"You don't have to be a genius to code, you just have to be persistent."





# **Graphs 1**



## Agenda:

Introduction to Graph
2. Types of Graphs

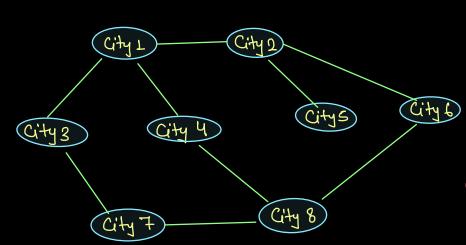
creation { 3. How to store data in Graph

{ 4. BFS (Breadth First Search) traveka!

Seanching { 5. Is Path Available from En Graph Source to Destination

#### Introduction to Graph

information of cities and their connectivity. store Warnt 40



with the help of graph data sinchre we com amplemed the storage of connection.

cities = vertex Links = Edge

No. of Verten ? -> 8

No , of Edga ? -> 9

Collection of vertex and logs one known as Graph.

neighbour of city 1 -> city 2, city 8, & city 4

neighbour of city 8 - City 6, City 4, City 7

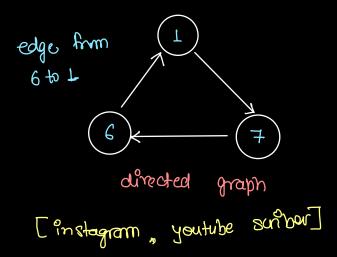
nbr(citys) = City2

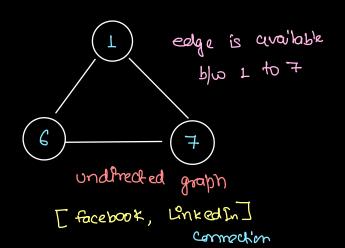
one know or neighbour. connect vertex Direct

#### ~~~~~~~~~

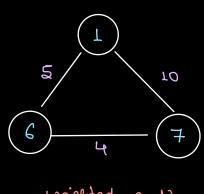
### Types of Graphs

I. Based on type of edges;

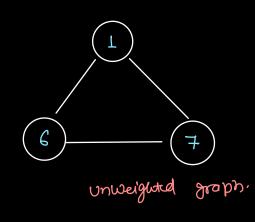




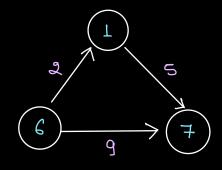
2. Based on Edge wt. present or not:



weighted graph



2. Combination of above types are also possible:



clirected weighted grapm

#### How to store data in Graph

There is too famous implement action available for graph.

- 1 Adjacency moutrix
- (2) Adjacency (16)

## 1. Adjacency matrix:

verter = 7, Edges = 8

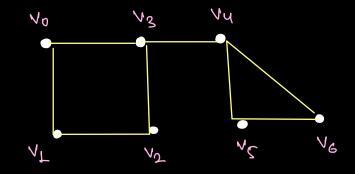
graph.

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			J						ろ ア	_ ପ	3	
	O	7	2	3	Ч	S	6	9>	3	3	7	  -
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7		0		D	0	0	G	Pnt u= edge[i][o];	5	4	5	
		O		)					6	Ч	6	
2	O		O	1	0	0	O	int V= edge[i][1];	7	5	6	レ
3		0		0	4	О	$\mathcal{C}$	//mate a connection	,	ondi	recte	d
٥	1	O		)	+	0		blo u fl 4				
4	೦	O	O	1	<u>ي</u>	1	1	u-v				
S	0	D	0	0		O		V - 1 U				
כ			O	O				- grap[u][v] = 1				
6	0	0	ව	0		1	0					
								graph[v][u] = 1				

	0	7	2	3	Ч	S	6			
0	0	T	0		0	0	0			
Τ	T	0		Q	0	0	G			
2	G		O	1	0	0	O			
3		0		0	1	O	O			
4	O	O	ð	1	ව	1	1			
S	00000									
6	000000000000000000000000000000000000000									
undirected + un weighted										
			8	rah)	h					

No. of vertex  $\rightarrow$  Row length or Column length

<u>-</u>1



## Weighted + Directed graph:

Vtx = 7, Edge = 8

Int[][] graph = new int[vtx][vtx];

	O	7	2	3	Ч	S	6
0	0	SZ SZ	0	10	O	O	0
Τ	O	0	4	Q	0	0	O
2	G	0	Ð	9	0	0	O
3	O	0	0	0	ळ	O	O
4	O	O	O	0	0	7	20
S	0	D	O	0	0	Ø	प्र
6	0	0	ව	0	0	0	0

int w= edge[i][o];

int w= edge[i][i];

int wt=edge[i][i];

Il Edge only from

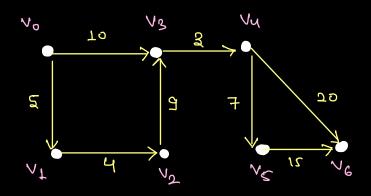
u to v with

weight with

graph[u][v]= wt;

	U	S	ωt
0	0	3 -	-10 /
$\tau$	0	7 -	-15 M
2	2	3 -	19 V
3	3	4 -	→3 ~
4	$\mathcal{T}$	ე -	741
S	4	5 -	77
6	Ч	6 -	→ 20 V
7	5	6	→ \\\\

	O	7	2	3	Ч	S	6
0	0	S	0	10	O	0	0
Τ	0	0	4	Q	0	0	G
2	G	0	Ð	9	0	0	O
3	O	0	0	0	ळ	O	O
4	O	O	O	0	0	7	20
S	0	D	0	0	0	D	15
6	0	0	ව	0	0	0	0



Major disadvandage of Adjacency moun's:

Major disadvantyr is wastage of memory, thatis why most of the time we will ameider Adjacency wist Prostead of Adjacency mentric

# 2. Adjacency 4875

undirected graph.



0	0	3
$\mathcal{T}$	0	1
2	<u>0</u>	3
3	3	4
Ч	$\mathcal{T}$	2
S	4	5
6	Ч	6
ユ	5	6

Vo	V3	٧٧	
VL	V2	VS	V6

Q	7	2	3	4	2	6	
							1
3	0 2	٦ ٦	0 2 4	3 5 6	4 6	Y (8	

Amay List < Amay List < Integer>> graph

Implementation PI: Understomeling

O	T		2	3	4		2		6	
		Ι.								7
2g 1	20		<u></u> ~	024	ન્ય િ	T	4		4 5	

Vtx=7, Edge=8

<del>ا</del> ا	0 _
4	0
<u>ე</u>	7 3
	~

int us edgelijloj; -- 2

int v= edge [i][i]; → 1

graph. get (2). add (1);

graph, get(1)-add(2);

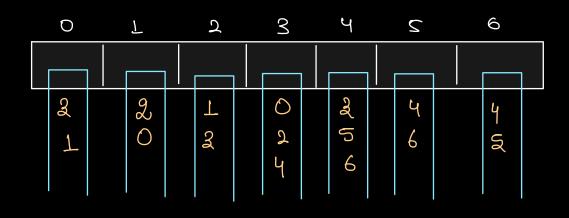
int uz edge [i][o];

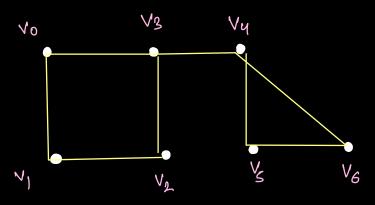
int += edge [i][1];

graph-get (u). odd(+);

graph get (V). add (W);

M: 2, V28



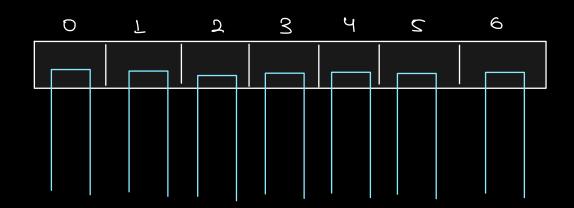


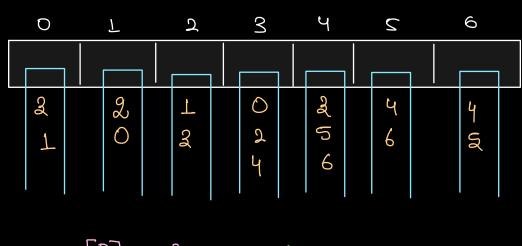
Array List < Intgu>> grapn = rrew AL<>();

for (int v=0; v < v+x; v++) {

 grapn.add (new AL<>());

}





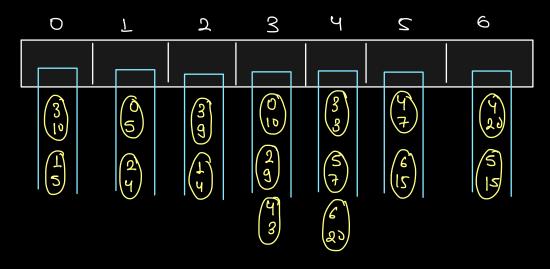
Screen Shot of implementation.

```
import java.util.*;
class Main {
    public static void display(ArrayList<ArrayList<Integer>> graph) {
        int n = graph.size(); // number of vertex
        for(int v = 0; v < n; v++) {
            System.out.print("[" + v + "] -> ");
            for(int nbr : graph.get(v)) {
                System.out.print(nbr + " ");
            System.out.println();
       public static void main(String args[]) {
           int vtx = 7;
           int[][] edges = {
               \{0, 3\}, \{0, 1\}, \{2, 3\}, \{3, 4\}, \{1, 2\}, \{4, 5\}, \{4, 6\}, \{5, 6\}
           };
           ArrayList<ArrayList<Integer>> graph = new ArrayList<>();
           for(int v = 0; v < vtx; v++) {
               graph.add(new ArrayList<>());
           }
           // container of graph is ready,
           for(int e = 0; e < edges.length; e++) {</pre>
               int u = edges[e][0];
               int v = edges[e][1];
               graph.get(u).add(v);
               graph.get(v).add(u);
           }
           int src = 5;
           bfs(graph, src);
           // display(graph);
   }
```

How to make welghted graph?

Ntx=7, Edge=8





	u	\$	ωt
0	0	3 -	-104
$\mathcal{T}$	0	7 -	<u> 15                                   </u>
2	2	3 -	19/
3	3	4 -	<b>→3</b> レ
Ч	$\mathcal{T}$	ე -	- 4 W
S	4	5 -	77 ~
6	Ч	6 -	→ 20 V
7	5	6	→ 15 V

```
'public class fair {

Int nbr;

Int wt;

Poir (int nbr, int wt) {

this nbr: nbr:

tws. wt = wt;

3
```

for lint u=0; V< V+x; V++) &

| graph.add (new AL<>());
}

for(int e=0; e< edgs. leyth; e+t) {

int u= edge[e](o);

int v= edge[e][i];

int wt= edge[e][i];

groph.get (u). add (new Pair (v, cot));
grapm.get (v). add (new Poir (u, wt));

Expected 0/3 :

Example

[0] - 1-1, 2-3

[1] - 0-1, 7-3

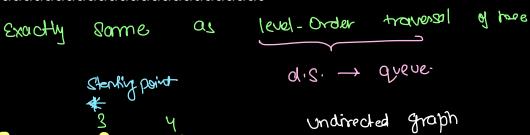
! ete,

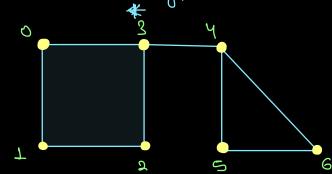
Topo

display.

BFS (Breadth First Search)

[Traversal of graph]





steps of BFS;

- Remove
- Remove
- work - Phint
- add Unvisited
neighbork.
- mark
- Add.

NOTE: Either starting point is given in problem or we can start traversal from only point.

```
public static void bfs(ArrayList<ArrayList<Integer>> graph, int src) {
   // number of vertex in graph
   int n = graph.size();
   // Make a visited array i.e. boolean array to mark visit of vertex
   boolean[] vis = new boolean[n];
   // Make a que to add vertex and starting of BFS
   Queue<Integer> qu = new ArrayDeque<>();
   // add source in queue and mark it true i.e. visited
   qu.add(src);
   vis[src] = true;
   while(qu.size() > 0) {
       // remove
       int rem = qu.remove();
       // work -> printing of vertex
       System.out.print(rem + " ");
       // add univisted neighbour
        for(int nbr : graph.get(rem)) {
            // if neighbour is univisted, mark it and add it
            if(vis[nbr] == false) {
                vis[nbr] = true;
                qu.add(nbr);
```

#### Is Path Available from Source to Destination

Given an undirected graph, source node and destination node. Cheek if there is a path Available from source to destination or not.

Src Ex1: Src=3, dst=6 ons: true (6) dstsrc Sre= 3, det; 6 Ex2: ons: Palso, 6) dst Ex 3: T ELC . 0 Sre = 1, dst = 4 on: false. source point, in End Finally Algo hom BAS solution: Start vis[det]: 8tadus of check SLC, 0 Remove Add unvisited 1 hbr

[ [fab] ziv