

Laziness in GHC Haskell

The features and principles

Presented by chip

ZJU Lambda
From here to World

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- 1 Appetizers
- 2 Thunk? What's it?
- 3 Why we need strictness?
- 4 Be more strict



```
possiblyBottom b =  
  case b of  
    True   → fst tup  
    False  → snd tup  
  where tup = (0, undefined)
```

If we apply **possiblyBottom** to **True**, we will get a 0.



Course 1: Outside in

A slightly arcane form:

```
possiblyBottom =  
  \f → f fst snd (0, undefined)
```

```
-- booleans in lambda form
```

```
true  :: a → a → a  
true = \a → (\b → a)
```

```
false :: a → a → a  
false = \a → (\b → b)
```



Course 1: Outside in

Nesting lambdas and reducing from the outside in:
(They are not in fact decomposed this way by the compiler)

```
(\f → f fst snd (0, undefined)) (\a → (\b → a))  
(\a → (\b → a)) fst snd (0, undefined)  
(\b → fst) snd (0, undefined)  
fst (0, undefined)  
0
```



Course 2: Evaluate to WHNF

```
length' :: [a] → Int
length' lst = go lst 0 where
    go [] acc      = acc
    go (x:xs) acc = go xs (acc+1)

main = let x = product [1..]
      in print $ length' [1, x]
```

It prints 2 !

What happened here?



Example 2: Evaluate to WHNF

The actual evaluation process:

```
length' [1, x]
= length' 1:(x:[])      -- 1:(x:[]) matches (x:xs)
= 1 + length' (x:[])    -- (x:[]), same with above
= 1 + 1 + length' []    -- [] matches []
= 1 + 1 + 0
= 2
```

Concept

In WHNF, we only evaluate the outermost constructor



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The Haskell Heap

The Haskell heap is a rather strange place.



Every item is wrapped up nicely in a box:
The Haskell heap is a heap of *presents* (thunks).



Present

When you actually want what's inside the present, you *open it up* (evaluate it).



Gift card

Sometimes you open a present, you get a *gift card* (data constructor). Gift cards have two traits.

- A name. (the **Just** gift card or **Right** gift card)
- And they tell you where the rest of your presents are.

There might be more than one (the tuple gift card), if you're a lucky duck!



Presents on the Haskell heap are rather mischievous.



Explode when you open it



Haunted by ghosts that open other presents when disturbed



What is a *thunk*?

`<thunk: expression-to-be-evaluated>`

- A box containing unevaluated expressions.
- Being evaluated when *needed*.
- Basically **anything** creates a thunk in (GHC) Haskell, by default



Example: Evaluate a thunk

How will this expression be evaluated?

map negate [1,2,3]



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Example: Evaluate a thunk

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<@: map negate <@: (1:2:3:[])>>

<@: negate <@: 1> : <@: map negate <@: [2,3]>>

—<@: 1> : <@: map negate <@: [2,3]>>

-1 : <@: map negate <@: [2,3]>>



Thunk brings us...

- On-demand data types.
- Call-by-need strategy.
- Memory reuse on CAF (Constant Applicative Forms).
- ...

```
fibs :: [Integer]  
fibs = 1 : 1 : zipWith (+) fibs (tail fibs)
```



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Thunks are good, but...


```
foldl (+) 0 (1:2:3:[])  
  = foldl (+) (0 + 1)                (2:3:[])  
  = foldl (+) ((0 + 1) + 2)          (3:[])  
  = foldl (+) (((0 + 1) + 2) + 3) []  
  =                               (((0 + 1) + 2) + 3)
```

What about **foldl** (+) 0 [1..1000000000] ?

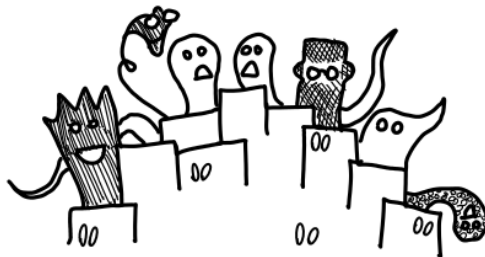


Memory leak

After executing **foldl** (+) 0 [1..1000000000]

Process Name	Status	% CPU	Nice	ID	Memory ▾
 ghc	Running	44	0	30047	4.0 GiB

A veritable ghost jamboree in our memory!



Example: Profiling

```
import System.Environment
import Text.Printf

main = do
    [d] ← map read `fmap` getArgs
    printf "%f\n" (mean [1..d])

mean :: [Double] → Double
mean xs = sum xs / fromIntegral (length xs)
```



RTS - Runtime System

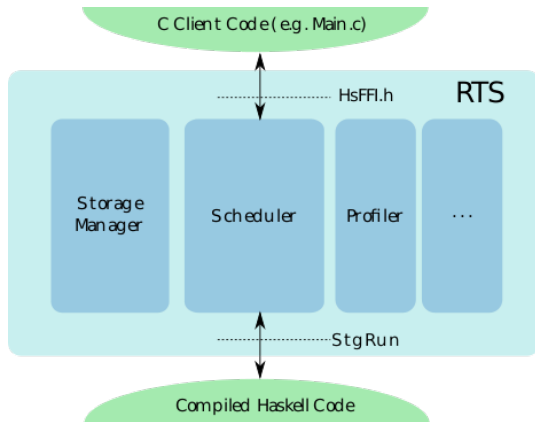


Figure: RTS Overview



Compile it.

```
ghc -make -rtsopts -O2 a.hs  
./a 1e7 +RTS -sstderr
```

Output:

```
...  
664 MB total memory in use
```

```
MUT      time      1.791s  (  1.804s elapsed)  
GC        time      2.255s  (  2.282s elapsed)  
Total     time      4.102s  (  4.146s elapsed)
```

```
%GC      time      55.0%  (55.0% elapsed)
```



Mark the cost centres

- SCC pragma
- Option **-auto-all**
- and **-caf-all**, if needed

Then, compile with option **-prof**

Run **./a 1e7 +RTS -p**, we get a file **a.prof**



Basic Profiling

```
1  Wed May 22 17:43 2019 Time and Allocation Profiling Report (Final)
2
3  a +RTS -p -RTS 1e7
4
5  total time =      2.57 secs  (2570 ticks @ 1000 us, 1 processor)
6  total alloc = 1,680,116,384 bytes  (excludes profiling overheads)
7
8  COST CENTRE MODULE %time %alloc
9
10 main          Main      87.2 100.0
11 mean          Main      12.8  0.0
12
13
14
15 COST CENTRE MODULE no. entries individual inherited
16                                %time %alloc %time %alloc
17 MAIN          MAIN          60      0      0.0  0.0 100.0 100.0
18 main          Main         121      0  87.2 100.0 100.0 100.0
19 mean          Main         124      1  12.8  0.0  12.8  0.0
20 CAF:main1     Main         118      0   0.0  0.0   0.0  0.0
21 main          Main         120      1   0.0  0.0   0.0  0.0
22 CAF:main2     Main         117      0   0.0  0.0   0.0  0.0
23 main          Main         122      0   0.0  0.0   0.0  0.0
24 CAF:main9     Main         114      0   0.0  0.0   0.0  0.0
25 main          Main         123      0   0.0  0.0   0.0  0.0
26 CAF           GHC.IO.Handle.FD 106      0   0.0  0.0   0.0  0.0
27 CAF           Text.Read.Lex    102      0   0.0  0.0   0.0  0.0
28 CAF           GHC.Conc.Signal  101      0   0.0  0.0   0.0  0.0
29 CAF           GHC.Float        100      0   0.0  0.0   0.0  0.0
30 CAF           GHC.IO.Encoding   99      0   0.0  0.0   0.0  0.0
31 CAF           GHC.IO.Encoding.Iconv 79      0   0.0  0.0   0.0  0.0
```

Figure: Profiling message generated by RTS



Heap Profiling

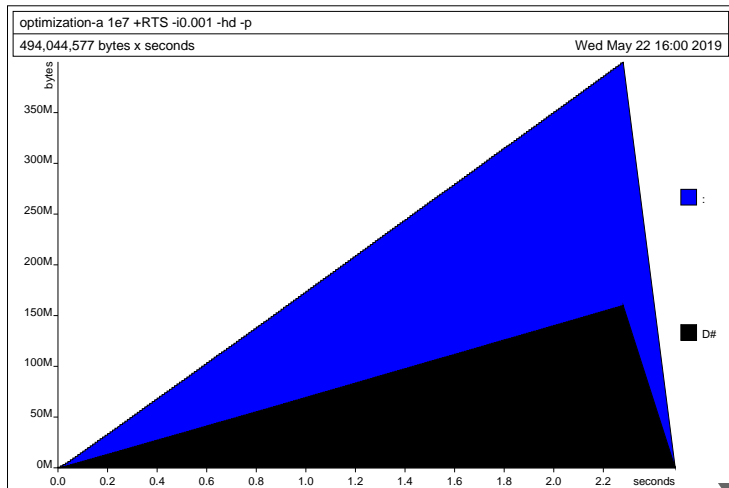


Figure: Break by constructor/closure



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Using *seq*

seq :: $a \rightarrow b \rightarrow b$

seq evaluates its first argument to **WHNF**, and return the second one.

$(\$!) :: (a \rightarrow b) \rightarrow a \rightarrow b$ -- |infixr 0|

$\$!$ is similar with $\$$, but evaluates its argument to **WHNF**.



deepseq :: **NFData** a \Rightarrow a \rightarrow b \rightarrow b

(\$!!) :: **NFData** a \Rightarrow (a \rightarrow b) \rightarrow a \rightarrow b -- |infixr 0|

force :: **NFData** a \Rightarrow a \rightarrow a

force x = x `deepseq` x

class **NFData** a **where**

rnf :: a \rightarrow ()



par :: a → b → b -- |infixr 0|

Indicates that it may be beneficial to evaluate the first argument in parallel with the second. Returns the value of the second argument.

pseq :: a → b → b -- |infixr 0|

Guarantee the order of evaluation in parallelism.



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pseq :: a → b → b -- |infixr 0|

Guarantee the order of evaluation in parallelism.

a `seq` b \iff b `seq` a `seq` b

