

CMPE 200
Computer Architecture & Design

Lecture 2.
Processor Instruction Set
Architecture & Language (1)

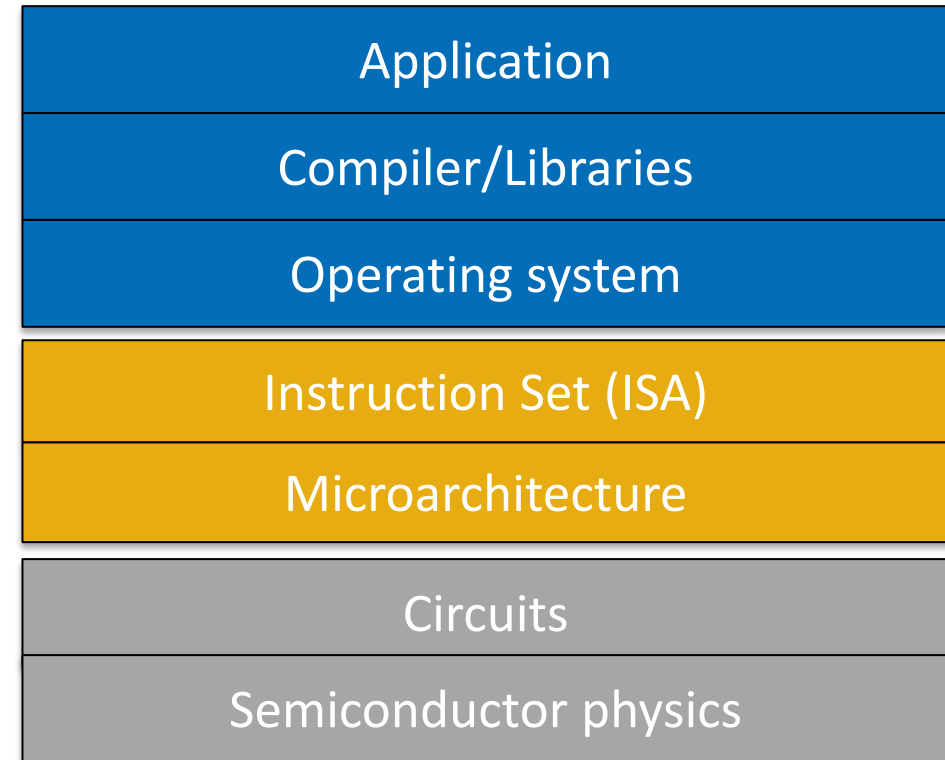
Haonan Wang



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Instruction Set Architecture & Microarchitecture

- **Instruction Set Architecture**
 - the programmer's view of the computer, defined by a set of instructions that each specifies an operation and its operand(s)
- **Microarchitecture**
 - the way that the ISA is implemented in hardware



What Does the ISA Deal With Specifically?

Example: a C program that reads two integer values from “file.txt” file and prints the sum of them.

```
#include <stdio.h>
#include <string.h>

int numbers[2];

void myfunction(void)
{
    FILE *fp;
    int size = 2;
    int sum = 0;

    /* Open file for reading */
    fp = fopen("mynumbers.txt", "r");

    /* Read and display data */
    fread(numbers, sizeof(int), size, fp);
    fclose(fp);
    sum = numbers[0] + numbers[1];
    printf("Sum = %d\n", sum);
}

int main (void) {
    myfunction();
    return(0);
}
```

Processor (CPU)

Understands
and executes
each line of the
code.

Uses fast on-
chip memories

Memory (DRAM)

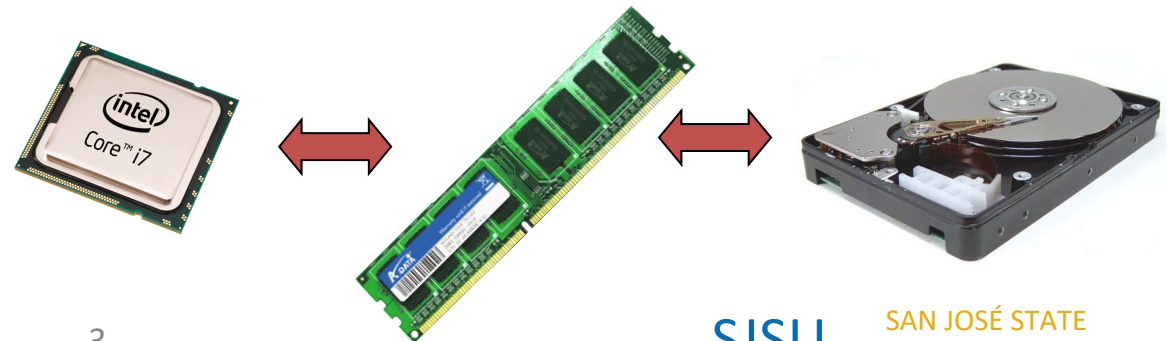
Provides
operands to
CPU

(*fp, size, sum,
numbers[2])

Storage (HDD)

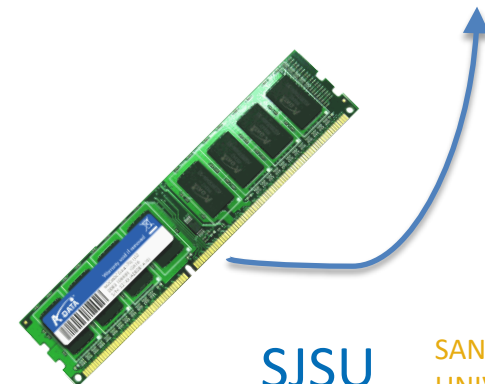
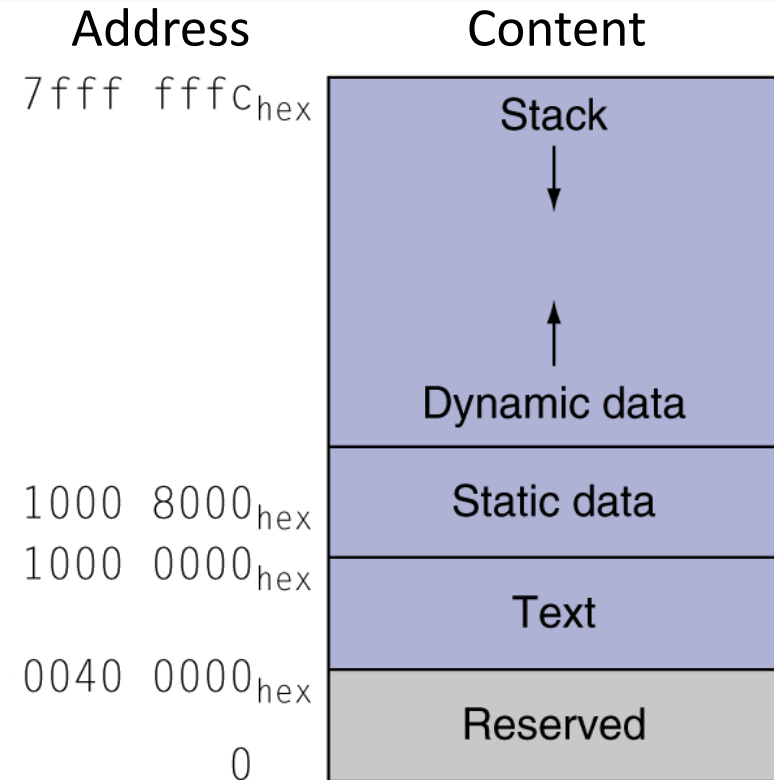
Provides file
inputs and
program code

(file.txt)



Memory Layout

- **Text:** program code
- **Static data:** global variables
 - e.g., static variables in C, constant arrays and strings
- **Dynamic data:** heap
 - e.g., malloc in C, new in Java
 - Grows from bottom (lower address) to top (higher address)
- **Stack:** temporal storage for functions
 - e.g., return address of sub-functions, local variables
 - Grows from top to bottom



Memory Layout

```
#include <stdio.h>
#include <string.h>

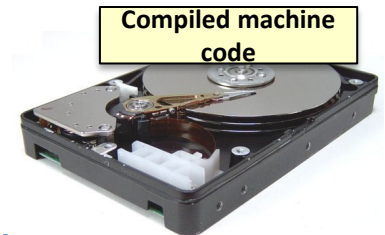
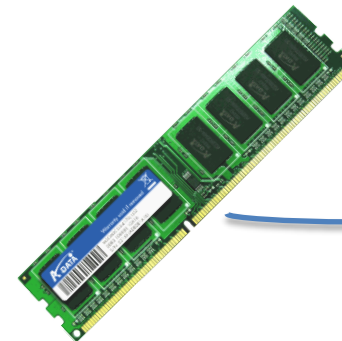
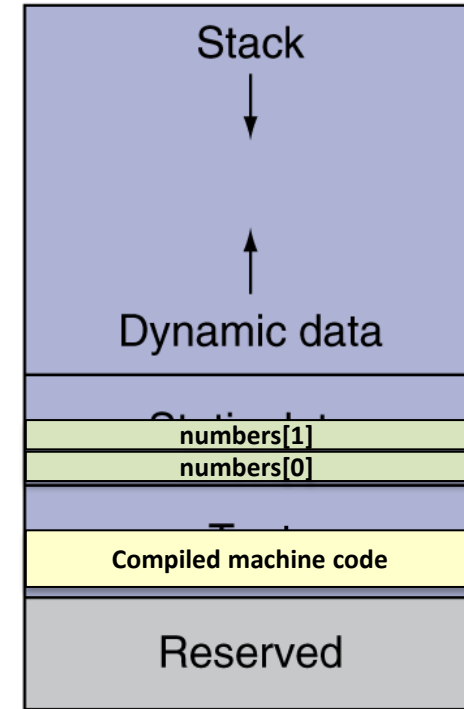
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int main (void) {
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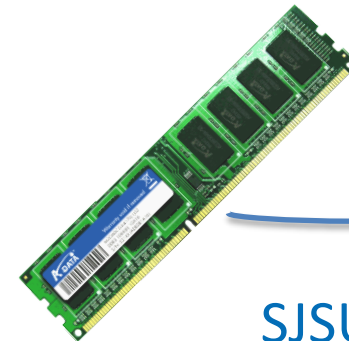
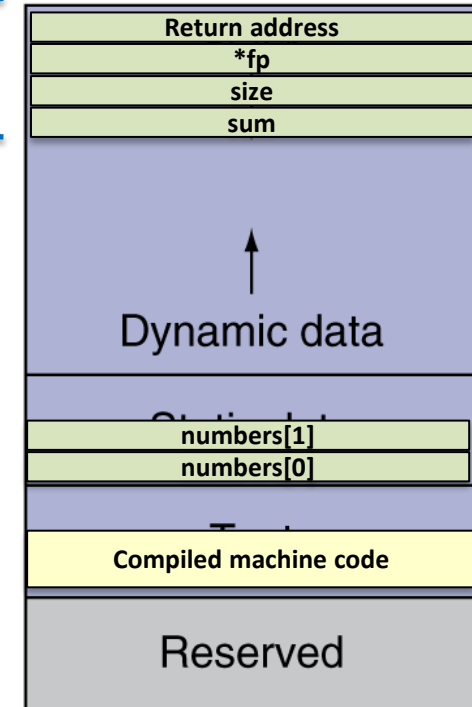
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int main (void) {
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```

Stack region for myfunction



Memory Layout

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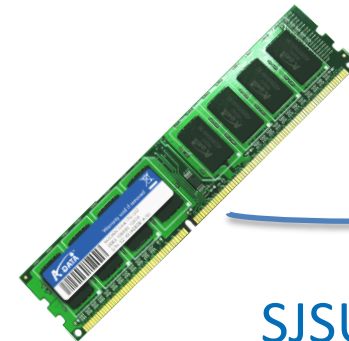
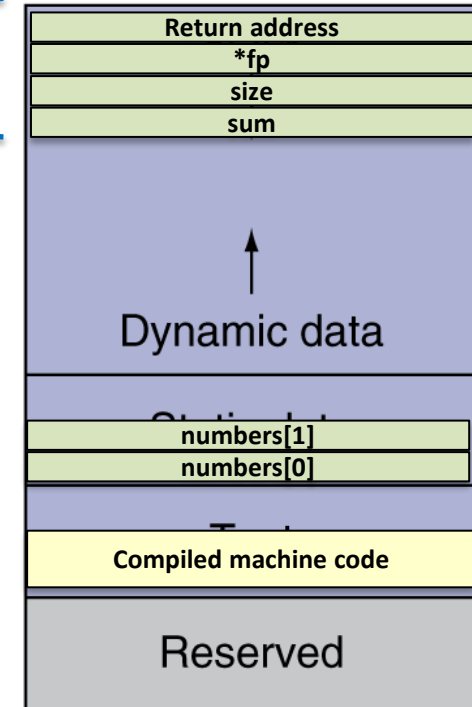
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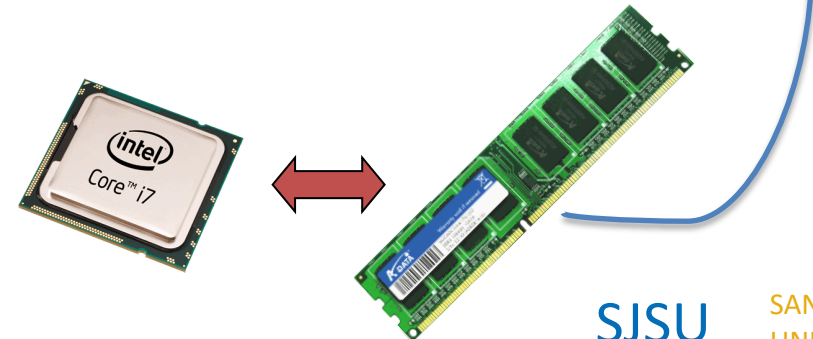
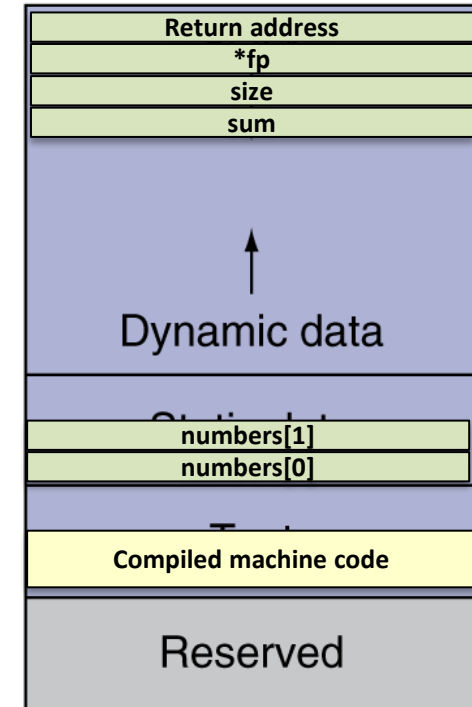
int main (void) {
    myfunction();
    return(0);
}
```

Stack region for myfunction {



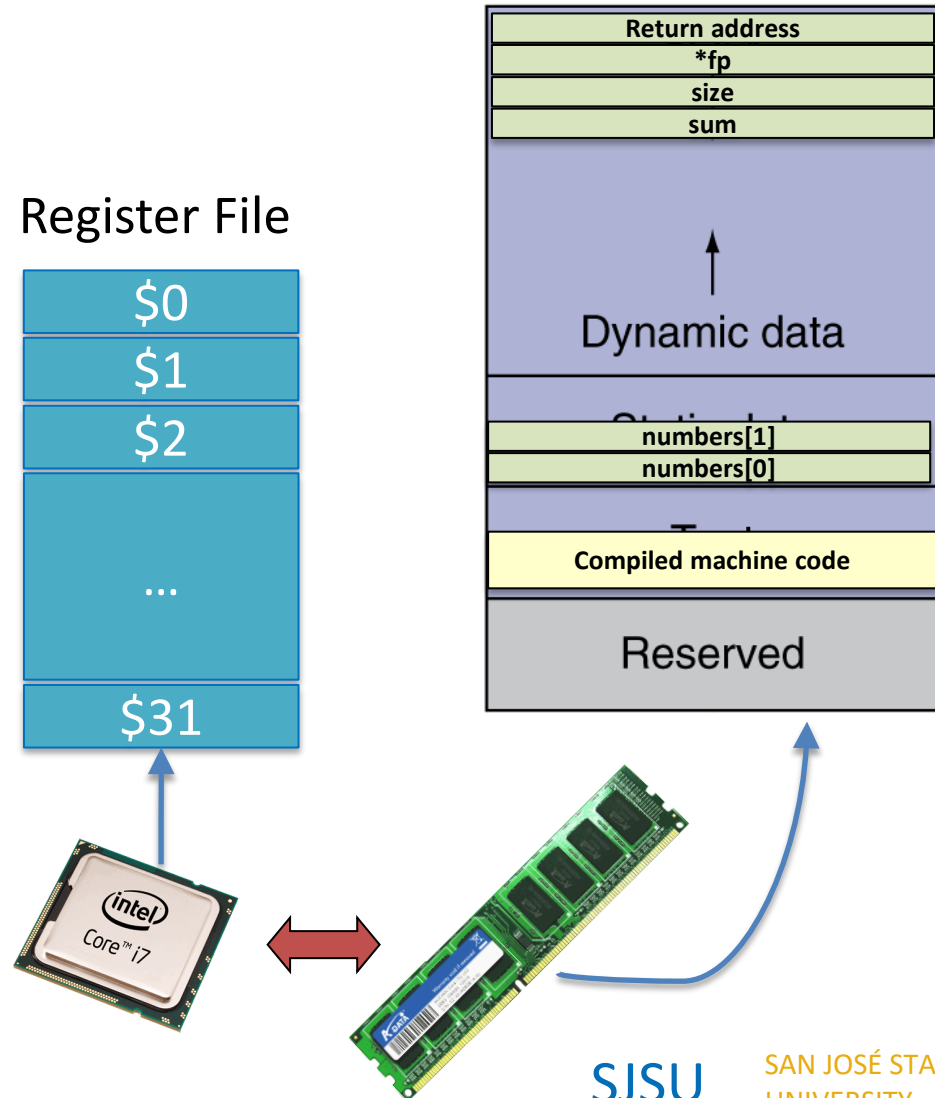
Memory Access Is Slow

- Variables are all stored in Memory, which is **outside of CPU**
 - Slow..
- How can we execute the operations faster?
 - Use **on-chip memory**



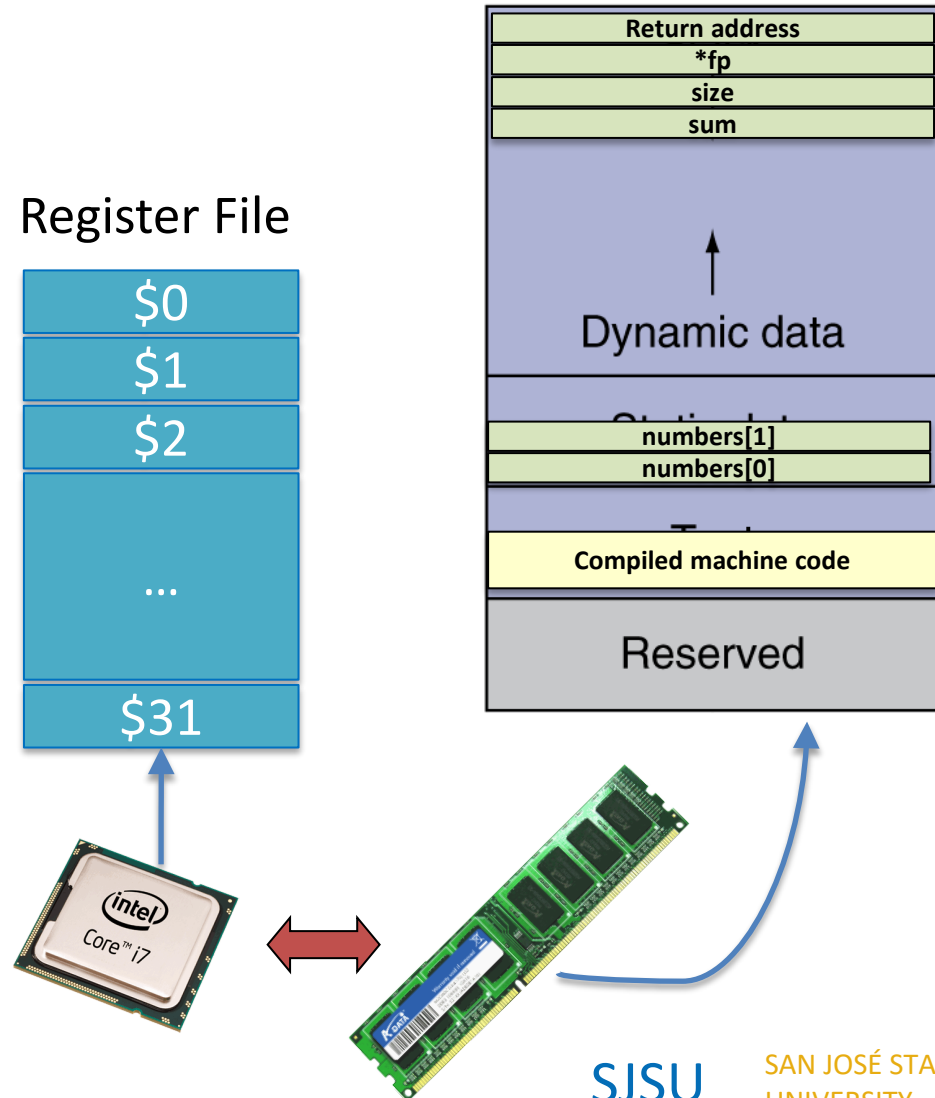
Register File

- **Register File**
 - On-chip memory that stores “Registers”
- **Registers**
 - Temporal space to maintain operand values and calculation results before storing back to memory
 - Presented with “\$” + “register id” in MIPS processor
 - i.e. \$0 : 0th register,
\$1 : 1st register



Typical Execution Steps

1. Load operand values from memory to registers
2. Do computation on registers
3. Move the results from register to memory

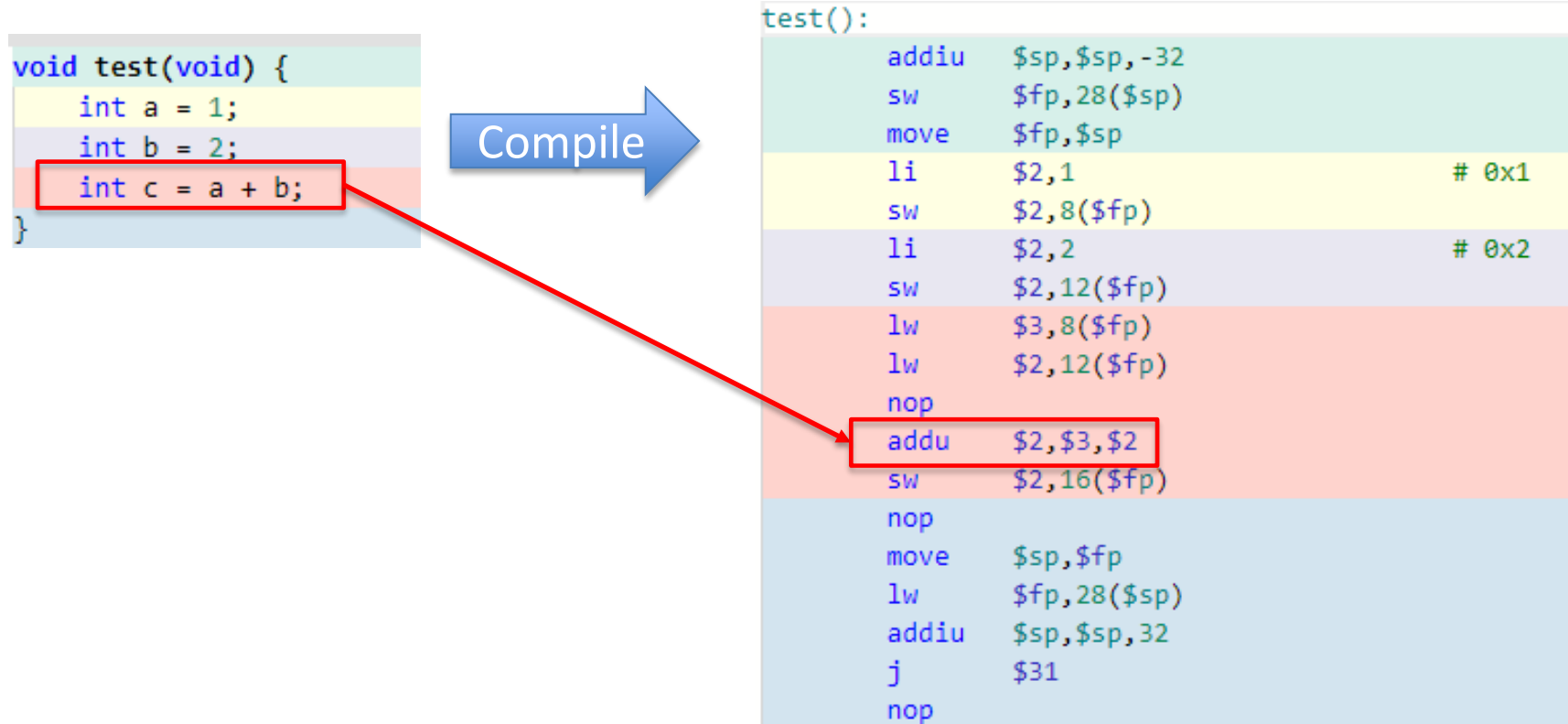


Operations in Hardware

- **Hardware can do one operation at a time**
 - Arithmetic operations
 - add, sub, mult, div, ...
 - Data movement
 - move, load data, store data, ...
 - Logical operations
 - shift, and, or, xor, ...
 - Conditional operations
 - jump, branch on condition, ...
- Format of assembly instructions that uses register operands
 - **Command Result, Operand 1, Operand 2**
 - **E.g. $C = A + B \rightarrow \text{Add } C, A, B$** (A, B, C should be replaced by register id)

HLL vs. Hardware Operations

- HLL are designed to be understood and programmed easily by programmers
- A HLL line may be a combination of multiple assembly instructions



Code Example: Memory & Registers

C code

```
void test(void) {  
    int a = 1;  
    int b = 2;  
    int c = a + b;  
}
```

Compile

MIPS Assembly code

```
test():  
    addiu    $sp,$sp,-32  
    sw       $fp,28($sp)  
    move     $fp,$sp  
    li       $2,1                # 0x1  
    sw       $2,8($fp)  
    li       $2,2                # 0x2  
    sw       $2,12($fp)  
    lw       $3,8($fp)  
    lw       $2,12($fp)  
    nop  
    addu     $2,$3,$2  
    sw       $2,16($fp)  
    nop  
    move     $sp,$fp  
    lw       $fp,28($sp)  
    addiu    $sp,$sp,32  
    j        $31  
    nop
```

Operations involved:

- Value of 'a' is loaded from mem to \$3
- Value of 'b' is loaded from mem to \$2
- **Add** operation done on \$2 and \$3
- Value of 'c' is stored to mem

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void test(void) {  
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}
```

Compile

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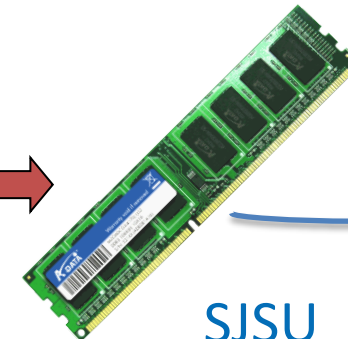
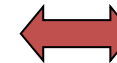
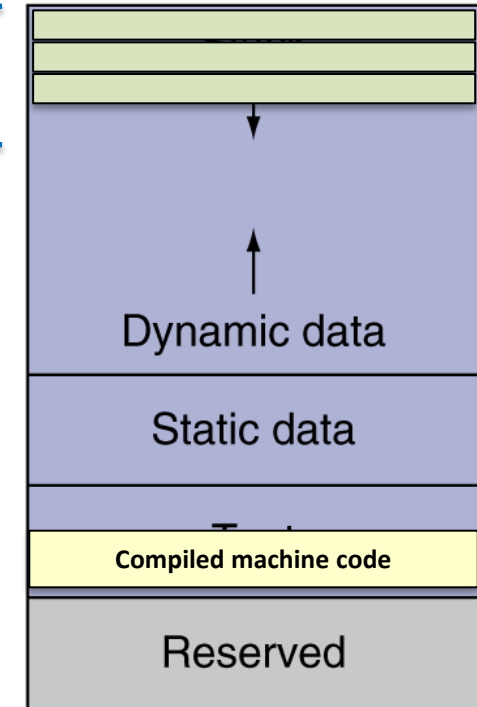
executing line



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```

Stack region
allocation for test()

Register File



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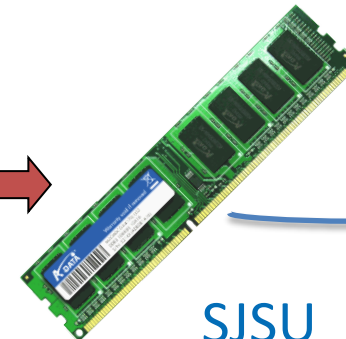
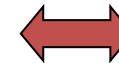
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executing line

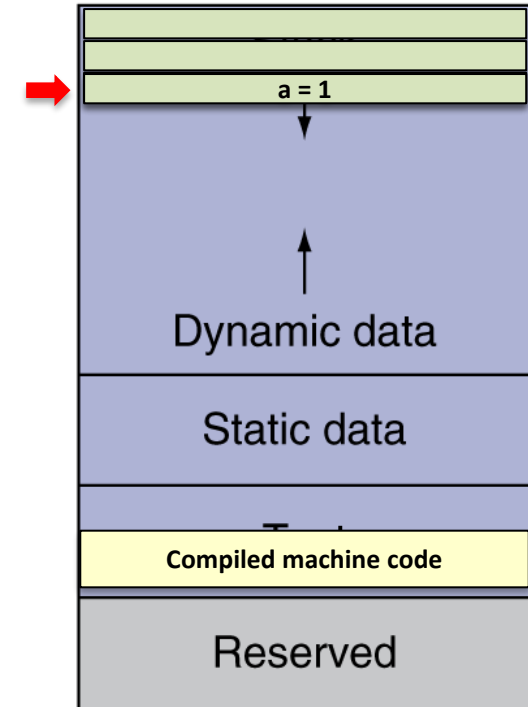


Register File



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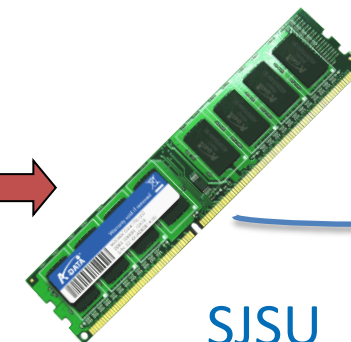
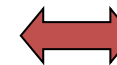
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executing line

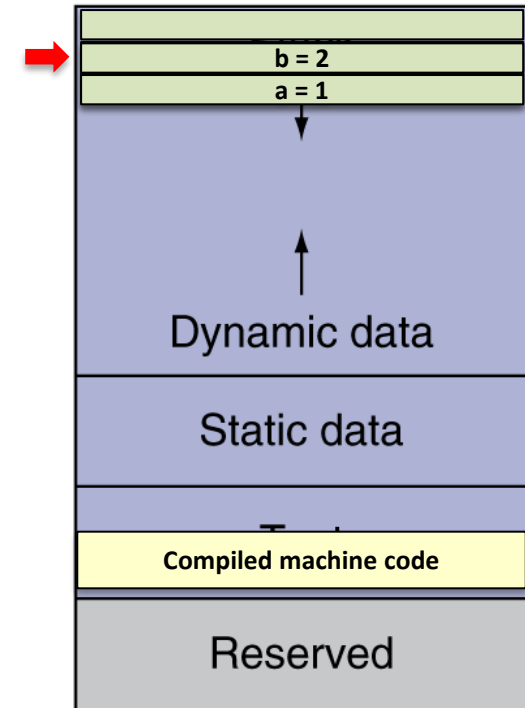


Register File



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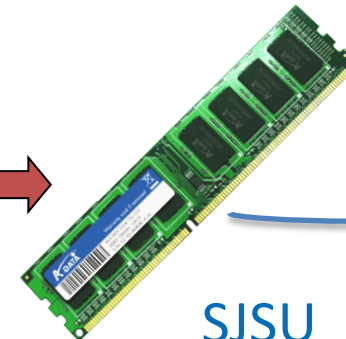
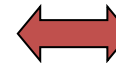
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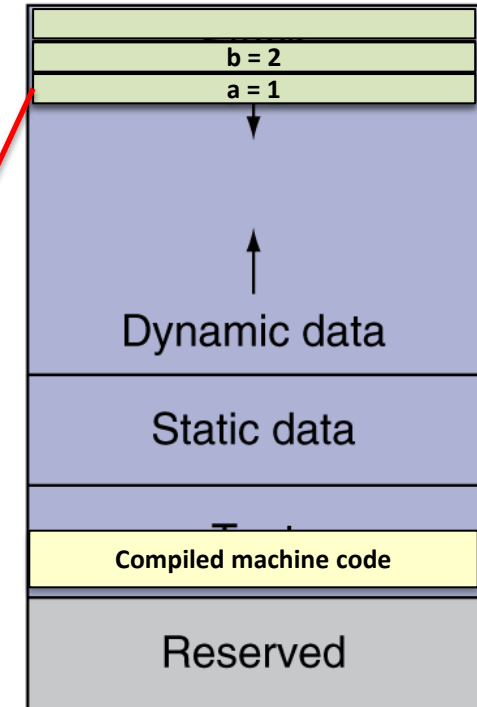


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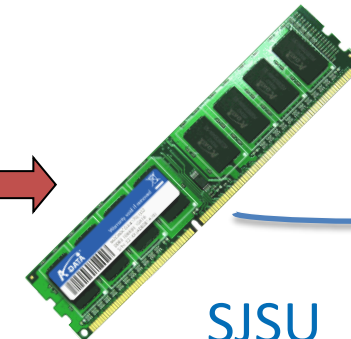
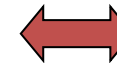
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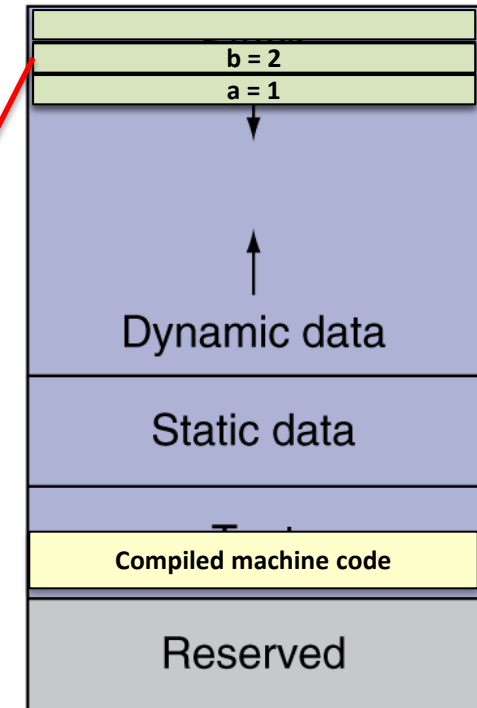


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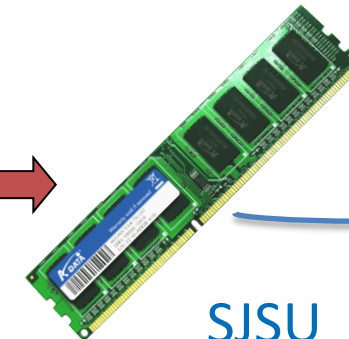
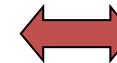
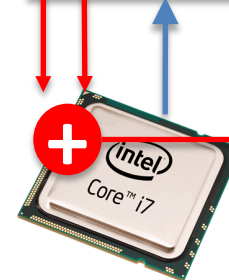
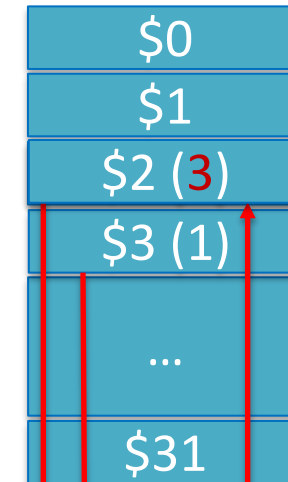
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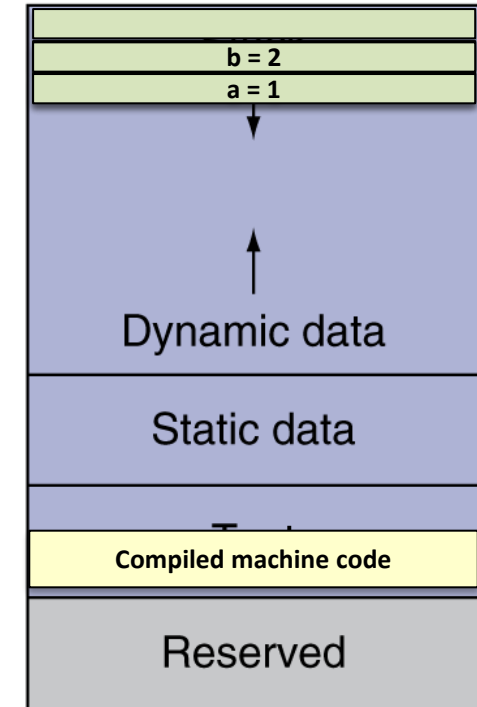


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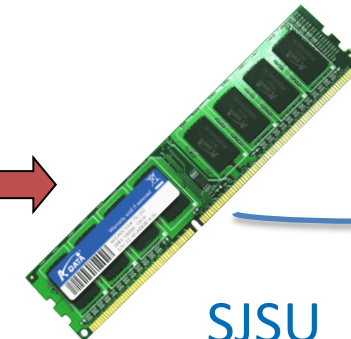
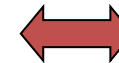
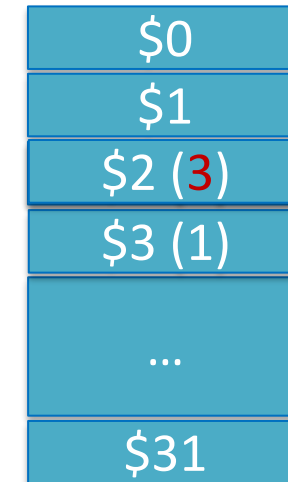
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executing line

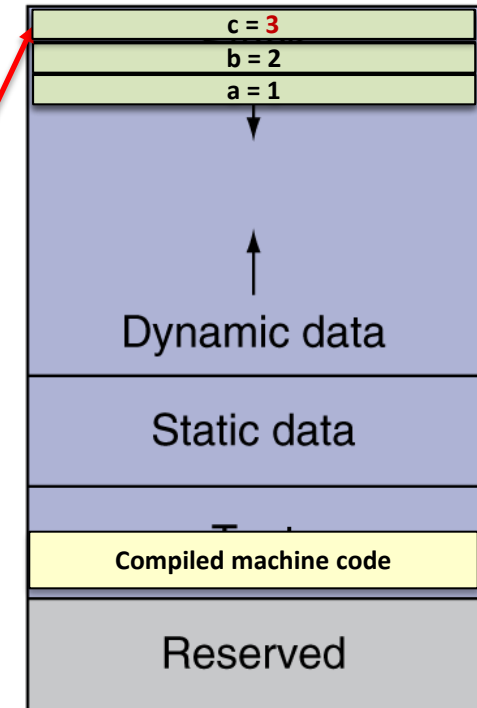


Register File



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Registers of MIPS CPUs

Assembler Name	Register Number	Description
\$zero	\$0	Constant 0 value
\$at	\$1	Assembler temporary
\$v0-\$v1	\$2-\$3	Function return values
\$a0-\$a3	\$4-\$7	Function Arguments
\$t0-\$t7	\$8-\$15	Temporaries
\$s0-\$s7	\$16-\$23	Saved Temporaries
\$t8-\$t9	\$24-\$25	Temporaries
\$k0-\$k1	\$26-\$27	Reserved for OS kernel
\$gp	\$28	Global Pointer (Global and static variables/data)
\$sp	\$29	Stack Pointer
\$fp	\$30	Frame Pointer
\$ra	\$31	Return Address

Exercise: Register Reuse

- **Example:**
 - Assume g, h, i and j are loaded to registers, \$t0, \$t1, \$t2, \$t3
 - HLL:
$$f = (g + h) - (i + j);$$
 - Assembly (with add & sub operations)?
add \$t0, \$t0, \$t1
add \$t2, \$t2, \$t3
sub \$t2, \$t0, \$t2
- **Register values are maintained unless overwritten**

A Brief Review

- **One line of HLL can involve:**

1. **Load operand values from memory to registers**

→ { Data load operations
move, load data, ...

2. **Do computation on registers**

→ { Arithmetic operations
add, sub, mult, div, ...
Logical operations
shift, and, or, xor, ...
Conditional operations
jump, branch on condition, ...

3. **Move the results from register to memory**

→ { Data store operations
move, store data, ...

Operations of MIPS CPUs

C operator	Assembly Operations	Comments
C = A + B	add C, A, B	Add two values
C = A - B	sub C, A, B	Subtract one from another
C = A * B	mul C, A, B	Multiply two values
C = A & B	and C, A, B	Logical AND operation
C = A B	or C, A, B	Logical OR operation
C = A ^ B	xor C, A, B	Logical XOR operation
C = A << shamt	sll C, A, shamt	Shift left by shamt
C = A >> shamt	srl C, A, shamt	Shift right by shamt
If (A < B) C = 1	slt C, A, B	Set if less than
C = Memory	lw C, Memory address	Load value from memory
Memory = C	sw C, Memory address	Store value to memory
If (A == B) go to Addr	beq A, B, address	Jump if A == B
Many more ...		

- Note: A, B, C in the table should be replaced by proper register ids

About MIPS Assembly

- **CPUs use their own assembly languages**
 - Assembly language of ARM, Pentium, Opteron... are all different
- **MIPS Assembly Features**
 - Very similar to ARM Assembly
 - One of the earliest RISC architectures that used pipelined instruction processing
- **Developed by MIPS Technologies**
 - Now maintained by Wave Computing
 - Various generations
 - MIPS I~V
 - MIPS32 (32-bit processor)
 - MIPS64 (64-bit processor)
 - microMIPS..

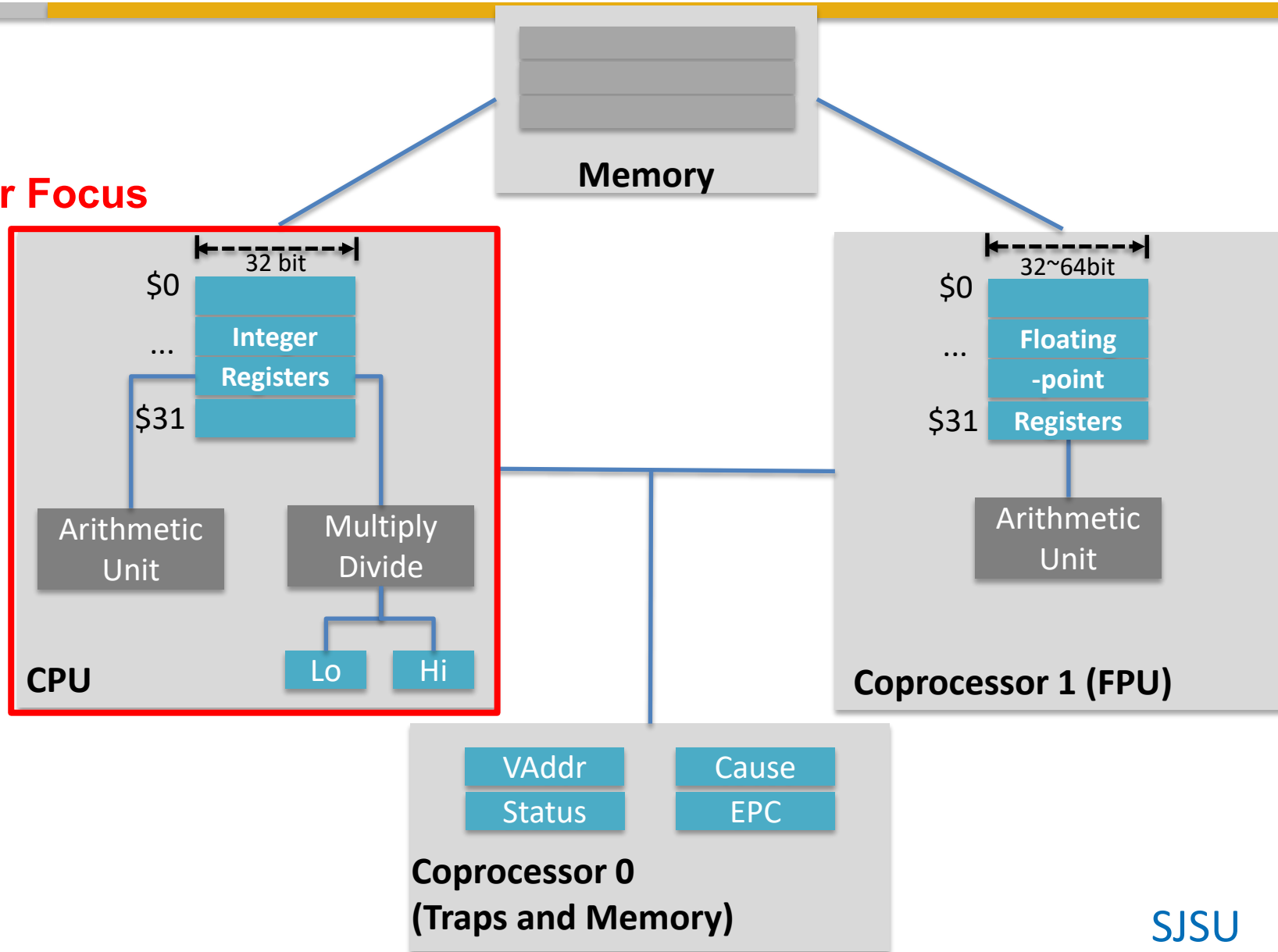


Two types of Computers

- **Reduced Instruction Set Computer (RISC):** MIPS, ARM, ...
 - Operands are in registers
 - Each instruction can do only one operation: Simple but longer codes
- **Complex Instruction Set Computer (CISC):** Intel, AMD, ...
 - Operands can be in registers, stacks, accumulators, in memories
 - Each instruction can do multiple operations: Complex but shorter codes
 - Preferred when memory was very expensive

MIPS Processor Organization

Our Focus



About MIPS32 Processor

- Instructions are 32-bit wide
- Registers and Computing Logic use 32-bit data
- Memory bus is logically 32-bit wide
- 32 general purpose registers (GPRs) for integer and address values
 - A few special ones (i.e. \$zero: constant 0, \$fp: frame pointer, \$sp: stack pointer..)
- 32 floating point registers for floating point operations (not our focus)

Conclusion Time

What is ISA?

Hardware operation definition

What is Microarchitecture?

ISA implementation

What is Register File?

On-chip memory

SAN JOSÉ STATE UNIVERSITY *powering* SILICON VALLEY

