

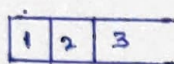
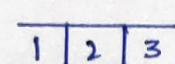
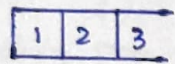
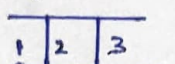
## Queue

- It is also a Linear Data Structure which follows FIFO.
- FIFO (First In First Out), for ex: If we are in a queue, the first person to come in, he will go out first.

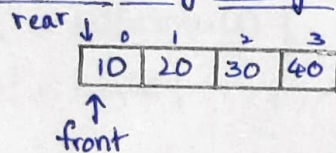
Syntax:-

Interface, not or a class  
`Queue <DataType> queue = new LinkedList<>();`

## Operations:

Stack	Queue
1. push()	1. add()
	
2. pop()	2. poll() → the first inserted element will be returned
 <u>ans: 3, 2, 1</u>	 <u>ans: 1, 2, 3.</u>
3. peek()	3. peek()
last element → 3	first Element → 1
4. isEmpty() → True/false	4. isEmpty() → true/false

## Queue Implementation Using Array:



rear = points to the last element  
front = points to the first element.

- Treat the array as Circular Array, even though the array is not circular array, rear will increment like it is circular.
- Take 2 pointers, front and rear, initially pointing to same location.
- Insert / add() / enqueue()
- a) check overflow condition ⇒ insert
- Delete / poll() / dequeue()
- b) check Underflow condition ⇒ f++



```
void Enqueue (int number) {
```

```
    rear ++;
```

```
    rear = (rear + 1) % arr.length;
```

```
    if (rear == front)
```

```
    { s.o.pln("Overflow");
```

```
        // we need to reset our rear pointer.
```

```
        if (rear == 0)
```

```
            rear = arr.length - 1;
```

```
        else
```

```
            rear = arr.length rear - 1;
```

```
    }
```

```
    else {
```

```
        queue[rear] = number;
```

```
    }
```

```
int dequeue () {
```

```
    if (front == rear)
```

```
    { s.o.pln("Underflow");
```

```
        return -1;
```

```
    }
```

```
    else {
```

```
        front = (front + 1) % arr.length;
```

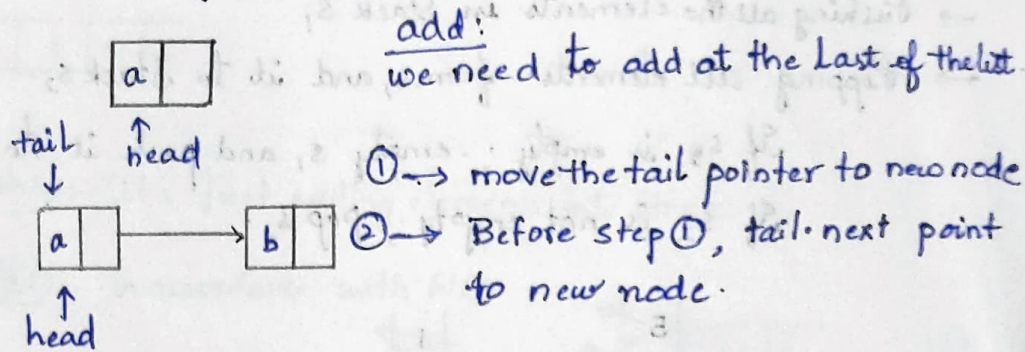
```
        number = queue[front];
```

```
        return number;
```

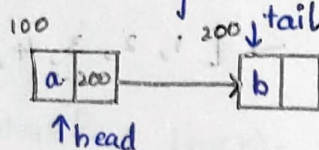
```
    }
```

```
}
```

# Queue Implementation Using Linked List:

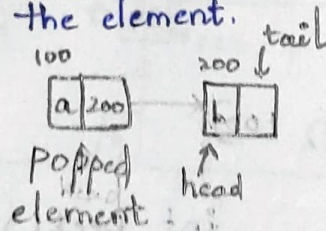


So, now tail points to b node.



Pop():-

- popping the head ; → pointing the head to head.next;
- when pop the element.



```
void push (int n){  
    Node node = new Node(n);  
    if (head == null || tail == null) {  
        head = tail = node;  
    }  
    else {  
        tail.next = node;  
        tail = node;  
    }  
}
```

```
int pop(){  
    if (head == null) {  
        return -1; // queue is empty  
    }  
    Node node = head;  
    head = head.next;  
    if (head == null) {  
        tail = null; → reaching  
                        the tail  
                        if queue  
                        becomes  
                        empty.  
    }  
    return node.data;  
}
```



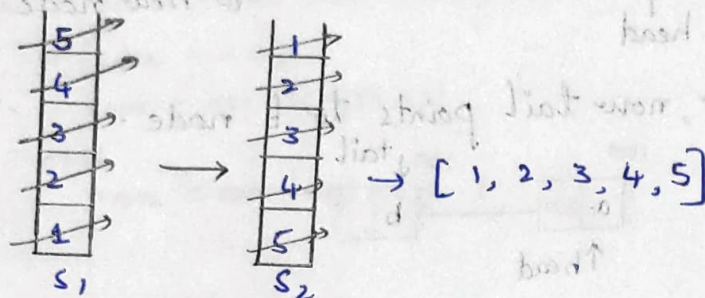
## Using Stacks:

→ Pushing all the elements in Stack  $S_1$

→ Popping all elements from  $S_1$  and push it to Stack  $S_2$ .

If  $S_2$  is empty; empty  $S_1$  and push it to  $S_2$  and then pop.

If  $S_2$  is not empty; pop  $S_2$ .

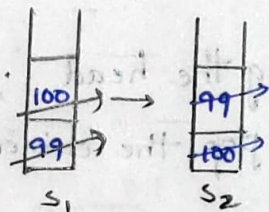


→ add, more elements to  $S_1$

→  $S_2$  is empty, so empty  $S_1$  and push to  $S_2$ .

→ pop  $S_2$ .

$[1, 2, 3, 4, 5, 99, 100]$



```
void push(int n){
```

```
    S1.push(n);
```

```
}
```

```
int pop(){
```

```
    peek();
```

```
    return S2.pop();
```

```
}
```

```
int peek(){
```

```
    if (S2.isEmpty()){
```

```
        while (S1.isEmpty()){
```

```
            S2.push(S1.pop());
```

```
        }
```

```
    }
```

```
}
```

```
boolean empty(){
```

```
    // check both S1 and S2
```

```
    return S1.isEmpty() && S2.isEmpty();
```

```
}
```

# Using One Stack: Reverse a Stack Using Recursion

1 2 3 4 5

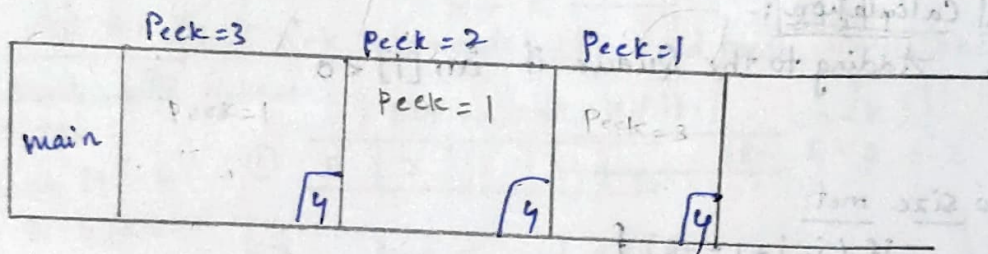
5  
4  
3  
2  
1

Push():- It's just adding elements into stack

Pop():- in accordance with FIFO.

1. int stackPeek = stack.peak();
2. stack.pop();
3. rev(stack);
4. insertAtBottom(stackPeek, stack);

1 2 3



2 2 1

→ Pop 1

2

→ Pop 2

⇒

3 2 1

insertAtBottom(int stackPeek)

{

if (stack.isEmpty())

{

stack.push(stackPeek);

}

else {

int peek = stack.peak();

stack.pop();

insertAtBottom(stackPeek);

stack.push(peek);

}

}



# \*\* First -ve Integer in Every Window of Size k:

I/P: -8, 2, 3, -6, 10, k=2

O/P: -8, 0, -6, -6

- For each window, check if there is any negative number.
- If negative number is available get the first negative number.
- If not result is zero.

Algorithm:-

j (right end)  
-8, 2, 3, -6, 10  
i (window start)

```
while (j < arr.length)
{
    // initial calculation
    // If window hits 'k', other calculation
    i++;
    j++;
}
```

## 1. Initial Calculation:-

Adding to the Queue if  $arr[i] < 0$

-8

## 2. Window size met:

```
if (j - i + 1 == k) {
    if (!queue.isEmpty()) {
        s.o.pln(queue.peek());
    }
}
```

K=2

-8, 2, 3, -6, 10

-8

j++,  $\Rightarrow$  j is at 2

$j - i + 1 \Rightarrow 1 - 0 + 1 = 2$

res[0] = -8

arr[0] = queue.peek

= -8 ✓

So queue.remove

1

$j - i + 1 \Rightarrow 2 - 1 + 1 = 2$

int res[] = new int[n-k+1];

int i=0; j=0;

while(j < n){

add to the queue if element < 0

```
if (arr[j] < 0) {
    queue.add(arr[j]);
}
```

if (j - i + 1 == k) {

if (!queue.isEmpty()) {

res[i] = queue.peek();

if (arr[i] == queue.peek()) {  
queue.remove();

else {  
res[i] = 0;

i++;

j++;

arr[1] = 0 ✓

-6

arr[2] = -6 ✓

arr[2] != peek, so

-6

arr[2] != 0

arr[3] = 6 ✓

If any element

is present add

add it to array

else add '0' to

the array



# \* Reverse First K Elements from Queue:

I/P:

1 2 3 4 5     $k=3$

O/P:

5 4 3 2 1

Algo:

1. Push first  $k$  elements to the stack.

3  
2  
1

2. Push the stack elements back to the Queue.

4 5 3 2 1

3. For the first  $n-k$  elements pop and push Simultaneously back into queue.

POP 4 and add it     $(n-k)$   
5 - 3 = 2

① 5 3 2 1 4

② POP 5 and add it

3 2 1 4 5

Stack < Integer > stack = new Stack < > ();

for (int i = 0; i < k; i++)

{ stack.add(q.poll());

}

adding first  $k$  elements to Stack

while (!stack.isEmpty()) {

q.add(stack.pop());

}

adding elements from Stack to Queue

for (int i = 0; i < q.size() - k; i++)

{ int element = q.poll();

q.add(element);

}

polling & adding back upto  $(n-k)$  to the Queue.

return q;

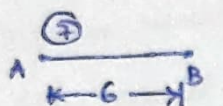


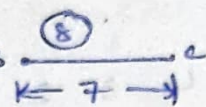
## First Circular Tour: (Petrol & Distance)

Petrol  $\Rightarrow$  <sup>A</sup>7, <sup>B</sup>8, <sup>C</sup>5, <sup>D</sup>11, <sup>E</sup>7, <sup>F</sup>6

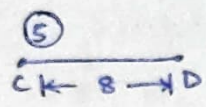
O/P: D

distance  $\Rightarrow$  6, 7, 8, 9, 7, 5

(a)  extra = 1L

b)  Petrol available = 1 + 8 = 9

extraFuel = 9 - 7

(c)  Petrol available = 2 + 5 = 7

extraFuel = 7 - 8 = -1 < 0

→ We need to find the point from which we can start such that we can start and cover the entire route.

→ Start at point D:

extraFuel = 11 - 9 = 2L

→ At point E: Petrol Available = 2 + 7 = 9

extraFuel = 9 - 7 = 2L

→ At point F: Petrol Available = 2 + 6 = 8

extraFuel = 8 - 5 = 3L

→ At point A: Petrol available = 3 + 7 = 10

extraFuel = 10 - 6 = 4L

→ At point B: Petrol: 4 + 8 = 12L


extraFuel = 12 - 7 = 5L

→ At point C: Petrol: 5 + 5 = 10L

extraFuel: 10 - 8 = 2L

→ We can each Reach point D:

Brute Force:

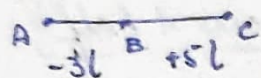
  
7, 8, 5, 11, 7, 6

6, 7, 8, 9, 7, 5

We need to check at each point, So

$O(n^2)$





-21 Storage if we start at B.

At A:

extra = -3

required = -3

At B:

next

extra = 0 + 5

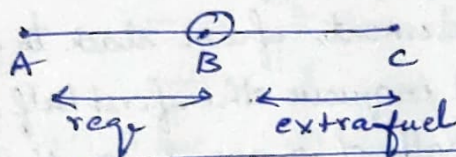
	A	B	C	D	E	F
P	7	8	5	11	7	6
d	6	7	8	9	7	5

$$eF = (7 - 6) + (8 - 7) + (5 - 8)$$

$$= 1 + 1 - 3$$

$$= -1 < 0$$

Start a Fresh



$$\boxed{\text{required} \leq \text{extraFuel}}$$

```
int start = 0;
```

```
int requiredPreviousFuel = 0;
```

```
int extraFuel = 0;
```

```
for(int i = 0; i < petrol.length; i++) {
```

```
    extraFuel += petrol[i] - distance[i];
```

```
    if (extraFuel < 0) {
```

```
        start = i + 1;
```

```
        requiredPreviousFuel = extraFuel;
```

```
        extraFuel = 0;
```

```
    }
```

```
}
```

```
int ans = -1;
```

```
if (extraFuel >= Math.abs(requiredPreviousFuel)) {
```

```
    ans = start;
```

```
}
```

```
return ans;
```



## \*\* Interleave the First Half of the Queue With Second Half:

→ Given a queue of integers of even length, rearrange the elements by interleaving the first half of the queue with 2<sup>nd</sup> half.

Input: 

4	3	2	1
---	---	---	---

1	2	3	4	5	6
---	---	---	---	---	---

4	3
---	---

2	1
---	---

1	2	3
---	---	---

4	5	6
---	---	---

4	2	3	1
---	---	---	---

1	4	2	5	3	6
---	---	---	---	---	---

### Algorithm:-

1. add first half of the queue to a stack.
2. Enqueue the elements from stack to queue.
3. Dequeue and enqueue the first half of queue.
4. add first half of queue to a stack.
5. add elements to the list from stack and then queue each time.

3
4

2	1
---	---

2	1	3	4
---	---	---	---

3	4	2	1
---	---	---	---

4
3

2	1
---	---

4	2	3	1
---	---	---	---

Stack < Integer > st = new Stack < > ();

① for (int i = 0; i < N/2; i++) {  
    st.push(q.poll());  
}

② while (!st.isEmpty()) {  
    q.add(st.pop());  
}

③ for (int i = 0; i < N/2; i++) {  
    int front = q.poll();  
    q.add(front);  
}

④ for (int i = 0; i < N/2; i++) {  
    st.add(q.poll());  
}

⑤ while (st.isEmpty() && !q.isEmpty()) {  
    list.add(st.pop());  
    list.add(q.poll());  
}