

# Find Minimum in Rotated Sorted Array

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10:40 PM

## 1. The Mental Model

A rotated sorted array isn't just a random mix. It is two sorted slopes separated by a drop.

- left slope: High numbers (e.g., 4, 5, 6)
- Right slope: Low numbers (e.g., 1, 2, 3)
- The goal: find the start of the Right slope.

## 2. The Golden Rule

Compare Mid vs. End

Never compare mid vs. Start. It is unreliable because the rotation might start at start. Always compare against the Rightmost element.

DECISION LOGIC:

1. If  $\text{nums}[\text{mid}] > \text{nums}[\text{end}]$ :

- meaning: mid is part of the Left slope (High)
- Implication: The min is somewhere to the Right.
- Action:  $\text{start} = \text{mid} + 1$  (we can safely discard mid).

2. If  $\text{nums}[\text{mid}] \leq \text{nums}[\text{end}]$ :

- meaning: mid is part of the Right slope (Low).
- Implication: The min is either at mid or to the left.

## 3. The "Trap" (Why $\text{end}=\text{mid}$ ?)

You will be tempted to write  $\text{end} = \text{mid} - 1$ . Don't!

- scenario:  $[3, 1, 2]$

- mid is 1

- check:  $1 \leq 2$  (True)

- If you do  $\text{end}=\text{mid}-1$  : You set end to index 0 ie 3.

You just threw away the 1 (the answer)!

- If you do  $\text{end}=\text{mid}$  : You keep the 1 in the search space.

## 4. Summary

- Complexity:  $O(\log n)$
- loop: while  $\text{start} < \text{end}$
- move Right:  $\text{start} = \text{mid} + 1$
- move Left:  $\text{end} = \text{mid}$
- final ans :  $\text{nums}[\text{start}]$