

De Morgan's Theorem 40

- Digital Circuit can be implemented by many ways.
- De Morgan's theorem is used to simplify boolean equation.
- Digital expression is simply made by three basic operation

1) Boolean AND \rightarrow  $Y = A.B$

2) Boolean OR \rightarrow  $Y = A+B$

3) Boolean NOT \rightarrow  \bar{A}

- De Morgan's Law

$$\overline{A.B} = \bar{A} + \bar{B} \quad \text{--- ①}$$

- Complement of product is Sum of Complement

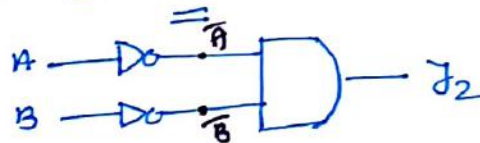
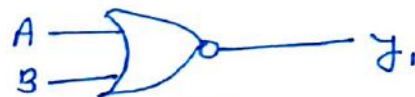


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$$\overline{A+B} = \bar{A} . \bar{B} \quad \text{--- ②}$$

- Complement of Sum is Product of Complement



Truth Table.

A	B	Y_1	\bar{A}	\bar{B}	Y_2
0	0	1	1	1	1
0	1	1	1	0	1
1	0	1	0	1	1
1	1	0	0	0	0

A	B	Y_1	\bar{A}	\bar{B}	Y_2
0	0	1	1	1	1
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	0

Ex $\overline{x + y.z} = \bar{x} . \bar{y.z}$
 $= \bar{x} . \overline{y.z}$
 $= \bar{x} . y.z$

Boolean Rules [Boolean Algebra Rules] 41

AND gate.

A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

- $A.A = A$
- $A.0 = 0$
- $A.1 = A$
- $A.\bar{A} = 0$

NOT Gate

A	Y
0	1
1	0

- $\bar{\bar{A}} = A$

Distributive Law

$$- A.(B+C) = A.B + A.C$$

$$\text{Imp} - A + (B.C) = (A+B).(A+C)$$

OR gate

A	B	Y
0	0	0
0	1	1
1	0	1
1	1	1

- $A+A = A$
- $A+0 = A$
- $A+1 = 1$
- $A+\bar{A} = 1$

$$\text{Imp} \left\{ \begin{array}{l} - A + \bar{A}B = (A + \bar{A}).(A + B) = A + B \\ - \bar{A} + AB = (\bar{A} + A).(\bar{A} + B) = \bar{A} + B \end{array} \right.$$

Commutative Law

$$A+B = B+A$$

&

$$A.B = B.A$$

Associative Law.

$$(A+B)+C = A+(B+C)$$

&

$$(A.B).C = A.(B.C).$$

De Morgan's Law

$$\overline{A+B} = \bar{A}.\bar{B} \text{ --- (1)}$$

$$\overline{A.B} = \bar{A} + \bar{B} \text{ --- (2)}$$

- Priority \rightarrow NOT \downarrow Highest
- \rightarrow AND \downarrow
- \rightarrow OR \downarrow Lowest

$$\begin{aligned} \text{Ex } P &= \underline{X}Y\bar{Z} + \underline{X}\bar{Y}\bar{Z} + Y\bar{Z} \\ &= X\bar{Z}(Y + \bar{Y}) + Y\bar{Z} \\ &= X\bar{Z} + Y\bar{Z} \\ &= (X+Y).\bar{Z} \end{aligned}$$

$$A \cdot 1 = A$$

Boolean Algebra Rules

NOT

$$\overline{\overline{A}} = A$$

AND

$$A \cdot A = A$$

$$A \cdot 0 = 0$$

$$A \cdot 1 = A$$

$$A \cdot \overline{A} = 0$$

OR

$$A + A = A$$

$$A + 0 = A$$

$$A + 1 = 1$$

$$A + \overline{A} = 1$$

Distributive Law

$$A + \overline{A}B = A + B$$

$$\overline{A} + AB = \overline{A} + B$$

$$A \cdot (B + C) = A \cdot B + A \cdot C$$

De-Morgan's Law

$$\overline{A + B} = \overline{A} \cdot \overline{B}$$

$$\overline{A \cdot B} = \overline{A} + \overline{B}$$

Boolean Algebra Examples 43

Consensus Theorem 42

Theorem.1 $AB + \bar{A}C + BC = AB + \bar{A}C$

$$\begin{aligned} \text{LHS} &= \underline{AB} + \underline{\bar{A}C} + \underline{BC(A + \bar{A})} \\ &= \underline{AB} + \underline{\bar{A}C} + \underline{BCA} + \underline{BC\bar{A}} \\ &= AB(1+C) + \bar{A}C(1+B) \\ &= AB + \bar{A}C \\ &= \text{RHS} \end{aligned}$$

Theorem.2 $(A+B)(\bar{A}+C)(B+C) = (A+B)(\bar{A}+C)$

$$\begin{aligned} \text{LHS} &= (\underline{A+B})(\underline{\bar{A}+C})(B+C) \\ &= (\underline{A\bar{A}} + AC + B\bar{A} + BC)(B+C) \quad [A\bar{A} = 0] \\ &= (AC + B\bar{A} + BC)(B+C) \\ &= \underline{ACB} + \underline{B\bar{A}B} + \underline{BCB} + \underline{ACC} + \underline{B\bar{A}C} + \underline{BCC} \\ &= \underline{ABC} + \underline{B\bar{A}} + \underline{BC} + AC + \underline{\bar{A}BC} + \underline{BC} \\ &= \underline{ABC} + \underline{B\bar{A}} + \underline{BC} + AC + \underline{\bar{A}BC} \\ &= BC[A + 1 + \bar{A}] + B\bar{A} + AC \\ &= BC + B\bar{A} + AC \end{aligned}$$

$$\begin{aligned} \text{RHS} &= (A+B)(\bar{A}+C) \\ &= \underline{A\bar{A}} + AC + B\bar{A} + BC \\ &= AC + B\bar{A} + BC \end{aligned}$$

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$$1) \overline{AB} + \bar{A} + AB$$

$$= \overline{\underline{X}} + \bar{A} + \underline{X} \quad [\because X = AB]$$

$$= \bar{A} + 1$$

$$= 1$$

$$= 0$$

↓ You may not see that sometimes.

$$= \overline{AB} + \bar{A} + AB$$

$$= \overline{AB} \cdot \bar{A} \cdot \overline{AB}$$

$$= AB \cdot A \cdot \overline{AB}$$

$$= AB \cdot \overline{AB}$$

$$= 0$$

$$2) A (B + \overline{C (AB + A\overline{C})})$$

$$= A (B + \overline{C} [\overline{AB} \cdot \overline{A\overline{C}}])$$

$$= A (B + \overline{C} [(\bar{A} + \bar{B}) \cdot (\bar{A} + \overline{\overline{C}})])$$

$$= A (B + \overline{C} [(\bar{A} + \bar{B}) \cdot (\bar{A} + C)])$$

$$= A (B + \overline{C} (\bar{A}\bar{A} + \bar{A}C + \bar{B}\bar{A} + \bar{B}C))$$

$$= A (B + \overline{C} (\underline{\bar{A}} + \underline{\bar{A}C} + \underline{\bar{A}\bar{B}} + \bar{B}C))$$

$$= A (B + \overline{C} (\bar{A} (1 + C + \bar{B}) + \bar{B}C))$$

$$= A (B + \overline{C} (\bar{A} + \bar{B}C))$$

$$= A (B + \overline{C}\bar{A} + \underline{\bar{B}C\overline{\overline{C}}})$$

$$= A (B + \overline{C}\bar{A})$$

$$= AB + \overline{C}\bar{A}A$$

$$= AB$$

3) If $\bar{A} + AB = 0$, then find values of A & B.

$$\Rightarrow \bar{A} + AB = 0$$

$$\Rightarrow \bar{A} + B = 0$$

$$\boxed{A=1} \text{ \& \ } \boxed{B=0}$$

Boolean Algebra Examples hh

1) $\overline{A}\overline{B}(C + BD) + \overline{A}\overline{B}$, Simplify given boolean eqn.

$$= \overline{A}\overline{B}(C + BD) + \overline{A}\overline{B}$$

$$= \overline{A}\overline{B}C + \overline{A}\overline{B}BD + \overline{A}\overline{B}$$

$$= \overline{A}\overline{B}C + \overline{A}\overline{B}$$

$$= \overline{B}(\overline{A} + AC)$$

$$= \boxed{\overline{B}(\overline{A} + C)}$$

2) $\overline{ABC}(A + B + C)$

$$= (\overline{A} + \overline{B} + \overline{C}) \cdot (\overline{A} \cdot \overline{B} \cdot \overline{C})$$

$$= \overline{A}\overline{A}\overline{B}\overline{C} + \overline{B}\overline{A}\overline{B}\overline{C} + \overline{C}\overline{A}\overline{B}\overline{C}$$

$$= \overline{A}\overline{B}\overline{C} + \overline{A}\overline{B}\overline{C} + \overline{A}\overline{B}\overline{C}$$

$$= \overline{A}\overline{B}\overline{C}$$

3) $\overline{(A + \overline{BC})}(A\overline{B} + \overline{ABC})$

$$= (\overline{A} \cdot \overline{\overline{BC}})(A\overline{B} + \overline{A} + \overline{B} + \overline{C})$$

$$= (\overline{A} \cdot BC)(\overline{A}\overline{B} + \overline{A} + \overline{B} + \overline{C})$$

$$= (\overline{A} \cdot BC)(\overline{B}(A+1) + \overline{A} + \overline{C})$$

$$= \overline{A}BC(\overline{A} + \overline{B} + \overline{C})$$

$$= \overline{A}\overline{A}BC + \overline{B}\overline{A}BC + \overline{C}\overline{A}BC$$

$$= \overline{A}BC$$

4) $A + \overline{B}C(A + \overline{BC})$

$$= A + \overline{B}C(A + \overline{B} + \overline{C})$$

$$= A + \overline{B}C(A + B + \overline{C})$$

$$= A + \overline{B}CA + \overline{B}CB + \overline{B}C\overline{C}$$

$$= A + \overline{B}CA$$

$$= A(1 + \overline{B}C)$$

$$= A$$

Boolean Algebra Examples 45

1) If $x=1$ in the logic eqn.

$$[x + z[\bar{y} + [\bar{z} + x\bar{y}]]][\bar{x} + \bar{z}(x+y)] = 1$$

then

A) $y=z$ B) $y=\bar{z}$ C) $z=1$ D) $z=0$

$$\Rightarrow [x + z[\bar{y} + [\bar{z} + x\bar{y}]]][\bar{x} + \bar{z}(x+y)] = 1$$

$$\Rightarrow [\underline{1 + z[\bar{y} + [\bar{z} + 1.\bar{y}]]}][\underline{0 + \bar{z}(1+y)}] = 1$$

$[1+A] = 1$ $[0+A=A]$

$$\Rightarrow 1.[\bar{z}.1] = 1$$

$$\Rightarrow \bar{z} = 1$$

$$\Rightarrow \boxed{z=0}$$

2) If we have 3 variables A, B & C. Find the output $y=1$ for majority of "1" in A, B & C. also minimise the function How 00

A	B	C	Y
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	← 1
1	0	0	0
1	0	1	← 1
1	1	0	← 1
1	1	1	← 1

$$\begin{aligned}
 Y &= \bar{A}BC + A\bar{B}C + \underline{ABC} + \underline{ABC} \\
 &= \bar{A}BC + A\bar{B}C + \underline{AB(\bar{C}+C)} \\
 &\quad [\bar{C}+C=1] \\
 &= \bar{A}BC + \underline{A\bar{B}C} + \underline{AB} \\
 &= \bar{A}BC + A[\bar{B}C + B] \\
 &\quad [\text{As distributive law}] \\
 &\quad \underline{BC + B = B+C} \\
 &= \bar{A}BC + A(B+C) \\
 &= \underline{\bar{A}BC} + \underline{AB} + AC \\
 &= B[\underline{\bar{A}C} + \underline{A}] + AC \\
 &\quad [\bar{A}C + A = A+C] \\
 &= B(A+C) + AC \\
 &= \underline{BA + BC + AC}
 \end{aligned}$$

Dual and Self Dual of Boolean expression 46

- To get dual of given expression, we need to replace

* OR with AND

* AND with OR

* 1 with 0

* 0 with 1

e.g.1 Find out dual of

$$\Rightarrow (A+B).(\bar{A}+C).(B+C) = (A+B).(\bar{A}+C)$$

$$\Rightarrow \boxed{A.B + \bar{A}.C + B.C = A.B + \bar{A}.C}$$

e.g.2 $F = AB + \bar{A}BC + \bar{A}\bar{C}$, Then find dual of F.

$$F_d = (A+B).(A+\bar{B}+C).(A+\bar{C})$$

- Self Dual - If dual of function is same function then it is referred as self dual.

e.g.3 $F = AB + BC + AC$

Find given function F is self dual or not.

$$F_d = (A+\underline{B}).(\underline{B}+C).(A+C)$$

$$= (B+A.C).(A+C) \quad \text{[As per distributive rule } (A+B)(B+C) = (B+AC)]$$

$$= BA + BC + \underline{A.C.A} + \underline{A.C.C}$$

$$= BA + BC + \underline{AC} + \underline{AC}$$

$$= BA + BC + AC$$

$$\boxed{F_d = F} \quad \text{So function F is self dual.}$$

- For n number of variable, total self dual = $2^{2^{n-1}}$

e.g.4 For $n=2$ variable find out total self dual.

$$= 2^{2^{n-1}} = 2^{2^{2-1}} = 2^2 = 4$$

→ we have variable A & B.

$$A \longleftrightarrow A$$

$$B \longleftrightarrow B$$

$$\bar{A} \longleftrightarrow \bar{A}$$

$$\bar{B} \longleftrightarrow \bar{B}$$

e.g.5 If $n=5$ variables, then find total self dual.

$$= 2^{2^{n-1}} = 2^{2^{5-1}} = 2^{2^4} = 2^{16} = 65536$$

SOP, POS & Canonical Form of Boolean function Representation

SOP = Sum of Product [DNF - Disjunctive Normal Form]

- It is a summation of product terms.

Eg. $Y = \underbrace{A.B + A.\bar{B}.C + A\bar{B}\bar{C}}_{\text{Product terms.}}$ Summation

POS = Product of sum [CNF - Conjunctive Normal Form].

- It is a product of sum terms.

Eg. $Y = \underbrace{(A+B) \cdot (A+C) \cdot (\bar{A}+\bar{B})}_{\text{Summation terms.}}$ Product.

Canonical Form

- Standard SOP (SSOP)

- Each product term contains all the variables of the function.

Eg. $F(A,B,C) = \underline{A.B.C} + \underline{A.\bar{B}.C} + \underline{A.B.\bar{C}}$

SSOP ✓
SOP X

Standard POS (SPOS)

- Each sum terms contains all the variables of the function.

Eg. $F(A,B,C) = \underline{(A+B+\bar{C})} \cdot \underline{(\bar{A}+B+C)}$

SPOS ✓
POS X

Minterms and Maxterms in Boolean function Representation

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Minterms

- Each Individual term in SSOP is called as minterms

Maxterms

- Each Individual term in SPOS is called as Maxterm

For 3 variables

F	A B C	minterms	maxterms
0	0 0 0	$\bar{A}\bar{B}\bar{C} \rightarrow m_0$	$(A+B+C) \rightarrow m_0$
1	0 0 1	$\bar{A}\bar{B}C \rightarrow m_1$	$(A+B+\bar{C}) \rightarrow m_1$
0	0 1 0	$\bar{A}B\bar{C} \rightarrow m_2$	$(A+\bar{B}+C) \rightarrow m_2$
1	0 1 1	$\bar{A}BC \rightarrow m_3$	$(A+\bar{B}+\bar{C}) \rightarrow m_3$
1	1 0 0	$A\bar{B}\bar{C} \rightarrow m_4$	$(\bar{A}+B+C) \rightarrow m_4$
0	1 0 1	$A\bar{B}C \rightarrow m_5$	$(\bar{A}+B+\bar{C}) \rightarrow m_5$
1	1 1 0	$AB\bar{C} \rightarrow m_6$	$(\bar{A}+\bar{B}+C) \rightarrow m_6$
1	1 1 1	$ABC \rightarrow m_7$	$(\bar{A}+\bar{B}+\bar{C}) \rightarrow m_7$

$$F = \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}\bar{C} + ABC + \bar{A}B\bar{C} + ABC$$

$$= \Sigma m(1, 3, 4, 6, 7)$$

$$F = (A+B+C) \cdot (A+\bar{B}+C) \cdot (\bar{A}+B+\bar{C})$$

$$= \Pi m(0, 2, 5)$$

- Max-terms and Min-terms are complement to each other.

$$F_{ABC} = \Sigma m(0, 2, 3, 5) \rightarrow F_{ABC} = \Pi m(1, 4, 6, 7)$$

↑

3 variables \rightarrow total min & max terms = $2^3 = 8$ (0, 1, 2, 3, 4, 5, 6, 7)

SOP to SSOP Conversion 49

Step 1 - Identify the missing variables in product terms.

$$F(A, B, C) = \underbrace{A.B}_{\substack{\uparrow \\ C \text{ is} \\ \text{missing}}} + \underbrace{A.\bar{B}.C}_{\substack{\uparrow \\ B \text{ is} \\ \text{missing}}} + \underbrace{A.\bar{C}}_{\substack{\uparrow \\ B \text{ is} \\ \text{missing}}}$$

Step 2 - Multiply [missing variable + it's complement]

$$\begin{aligned} F(A, B, C) &= A.B.[C + \bar{C}] + A.\bar{B}.C + A.\bar{C}.[B + \bar{B}] \\ &= \underline{A.B.C} + \underline{A.B.\bar{C}} + \underline{A.\bar{B}.C} + \underline{A.\bar{B}.\bar{C}} + \underline{A.B.C} + \underline{A.\bar{B}.C} \end{aligned}$$

Step 3 - Neglect the repeated terms.

$$\begin{aligned} F(A, B, C) &= \underbrace{A.B.C}_{m_7} + \underbrace{A.B.\bar{C}}_{m_6} + \underbrace{A.\bar{B}.C}_{m_5} \\ &= \Sigma_m(5, 6, 7) \end{aligned}$$

$$\begin{aligned} F(A, B, C, D) &= \underbrace{A.B}_{\substack{\downarrow \\ C \neq D \\ \text{is missing}}} + \underbrace{A.C}_{\substack{\downarrow \\ B \neq D \\ \text{is missing}}} + \underbrace{A.B.C.\bar{D}} \\ &= A.B.(C + \bar{C})(D + \bar{D}) + A.C.(B + \bar{B})(D + \bar{D}) + A.B.C.\bar{D} \\ &= \underline{A.B.C.D} + \underline{A.B.C.\bar{D}} + \underline{A.B.\bar{C}.D} + \underline{A.B.\bar{C}.\bar{D}} + \underline{A.B.C.D} + \underline{A.B.C.\bar{D}} + \underline{A.\bar{B}.C.D} + \underline{A.\bar{B}.C.\bar{D}} \\ &= \underline{A.B.C.D} + \underline{A.B.C.\bar{D}} + \underline{A.B.\bar{C}.D} + \underline{A.B.\bar{C}.\bar{D}} + \underline{A.\bar{B}.C.D} + \underline{A.\bar{B}.C.\bar{D}} \\ &= \underbrace{m_{15}}_{\downarrow} + \underbrace{m_{14}}_{\downarrow} + \underbrace{m_{13}}_{\downarrow} + \underbrace{m_{12}}_{\downarrow} + \underbrace{m_{11}}_{\downarrow} + \underbrace{m_{10}}_{\downarrow} \\ &= \Sigma_m(10, 11, 12, 13, 14, 15) \end{aligned}$$

How

$$F(A, B, C, D) = A + C.D$$

↑ solve this in SSOP

POS to SPOS Conversion SO

Step 1 - Identify the missing Variable.

$$F(A, B, C) = (\bar{A} + \bar{B}) \cdot \underline{A} \cdot (\underline{A + B + \bar{C}})$$

\downarrow C is missing \downarrow B & C is missing

Step 2 - Add with that Variable & its complement Separately.

$$F(A, B, C) = (\bar{A} + \bar{B} + C) \cdot (\bar{A} + \bar{B} + \bar{C})$$

$$(\underline{A + B + C}) \cdot (\underline{A + B + \bar{C}}) \cdot (\underline{A + \bar{B} + C}) \cdot (\underline{A + \bar{B} + \bar{C}})$$

x

Step 3 - Neglect repeated terms.

$$F(A, B, C) = (\bar{A} + \bar{B} + C) (\bar{A} + \bar{B} + \bar{C}) (A + B + C) (A + B + \bar{C})$$

m_6 m_7 m_0 m_1

$$(\underline{A + \bar{B} + C}) (\underline{A + \bar{B} + \bar{C}})$$

m_2 m_3

$$= \Pi M(0, 1, 2, 3, 6, 7)$$

eg

$$F(A, B, C) = (\underline{A + \bar{B}}) (\underline{\bar{A} + C}) (\underline{\bar{A} + \bar{B} + C})$$

\downarrow C is missing \downarrow B is missing

$$= (\underline{A + \bar{B} + C}) (\underline{A + \bar{B} + \bar{C}}) (\underline{\bar{A} + B + C}) (\underline{\bar{A} + \bar{B} + C})$$

$$= (\underline{A + \bar{B} + C}) (\underline{A + \bar{B} + \bar{C}}) (\underline{\bar{A} + B + C}) (\underline{\bar{A} + \bar{B} + C})$$

$\uparrow m_2$ $\uparrow m_3$ $\uparrow m_4$ $\uparrow m_6$

$$= \Pi M(2, 3, 4, 6)$$

Hw $F(A, B, C) = B(A + C)$

↑ Convert into SPOS.

SSOP to SPOS Conversion & 5
 SPOS to SSOP Conversion

1) $F(A, B, C) = \sum_m (0, 1, 4, 7)$. (convert SSOP to SPOS)
 \downarrow
 $n=3$, $[0, 1, 2, 3, 4, 5, 6, 7]$
 $= \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C + A\bar{B}\bar{C} + ABC$

$F(A, B, C) = \prod M [2, 3, 5, 6]$
 $= (A + \bar{B} + C)(A + \bar{B} + \bar{C})(\bar{A} + B + \bar{C})(\bar{A} + \bar{B} + C)$

2) $F(A, B, C) = \prod M (1, 2, 6)$. (convert SPOS to SSOP)
 \downarrow
 $n=3$, $[0, 1, 2, 3, 4, 5, 6, 7]$
 $= (A + B + \bar{C})(A + \bar{B} + C)(\bar{A} + \bar{B} + C)$

$F(A, B, C) = \sum_m (0, 5, 4, 5, 7)$
 $= \bar{A}\bar{B}\bar{C} + \bar{A}BC + A\bar{B}\bar{C} + A\bar{B}C + ABC$

3) $F(A, B, C) = ABC + AB\bar{C} + \bar{A}BC$, (convert SSOP to SPOS)
 $\downarrow \quad \quad \downarrow \quad \quad \downarrow$
 $m_7 \quad \quad m_6 \quad \quad m_3$
 \downarrow
 $n=3$, $[0, 1, 2, 3, 4, 5, 6, 7]$
 $= \sum_m (3, 6, 7)$

$F(A, B, C) = \prod M (0, 1, 2, 4, 5)$
 $= (A + B + C)(A + B + \bar{C})(A + \bar{B} + C)(\bar{A} + B + C)(\bar{A} + B + \bar{C})$

4) $F(A, B, C) = (A + \bar{B} + C)(\bar{A} + B + \bar{C})(A + B + \bar{C})$, (convert SPOS to SSOP)
 $\downarrow \quad \quad \downarrow \quad \quad \downarrow$
 $m_2 \quad \quad m_5 \quad \quad m_1$

$F(A, B, C) = \prod M (2, 2, 5)$
 \uparrow
 $n=3$, $[0, 1, 2, 3, 4, 5, 6, 7]$

$F(A, B, C) = \sum_m (0, 3, 4, 6, 7)$
 $= \bar{A}\bar{B}\bar{C} + \bar{A}BC + A\bar{B}\bar{C} + A\bar{B}C + ABC$

Example on SOP & POS 52

1) If $n=3$ variables then total minterms is 8
 total maxterms is 8
 total terms is 256

- total minterms = $2^n = 2^3 = 8$
 - total maxterms = $2^n = 2^3 = 8$
 - total terms = $2^{2^n} = 2^{2^3} = 2^8 = 256$
- total self dual = $2^{2^{n-1}} = 2^{2^{3-1}} = 2^4 = 16$

2) $f(A, B, C) = \underline{A} + \underline{\overline{B}C}$. Find SSOP & SPOS.

\downarrow \downarrow
 $B \& C$ A is
 is missing missing.

$$\begin{aligned}
 &= A.(B+\overline{B}).C.(1) + \overline{B}.C.(A+\overline{A}) \\
 &= \underline{A.B.C} + \underline{A.B.\overline{C}} + \underline{A.\overline{B}.C} + \underline{A.\overline{B}.\overline{C}} + \underline{\overline{A}.\overline{B}.C} + \underline{\overline{A}.\overline{B}.\overline{C}} \\
 &= ABC + AB\overline{C} + A\overline{B}C + A\overline{B}\overline{C} + \overline{A}\overline{B}C \\
 &\quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \\
 &\quad m_7 \quad m_6 \quad m_5 \quad m_4 \quad m_1
 \end{aligned}$$

$$= \Sigma m(1, 4, 5, 6, 7)$$

$$n=3, [0, 1, 2, 3, \underline{4}, \underline{5}, \underline{6}, \underline{7}]$$

$$f(A, B, C) = \Pi m[0, 2, 3]$$

$$= (A+B+C).(A+\overline{B}+C).(A+\overline{B}+\overline{C})$$

3) $f = \underline{(A+B)} \underline{(A+C)}$. find total minterms & maxterms.

\uparrow \uparrow
 C is missing B is missing.

$$\begin{aligned}
 &= (A+B+C).(A+B+\overline{C}).(A+\overline{B}+C).(A+\overline{B}+\overline{C}) \\
 &= (A+B+C).(A+B+\overline{C}).(A+\overline{B}+C) \\
 &\quad \downarrow \quad \downarrow \quad \downarrow \\
 &\quad m_0 \quad m_1 \quad m_2
 \end{aligned}$$

$$= \Pi m(0, 1, 2)$$

$$n=3, [\underline{0}, \underline{1}, \underline{2}, 3, 4, 5, 6, 7]$$

$$f = \Sigma m(3, 4, 5, 6, 7)$$

$$= \overline{A}BC + A\overline{B}C + A\overline{B}\overline{C} + ABC$$