

## Today's Content

- a) Search  $k$  in BT
- b) LCA in Binary Tree
- c) No: of Nodes at  $k$  from given Node
- d) Fill right Pointer

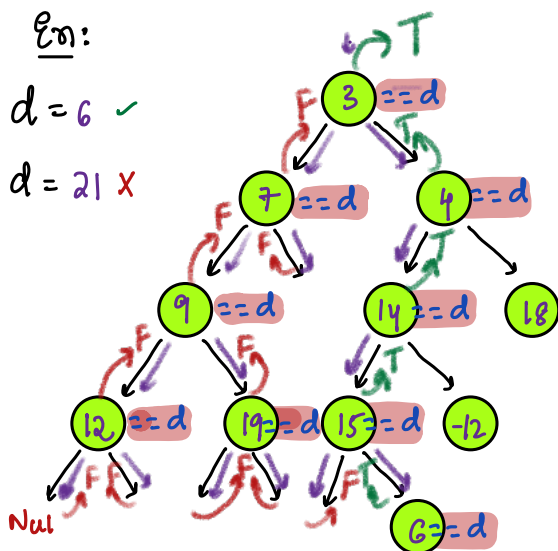
Q1: Given a B.T which contains all unique values.

↓ Search if there exists a  $d$  in Binary Tree

Ex:

$d = 6$  ✓

$d = 21$  ✗



Ass: Given Tree, Search for  $d$  return T/F

bool check(Node root, int d) {

if (root == null) { return false; }

if (root.data == d) { return true; }

if (check(root.left, d) || check(root.right, d))

{ return true;

} else { return false; }

Path  $d: 6 \Rightarrow 3 \ 4 \ 14 \ 15 \ 6$

Obs: If we store all nodes, which returns true, it gives Path between source & root node.

Note: According to above logic: 6 15 14 4 3 : Source node to root node

ArrayList<Node> al;

boolean check(Node root, int s) { TC:  $O(N)$  SC:  $O(H)$

↳ // height of Tree

if (root == null) { return false; }

if (root.data == s) { al.add(root); return true; }

if (check(root.left, s) || check(root.right, s))

{ al.add(root); return true;

} else { return false; }

ArrayList<Node> getPath(Node root, int s) {

al = new ArrayList<>(); // Only before function we initialize

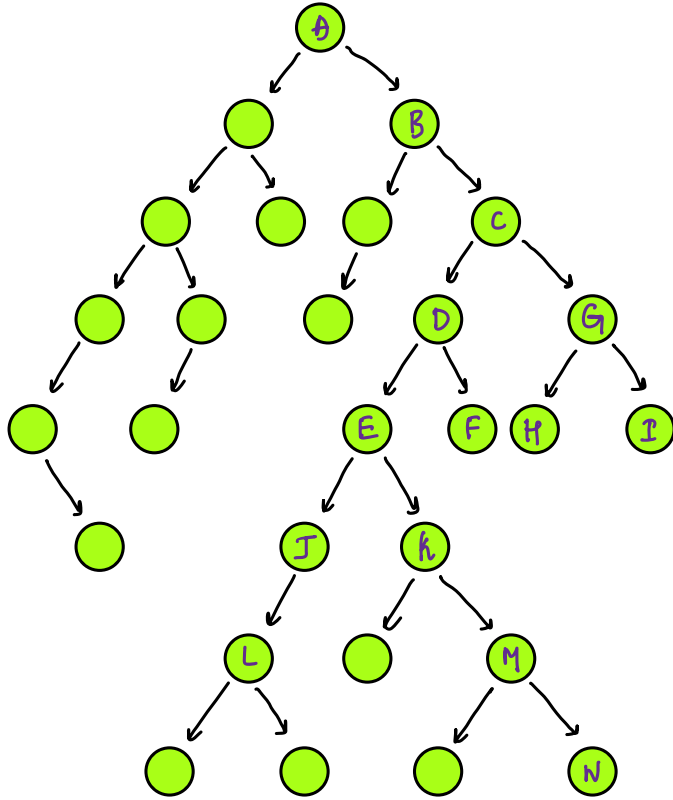
check(root, s);

// reverse arraylist al

return al;

LCA: Lowest Common Ancestor: lowest common ancestor is defined between 2 nodes  $p$  &  $q$  as the lowest level node in Tree, which is an ancestor to both  $p$  &  $q$ . Ancestor: itself/parents/grandparent.../root node

Note: We allow a node to be ancestor to itself.



Ex:  $LCA(K, L)$ : Node E is lowest common Ancestor

K: A B C D E K

L: A B C D E J L

Ex:  $LCA(J, M)$ :

J: A B C D E J

M: A B C D E K

Ex:  $LCA(D, C)$ :

D: A B C D

C: A B C

Node  $LCA(Node\ root, int\ p, int\ q)$  { Tc:  $O(N)$

//Step1:

1. get path of  $p$  from root

2. get path of  $q$  from root

//Step2:

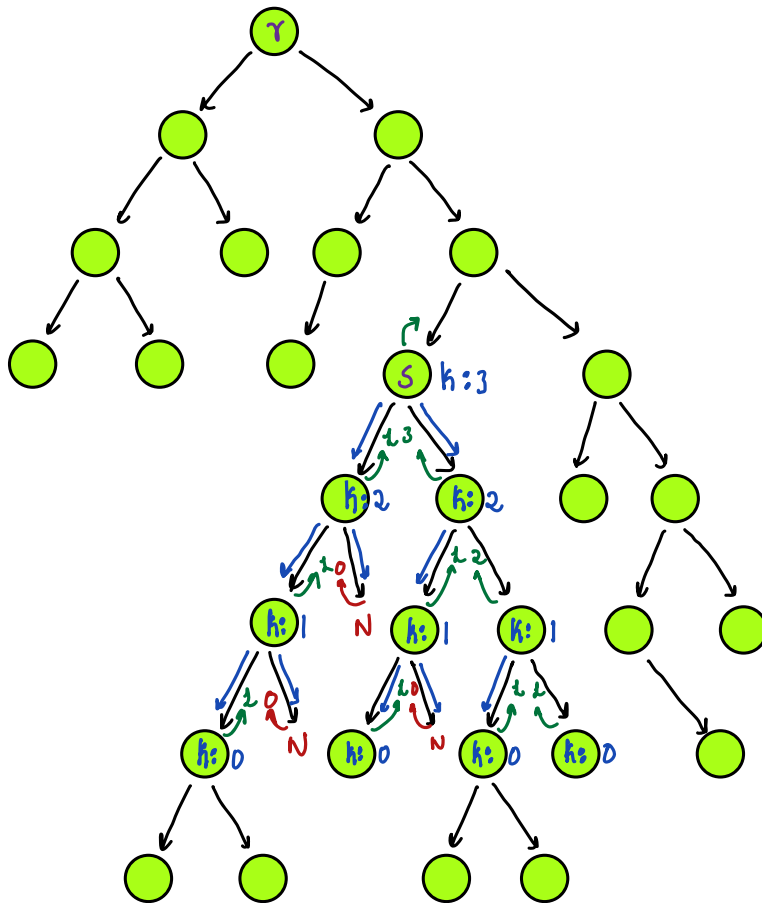
Iterate on path & get last common node.

Q2: Given a source node how many nodes are there distance exactly k.

Note1: All Nodes be below source.

Note2: Distance between nodes calculate, no: of edges between them

Ex:



```
int below(Node s, int k) {
```

```
    if (s == null) { return 0; }
```

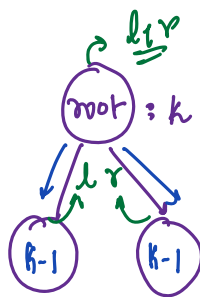
```
    if (k == 0) { return 1; }
```

```
    int l = below(s.left, k-1);
```

```
    int r = below(s.right, k-1);
```

```
    return l+r;
```

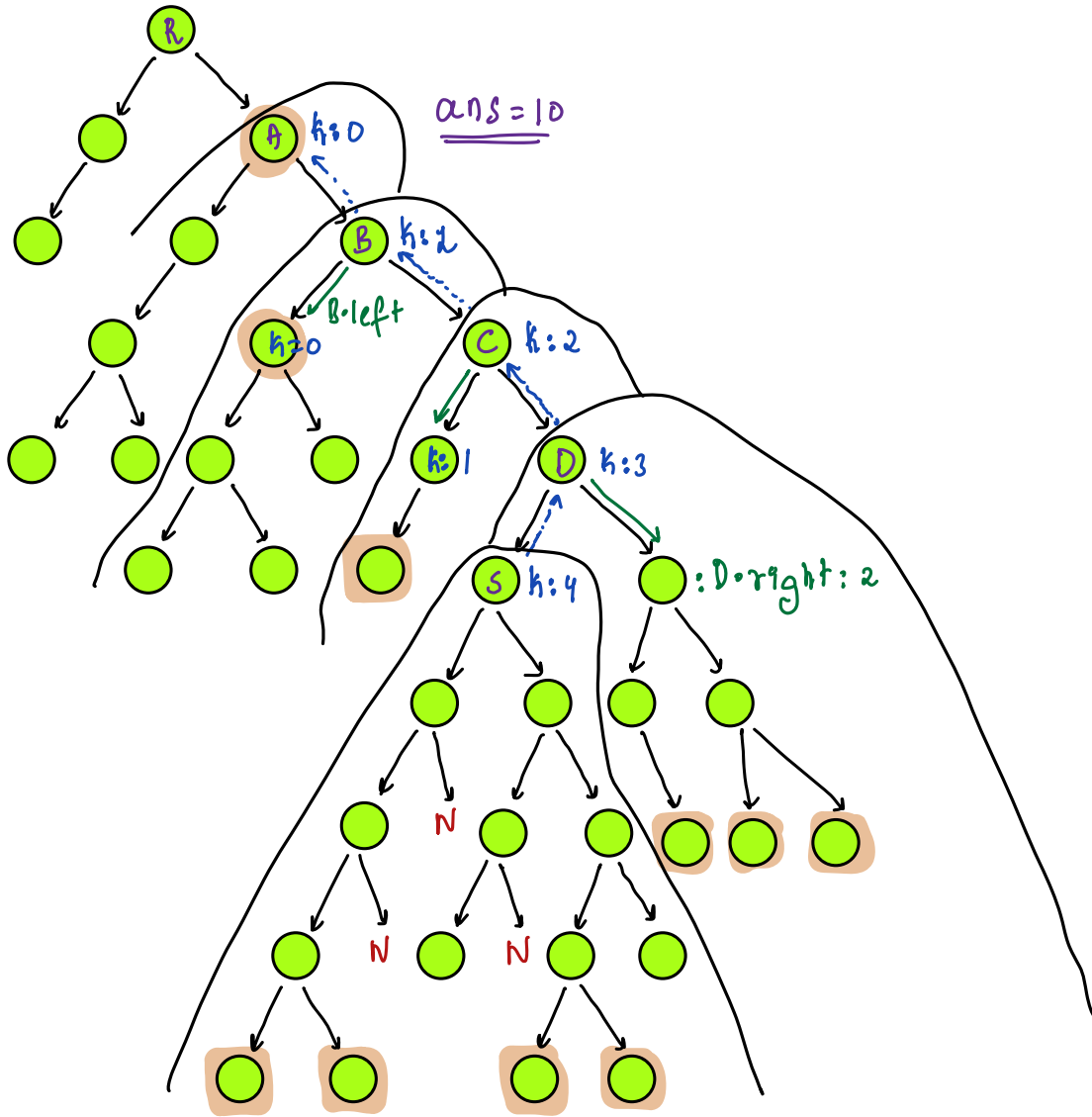
```
}
```



39 Calculate no: of nodes are at distance  $k$  from source.

Note: Only source node value is given, we need to search it first

Note: Binary Tree contains only distinct values



Step 1: Get path from source  $\rightarrow$  root node

Note: In path for node at index  $i$  's child is at index  $i-1$

Steps: Path  $k = \text{below}(s, k)$   
 $\text{below}(s, k)$

0 1 2 3 4 5

$s$   $D, k-1$   $C, k-2$   $B, k-3$   $A, k-4$   $R$

$s, k$   $D, \text{right } k-2$   $C, \text{left } k-3$   $B, \text{left } k-4$   $C, k-2$   $B, k-3$

$\text{ans} += 4$   $\text{ans} += 3$   $\text{ans} += 1$   $\text{ans} += 1$   $\text{pf } k-4 == 0: \text{ans} += 1 \text{ break;}$

0 1 2 3 4

$p: s D C B A$

```
int countNodes(Node r, int s, int d) { Tc: O(N) Sc: O(H)
```

// Step 1: Get path from source  $\rightarrow$  root.

ArrayList<Node> p = getPath(r, s); // Path from root  $\rightarrow$  source

p.reverse() // reverse it

int ans = 0

ans = ans + below(p.get(0), h);

int N = p.length();

$\circ [i-1]$

for (int i = 1; i < N; i++) {

int dist = d - i;

if (dist == 0) { ans = ans + 1; break; }

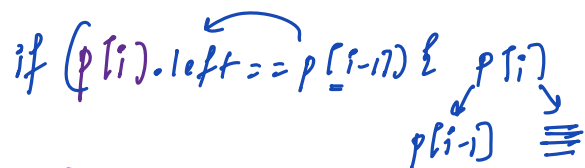
// For  $i^{\text{th}}$  index, child is at  $i-1^{\text{th}}$  index

if (p.get(i).left == p.get(i-1)) {

ans = ans + below(p.get(i).right, dist - 1)

else {

ans = ans + below(p.get(i).left, dist - 1)



if (p[i].right == p.get(i-1))



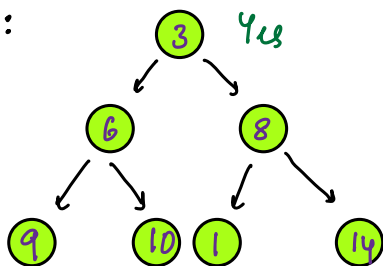
return ans;

}

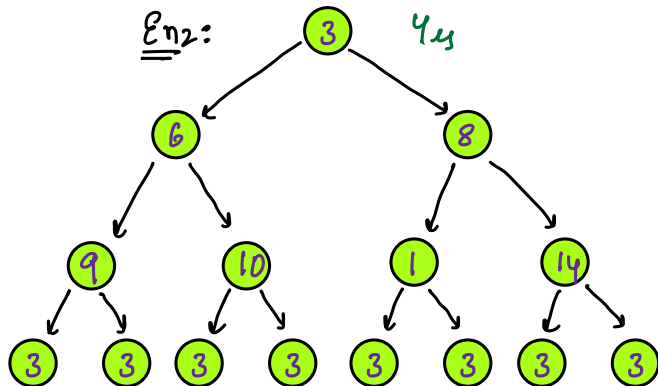
## Perfect Binary Tree:

A binary tree in which all the leaf nodes are at same depth and all non leaf nodes have exactly 2 children

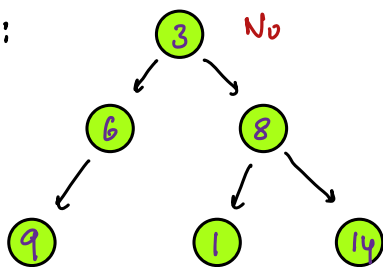
Ex1:



Ex2:



Ex3:



Ex4:

