testing of the SoC address decoder involved writing a testbench which takes address as input and gives out the outputs (RdSel) with an eyeball test. This decoder was also relatively simple to design and test. To test, we gave the address as input value and examined the outputs of RdSel as we changed the input address (A) between all possible values i.e., (0x0000, 0x0800, 0x0900).

To safely confirm that our design works as intended, RdSel must output the following: the value when A[3:2] = 00 must be decoded to RdSel = 0, the value when A[3:2] = 10 must be decoded to RdSel = 2, the value when A[3:2] = 11 must be decoded to RdSel = 3. The waveform of this functional verification is shown in Figure 1.20.



Figure 1.20 – Waveform of System Address Decoder after testing

d) Testbench for the Pipelined System-on-Chip (SoC) Design:

The functional verification of the pipelined system-on-chip proved to be the most challenging. There were so many modifications to the design of the CPU that it became extremely cumbersome to debug issues as they continued to pop up. Nevertheless, we created a working pipelined MIPS processor. We gave the pipelined SoC an input for n on gpI1 and Go signal on gpI2. We also added and removed new probes to debug the processor along the way. The GPIO output ports (gpO1 and gpO2) were examined for the correct result of the factorial calculation and Error and Done bits, respectively. The memory file for this simulation was lightly modified with NOPs between instructions with data dependencies as described earlier in the Design Methodology section. Execution stops when the processor reaches the final instruction. The following waveforms show the output of the whole SoC with different 'n' values.

The below waveform in *Figure 1.21* output shows the entire execution process of the pipelined MIPS processor. Of note is the input port gpI1 = 0x4 with $n = 4$ and the output port gpO2 with $n! = 0x18 = 24$. Apart from that, the Go signal is being given to the system by the user to start the factorial computation.



Figure1.21 – Waveform of SoC when $n=4$ after testing

The below waveform in *Figure 1.22* output shows the entire execution process of the pipelined MIPS processor. Of note is the input port $gpi1 = 0x05$ with $n = 5$ and the output port $gpO2$ with $n! = 0x78 = 24$. Apart from that, the Go signal is being given to the system by the user to start the factorial computation.

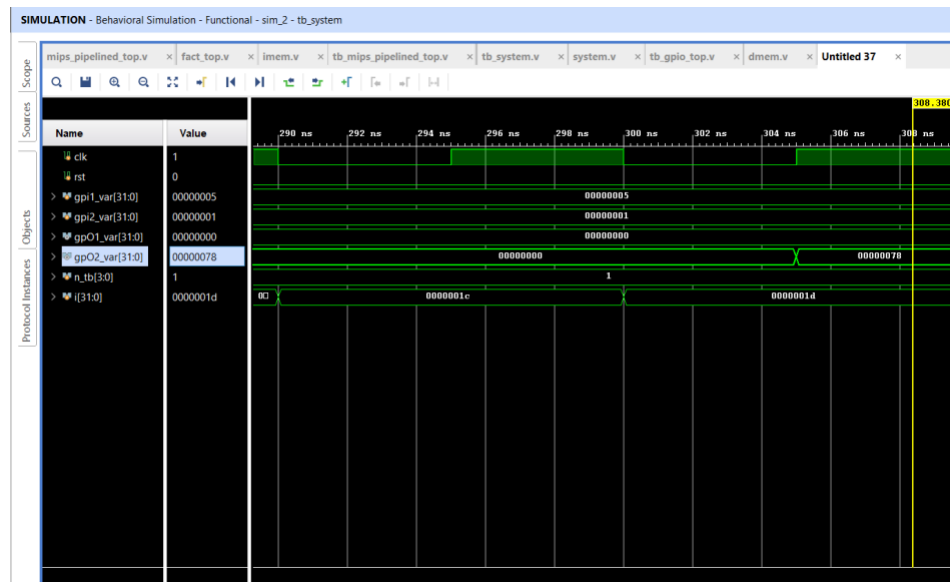


Figure1.22 – Waveform of SoC when $n=5$ after testing

COLLABORATION SECTION:

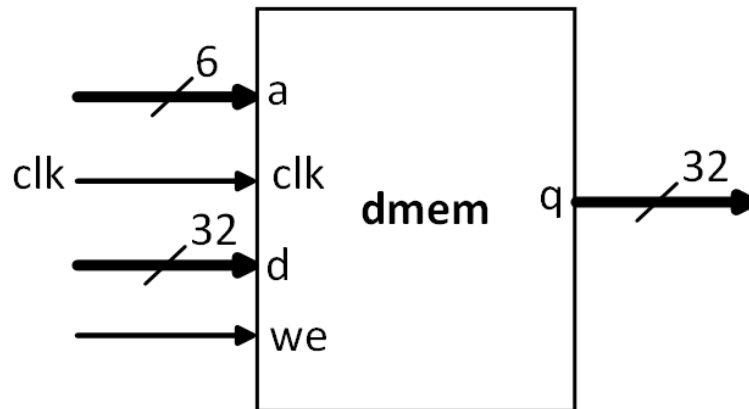
1. Divided the work of individual module design such as Harish worked on the design of Interface Wrappers, and I worked on the design of the Pipelined MIPS processor.
2. Worked together on the pipelined MIPS Datapath block diagram designs using Microsoft Visio tool.
3. Worked together to divide the control unit decoder truth tables into stages and inserted values into them using an excel sheet.
4. By collaborating with each other, started implementing the discussed logic in Vivado and observed the output waveforms after the simulation by writing a testbench for all major components used in the SoC design.
5. Worked together to devise the assembly code to let the pipelined processor call the corresponding interface wrappers and the factorial accelerator.
6. Collaborated in extracting the memfile.dat file from the MARS simulator using Dump to memory option.

CONCLUSION:

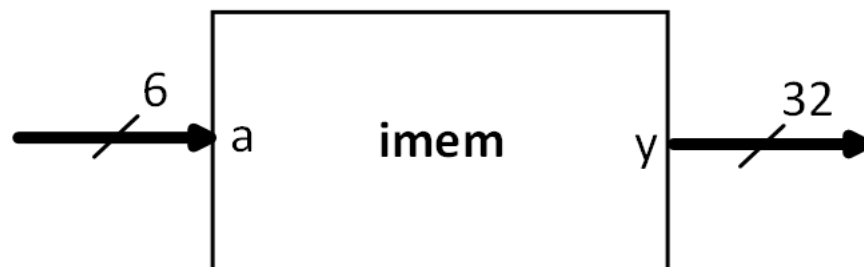
In conclusion, we end the report by converting the fifteen-instruction 32-bit single-cycle MIPS processor into a five-stage, pipelined design. We also were able to interface the pipelined processor with a factorial accelerator and a GPIO unit using the memory-mapped interface registers. Apart from that, we also attempted to use data forwarding technique to overcome few data hazards but were unsuccessful in doing that. A performance analysis is also done between the runtime cycles of single-cycle, pipelined and factorial accelerator while performing the factorial computation using recursive assembly code. Finally, implemented few testbenches to test all the major components in the SoC and completed the overall design of the system. Even-though it was a tough assignment, we stuck to the tasks and the lecture notes with some help from the internet (YouTube) to successfully complete the assignment.

APPENDIX

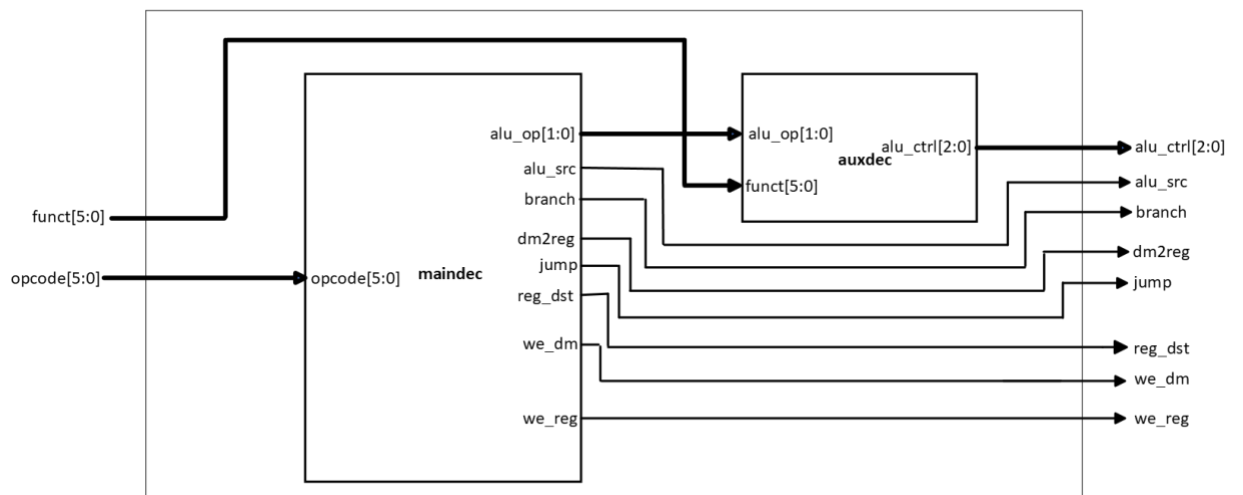
8) Schematic of the Data Memory module used in the pipelined processor:



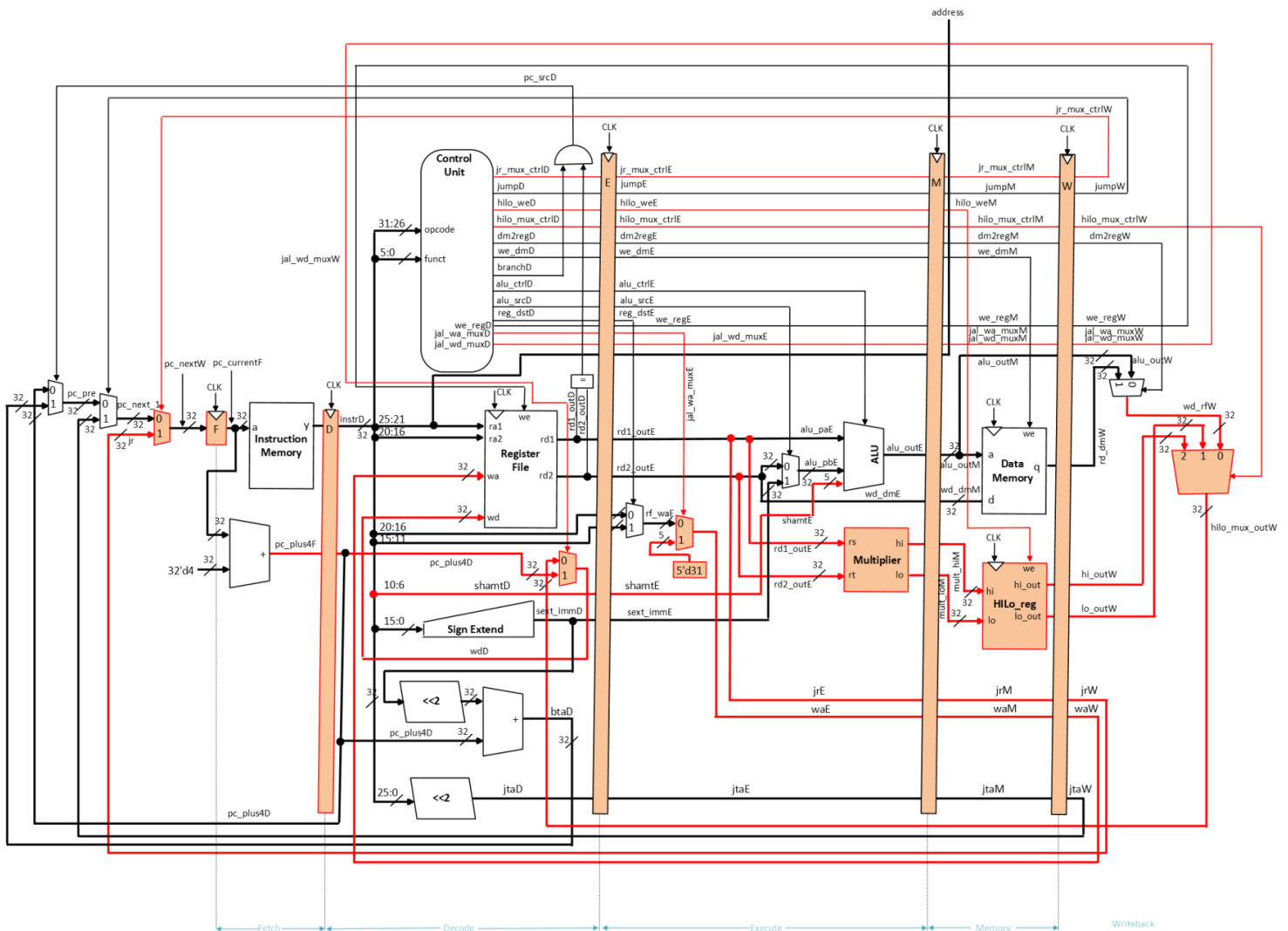
9) Schematic of the Instruction Memory module used in the pipelined processor:



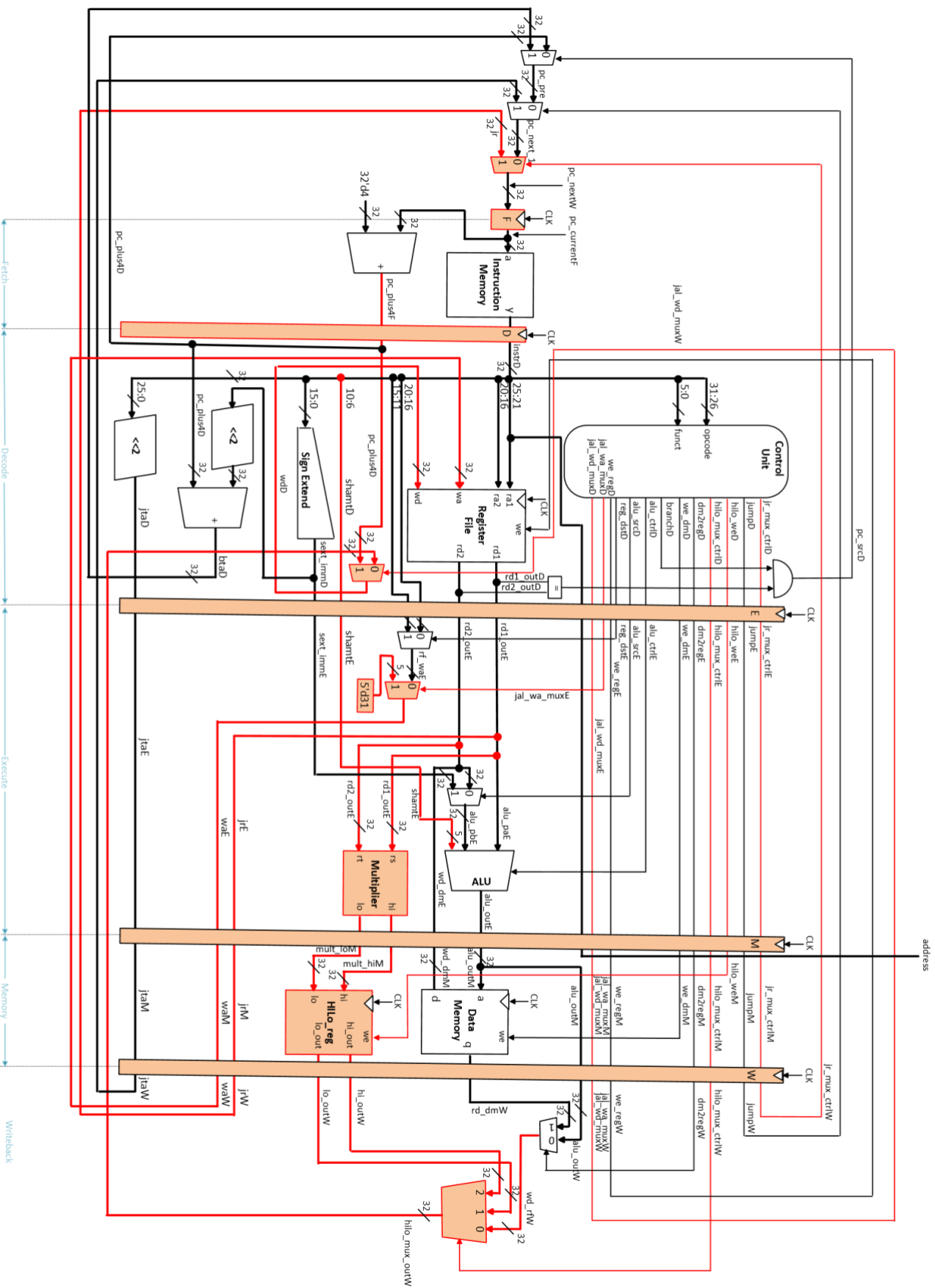
10) Schematic of the Control unit used in the pipelined MIPS processor:



11) Datapath of the pipelined MIPS processor:



Final schematic of pipelined MIPS processor. Modules in red are modules added to support new instructions in Assignment 7. Vertical rectangle modules are pipeline registers which are shown in the above diagram with yellow fill. A larger, horizontal view is shown in the next page.



12) Control unit Truth table for the pipelined processor (Decode Stage):

Type	Instruction	Input		Decode
		<i>opcode [31:26]</i>	<i>funct [5:0]</i>	<i>branchD</i>
R	MULTU	00_0000	01_1000	0
	MFHI	00_0000	01_0000	0
	MFLO	00_0000	01_0010	0
	JR	00_0000	00_1000	0
	SLL	00_0000	00_0000	0
	SRL	00_0000	00_0010	0
J	JAL	00_0011	X	0
R	ADD	00_0000	10_0000	0
	SUB	00_0000	10_0010	0
	AND	00_0000	10_0100	0
	OR	00_0000	10_0101	0
	SLT	00_0000	10_1010	0
I	LW	10_0011	X	0
	SW	10_1011	X	0
	BEQ	00_0100	X	1
	ADDI	00_1000	X	0
J	J	00_0010	X	0

13) Control unit Truth table for the pipelined processor (Execute Stage):

Type	Instruction	Input		Execute				
		<i>opcode [31:26]</i>	<i>funct [5:0]</i>	<i>alu_ctrl</i> <i>E</i>	<i>alu_op</i>	<i>alu_srcE</i>	<i>reg_dst</i> <i>E</i>	<i>jal_wa_muxE</i>
R	MULTU	00_0000	01_1000	xxxx	10	0	1	0
	MFHI	00_0000	01_0000	xxxx	10	0	1	0
	MFLO	00_0000	01_0010	xxxx	10	0	1	0
	JR	00_0000	00_1000	xxxx	10	0	1	0
	SLL	00_0000	00_0000	1001	10	0	1	0
	SRL	00_0000	00_0010	1010	10	0	1	0
J	JAL	00_0011	X	X	00	0	0	1
R	ADD	00_0000	10_0000	0010	10	0	1	0
	SUB	00_0000	10_0010	0110	10	0	1	0
	AND	00_0000	10_0100	0010	10	0	1	0
	OR	00_0000	10_0101	0001	10	0	1	0
	SLT	00_0000	10_1010	0111	10	0	1	0
I	LW	10_0011	X	X	00	1	0	0
	SW	10_1011	X	X	00	1	0	0
	BEQ	00_0100	X	X	01	0	0	0
	ADDI	00_1000	X	X	00	1	0	0
J	J	00_0010	X	X	00	0	0	0

14) Control unit Truth table for the pipelined processor (Memory Stage):

Type	Instruction	Input		Memory	
		<i>opcode [31:26]</i>	<i>funct [5:0]</i>	<i>hilo_weM</i>	<i>we_dmM</i>
R	MULTU	00_0000	01_1000	1	0
	MFHI	00_0000	01_0000	0	0
	MFLO	00_0000	01_0010	0	0
	JR	00_0000	00_1000	0	0
	SLL	00_0000	00_0000	0	0
	SRL	00_0000	00_0010	0	0
J	JAL	00_0011	X	0	0
R	ADD	00_0000	10_0000	0	0
	SUB	00_0000	10_0010	0	0
	AND	00_0000	10_0100	0	0
	OR	00_0000	10_0101	0	0
	SLT	00_0000	10_1010	0	0
I	LW	10_0011	X	0	0
	SW	10_1011	X	0	1
	BEQ	00_0100	X	0	0
	ADDI	00_1000	X	0	0
J	J	00_0010	X	0	0

15) Control unit Truth table for the pipelined processor (Writeback Stage):

Type	Instr	Input		Writeback					
		<i>opcode [31:26]</i>	<i>funct [5:0]</i>	<i>jr_mux_ctrlW</i>	<i>jumpW</i>	<i>hilo_mux_ctrlW</i>	<i>dm2regW</i>	<i>we_regW</i>	<i>jal_wd_muxW</i>
R	MULTU	00_0000	01_1000	0	0	00	0	1	0
	MFHI	00_0000	01_0000	0	0	11	0	1	0
	MFLO	00_0000	01_0010	0	0	01	0	1	0
	JR	00_0000	00_1000	1	0	00	0	0	0
	SLL	00_0000	00_0000	0	0	00	0	1	0
	SRL	00_0000	00_0010	0	0	00	0	1	0
J	JAL	00_0011	X	0	1	00	0	1	1
R	ADD	00_0000	10_0000	0	0	00	0	1	0
	SUB	00_0000	10_0010	0	0	00	0	1	0
	AND	00_0000	10_0100	0	0	00	0	1	0
	OR	00_0000	10_0101	0	0	00	0	1	0
	SLT	00_0000	10_1010	0	0	00	0	1	0
I	LW	10_0011	X	0	0	00	1	1	0
	SW	10_1011	X	0	0	00	0	0	0
	BEQ	00_0100	X	0	0	00	0	0	0
	ADDI	00_1000	X	0	0	00	0	1	0
J	J	00_0010	X	0	1	00	0	0	0

16) Source Code of tb_system.v:

```
module tb_system;

reg clk;
reg rst;

reg [31:0] gpi1_var;
reg [31:0] gpi2_var;
wire [31:0] gpO1_var;
wire [31:0] gpO2_var;
reg [3:0] n_tb;
integer i;

system DUT(
    .clk(clk),
    .reset(rst),
    .gpi1(gpi1_var),
    .gpi2(gpi2_var),
    .gpO1(gpO1_var),
    .gpO2(gpO2_var)
);

task tick;
    begin
        clk = 1'b0; #5;
        clk = 1'b1; #5;
    end
endtask

task reset;
    begin
        rst = 1'b0; #5;
        rst = 1'b1; #5;
        rst = 1'b0;
    end
endtask

initial begin
    reset;
    n_tb = 4'd1;
```

```

gpi1_var = 32'h4;
gpi2_var = 32'h1;

for(i=0;i<30;i=i+1)
    begin
        tick;
    end

$finish;
end

endmodule

```

17) Source Code of tb_fact_top.v:

```

module tb_fact_top;

    reg    [1:0]  A_tb;
    reg      WE_tb;
    reg    [3:0]  WD_tb;
    reg      Rst_tb;
    reg      Clk_tb;
    wire  [31:0]  RD_tb;
    integer i;

    task tick;
    begin
        Clk_tb = 0; #1;
        Clk_tb = 1; #1;
    end
endtask

    fact_top DUT1 (
        .A(A_tb),
        .WE(WE_tb),
        .WD(WD_tb),
        .Rst(Rst_tb),
        .clk(Clk_tb),
        .RD(RD_tb)
    );

    initial begin
        // Reset fact_top module

```

```

Rst_tb = 1;
tick;
Rst_tb = 0;

for (i=0; i<14; i=i+1) begin
    // Load n and assert WE
    WD_tb = i;
    WE_tb = 1;
    A_tb = 2'b00;
    tick;

    // Give go signal
    WD_tb = 1;
    A_tb = 2'b01;
    tick;

    // Monitor Control Output (Done, Err)
    A_tb = 2'b10;
    tick;
    while (RD_tb == 0) begin
        tick;
    end

    // Examine result if Done = 1
    if(RD_tb == 1) begin
        // Examine Result
        A_tb = 2'b11;
        tick;
    end else if(RD_tb == 2) begin
        $display("Error bit triggered for n = %d", i);
        tick;
    end

    // Reset factorial module
    Rst_tb = 0;
    WD_tb = 0;
    A_tb = 0;
    tick;
end
end
endmodule

```

18) Source Code of tb_gpio_top.v:

```
`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 12/01/2022 12:10:16 AM
// Design Name:
// Module Name: tb_gpio_top
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

module tb_gpio_top;

    //inputs
    reg  [1:0]  A_tb;
    reg        WE_tb;
    reg  [31:0] gpI1_tb;
    reg  [31:0] gpI2_tb;
    reg  [31:0] WD_tb;
    reg        rst_tb;
    reg        clk_tb;

    //outputs
    wire [31:0] RD_tb;
    wire [31:0] gpO1_tb;
    wire [31:0] gpO2_tb;

    gpio_top DUT(
        .A(A_tb),
        .WE(WE_tb),
```

```

.gpi1(gpI1_tb),
.gpi2(gpI2_tb),
.WD(WD_tb),
.RST(rst_tb),
.CLK(clk_tb),
.RD(RD_tb),
.gpo1(gpO1_tb),
.gpo2(gpO2_tb)
);

task tick;
begin
    clk_tb = 0; #5; // Set bit 1 --> 0 after 5 time units
    clk_tb = 1; #5;
end
endtask

initial begin
begin
    // Reset GPIO module
    rst_tb = 1;
    tick;
    rst_tb = 0;

    // Set up GPIO module
    WE_tb = 1'b1;

    // Give test values to input
    WD_tb = 3;
    gpI1_tb = 5;
    gpI2_tb = 7;

    // Test all values for A
    A_tb = 2'b00;
    tick;

    A_tb = 2'b01;
    tick;

    A_tb = 2'b10;
    tick;

    A_tb = 2'b11;

```

```

        tick;
    end
end
endmodule

```

19) Source Code of tb_system_ad.v:

```

module tb_system_ad;
    reg        WE_tb;
    reg  [31:0] A_tb;
    wire        WE1_tb, WE2_tb, WEM_tb;
    wire  [1:0] RdSel_tb;

    system_AD DUT1 (
        .WE(WE_tb),
        .A(A_tb),
        .WE1(WE1_tb),
        .WE2(WE2_tb),
        .WEM(WEM_tb),
        .RdSel(RdSel_tb)
    );

    initial begin
        WE_tb = 0;
        A_tb = 32'h000; #100;
        A_tb = 32'h800; #100;
        A_tb = 32'h900; #100;
    end
endmodule

```

20) Source Code of tb_mips_pipelined_top.v:

```

module tb_mips_pipelined_top;

    reg clk;
    reg rst;
    wire we_dm;
    wire [31:0] pc_current;
    wire [31:0] instr;
    wire [31:0] alu_out;
    wire [31:0] wd_dm;
    wire [31:0] rd_dm;

```

```

wire [31:0] DONT_USE;

integer i;

mips_pipelined_top DUT(
    .clk(clk),
    .rst(rst),
    .ra3(5'h0),
    .we_dm(we_dm),
    .pc_current(pc_current),
    .instr(instr),
    .alu_out(alu_out),
    .wd_dm(wd_dm),
    .rd_dm(rd_dm),
    .rd3(DONT_USE)
);

task tick;
begin
    clk = 1'b0; #5;
    clk = 1'b1; #5;
end
endtask

task reset;
begin
    rst = 1'b0; #5;
    rst = 1'b1; #5;
    rst = 1'b0;
end
endtask

initial begin
    reset;
    for(i=0;i<300;i=i+1)
        begin
            tick;
        end
    reset;
    while(pc_current != 32'hb4) tick;
    $finish;
end
endmodule

```

21) Source Code of system.v:

```

module system(
    input clk,

```

```

    input reset,
    input [31:0] gpi1,
    input [31:0] gpi2,
    output [31:0] gpO1,
    output [31:0] gpO2
);

wire [31:0] PC;
wire [31:0] Instruction;
wire MemWrite;
//wire [31:0] address;
wire [31:0] WriteData;
wire [31:0] DMemData;
wire [31:0] FactData;
wire [31:0] GPIOData;
wire [31:0] ReadData;

wire WE1;
wire WE2;
wire WEM;
wire [1:0] RdSel;

imem iMemory(
    .a(PC[7:2]),
    .y(Instruction)
);

wire [4:0] dummy5;
wire [31:0] alu_out;
wire [31:0] dummy;
//wire [31:0] address;

mips_pipelined MIPS(
    .clk(clk),
    .rst(reset),
    .ra3(dummy5),
    .we_dm(MemWrite),
    .pc_current(PC),
    .instr(Instruction),
    .alu_out(alu_out),
    .wd_dm(WriteData),
    .rd_dm(ReadData),
    .rd3(dummy)
    //.address(address)
);

```



```

dmem dMemory(
    .clk(clk),
    .we(WEM),
    .a(alu_out[7:2]),
    .d(WriteData),
    .q(DMemData)
);

fact_top fact(
    .A(alu_out[3:2]),
    .WE(WE1),
    .WD(WriteData[3:0]),
    .Rst(reset),
    .clk(clk),
    .RD(FactData)
);

gpio_top gpio(
    .A(alu_out[3:2]),
    .WE(WE2),
    .gpi1(gpi1),
    .gpi2(gpi2),
    .WD(WriteData),
    .RST(reset),
    .CLK(clk),
    .RD(GPIOData),
    .gpo1(gpO1),
    .gpo2(gpO2)
);

system_AD AD(
    .WE(MemWrite),
    .A(alu_out),
    .WEM(WEM),
    .WE1(WE1),
    .WE2(WE2),
    .RdSel(RdSel)
);

mux4 #(32) rdDM_mux(
    .sel(RdSel),
    .a(DMemData),
    .b(DMemData),
    .c(FactData),
    .d(GPIOData),
    .y(ReadData)

```

```
);  
  
endmodule
```

22) Source Code of system_AD.v:

```
module system_AD(  
    input WE,  
    input [31:0] A,  
    output reg WEM,  
    output reg WE1,  
    output reg WE2,  
    output reg [1:0] RdSel  
);  
  
always @ (*) begin  
    // Address range in if statements based on memory map table  
    if (A >= 32'h0 && A <= 32'hFC) begin  
        WEM = WE;  
        WE1 = 0;  
        WE2 = 0;  
        RdSel = 2'b00;  
    end  
    else if (A >= 32'h800 && A <= 32'h80C) begin  
        WEM = 0;  
        WE1 = WE;  
        WE2 = 0;  
        RdSel = 2'b10;  
    end  
    else if (A >= 32'h900 && A <= 32'h90C) begin  
        WEM = 0;  
        WE1 = 0;  
        WE2 = WE;  
        RdSel = 2'b11;  
    end  
    else begin  
        WEM = 0;  
        WE1 = 0;  
        WE2 = 0;  
        RdSel = 2'b00;  
    end  
end  
end  
endmodule
```

23) Source Code of fact_top.v:

```
`timescale 1ns / 1ps
```

```

/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 10:49:51 PM
// Design Name:
// Module Name: fact_top
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

```

```

module fact_top(
    input [1:0] A,
    input      WE,
    input [3:0] WD,
    input      Rst,
    input      clk,
    output [31:0] RD
);

    wire [1:0] RdSel;
    wire WE1, WE2;
    wire [3:0] n;
    wire GoPulseCmb, GoPulse;
    wire Go;

    wire [31:0] nf;
    wire fact_done, fact_err;
    wire [31:0] result;

    wire ResDone, ResErr;

    and_fact #(1) and_fact (
        .a      (WE2),
        .b      (WD[0]),
        .c      (GoPulseCmb)
    );
    fact_ad address_decoder (
        .A      (A),

```

```

        .WE      (WE),
        .WE2     (WE2),
        .WE1     (WE1),
        .RdSel   (RdSel)
    );

    fact_reg #(4) n_reg (
        .Rst      (Rst),
        .Clk      (clk),
        .Load_Reg  (WE1),
        .D        (WD),
        .Q        (n)
    );

    fact_reg #(1) go_reg (
        .Rst      (Rst),
        .Clk      (clk),
        .Load_Reg  (WE2),
        .D        (WD[0]),
        .Q        (Go)
    );

    fact_reg #(1) go_pulse_reg (
        .Rst      (Rst),
        .Clk      (clk),
        .Load_Reg  (1'b1),
        .D        (GoPulseCmb),
        .Q        (GoPulse)
    );

    //fact fact (); // TODO
    FSMmult #(32) fact (
        .D        ({28'b0,n}),
        .GO       (GoPulse),
        .RST      (Rst),
        .CLK      (clk),
        .doneF    (fact_done),
        .ERROR    (fact_err),
        .out      (nf)
    );

    fact_reg #(32) result_reg (
        .Rst      (Rst),
        .Clk      (clk),
        .Load_Reg  (fact_done),
        .D        (nf),
        .Q        (result)
    );

    fact_res_done_reg fact_res_done_reg(
        .Clk      (clk),
        .Rst      (Rst),

```

```

        .GoPulseCmb (GoPulseCmb),
        .Done      (fact_done),
        .ResDone   (ResDone)
    );
    fact_res_err_reg fact_res_err_reg (
        .Clk      (clk),
        .Rst      (Rst),
        .GoPulseCmb (GoPulseCmb),
        .Err      (fact_err),
        .ResErr   (ResErr)
    );
    mux4 #(32) fact_mux (
        .sel      (RdSel),
        .a        ({28'b0,n}),
        .b        ({31'b0,Go}),
        .c        ({30'b0,ResErr,ResDone}),
        .d        (result),
        .y        (RD)
    );
endmodule

```

24) Source Code of the gpio_top.v:

```

module gpio_top #(parameter WIDTH=32)(
    input wire [1:0] A,
    input wire      WE,
    input wire [WIDTH-1:0] gpi1, //gpi1 not L or I
    input wire [WIDTH-1:0] gpi2,
    input wire [WIDTH-1:0] WD,
    input wire      RST,
    input wire      CLK,
    output wire [WIDTH-1:0] RD,
    output wire [WIDTH-1:0] gpo1,
    output wire [WIDTH-1:0] gpo2
);

wire WE1;
wire WE2;
wire [1:0] RdSel;

gpio_ad gpio_ad(
    .A (A),
    .WE (WE),
    .WE1 (WE1),
    .WE2 (WE2),
    .RdSel (RdSel)
);

```

```

//gpo1
fact_reg #(32) gpo1_reg(
    .Clk(CLK),
    .Rst(RST),
    .D(WD),
    .Load_Reg(WE1),
    .Q(gpo1)
);
//gpo2
fact_reg #(32) gpo2_reg(
    .Clk(CLK),
    .Rst(RST),
    .D(WD),
    .Load_Reg(WE2),
    .Q(gpo2)
);
mux4 #(32) mux_out(
    .sel(RdSel),
    .a (gpi1),
    .b (gpi2),
    .c (gpo1),
    .d (gpo2),
    .y (RD)
);

endmodule

```

25) Source Code of mux4.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer: Tirumala Saiteja Goruganthu and Harish Marepalli
//
// Create Date: 11/30/2022 10:53:29 PM
// Design Name:
// Module Name: mux4
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//

```

```

////////////////////////////////////
module mux4 #(parameter WIDTH = 8) (
    input wire [1:0] sel,
    input wire [WIDTH-1:0] a,
    input wire [WIDTH-1:0] b,
    input wire [WIDTH-1:0] c,
    input wire [WIDTH-1:0] d,
    output reg [WIDTH-1:0] y
);

always @ (*)
begin
    case (sel)
        2'b00: y <= a;
        2'b01: y <= b;
        2'b10: y <= c;
        2'b11: y <= d;
        default: y <= {WIDTH{1'bx}};
    endcase
// assign y = (sel) ? b : a;
end

endmodule

```

26) Source Code of mips_pipelined_top.v:

```

`timescale 1ns / 1ps
////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 12/01/2022 04:11:53 PM
// Design Name:
// Module Name: mips_pipelined_top
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////

```



```
// Company:
// Engineer: Tirumala Saiteja Goruganthu and Harish Marepalli
//
// Create Date: 12/01/2022 04:26:16 PM
// Design Name:
// Module Name: mips_pipelined
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////
```

```
module mips_pipelined(
    input wire    clk,
    input wire    rst,
    input wire [4:0] ra3,
    input wire [31:0] instr,
    input wire [31:0] rd_dm,
    output wire    we_dm,
    output wire [31:0] pc_current,
    output wire [31:0] alu_out,
    output wire [31:0] wd_dm,
    output wire [31:0] rd3,
    output wire [31:0] address
);
```

```
/* Wires */
```

```
/* FETCH STAGE */
```

```
    wire [31:0] pc_currentF;
    wire [31:0] pc_plus4F;
```

```
/* DECODE STAGE */
```

```
    wire    pc_srcD;
    wire [31:0] instrD;
    wire [31:0] pc_plus4D;
```

```
    wire [31:0] btaD;
```

```
wire [31:0] rd1D;  
wire [31:0] rd2D;
```

```
wire [31:0] wdD;  
wire [31:0] jtaD;
```

```
/* CU */
```

```
    wire    branchD;  
    wire    jumpD;  
    wire    reg_dstD;  
    wire    we_regD;  
    wire    alu_srcD;  
    wire    we_dmD;  
    wire    dm2regD;  
    wire [3:0] alu_ctrlD;  
    wire [1:0] hilo_mux_ctrlD;  
    wire    hilo_weD;  
    wire    jr_mux_ctrlD;  
    wire    jal_wd_muxD;  
    wire    jal_wa_muxD;
```

```
wire [31:0] sext_immD;
```

```
wire [4:0] instrD_20_16;  
wire [4:0] instrD_15_11;
```

```
wire [4:0] shamtD;
```

```
wire [31:0] addressD;
```

```
/* EXECUTE STAGE */
```

```
wire [31:0] wdE;
```

```
    wire    jr_mux_ctrlE;  
    wire    jumpE;  
    wire    hilo_weE;  
    wire [1:0] hilo_mux_ctrlE;  
    wire    dm2regE;  
    wire    we_dmE;  
    wire    branchE; //doubt  
    wire [3:0] alu_ctrlE;  
    wire    alu_srcE;  
    wire    reg_dstE;  
    wire    we_regE;  
    wire    jal_wa_muxE;
```

```

wire    jal_wd_muxE;

wire [31:0] rd1_outE;
wire [31:0] rd2_outE;

wire [4:0] instrE_20_16;
wire [4:0] instrE_15_11;

wire [31:0] pc_plus4E; //doubt

wire [4:0] shamtE;
wire [31:0] sext_immE;

wire [31:0] alu_outE;

wire [31:0] addressE;

wire [31:0] jtaE;

wire [31:0] jrE;
wire [4:0] waE;

wire [31:0] mult_hiE;
wire [31:0] mult_loE;

wire [31:0] wd_dmE;

assign wd_dmE = rd2_outE;
//assign wd_dm = wd_dmE;

assign alu_out = alu_outM;

/* MEMORY STAGE */
//wire    pc_srcM;
// wire [31:0] btaM;
wire    jr_mux_ctrlM;
wire    jumpM;
wire    hilo_weM;
wire [1:0] hilo_mux_ctrlM;
wire    dm2regM;
wire    we_dmM;
wire    we_regM;
wire [31:0] alu_outM;

wire [31:0] wd_dmM;
wire [31:0] mult_hiM;
wire [31:0] mult_loM;

```

```

wire [31:0] jrM;
wire [4:0] waM;
wire      jal_wa_muxM;
wire      jal_wd_muxM;
wire [31:0] rd_dmM;
wire [31:0] hi_outM;
wire [31:0] lo_outM;
wire [31:0] jtaM;

/* WRITEBACK STAGE */
wire [31:0] pc_nextW;
wire      jr_mux_ctrlW;
wire      jumpW;
wire [31:0] jrW;
wire      we_regW;
wire [4:0] waW;
wire [31:0] jtaW;
wire      jal_wd_muxW;
wire [31:0] hilo_mux_outW;
wire [1:0] hilo_mux_ctrlW;
wire      dm2regW;
wire [31:0] alu_outW;
wire [31:0] hi_outW;
wire [31:0] lo_outW;
wire [31:0] wd_rfW;
wire [31:0] rd_dmW;

assign pc_current = pc_currentF;

/* FETCH STAGE */

wire [31:0] pc_pc_src_mux;
wire [31:0] pc_jump_mux;

                                //for branches
mux2 #(32) pc_src_mux (
    .sel      (pc_srcD),
    .a        (pc_plus4D), //TODO: double check that this shouldn't be 27:0
    .b        (btaD),

    .y        (pc_pc_src_mux)
);

mux2 #(32) jump_mux (
    .sel      (jumpW),
    .a        (pc_pc_src_mux),

```

```

        .b          (jtaW),

        .y          (pc_jump_mux)
    );

    mux2 #(32) jr_mux    (
        .sel         (jr_mux_ctrlW),
        .a           (pc_jump_mux),
        .b           (jrW),

        .y           (pc_nextW)
    );

    dreg #(32) pc_reg    (
        .clk         (clk),
        .rst         (rst),
        .d           (pc_nextW),

        .q           (pc_currentF)
    );
    //imem

    adder pc_plus4      (
        .a           (pc_currentF),
        .b           (32'd4),

        .y           (pc_plus4F)
    );

    /* D Stage Reg Interface */
    D_Stage_Reg D_Stage_Reg (
        .clk         (clk),
        .rst         (rst),
        .instrF      (instr),
        .pc_plus4F   (pc_plus4F),

        .instrD      (instrD),
        .pc_plus4D   (pc_plus4D)
    );

    /* DECODE STAGE */
    controlunit cu (
        .opcode      (instrD[31:26]),
        .funct       (instrD[5:0]),

        .branch      (branchD),
        .jump        (jumpD),

```

```

.reg_dst      (reg_dstD),
.we_reg       (we_regD),
.alu_src      (alu_srcD),
.we_dm        (we_dmD),
.dm2reg       (dm2regD),
.alu_ctrl     (alu_ctrlD),
.hilo_mux_ctrl (hilo_mux_ctrlD),
.hilo_we      (hilo_weD),
.jr_mux_ctrl  (jr_mux_ctrlD),
.jal_wd_mux_sel (jal_wd_muxD),
.jal_wa_mux_sel (jal_wa_muxD)

);
regfile rf (
.clk      (clk),
.we       (we_regW),
.ra1      (instrD[25:21]),
.ra2      (instrD[20:16]),
.ra3      (ra3),
.wa       (waW),
.wd       (wdD),
.rd1      (rd1D),
.rd2      (rd2D),
.rd3      () //TODO not sure if we're using this
);
signext se (
.a        (instrD[15:0]),
.y        (sext_immD)
);

mux2 #(32) jal_wd_mux (
.sel      (jal_wd_muxW),
.a        (hilo_mux_outW),
.b        (pc_plus4D),

.y        (wdD)
);

assign shamtD = instrD[10:6];

assign instrD_20_16 = instrD[20:16];
assign instrD_15_11 = instrD[15:11];

assign jtaD = {pc_plus4D[31:28], instrD[25:0], 2'b00}; //TODO not sure about this
one.

addner pc_plus_br (
.a      ({sext_immD[29:0], 2'b00}),
.b      (pc_plus4D),

```

```

.y          (btaD)
);

assign pc_srcD = ((rd1D == rd2D) && branchD) ? 1 : 0;

assign addressD = instrD;

/* E Stage Reg Interface */
E_Stage_Reg E_Stage_Reg (
.clk          (clk),
.rst          (rst),
.jr_mux_ctrlD (jr_mux_ctrlD),
.jumpD        (jumpD),
.hilo_weD     (hilo_weD),
.hilo_mux_ctrlD (hilo_mux_ctrlD),
.dm2regD      (dm2regD),
.we_dmD       (we_dmD),
//.branchD    (branchD),
.alu_ctrlD    (alu_ctrlD),
.alu_srcD     (alu_srcD),
.reg_dstD     (reg_dstD),
.we_regD      (we_regD),
.jal_wa_muxD  (jal_wa_muxD),
.jal_wd_muxD  (jal_wd_muxD),

.rd1D         (rd1D),
.rd2D         (rd2D),

.instrD_20_16  (instrD_20_16),
.instrD_15_11  (instrD_15_11),

.pc_plus4D     (pc_plus4D),

.shamtD        (shamtD),
.sext_immD     (sext_immD),

.addressD      (addressD),
.jtaD          (jtaD),

.jr_mux_ctrlE  (jr_mux_ctrlE),
.jumpE         (jumpE),
.hilo_weE      (hilo_weE),
.hilo_mux_ctrlE (hilo_mux_ctrlE),
.dm2regE       (dm2regE),
.we_dmE        (we_dmE),
//.branchE     (branchE),
.alu_ctrlE     (alu_ctrlE),

```

```

.alu_srcE      (alu_srcE),
.reg_dstE      (reg_dstE),
.we_regE       (we_regE),
.jal_wa_muxE   (jal_wa_muxE),
.jal_wd_muxE   (jal_wd_muxE),

.rd1_outE      (rd1_outE),
.rd2_outE      (rd2_outE),

.instrE_20_16  (instrE_20_16),
.instrE_15_11  (instrE_15_11),

.pc_plus4E     (pc_plus4E),

.shamtE        (shamtE),
.sext_immE     (sext_immE),

.addressE      (addressE),
.jtaE          (jtaE)

);

/* EXECUTE STAGE */

wire [31:0] alu_paE;
wire [31:0] alu_pbE;
wire      zeroE;

wire [4:0] rf_wa;

assign alu_paE = rd1_outE;

mux2 #(32) alu_pb_mux (
    .sel      (alu_srcE),
    .a        (rd2_outE),
    .b        (sext_immE),
    .y        (alu_pbE)
);
alu alu      (
    .op       (alu_ctrlE),
    .a        (alu_paE),
    .b        (alu_pbE),
    //zero     (zeroE),
    .y        (alu_outE),
    .shamt    (shamtE)
);

mux2 #(5) rf_wa_mux ( //TODO move to D
    .sel      (reg_dstE),

```



```

        .a      (instrE_20_16),
        .b      (instrE_15_11),
        .y      (rf_wa)
    );

    mux2 #(5) jal_wa_mux (
        .sel      (jal_wa_muxE),
        .a      (rf_wa),
        .b      (5'd31),

        .y      (waE)

    );

    mult_inf #(32) mult (
        .a      (rd1_outE),
        .b      (rd2_outE),

        .out      ({mult_hiE, mult_loE})
    );

    /* M Stage Interface */
    M_Stage_Reg M_Stage_Reg (
        .clk      (clk),
        .rst      (rst),
        .jr_mux_ctrlE (jr_mux_ctrlE),
        .jumpE      (jumpE),
        .hilo_weE      (hilo_weE),
        .hilo_mux_ctrlE (hilo_mux_ctrlE),
        .dm2regE      (dm2regE),
        .we_dmE      (we_dmE),
        .we_regE      (we_regE),
        .alu_outE      (alu_outE),

        .wd_dmE      (wd_dmE),
        .mult_hiE      (mult_hiE),
        .mult_loE      (mult_loE),

        .jrE      (jrE),
        .waE      (waE),
        .jal_wa_muxE (jal_wa_muxE),
        .jal_wd_muxE (jal_wd_muxE),
        .addressE      (addressE),
        .jtaE      (jtaE),

        .jr_mux_ctrlM (jr_mux_ctrlM),
        .jumpM      (jumpM),
        .hilo_weM      (hilo_weM),

```

```

.hilo_mux_ctrlM (hilo_mux_ctrlM),
.dm2regM      (dm2regM),
.we_dmM       (we_dmM),
.we_regM      (we_regM),
.alu_outM     (alu_outM),

.wd_dmM       (wd_dm),
.mult_hiM     (mult_hiM),
.mult_loM     (mult_loM),

.jrM          (jrM),
.waM          (waM),
.jal_wa_muxM  (jal_wa_muxM),
.jal_wd_muxM  (jal_wd_muxM),
.addressM     (address),
.jtaM         (jtaM)
);

/* MEMORY STAGE */

assign alu_out = alu_outM;
assign wd_dm  = wd_dmM;
assign we_dm  = we_dmM;
assign rd_dmM = rd_dm;
//inferred and gate
//connections to data mem
HiLo_reg #(32) hi_lo_reg (
.clk      (clk),
.rst      (rst),
.we       (hilo_weM),
.hi       (mult_hiM),
.lo       (mult_loM),

.hi_out   (hi_outM),
.lo_out   (lo_outM)
);

/* W Stage Reg Interface */
W_Stage_Reg W_Stage_Reg (
.clk      (clk),
.rst      (rst),
.jr_mux_ctrlM (jr_mux_ctrlM),
.jumpM     (jumpM),
.hilo_mux_ctrlM (hilo_mux_ctrlM),
.dm2regM    (dm2regM),
.we_regM    (we_regM),
.alu_outM    (alu_outM),
.rd_dmM     (rd_dmM),

```

```

        .hi_outM      (hi_outM),
        .lo_outM      (lo_outM),
        .jrM          (jrM),
        .waM          (waM),
        .jtaM         (jtaM),
        .jal_wd_muxM  (jal_wd_muxM),

        .jr_mux_ctrlW (jr_mux_ctrlW),
        .jumpW        (jumpW),
        .hilo_mux_ctrlW (hilo_mux_ctrlW),
        .dm2regW      (dm2regW),
        .we_regW      (we_regW),

        .alu_outW     (alu_outW),
        .rd_dmW       (rd_dmW),
        .hi_outW      (hi_outW),
        .lo_outW      (lo_outW),

        .jrW          (jrW),
        .waW          (waW),
        .jtaW         (jtaW),
        .jal_wd_muxW  (jal_wd_muxW)
    );

/* WRITEBACK STAGE */

    mux2 #(32) rf_wd_mux (
        .sel      (dm2regW),
        .a        (alu_outW),
        .b        (rd_dmW),

        .y        (wd_rfW)
    );
    mux4 #(32) hilo_mux (
        .sel      (hilo_mux_ctrlW),
        .a        (wd_rfW),
        .b        (lo_outW),
        .c        (hi_outW),
        .d        (32'd0),

        .y        (hilo_mux_outW)
    );

endmodule

```

28) Source Code of imem.v:

```
module imem (
    input wire [5:0] a,
    output wire [31:0] y
);

    reg [31:0] rom [0:63];

    initial begin
        $readmemh
("H:\\Masters\\Semester1\\CMPE200\\Assignment7\\memfile_withnops_updated2.dat",
rom);
    end

    assign y = rom[a];

endmodule
```

29) Source Code of dmem.v:

```
module dmem (
    input wire    clk,
    input wire    we,
    //input wire    rst,
    input wire [5:0] a,
    input wire [31:0] d,
    output wire [31:0] q
);

    reg [31:0] ram [0:63];

    integer n;

    initial begin
        for (n = 0; n < 64; n = n + 1) ram[n] = 32'hFFFFFFFF;
    end

    always @ (posedge clk) begin
        if (we) ram[a] <= d;
    end

    assign q = ram[a];

endmodule
```

30) Source Code of mux2.v:

```
module mux2 #(parameter WIDTH = 8) (
    input wire      sel,
    input wire [WIDTH-1:0] a,
    input wire [WIDTH-1:0] b,
    output wire [WIDTH-1:0] y
);

    assign y = (sel) ? b : a;

endmodule
```

31) Source Code of dreg.v:

```
module dreg # (parameter WIDTH = 32) (
    input wire      clk,
    input wire      rst,
    input wire [WIDTH-1:0] d,
    output reg [WIDTH-1:0] q
);

    always @ (posedge clk, posedge rst) begin
        if (rst) q <= 0;
        else    q <= d;
    end
endmodule
```

32) Source Code of adder.v:

```
module adder (
    input  wire [31:0] a,
    input  wire [31:0] b,
    output wire [31:0] y
);

    assign y = a + b;

endmodule
```

33) Source Code of D_Stage_Reg.v:

[illegible]

```

// Engineer:
//
// Create Date: 11/30/2022 10:58:47 PM
// Design Name:
// Module Name: D_Stage_Reg
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////

module D_Stage_Reg(
    input    clk, rst,
    input [31:0] instrF,
    input [31:0] pc_plus4F,

    output reg [31:0] instrD,
    output reg [31:0] pc_plus4D
);

always @ (negedge clk, posedge rst) begin
    if (rst) begin
        instrD    <= 0;
        pc_plus4D <= 0;
    end
    else begin
        instrD    <= instrF;
        pc_plus4D <= pc_plus4F;
    end
end

endmodule

```

34) Source Code of controlunit.v:

```

module controlunit (
    input wire [5:0] opcode,
    input wire [5:0] funct,
    output wire      branch,
    output wire      jump,

```

```

        output wire    reg_dst,
        output wire    we_reg,
        output wire    alu_src,
        output wire    we_dm,
        output wire    dm2reg,
        output wire [3:0] alu_ctrl,
        output wire [1:0] hilo_mux_ctrl,
        output wire    hilo_we,
        output wire    jr_mux_ctrl,
        output wire    jal_wd_mux_sel,
        output wire    jal_wa_mux_sel
    );

    wire [1:0] alu_op;
    wire [1:0] hilo_mux_internal;

    maindec md (
        .opcode      (opcode),
        .branch      (branch),
        .jump        (jump),
        .reg_dst     (reg_dst),
        .we_reg      (we_reg),
        .alu_src     (alu_src),
        .we_dm       (we_dm),
        .dm2reg      (dm2reg),
        .alu_op      (alu_op),
        .jal_wa_mux_sel (jal_wa_mux_sel),
        .jal_wd_mux_sel (jal_wd_mux_sel)
    );

    auxdec ad (
        .alu_op      (alu_op),
        .funct       (funct),
        .alu_ctrl    (alu_ctrl),
        .hilo_mux_ctrl (hilo_mux_internal),
        .hilo_we     (hilo_we),
        .jr_mux_ctrl  (jr_mux_ctrl)
    );

    assign hilo_mux_ctrl = (hilo_mux_internal) ? hilo_mux_internal : 2'b0;

endmodule

```

35) Source Code of regfile.v:

```

module regfile (
    input wire    clk,
    input wire    we,

```



```
// Create Date: 11/30/2022 10:58:47 PM
// Design Name:
// Module Name: E_Stage_Reg
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////////////////////////////////
```

```
module E_Stage_Reg(
    input    clk, rst,
    input    jr_mux_ctrlD,
    input    jumpD,
    input    hilo_weD,
    input [1:0] hilo_mux_ctrlD,
    input    dm2regD,
    input    we_dmD,

    //input    branchD,
    input [3:0] alu_ctrlD,
    input    alu_srcD,
    input    reg_dstD,
    input    we_regD,
    input    jal_wa_muxD,
    input    jal_wd_muxD,

    input [31:0] rd1D,
    input [31:0] rd2D,

    input [4:0] instrD_20_16,
    input [4:0] instrD_15_11,

    input [31:0] pc_plus4D,

    input [4:0] shamtD,
    input [31:0] sext_immD,

    input [31:0] addressD,
    input [31:0] jtaD,
```

```

output reg    jr_mux_ctrlE,
output reg    jumpE,
output reg    hilo_weE,
output reg [1:0] hilo_mux_ctrlE,
output reg    dm2regE,
output reg    we_dmE,
//output reg    branchE,
output reg [3:0] alu_ctrlE,
output reg    alu_srcE,
output reg    reg_dstE,
output reg    we_regE,
output reg    jal_wa_muxE,
output reg    jal_wd_muxE,

```

```

output reg [31:0] rd1_outE,
output reg [31:0] rd2_outE,

```

```

output reg [4:0] instrE_20_16,
output reg [4:0] instrE_15_11,

```

```

output reg [31:0] pc_plus4E,

```

```

output reg [4:0] shamtE,
output reg [31:0] sext_immE,

```

```

output reg [31:0] addressE,
output reg [31:0] jtaE
);

```

```

always @ (negedge clk, posedge rst) begin

```

```

    if (rst) begin
        jr_mux_ctrlE  <= 0;
        jumpE         <= 0;
        hilo_weE      <= 0;
        hilo_mux_ctrlE <= 0;
        dm2regE       <= 0;
        we_dmE        <= 0;
        //branchE      <= 0;
        alu_ctrlE      <= 0;
        alu_srcE       <= 0;
        reg_dstE       <= 0;
        we_regE        <= 0;
        jal_wa_muxE    <= 0;
        jal_wd_muxE    <= 0;
    end

```

```

    rd1_outE  <= 0;
    rd2_outE  <= 0;

```

```

instrE_20_16 <= 0;
instrE_15_11 <= 0;

pc_plus4E    <= 0;

shamtE       <= 0;
sext_immE    <= 0;

addressE     <= 0;
jtaE         <= 0;
end

else begin
    jr_mux_ctrlE <= jr_mux_ctrlD;
    jumpE        <= jumpD;
    hilo_weE     <= hilo_weD;
    hilo_mux_ctrlE <= hilo_mux_ctrlD;
    dm2regE      <= dm2regD;
    we_dmE       <= we_dmD;
    //branchE    <= branchD;
    alu_ctrlE    <= alu_ctrlD;
    alu_srcE     <= alu_srcD;
    reg_dstE     <= reg_dstD;
    we_regE      <= we_regD;
    jal_wa_muxE  <= jal_wa_muxD;
    jal_wd_muxE  <= jal_wd_muxD;

    rd1_outE     <= rd1D;
    rd2_outE     <= rd2D;

    instrE_20_16 <= instrD_20_16;
    instrE_15_11 <= instrD_15_11;

    pc_plus4E    <= pc_plus4D;

    shamtE       <= shamtD;
    sext_immE    <= sext_immD;

    addressE     <= addressD;
    jtaE         <= jtaD;
end

end

endmodule

```

38) Source Code of alu.v:

```
`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 10:45:02 PM
// Design Name:
// Module Name: alu
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

module alu (
    input wire [3:0] op,
    input wire [31:0] a,
    input wire [31:0] b,
    input wire [4:0] shamt,
    output wire      zero,
    output reg  [31:0] y, hi, lo
);

    assign zero = (y == 0);

    always @ (op, a, b) begin
        case (op)
            4'b0000: y <= a & b;
            4'b0001: y <= a | b;
            4'b0010: y <= a + b;
            4'b0110: y <= a - b;
            4'b0111: y <= (a < b) ? 1 : 0;
            4'b1000: {hi,lo} <= a * b;
            4'b1001: y <= b << shamt;
            4'b1010: y <= b >> shamt;
        endcase
    end

endmodule
```

39) Source Code of mult_inf.v:

```
`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:01:17 PM
// Design Name:
// Module Name: mult_inf
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

module mult_inf #(parameter SIZE = 32)(

    input wire [SIZE - 1:0] a, b,
    output reg [2*SIZE - 1:0] out

);

    always @ (a, b) begin
        out <= a * b;
    end
endmodule
```

40) Source Code of M_Stage_Reg.v:

```
`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 10:58:47 PM
// Design Name:
// Module Name: M_Stage_Reg
```

```
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////
```

```
module M_Stage_Reg(
    input      clk, rst,
    input      jr_mux_ctrlE,
    input      jumpE,
    input      hilo_weE,
    input [1:0] hilo_mux_ctrlE,
    input      dm2regE,
    input      we_dmE,
    //input     branchE,
    input      we_regE,

    //input     zeroE,

    input [31:0] alu_outE,

    input [31:0] wd_dmE,

    input [31:0] mult_hiE,
    input [31:0] mult_loE,

    //input [31:0] btaE,

    input [31:0] jrE,
    input [4:0]  waE,
    input      jal_wa_muxE,
    input      jal_wd_muxE,

    input [31:0] addressE,
    input [31:0] jtaE,

    output reg   jr_mux_ctrlM,
    output reg   jumpM,
    output reg   hilo_weM,
    output reg [1:0] hilo_mux_ctrlM,
    output reg   dm2regM,
```

```

output reg    we_dmM,
//output reg   branchM,
output reg    we_regM,

//output reg    zeroM,

output reg [31:0] alu_outM,

output reg [31:0] wd_dmM,

output reg [31:0] mult_hiM,
output reg [31:0] mult_loM,

//output reg [31:0] btaM,

output reg [31:0] jrM,
output reg [4:0]  waM,
output reg    jal_wa_muxM,
output reg    jal_wd_muxM,
output reg [31:0] addressM,
output reg [31:0] jtaM
);

always @ (negedge clk, posedge rst) begin
  if (rst) begin
    jr_mux_ctrlM  <= 0;
    jumpM        <= 0;
    hilo_weM     <= 0;
    hilo_mux_ctrlM <= 0;
    dm2regM      <= 0;
    we_dmM       <= 0;
    //branchM     <= 0;
    we_regM      <= 0;

    //zeroM       <= 0;

    alu_outM     <= 0;

    wd_dmM       <= 0;

    mult_hiM     <= 0;
    mult_loM     <= 0;

    //btaM        <= 0;

    jrM          <= 0;
    waM          <= 0;

    jal_wa_muxM  <= 0;

```

```

    jal_wd_muxM    <= 0;
    addressM       <= 0;
    jtaM           <= 0;
end

else begin
    jr_mux_ctrlM   <= jr_mux_ctrlE;
    jumpM          <= jumpE;
    hilo_weM       <= hilo_weE;
    hilo_mux_ctrlM <= hilo_mux_ctrlE;
    dm2regM        <= dm2regE;
    we_dmM         <= we_dmE;
    //branchM       <= branchE;
    we_regM        <= we_regE;

    //zeroM         <= zeroE;

    alu_outM        <= alu_outE;

    wd_dmM          <= wd_dmE;

    mult_hiM        <= mult_hiE;
    mult_loM        <= mult_loE;

    //btaM          <= btaE;

    jrM             <= jrE;
    waM             <= waE;

    jal_wa_muxM     <= jal_wa_muxE;
    jal_wd_muxM     <= jal_wd_muxE;
    addressM        <= addressE;
    jtaM            <= jtaE;
end
end
endmodule

```

41) Source Code of HiLo_reg.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:01:17 PM
// Design Name:
// Module Name: HiLo_reg
// Project Name:

```



```

// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////

module HiLo_reg #(parameter WIDTH = 32) (
    input wire clk, rst, we,
    input wire [WIDTH - 1:0] hi, lo,
    output reg [WIDTH - 1:0] hi_out, lo_out
);

    always @ (posedge clk, posedge rst)
    begin
        if (rst) {hi_out, lo_out} <= 0;
        else if (we) {hi_out, lo_out} <= {hi,lo};
        else {hi_out, lo_out} <= {hi_out, lo_out};
    end
endmodule

```

42) Source Code of W_Stage_Reg.v:

```

`timescale 1ns / 1ps
////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 10:58:47 PM
// Design Name:
// Module Name: W_Stage_Reg
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//

```

//

```
module W_Stage_Reg(
    input      clk, rst,
    input      jr_mux_ctrlM,
    input      jumpM,
    input [1:0] hilo_mux_ctrlM,
    input      dm2regM,
    input      we_regM,

    input [31:0] alu_outM,

    input [31:0] rd_dmM,

    input [31:0] hi_outM,
    input [31:0] lo_outM,

    input [31:0] jrM,
    input [4:0]  waM,
    input [31:0] jtaM,
    input      jal_wd_muxM,

    output reg   jr_mux_ctrlW,
    output reg   jumpW,
    output reg [1:0] hilo_mux_ctrlW,
    output reg   dm2regW,
    output reg   we_regW,

    output reg [31:0] alu_outW,

    output reg [31:0] rd_dmW,

    output reg [31:0] hi_outW,
    output reg [31:0] lo_outW,
    output reg      jal_wd_muxW,

    output reg [31:0] jrW,
    output reg [4:0]  waW,
    output reg [31:0] jtaW
);
always @ (negedge clk, posedge rst) begin
    if (rst) begin
        jr_mux_ctrlW <= 0;
        jumpW        <= 0;
        hilo_mux_ctrlW <= 0;
        dm2regW       <= 0;
        we_regW       <= 0;
    end
end
```

```

    alu_outW    <= 0;
    jal_wd_muxW <= 0;
    rd_dmW      <= 0;

    hi_outW     <= 0;
    lo_outW     <= 0;

    jrW         <= 0;
    waW         <= 0;
    jtaW        <= 0;
end

else begin
    jr_mux_ctrlW <= jr_mux_ctrlM;
    jumpW        <= jumpM;
    hilo_mux_ctrlW <= hilo_mux_ctrlM;
    dm2regW      <= dm2regM;
    we_regW      <= we_regM;

    alu_outW     <= alu_outM;

    rd_dmW       <= rd_dmM;

    hi_outW      <= hi_outM;
    lo_outW      <= lo_outM;

    jrW          <= jrM;
    waW          <= waM;

    jtaW         <= jtaM;

    jal_wd_muxW  <= jal_wd_muxM;
end
end
endmodule

```

43) Source Code of maindec.v:

```

module maindec (
    input  wire [5:0] opcode,
    output wire      branch,
    output wire      jump,
    output wire      reg_dst,
    output wire      we_reg,
    output wire      alu_src,
    output wire      we_dm,
    output wire      dm2reg,
    output wire [1:0] alu_op,

```

```

        output wire    jal_wa_mux_sel,
        output wire    jal_wd_mux_sel
    );

    reg [10:0] ctrl;

    assign {branch, jump, reg_dst, we_reg, alu_src, we_dm, dm2reg, alu_op,
jal_wa_mux_sel, jal_wd_mux_sel} = ctrl;

    always @ (opcode) begin
        case (opcode)
            6'b00_0000: ctrl = 11'b0_0_1_1_0_0_0_10_0_0; // R-type
            6'b00_1000: ctrl = 11'b0_0_0_1_1_0_0_00_0_0; // ADDI
            6'b00_0100: ctrl = 11'b1_0_0_0_0_0_0_01_0_0; // BEQ
            6'b00_0010: ctrl = 11'b0_1_0_0_0_0_0_00_0_0; // J
            6'b00_0011: ctrl = 11'b0_1_0_1_0_0_0_00_1_1; // JAL //TODO
            6'b10_1011: ctrl = 11'b0_0_0_0_1_1_0_00_0_0; // SW
            //6'b10_0011: ctrl = 11'b0_0_0_1_1_0_1_00_0_0; // LW
            6'b10_0011: ctrl = 11'b0_0_0_1_1_0_1_00_0_0; // LW
            default:    ctrl = 11'bx_x_x_0_x_0_x_xx_x_x;
            //default:    ctrl = 11'b0_0_0_0_0_0_0_00_0_0;
        endcase
    end

endmodule

```

44) Source Code of auxdec.v:

```

module auxdec (
    input wire [1:0] alu_op,
    input wire [5:0] funct,
    output wire [3:0] alu_ctrl,
    output wire [1:0] hilo_mux_ctrl,
    output wire hilo_we,
    output wire jr_mux_ctrl
);

    reg [7:0] ctrl;

    assign {alu_ctrl, hilo_mux_ctrl, hilo_we, jr_mux_ctrl} = ctrl;

    always @ (alu_op, funct) begin
        case (alu_op)
            2'b00: ctrl = 8'b0010_00_0_0;    // ADD
            2'b01: ctrl = 8'b0110_00_0_0;    // SUB
            default: case (funct)
                6'b10_0100: ctrl = 8'b0000_00_0_0; // AND
                6'b10_0101: ctrl = 8'b0001_00_0_0; // OR
            endcase
        endcase
    end
endmodule

```

```

        6'b10_0000: ctrl = 8'b0010_00_0_0; // ADD
        6'b10_0010: ctrl = 8'b0110_00_0_0; // SUB
        6'b10_1010: ctrl = 8'b0111_00_0_0; // SLT
        6'b01_1001: ctrl = 8'bxxxx_00_1_0; // MULTU
        6'b01_0000: ctrl = 8'b0000_11_0_0; // MFHI
        6'b01_0010: ctrl = 8'b0000_01_0_0; // MFLO
        6'b00_0000: ctrl = 8'b1001_00_0_0; // SLL
        6'b00_0010: ctrl = 8'b1010_00_0_0; // SRL
        6'b00_1000: ctrl = 8'bxxxx_00_0_1; // JR
        default: ctrl = 8'bxxxx_xx_x_x;
    endcase
endcase
end
endmodule

```

45) Source Code of and_fact.v:

```

module and_fact #(parameter WIDTH = 32)(
    input [WIDTH - 1:0] a, b,
    output wire c
);

    assign c = a & b;
endmodule

```

46) Source Code of fact_ad.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:02:37 PM
// Design Name:
// Module Name: fact_ad
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//

```

```

////////////////////////////////////
module fact_ad(
    input wire [1:0] A,
    input wire WE,
    output reg WE1, WE2,
    output wire [1:0] RdSel
);

always @ (*) begin
    case(A)
        2'b00: begin
            WE1 <= WE;
            WE2 <= 1'b0;
        end
        2'b01: begin
            WE1 <= 1'b0;
            WE2 <= WE;
        end
        2'b10: begin
            WE1 <= 1'b0;
            WE2 <= 1'b0;
        end
        2'b11: begin
            WE1 <= 1'b0;
            WE2 <= 1'b0;
        end
        default: begin
            WE1 <= 1'bx;
            WE1 <= 1'bx;
        end
    endcase
end
assign RdSel = A;
endmodule

```

47) Source Code of fact_reg.v:

```

`timescale 1ns / 1ps
////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:11:45 PM
// Design Name:
// Module Name: fact_reg

```

```

// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////

module fact_reg #(parameter w = 32)(
    input wire Clk, Rst,
    input wire [w-1:0] D,
    input wire Load_Reg,
    output reg [w-1:0] Q
);

    always @ (posedge Clk, posedge Rst) begin
        if (Rst)
            Q <= 0;
        else if (Load_Reg)
            Q <= D;
        else
            Q <= Q;
        end
    endmodule

```

48) Source Code of FSMult.v:

```

module FSMmult #(parameter DATA_WIDTH = 32)(
    input [DATA_WIDTH-1:0] D,
    input GO,
    input RST,
    input CLK,
    output wire doneF,
    output wire ERROR,
    output wire [DATA_WIDTH-1:0] out
);

    wire load_cnt, cnt_en, mux_sel, load_reg, buf_oe, gt, ec;

    CU control(
        .GO(GO),

```

```

        .GT(gt),
        .CLK(CLK),
        .RST(RST),
        .EC(ec),
        .load_cnt(load_cnt),
        .cnt_en(cnt_en),
        .mux_sel(mux_sel),
        .load_reg(load_reg),
        .buf_oe(buf_oe),
        .done(doneF),
        .error(ERROR)
    );
    DP dataP(
        .clk(CLK),
        .load_cnt(load_cnt),
        .en_cnt(cnt_en),
        .sel_mux(mux_sel),
        .load_reg(load_reg),
        .oe_buf(buf_oe),
        .gt(gt),
        .ec(ec),
        .in(D),
        .out(out)
    );

endmodule

```

49) Source Code of fact_res_done_reg.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:19:40 PM
// Design Name:
// Module Name: fact_res_done_reg
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:

```



```
//
////////////////////////////////////////////////////////////////

module fact_res_done_reg(
    input Clk, Rst, GoPulseCmb, Done,
    output reg ResDone
);

    always @ (posedge Clk, posedge Rst) begin
        if (Rst)
            ResDone <= 1'b0;
        else
            ResDone <= (~GoPulseCmb) & (Done | ResDone);
        end
    endmodule
```

50) Source Code of fact_res_err_reg.v:

```
`timescale 1ns / 1ps
////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:34:29 PM
// Design Name:
// Module Name: fact_res_err_reg
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////////////////////////////////

module fact_res_err_reg(
    input Clk, Rst, GoPulseCmb, Err,
    output reg ResErr
);

    always @ (posedge Clk, posedge Rst) begin
        if (Rst)
```

```

        ResErr <= 1'b0;
    else
        ResErr <= (~GoPulseCmb) & (Err | ResErr);
    end
endmodule

```

51) Source Code of DP.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:08:41 PM
// Design Name:
// Module Name: DP
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

module DP #(parameter DATA_WIDTH = 32)(
    input wire clk,
    input wire load_cnt,
    input wire en_cnt,
    input wire sel_mux,
    input wire load_reg,
    input wire oe_buf,
    output wire gt,
    output wire ec,
    input wire [DATA_WIDTH-1:0] in,
    output wire [DATA_WIDTH-1:0] out);

    wire [DATA_WIDTH-1:0] mux_to_reg;
    wire [DATA_WIDTH-1:0] mul_to_mux;
    wire [DATA_WIDTH-1:0] cnt_to_cmp_and_mul;
    wire [DATA_WIDTH-1:0] reg_to_mul_and_buf;

```

```

        CNT #(DATA_WIDTH) down_counter(
            .in(in),
            .load_cnt(load_cnt),
            .en(en_cnt),
            .clk(clk),
            .out(cnt_to_cmp_and_mul)
        );

        Multiplexer #(DATA_WIDTH) product_multi(sel_mux,
mul_to_mux, 1, mux_to_reg);

        Register #(DATA_WIDTH) product_reg(clk, load_reg, mux_to_reg,
reg_to_mul_and_buf);

        Multiplexer #(DATA_WIDTH) product_buf(
            oe_buf,
            reg_to_mul_and_buf,
            0,
            out);

        CMP #(DATA_WIDTH) cmp_gt(
            cnt_to_cmp_and_mul, 1, gt);

        Multiplier #(DATA_WIDTH) multiply(
            reg_to_mul_and_buf, cnt_to_cmp_and_mul, mul_to_mux);
        CMP #(DATA_WIDTH) cmp_ec(
            in, 12, ec);

endmodule

```

52) Source Code of CNT.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:08:41 PM
// Design Name:
// Module Name: CNT
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//

```



```

module Multiplexer #(parameter DATA_WIDTH = 32)(
    input wire sel,
    input wire [DATA_WIDTH-1:0] a,
    input wire [DATA_WIDTH -1:0] b,
    output wire [DATA_WIDTH-1:0] out);

    assign out = (sel == 0) ? a : b;

endmodule

```

54) Source Code of Register.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:08:41 PM
// Design Name:
// Module Name: Register
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

module Register #(parameter DATA_WIDTH = 32)(
    input wire clk,
    input wire load_reg,
    input wire [DATA_WIDTH-1:0] in,
    output reg [DATA_WIDTH-1:0] out);

    always@(posedge clk) begin
        if(load_reg) begin
            out = in;
        end
    end

```

```

// else out = out
end

endmodule

```

55) Source Code of CMP.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:08:41 PM
// Design Name:
// Module Name: CMP
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

module CMP #(parameter DATA_WIDTH = 32)(

    input wire [DATA_WIDTH-1:0] a,
    input wire [DATA_WIDTH-1:0] b,
    output wire gt);

    assign gt = (a > b) ? 1 : 0;

endmodule

```

56) Source Code of CU.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:08:41 PM

```

```

// Design Name:
// Module Name: CU
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
////////////////////////////////////

module CU #(parameter DATA_WIDTH = 32)(
    input GO,
    input GT,
    input RST,
    input CLK,
    input EC,
    output reg load_cnt,
    output reg cnt_en,
    output reg mux_sel,
    output reg load_reg,
    output reg buf_oe,
    output reg done,
    output reg error
);

    parameter  S0 = 3'b000,
               S1 = 3'b001,
               S2 = 3'b010,
               S3 = 3'b011,
               S4 = 3'b100;
    reg [2:0] CS, NS;
    //NS driver
    always @ (posedge CLK, posedge RST)
    begin
        if(RST == 1) CS <= S0;
        else      CS <= NS;

    end
    //Output logic
    always @ (CS, GT, EC)
    begin
        case(CS)

```

```

S0: begin
    load_cnt = 0;
    cnt_en = 0;
    mux_sel = 0;
    load_reg = 0;
    buf_oe = 0;
    done = 0;
    error = 0;
end
//Load
S1: begin

    load_cnt = 1;
    cnt_en = 0;
    mux_sel = 1;
    load_reg = 1;
    buf_oe = 0;
    done = 0;
    error = 0;

end
//Finished
S2: begin

    if(EC)begin
        load_cnt <= 0;
        cnt_en <= 0;
        mux_sel <= 0;
        load_reg <= 0;
        buf_oe <= 0;
        done <= 1;
        error <= 1;
    end
    else if(!EC && !GT)begin
        load_cnt = 0;
        cnt_en = 0;
        mux_sel = 0;
        load_reg = 0;
        buf_oe = 0;
        done = 1;
        error = 0;
    end
    else begin
        load_cnt = 0;
        cnt_en = 1;
        mux_sel = 0;
        load_reg = 1;
        buf_oe = 0;
    end
end

```



```

        done = 0;
        error = 0;

        end
    end

    endcase
end
always @ (CS, GO, GT, EC)
begin
    case(CS)
        S0: begin
            if(GO == 1) NS <= S1;    // Go = 1 and D <= 12, go state 1
            else NS <= S0;           // Go = 0
        end
        S1: begin
            NS <= S2;
        end
        //Finished
        S2:
            if(GT == 0 || EC == 1) NS <= S0;
            else NS <= S2;

    endcase
end
endmodule

```

57) Source Code of gpio_ad.v:

```

module gpio_ad(
    input wire [1:0] A,
    input wire      WE,
    output reg      WE1,
    output reg      WE2,
    output wire [1:0] RdSel
);
always @ (*) begin
    case(A)
        2'b00:begin
            WE1 = 1'b0;
            WE2 = 1'b0;
        end
        2'b01: begin
            WE1 = 1'b0;
            WE2 = 1'b0;
        end
    end
end

```

```

        2'b10: begin
            WE1 = WE;
            WE2 = 1'b0;
        end
        2'b11: begin
            WE1 = 1'b0;
            WE2 = WE;
        end

        default: begin
            WE1 = 1'bx;
            WE2 = 1'bx;
        end
    endcase
end

assign RdSel = A;

endmodule

```

58) Source Code of Multiplier.v:

```

`timescale 1ns / 1ps
/////////////////////////////////////////////////////////////////
// Company:
// Engineer:
//
// Create Date: 11/30/2022 11:08:41 PM
// Design Name:
// Module Name: Multiplier
// Project Name:
// Target Devices:
// Tool Versions:
// Description:
//
// Dependencies:
//
// Revision:
// Revision 0.01 - File Created
// Additional Comments:
//
/////////////////////////////////////////////////////////////////

module Multiplier #(parameter DATA_WIDTH = 32)(
    input wire [DATA_WIDTH-1:0] a,

```

```

        input wire [DATA_WIDTH-1:0] b,
        output wire [DATA_WIDTH-1:0] out);

        assign out = a * b;

endmodule

```

59) Assembly Code in pipelined MIPS to trigger the SoC:

```

main:      addi $t0, $0, 0x0F
           lw $t1, 0x0900($0) #Take n value from GPIO
           lw $t2, 0x0904($0) #Take Go value from GPIO
           sll $0, $0, 0
           sll $0, $0, 0
           sw $t1, 0x0800($0) #Give this n value to Factorial Accelerator
           sw $t2, 0x0804($0) #Give the Go signal to Factorial Accelerator
poll:      lw $t3, 0x0808($0) #Continuosuly poll the Done/Err bit
           sll $0, $0, 0
           sll $0, $0, 0
           sll $0, $0, 0
           beq $t3, $0, poll    #Continuosuly poll the Done/Err bit
           srl $t4, $t3, 1      #Shift right to access the Err bit
           lw $t5, 0x080C($0) #Get the n! value into processor
           sll $0, $0, 0
           sll $0, $0, 0
           sw $t4, 0x0908($0) #Processor gives Err bit to GPIO1
           sw $t5, 0x090C($0) #Processor gives n! value to GPIO2

```