

1 Advance Encryption Standard AES

→ It is a block cipher.

⇒ Round function in AES should be invertable.

- We have understand the following.

1. Round Function.
2. Key Scheduling Algorithm.

1.1 Round Function of AES - 128 bit

f_1, f_2, \dots, f_{10}

(A) $f_1 = f_2 = f_3 = \dots = f_9$

(B) f_{10} is different from $f_i, i = 1, 2, 3, \dots, 9$.

First 9 round functions are exactly same and 10th round function is different from other 9 round functions.

- The first 9 round functions (i.e., f_1, f_2, \dots, f_9) are based on the following.

i. Sub bytes.

$$f_i : \{0, 1\}^{128} \rightarrow \{0, 1\}^{128}$$

ii. Shift rows.

iii. Mix columns.

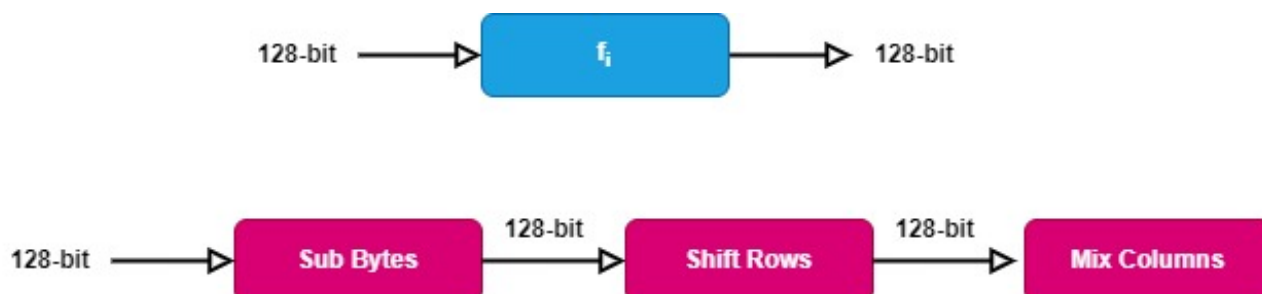
- The 10th round function (i.e., f_{10}) is based on the following.

i. Sub bytes.

$$f_{10} : \{0, 1\}^{128} \rightarrow \{0, 1\}^{128}$$

ii. Shift rows.

$f_i(X) = \text{Mix Column}(\text{Shift Row}(\text{Sub Bytes}(X)))$



1.2 Sub bytes

$$\text{Sub bytes} : \{0, 1\}^{128} \rightarrow \{0, 1\}^{128}$$

$S \rightarrow \text{input.}$

$$S \rightarrow \begin{vmatrix} S_{00} & S_{01} & S_{02} & S_{03} \\ S_{10} & S_{11} & S_{12} & S_{13} \\ S_{20} & S_{21} & S_{22} & S_{23} \\ S_{30} & S_{31} & S_{32} & S_{33} \end{vmatrix}$$

$$S_{ij} \rightarrow 8 \text{ bit.}$$

$P \rightarrow \text{Plain text 128 bit.}$

$$P \rightarrow \begin{vmatrix} P_0 & P_4 & P_8 & P_{12} \\ P_1 & P_5 & P_9 & P_{13} \\ P_2 & P_6 & P_{10} & P_{14} \\ P_3 & P_7 & P_{11} & P_{15} \end{vmatrix} + K_1 \rightarrow (S_{ij})_{4 \times 4} \quad \text{len}(P_i) = 8 \text{ bit}$$

$K_1 \rightarrow 128 \text{ bit. } K_0, K_1, \dots, K_{15}$

$$K_1 \rightarrow \begin{vmatrix} K_0 & K_4 & K_8 & K_{12} \\ K_1 & K_5 & K_9 & K_{13} \\ K_2 & K_6 & K_{10} & K_{14} \\ K_3 & K_7 & K_{11} & K_{15} \end{vmatrix}$$

1.3 Sub bytes

$$S = (S_{ij})_{4 \times 4}$$

$$S : \{0, 1\}^8 \rightarrow \{0, 1\}^8 \quad S(0) = 0$$

$$\text{i. } (c_7 c_6 c_5 c_4 c_3 c_2 c_1 c_0) = (01100011) = (63)_{16}$$

$$\text{ii. } S(S_{ij}) = (a_7 a_6 a_5 a_4 a_3 a_2 a_1 a_0)$$

iii. For $i = 0$ to 7

$$b_i = (a_i + a_{(i+4)\%8} + a_{(i+5)\%8} + a_{(i+6)\%8} + a_{(i+7)\%8} + c_i) \bmod 2$$

$$\text{iv. } (b_7 b_6 b_5 b_4 b_3 b_2 b_1 b_0)$$

$$\text{v. } S^1_{ij} = (b_7 b_6 b_5 b_4 b_3 b_2 b_1 b_0)$$

$$\bullet S : \{0, 1\}^8 \rightarrow \{0, 1\}^8 \quad S(0) = 0$$

$$X \neq 0 \in \{0, 1\}^8$$

$$S(X) = Y \in \{0, 1\}^8$$

$$X = (a_7 a_6 a_5 a_4 a_3 a_2 a_1 a_0), a_i \in \{0, 1\}$$

$$P(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_7 x^7$$

$$\deg(P(x)) \leq 7$$

$$P(x) \in \mathbb{F}_2[x]$$

$$g(x) = x^8 + x^4 + x^3 + x + 1$$

$g(x)$ is a primitive polynomial.

$$(\mathbb{F}_2[x]/\langle g(x) \rangle, +, *) \rightarrow \text{Field.}$$

$$\text{Find the multiplicative inverse of } P(x) \text{ under modulo } (x^8 + x^4 + x^3 + x + 1)$$

$$p(x).q(x) = 1 \bmod (x^8 + x^4 + x^3 + x + 1)$$

$$\Rightarrow p(x).q(x) - 1 = h(x).(x^8 + x^4 + x^3 + x + 1)$$

$$1 = p(x).q(x) + h_1(x).(x^8 + x^4 + x^3 + x + 1)$$

$$\gcd(a, b) = as + bt$$

$$\gcd(P(x), x^8 + x^4 + x^3 + x + 1) = 1$$

How to find $g(x)$?

\Rightarrow Extended Euclidean Algorithm finds $q(x)$

$$q(x) \rightarrow \deg(q(x)) \leq 7$$

$$q(x) = r_0 + r_1x + r_2x^2 + \dots + r_7x^7$$

$$q(x) \rightarrow (r_7r_6r_5r_4r_3r_2r_1r_0) \in \{0, 1\}^8$$

$$S(x) = Y = (r_7r_6r_5r_4r_3r_2r_1r_0)$$

Example

Find $S(01010011) = ?$

$$p(x) = x^6 + x^4 + x + 1$$

$$g(x) = x^8 + x^4 + x^3 + x + 1$$

$$\begin{array}{r} x^6 + x^4 + x + 1 \quad x^8 + x^4 + x^3 + x + 1 \quad (x^2 + 1) \\ \underline{x^2 + x + 1} \end{array}$$

$$\begin{array}{r} x^6 + x^4 + x^2 + x + 1 \\ \underline{x^6 + x^4 + x + 1} \end{array}$$

$$\begin{array}{r} x^2) \quad x^6 + x^4 + x + 1 \quad (x^4 + x^2) \\ \underline{x^6} \end{array}$$

$$\begin{array}{r} x^4 + x + 1 \\ \underline{x^4} \end{array}$$

$$\begin{array}{r} x + 1) \quad x^2 (x + 1) \\ \underline{x^2 + x} \end{array}$$

$$\begin{array}{r} x \\ \underline{x + 1} \end{array}$$

$$\begin{array}{r} 1 \\ \underline{} \end{array}$$

$$\begin{aligned} 1 &= x^2 + (x + 1)(x + 1) \\ &= x^2 + (x + 1)[(x^6 + x^4 + x + 1) + x^2(x^4 + x^2)] \\ &= x^2[x^5 + x^4 + x^3 + x^2 + 1] + (x + 1)[(x^6 + x^4 + x + 1)] \\ &= [(x^8 + x^4 + x^3 + x + 1) + (x^6 + x^4 + x^2 + x + 1)(x^2 + 1)][x^5 + x^4 + x^3 + x^2 + 1] + \\ &\quad (x + 1)[(x^6 + x^4 + x + 1)] \\ 1 &= [(x^6 + x^4 + x + 1)][(x^2 + 1)(x^5 + x^4 + x^3 + x^2 + 1) + (x + 1)] + [(x^8 + x^4 + x^3 + x + 1)][(x^5 + x^4 + x^3 + x^2 + 1)] \end{aligned}$$

Co-efficient of $(x^6 + x^4 + x + 1)$ is multiplicative inverse of $p(x)$ (i.e., $q(x)$)

$$\begin{aligned} \text{So, } q(x) &= (x^2 + 1)(x^5 + x^4 + x^3 + x^2 + 1) + (x + 1) \\ &= (x^7 + x^6 + x^5 + x^4 + x^2 + x^5 + x^4 + x^3 + x^2 + 1) + (x + 1) \end{aligned}$$

$$= x^7 + x^6 + x^3 + x$$

(11001010) → output.

c = (01100011)

$$b_i = (a_i + a_{(i+4)\%8} + a_{(i+5)\%8} + a_{(i+6)\%8} + a_{(i+7)\%8} + c_i) \bmod 2$$

$$b_0 = (a_0 + a_4 + a_5 + a_6 + a_7 + c_0) \bmod 2$$

$$= (0 + 0 + 0 + 1 + 1 + 1) \bmod 2$$

$$= 1$$

$$b_1 = (a_1 + a_5 + a_6 + a_7 + a_0 + c_1) \bmod 2$$

$$= (1 + 0 + 1 + 1 + 0 + 1) \bmod 2$$

$$= 0$$

$$b_2 = 1$$

$$b_3 = 1$$

$$b_4 = 0$$

$$b_5 = 1$$

$$b_6 = 1$$

$$b_7 = 1$$

Output of Sub bytes (1110 1101)

E D

Subbytes (0101 0011) = (1110 1101)

Subbytes (53) = (ED)

1.4 AES S-Boxes

	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	63	7C	77	7B	F2	6B	6F	C5	30	01	67	2B	FE	D7	AB	76
1	CA	82	C9	7D	FA	59	47	F0	AD	D4	A2	AF	9C	A4	72	C0
2	B7	FD	93	26	36	3F	F7	CC	34	A5	E5	F1	71	D8	31	15
3	04	C7	23	C3	18	96	05	9A	07	12	80	E2	EB	27	B2	75
4	09	83	2C	1A	1B	6E	5A	A0	52	3B	D6	B3	29	E3	2F	84
5	53	D1	00	ED	20	FC	B1	5B	6A	CB	BE	39	4A	4C	58	CF
6	D0	EF	AA	FB	43	4D	33	85	45	F9	02	7F	50	3C	9F	A8
7	51	A3	40	8F	92	9D	38	F5	BC	B6	DA	21	10	FF	F3	D2
8	CD	0C	13	EC	5F	97	44	17	C4	A7	7E	3D	64	5D	19	73
9	60	81	4F	DC	22	2A	90	88	46	EE	B8	14	DE	5E	0B	DB
A	E0	32	3A	0A	49	06	24	5C	C2	D3	AC	62	91	95	E4	79
B	E7	C8	37	6D	8D	D5	4E	A9	6C	56	F4	EA	65	7A	AE	08
C	BA	78	25	2E	1C	A6	B4	C6	E8	DD	74	1F	4B	BD	8B	8A
D	70	3E	B5	66	48	03	F6	0E	61	35	57	B9	86	C1	1D	9E
E	E1	F8	98	11	69	D9	8E	94	9B	1E	87	E9	CE	55	28	DF
F	8C	A1	89	0D	BF	E6	42	68	41	99	2D	0F	B0	54	BB	16

(a) S-box

	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	52	09	6A	D5	30	36	A5	38	BF	40	A3	9E	81	F3	D7	FB
1	7C	E3	39	82	9B	2F	FF	87	34	8E	43	44	C4	DE	E9	CB
2	54	7B	94	32	A6	C2	23	3D	EE	4C	95	0B	42	FA	C3	4E
3	08	2E	A1	66	28	D9	24	B2	76	5B	A2	49	6D	8B	D1	25
4	72	F8	F6	64	86	68	98	16	D4	A4	5C	CC	5D	65	B6	92
5	6C	70	48	50	FD	ED	B9	DA	5E	15	46	57	A7	8D	9D	84
6	90	D8	AB	00	8C	BC	D3	0A	F7	E4	58	05	B8	B3	45	06
7	D0	2C	1E	8F	CA	3F	0F	02	C1	AF	BD	03	01	13	8A	6B
8	3A	91	11	41	4F	67	DC	EA	97	F2	CF	CE	F0	B4	E6	73
9	96	AC	74	22	E7	AD	35	85	E2	F9	37	E8	1C	75	DF	6E
A	47	F1	1A	71	1D	29	C5	89	6F	B7	62	0E	AA	18	BE	1B
B	FC	56	3E	4B	C6	D2	79	20	9A	DB	C0	FE	78	CD	5A	F4
C	1F	DD	A8	33	88	07	C7	31	B1	12	10	59	27	80	EC	5F
D	60	51	7F	A9	19	B5	4A	0D	2D	E5	7A	9F	93	C9	9C	EF
E	A0	E0	3B	4D	AE	2A	F5	B0	C8	EB	BB	3C	83	53	99	61
F	17	2B	04	7E	BA	77	D6	26	E1	69	14	63	55	21	0C	7D

(b) Inverse S-box