# ${\bf IUT\ Pounopunik,\ Islamic\ University\ of\ Technology}$

Contents		6		10
1	All Macros	1	6.2 Combi	10 10
2	Data Structure	1		11 11
	2.1 Sparse Table	1	6.5 Pollard Rho	11
	2.2 BIT	$_{1}$	6.6 Chinese Remainder Theorem	11
		2		12
	v e	$\frac{1}{2}$	6.8 FFT	12
	9	3	6.9 NTT	12
			6.10 ModInverse	13
	9	3	6.11 Diophantine	13
		3	- a	
	•	3 7	· ·	14
		4		14
	2.10 SQRT Decomp	4		14
				$\frac{15}{16}$
3	Graph	5		16
	3.1 DSU BySize	5	, O	18
	3.2 MST Kruskal	5	7.6 Half Plane	19
	3.3 Dijkstra	5 8	B Equations and Formulas 2	20
	3.4 Bellman Ford	6	1	<b>2</b> 0
		6		$\frac{1}{20}$
	·	6	9	$\frac{1}{20}$
		6	<u> </u>	$\frac{1}{20}$
		$\frac{3}{7}$		20
		<u>'</u>		20
	3.9 BFS	"		20
4	String	7		
	4.1 Hashing	7		
		7		
		8		
		8		
		8		
		8		
		9		
		9		
	4.9 Z Algo			
5	DP 10	0		
	5.1 Bitmask	0		
	5.2 LIS			
	5.3 Divide and Conquer DP			
	5.5 Knuth Optimization			

```
Sublime Build

{
    "shell_cmd": "g++ -std=c++17 -o ${file_path}/${
        file_base_name} ${file} && ${file_path}/${
        file_base_name} < input.txt > output.txt",
    "working_dir": "${file_path}",
    "selector": "source.cpp"
}
```

# 

#### Stress-tester

```
#!/bin/bash
# Call as stresstester ITERATIONS [--count]
g++ gen.cpp -o gen
g++ sol.cpp -o sol
g++ brute.cpp -o brute
for i in $(seq 1 "$1"); do
   echo "Attempt $i/$1"
   ./gen > in.txt
    ./sol < in.txt > out1.txt
   ./brute < in.txt > out2.txt
   diff -y out1.txt out2.txt > diff.txt
   if [ $? -ne 0 ] ; then
       echo "Differing Testcase Found:"; cat in.txt
       echo -e "\nOutputs:"; cat diff.txt
       break
   fi
```

# 1 All Macros

```
/*--- DEBUG TEMPLATE STARTS HERE ---*/
void show(int x) {cerr << x;}
void show(long long x) {cerr << x;}
void show(double x) {cerr << x;}
void show(char x) {cerr << '\'' << x << '\'';}
void show(const string &x) {cerr << '\"' << x << '\"';}
void show(bool x) {cerr << (x ? "true" : "false");}
```

```
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\f'; show(x.first);</pre>
    cerr << ", "; show(x.second); cerr << '}'; }</pre>
template<typename T>
void show(T x) {int f = 0; cerr << "{"; for (auto &i: x)</pre>
    cerr << (f++ ? ", " : ""), show(i); cerr << "}";}</pre>
void debug_out(string s) {
 s.clear();
 cerr << s << '\n';
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
 s.erase(remove(s.begin(), s.end(), ''), s.end());
                "; // 8 spaces
  cerr << s.substr(0, s.find(','));</pre>
  s = s.substr(s.find(',') + 1);
 cerr << " = ":
  show(t):
  cerr << endl;
  if(sizeof...(v)) debug_out(s, v...);
#define dbg(x...) cerr << "LINE: " << __LINE__ << endl;</pre>
    debug_out(#x, x); cerr << endl;</pre>
/*--- DEBUG TEMPLATE ENDS HERE ---*/
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm
    .mmx.avx.tune=native")
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
 //find_by_order(k) --> returns iterator to the kth
      largest element counting from 0
 //order_of_key(val) --> returns the number of items in
      a set that are strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type, less<DT>,
    rb_tree_tag,tree_order_statistics_node_update>;
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
#ifdef LOCAL
#include "dbg.h"
```

```
#else
#define dbg(x...)
#endif
int main() {
  cin.tie(0) -> sync_with_stdio(0);
}
```

# 2 Data Structure

# 2.1 Sparse Table

#### 2.2 BIT

```
template <typename T> class BIT
{
  public:
    int n; vector<T> tree;

  BIT(int size) // 1-indexed
  {
      n = size; tree.assign(n+1, 0);
  }

  BIT(const vector<T> &a) : BIT(a.size())
  {
      for(int i = 1; i <= n; i++) update(i, a[i-1]);
  }

  T query(int i)</pre>
```

```
{
    T ans = 0;
    for(; i >= 1; i-= (i & -i)) ans+= tree[i];
    return ans;
}

T query(int 1, int r)
{
    return query(r) - query(1-1);
}

void update(int i, T delta)
{
    for(; i <= n; i+= (i & -i)) tree[i]+= delta;
}
};</pre>
```

# 2.3 Lazy SegmentTree

```
11 tree[4*MAX], lazy[4*MAX]; // 1-indexed
void build(vector<ll>&a, int b = 0, int e = -1, int v=1)
 if(e == -1) e = a.size()-1;
 if(b == e)
   tree[v] = a[b];
   return;
 int mid = (b+e)/2;
  build(a, b, mid, 2*v);
 build(a, mid+1, e,2*v+1);
 tree[v] = tree[2*v] + tree[2*v+1]:
ll query(int l, int r, int b, int e, int v=1, ll carry =
    0)
 if(b > r || e < 1) return 0;
 if(b >= 1 && e <= r) return tree[v]+carry*(e-b+1);</pre>
 int mid = (b+e)/2;
 11 lseg = query(1, r, b, mid, 2*v, carry+lazy[v]);
 ll rseg = query(l, r, mid+1, e, 2*v+1, carry+lazy[v]);
 return lseg + rseg;
void update(int 1, int r, 11 val, int b, int e, int v =
    1)
```

```
if(b > r || e < 1) return;
if(b >= 1 && e <= r)
{
   tree[v]+= (e-b+1)*val;
   lazy[v]+= val;
   return;
}

int mid = (b+e)/2;
update(1, r, val, b, mid, 2*v);
update(1, r, val, mid+1, e, 2*v+1);
tree[v] = tree[2*v] + tree[2*v+1]+ (e-b+1)*lazy[v];
}</pre>
```

# 2.4 Generic SegmentTree

```
template<typename ST, typename LZ>
class SegmentTree {
private:
 int n;
 ST *tree, identity;
 ST (*merge) (ST, ST);
 LZ *lazy, unmark;
 void (*mergeLazy)(int, int, LZ&, LZ);
 void (*applyLazy)(int, int, ST&, LZ);
 void build(vector<ST> &arr, int lo, int hi, int cur=1)
   if(lo == hi)
     tree[cur] = arr[lo-1];
     return;
   int mid = (hi+lo)/2, left = 2*cur, right = 2*cur+1;
   build(arr, lo, mid, left);
   build(arr, mid+1, hi, right);
   tree[cur] = merge(tree[left], tree[right]);
 void propagate(int lo, int hi, int cur)
   applyLazy(lo, hi, tree[cur], lazy[cur]);
   if(lo < hi)</pre>
     int mid = (lo+hi)/2, left = 2*cur, right = 2*cur+1;
     mergeLazy(lo, mid, lazy[left], lazy[cur]);
     mergeLazy(mid+1, hi, lazy[right], lazy[cur]);
   lazy[cur] = unmark;
```

```
}
 void update(int from, int upto, LZ delta, int lo, int
     hi, int cur=1)
  if(lo>hi) return;
  propagate(lo, hi, cur);
  if(from > hi or upto < lo) return;</pre>
  if(from<= lo and upto >= hi)
    mergeLazy(lo, hi, lazy[cur], delta);
    propagate(lo, hi, cur);
    return;
   int mid = (lo+hi)/2, left = 2*cur, right = 2*cur+1;
   update(from, upto, delta, lo, mid, left);
   update(from, upto, delta, mid+1, hi, right);
   tree[cur] = merge(tree[left], tree[right]);
}
 ST query(int from, int upto, int lo, int hi, int cur=1)
  if(lo>hi) return identity;
  propagate(lo, hi, cur);
   if(from > hi or upto < lo) return identity;</pre>
   if(from<= lo and upto >= hi) return tree[cur];
   int mid = (lo+hi)/2, left = 2*cur, right = 2*cur+1;
   ST lseg = query(from, upto, lo, mid, left);
  ST rseg = query(from, upto, mid+1, hi, right);
  return merge(lseg, rseg);
}
public:
 SegmentTree(
  vector<ST> arr, ST (*merge) (ST, ST), ST identity,
  void (*mergeLazy)(int, int, LZ&, LZ),
   void (*applyLazy)(int, int, ST&, LZ), LZ unmark
):
  n(arr.size()), merge(merge), identity(identity),
   mergeLazy(mergeLazy), applyLazy(applyLazy), unmark(
       unmark)
  tree = new ST[n*4];
  lazy = new LZ[n*4];
  build(arr, 1, n);
  fill(lazy, lazy+n*4, unmark);
```

```
void update(int from, int upto, LZ delta)
   update(from, upto, delta, 1, n);
  ST query(int from, int upto)
   return query(from, upto, 1, n);
  ~SegmentTree()
   delete[] tree;
   delete[] lazy;
};
11 add(11 1, 11 r) { return 1+r;}
void mergeAdd(int lo, int hi, ll &cur, ll pending) { cur
    += pending;}
void applyAdd(int lo, int hi, ll &cur, ll pending) { cur
    += pending*(hi-lo+1);}
void solve(int tcase)
 vector<ll> a(n);
 SegmentTree<11, 11> st(a, add, 0, mergeAdd, applyAdd,
      0);
```

#### 2.5 MO

```
struct node {
   LL l, r, idx;
};
bool cmp(const node &x, const node &y) {
   return x.r < y.r;
}

void add(LL x) {
   if(mp[x] % 2) curr++;
   mp[x]++;
}

void diminish(LL x) {
   if(mp[x] % 2 == 0) curr--;
   mp[x]--;
}

void solve()
{
   BLOCK_SIZE = sqrt(n) + 1;
   rep(i, 0, q-1) {</pre>
```

```
LL x, y; cin >> x >> y;
x--; y--;
query[x / BLOCK_SIZE].pb({x, y, i});
m = max(m, x / BLOCK_SIZE);
}
rep(i, 0, m) sort(all(query[i]), cmp);
LL mo_left = 0, mo_right = -1;
rep(i, 0, m) {
  for(auto [left, right, id] : query[i]) {
    while(mo_right < right) add(v[++mo_right]);
    while(mo_right > right) diminish(v[mo_right--]);
    while(mo_left < left) diminish(v[mo_left++]);
    while(mo_left > left) add(v[--mo_left]);
    answer[id] = curr;
}
rep(i, 0, q-1) cout << answer[i] << endl;</pre>
```

# 2.6 MergeSort Tree

```
vector<LL> tree[5*MAXN];
LL A[N];
void build_tree(LL now , LL curLeft, LL curRight) {
         if(curLeft == curRight) {
              tree[now].push_back(A[curLeft]);
              return;
         LL mid = (curLeft + curRight) / 2;
         build tree(2 * now, curLeft, mid);
         build_tree(2 * now + 1, mid + 1 , curRight);
         tree[now] = merge(tree[2 * now] , tree[2 * now +
              1]);
LL query(LL now, LL curLeft, LL curRight, LL 1, LL r, LL
       if(curRight < 1 || curLeft > r) return 0;
       if(curLeft >= 1 && curRight <= r)</pre>
              Return lower_bound(tree[now].begin(), tree
                   [now].end(), k) - tree[now].begin();
       LL mid = (curLeft + curRight) / 2;
       return query(2 * now, curLeft, mid, 1, r, k) +
           query(2 * now + 1, mid + 1, curRight, 1, r, k)
           );
```

# 2.7 BIT2d

```
const int N = 1008;
int bit[N][N], n, m;
int a[N][N], q;
void update(int x, int y, int val) {
```

## 2.8 SparseTable2d

```
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 505:
const int LOGN = 9;
// O(n^2 (logn)^2
// Supports Rectangular Query
int A[MAXN][MAXN];
int M[MAXN] [MAXN] [LOGN] [LOGN];
void Build2DSparse(int N) {
   for (int i = 1; i <= N; i++) {</pre>
       for (int j = 1; j <= N; j++) {</pre>
           M[i][j][0][0] = A[i][j];
       for (int q = 1; (1 << q) <= N; q++) {
           int add = 1 << (q - 1);
           for (int j = 1; j + add <= N; j++) {</pre>
               M[i][j][0][q] = max(M[i][j][0][q - 1], M[i]
                   [j + add][0][q - 1];
           }
   }
   for (int p = 1; (1 << p) <= N; p++) {</pre>
       int add = 1 << (p - 1);
       for (int i = 1; i + add <= N; i++) {</pre>
           for (int q = 0; (1 << q) <= N; q++) {
              for (int j = 1; j <= N; j++) {</pre>
                  M[i][j][p][q] = max(M[i][j][p - 1][q],
                       M[i + add][j][p - 1][q]);
              }
```

# 2.9 SegmentTree

```
#include <bits/stdc++.h>
using namespace std;
using 11 = long long;
using pii = pair<11, 11>;
const 11 N=5e5+5, mod=998244353;
using treenode = 11;
using lazynode = pii;
#define fir(a) for(int i=0; i<a; i++)</pre>
treenode treeidn = 0;
lazynode lazyidn = {1, 0};
vector<ll> v(N);
vector<treenode> tree(4*N, treeidn);
vector<lazynode> lazy(4*N, lazyidn);
treenode merge(treenode &a, treenode &b){
 return ;//add merge function here
void lazyapply(treenode &to, ll l, ll r, lazynode &fr){
 //apply lazy to treenode here
void lazymerge(lazynode &to, lazynode &fr){
 //combine to lazy updates here
```

```
//----dont touch: start
void build(ll id, ll l, ll r){
 if(l==r){
   tree[id]=v[l]:
   return;
 11 m=(1+r)/2;
 build(id*2+1, 1, m);
 build(id*2+2, m+1, r);
 tree[id] = merge(tree[id*2+1], tree[id*2+2]);
void push(ll id, ll l, ll r){
 if(1-r){
   11 m=(1+r)/2;
   lazyapply(tree[2*id+1], 1, m, lazy[id]);
   lazymerge(lazy[2*id+1], lazy[id]);
   lazyapply(tree[2*id+2], m+1, r, lazy[id]);
   lazymerge(lazy[2*id+2], lazy[id]);
   lazy[id] = lazyidn;
treenode query(ll id, ll l, ll r, ll ql, ll qr){
 push(id, 1, r);
 if(ql<=l && r<=qr) return tree[id];</pre>
 if(ql>r || qr<l) return treeidn;</pre>
 11 m=(1+r)/2;
 treenode tl=query(id*2+1, 1, m, ql, qr);
 treenode tr=query(id*2+2, m+1, r, ql, qr);
 return merge(tl, tr);
void update(ll id, ll l, ll r, ll ul, ll ur, lazynode uv)
 push(id, 1, r);
 if(ul<=1 && r<=ur){</pre>
   lazyapply(tree[id], 1, r, uv);
   lazymerge(lazy[id], uv);
   return;
 if(ul>r || ur<l) return;</pre>
 11 m=(1+r)/2;
 update(id*2+1, 1, m, ul, ur, uv);
 update(id*2+2, m+1, r, ul, ur, uv);
 tree[id]=merge(tree[id*2+1], tree[id*2+2]);
 return;
   -----dont touch: end
```

```
void solve(){
    ll n, q; cin>>n>q;
    fir(n) cin>>v[i];

build(0, 0, n-1);

while(q--){
    ll t; cin>>t;
    if(t){
        ll l, r; cin>>l>>r;
        cout<<query(0, 0, n-1, l, r-1)<<"\n";
    }else{
        ll l, r, a, b; cin>>l>>r>>a>b;
        update(0, 0, n-1, l, r-1, {a, b});
    }
}
return;
}
```

#### 2.10 SQRT Decomp

```
// input data
int n;
vector<int> a (n);
// preprocessing
int len = (int) sqrt (n + .0) + 1; // size of the block
    and the number of blocks
vector<int> b (len):
for (int i=0; i<n; ++i)</pre>
       b[i / len] += a[i];
// answering the queries
for (;;) {
       int 1, r;
   // read input data for the next query
       int sum = 0;
       for (int i=1; i<=r; )</pre>
               if (i % len == 0 && i + len - 1 <= r) {</pre>
                      // if the whole block starting at i
                            belongs to [1, r]
                      sum += b[i / len];
                      i += len;
              }
               else {
                      sum += a[i];
                      ++i;
int c_1 = 1 / len, c_r = r / len;
```

```
if (c_1 == c_r)
       for (int i=1; i<=r; ++i)</pre>
               sum += a[i];
else {
       for (int i=1, end=(c_l+1)*len-1; i<=end; ++i)</pre>
               sum += a[i];
       for (int i=c_l+1; i<=c_r-1; ++i)</pre>
               sum += b[i];
       for (int i=c_r*len; i<=r; ++i)</pre>
               sum += a[i];
}
void remove(idx); // TODO: remove value at idx from data
    structure
void add(idx);
                  // TODO: add value at idx from data
    structure
int get_answer(); // TODO: extract the current answer of
    the data structure
int block_size;
struct Query {
       int 1, r, idx;
       bool operator<(Query other) const</pre>
       {
               return make_pair(l / block_size, r) <</pre>
                            make_pair(other.l / block_size
                                 , other.r);
       }
};
vector<int> mo_s_algorithm(vector<Query> queries) {
       vector<int> answers(queries.size());
       sort(queries.begin(), queries.end());
       // TODO: initialize data structure
       int cur_1 = 0;
       int cur r = -1;
       // invariant: data structure will always reflect
            the range [cur_1, cur_r]
       for (Query q : queries) {
               while (cur_1 > q.1) {
                       cur 1--;
                       add(cur_1);
               }
               while (cur_r < q.r) {</pre>
                       cur_r++;
                       add(cur_r);
               while (cur_1 < q.1) {</pre>
                       remove(cur 1);
```

```
cur_1++;
              }
              while (cur_r > q.r) {
                      remove(cur r):
                      cur_r--;
              answers[q.idx] = get_answer();
       }
       return answers:
3 Graph
3.1 DSU BySize
vector<int> parent, setSize;
void make_set(int v) {
 parent[v] = v;
  setSize[v] = 1;
int find set(int v) {
 if (v == parent[v]) return v;
 return parent[v] = find_set(parent[v]);
void union_sets(int a, int b) {
 a = find set(a);
 b = find_set(b);
if (a != b) {
   if (setSize[a] < setSize[b]) swap(a, b);</pre>
   parent[b] = a:
   setSize[a] += setSize[b];
int main() {
 int n;
 cin >> n;
 parent.resize(n);
 setSize.resize(n);
 for (int i = 0; i < n; i++) make_set(i);</pre>
3.2 MST Kruskal
```

```
const ll sz = 1e5 + 7;
vector<ll> pr(sz);

ll find(ll x) {
   if (pr[x] == x) return x;
   return pr[x] = find(pr[x]);
}

void _union(ll x, ll y) {
   pr[find(y)] = find(x);
```

```
signed main() {
 ll n, m, i;
 cin >> n >> m;
 vector<tuple<11, 11, 11>> edg(m);
 iota(pr.begin(), pr.begin() + n + 1, 0);
 for (auto &[w, u, v] : edg) cin >> u >> v >> w;
 sort(edg.begin(), edg.end());
 11 cost = 0;
 for (auto [w, u, v] : edg) {
   if (find(u) != find(v)) {
     _union(u, v);
     cost += w;
 }
 for (i = 1; i < n; i++) {
   if (find(i) != find(i + 1)) {
     cout << "IMPOSSIBLE\n";</pre>
     return 0;
   }
 }
 cout << cost << "\n";
```

#### 3.3 Dijkstra

```
using pll = pair<ll, ll>;
vector<pll> adj[MAX];
vector<ll> dist(MAX, INF);
vector<ll> par(MAX, -1);

void dijkstra(int src)
{
    dist[src] = 0;
    priority_queue<pll, vector<pll>, greater<pll>> pq;
    pq.push({0, src});

    while(!pq.empty())
    {
        auto [d, u] = pq.top();
        pq.pop();

        if(d > dist[u]) continue;

        for(auto &[v, w]: adj[u])
        {
             if(dist[u]+w < dist[v])
            {
                 dist[v] = dist[u]+w;
            }
}</pre>
```

```
par[v] = u;
    pq.push({dist[v], v});
    }
}
```

#### 3.4 Bellman Ford

```
#define sz 100007
ll INF = 1e18;
vector<tuple<11, 11, 11>> edg;
vector<ll> dis(sz, INF);
void bellman_ford(ll n) {
 ll i, brk;
  dis[1] = 011;
 for (i = 1: i <= n: i++) {
   brk = 0;
   for (auto [u, v, w] : edg) {
     if (dis[v] > dis[u] + w)
       dis[v] = dis[u] + w; // for directional graph
     else
       brk++;
   }
   if (brk == n)
     break; // optimization
```

# 3.5 Floyd Warshall

bellman\_ford(n);

```
floyd_warshall(n);
```

#include <bits/stdc++.h>

for (i = 1; i <= n; i++) {</pre>

for (i = 1; i <= n; i++)

for (auto e : ord) {

if (!vis[e]) {

reverse(ord.begin(), ord.end());

if (!vis[i])

dfs1(i);

vis[i] = 0:

vector<ll> scc;

# 3.6 SCC

```
using namespace std;
#define ll long long
const 11 sz = 1e5 + 7;
vector<ll> adj[sz];
vector<ll> Radj[sz];
vector<bool> vis(sz);
vector<ll> ord;
void dfs1(ll cur) {
 vis[cur] = 1;
 for (auto nxt : adj[cur]) {
   if (!vis[nxt])
     dfs1(nxt);
 ord.push_back(cur);
void dfs2(11 cur) {
 vis[cur] = 1;
 for (auto nxt : Radj[cur]) {
   if (!vis[nxt])
     dfs2(nxt);
 }
signed main() {
 // strongly connected component kosaraju's algorithm
 ll n, m, a, b, i;
 cin >> n >> m;
 for (i = 0; i < m; i++) {</pre>
   cin >> a >> b;
   adj[a].push_back(b);
   Radj[b].push_back(a);
```

```
scc.push_back(e);
dfs2(e);
}

if (scc.size() == 1) {
  cout << "YES\n";
  return 0;
}

cout << "NO\n";

for (i = 1; i <= n; i++)
  vis[i] = 0;
dfs2(scc[0]);
if (vis[scc[1]])
  cout << scc[0] << " " << scc[1] << "\n";
else
  cout << scc[1] << " " " << scc[0] << "\n";
return 0;</pre>
```

#### 3.7 LCA

```
LL n, 1, timer;
vector<vector<LL>> adj;
vector<LL> tin, tout;
vector<vector<LL>> up;
void dfs(LL v, LL p) {
 tin[v] = ++timer;
 up[v][0] = p;
 for (LL i = 1; i <= 1; ++i)
   up[v][i] = up[up[v][i-1]][i-1];
 for (LL u : adj[v]) {
   if (u != p) dfs(u, v);
 tout[v] = ++timer;
bool is ancestor(LL u, LL v) {
 return tin[u] <= tin[v] && tout[u] >= tout[v];
LL lca(LL u, LL v) {
 if (is_ancestor(u, v)) return u;
 if (is_ancestor(v, u)) return v;
 for (LL i = 1; i >= 0; --i) {
   if (!is_ancestor(up[u][i], v)) u = up[u][i];
 return up[u][0];
void preprocess(LL root) {
 tin.resize(n);
 tout.resize(n);
 timer = 0;
```

```
1 = ceil(log2(n));
 up.assign(n, vector<LL>(1 + 1));
 dfs(root, root);
3.8 EulerTourTree
using ll = long long;
using vi = vector<11>;
using grid = vector<vi>;
void et(grid &edg, ll at, ll pt, grid &tr, ll &id){
 tr[0][id]=at; //val[at];
 tr[1][at]=id++;
 for(ll to: edg[at]) if(to-pt){
   et(edg, to, at, tr, id);
 tr[0][id]=at: //val[at]:
 tr[2][at]=id++;
 return:
grid etour(grid &edg, ll rt){
 ll cn=edg.size(), id=1;
 grid tour={vi(2*cn, 0), vi(cn), vi(cn)};
 et(edg, rt, 0, tour, id);
 return tour;
3.9 BFS
```

```
ll bfs(grid &edg, ll sn){
 ll cn=edg.size(), lv=-1, cl=0, nl=1, at, ls;
 vi vst(cn+1, 0), prt(cn+1, -1);
  queue<11> call;
  call.push(sn); vst[sn]++;
  while(!call.empty()){
   if(!cl){
     lv++; cl=nl; nl=0;
   }
   at=call.front():
   //if(at==en) return lv;
   call.pop(); cl--; ls=at;
   for(ll to:edg[at]){
     if(!vst[to]){
       prt[to] = at;
       call.push(to);
```

```
vst[to]++:
      nl++;
   }
 }
 return 0;
 //return ls; //for deepest.
4 String
4.1 Hashing
```

```
class HashedString {
private:
 static const long long M = 1e9 + 7;
 static const long long B = 256;
 static vector<long long> pow;
 vector<long long> p_hash;
public:
 HashedString(const string& s) : p_hash(s.size() + 1) {
   while (pow.size() < s.size()) {</pre>
     pow.push_back((pow.back() * B) % M);
   p hash[0] = 0;
   for (int i = 0; i < s.size(); i++) {</pre>
     p hash[i + 1] = ((p hash[i] * B) % M + s[i]) % M;
   }
 }
 long long getHash(int start, int end) {
   long long raw_val = (
     p_hash[end + 1] - (p_hash[start] * pow[end - start
          + 1])
   );
   return (raw_val % M + M) % M;
};
vector<long long> HashedString::pow = {1};
```

```
4.2 Double hash
// define +, -, * for (PLL, LL) and (PLL, PLL), % for (
    PLL. PLL):
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
PLL power(PLL a, LL p) {
 PLL ans = PLL(1, 1);
 for(; p; p >>= 1, a = a * a % mod) {
     if(p \& 1) ans = ans * a % mod;
```

```
return ans;
PLL inverse(PLL a) { return power(a, (mod.ff - 1) * (mod.
    ss - 1) - 1); }
PLL inv_base = inverse(base);
PLL val;
vector<PLL> P:
void hash_init(int n) {
 P.resize(n + 1);
 P[0] = PLL(1, 1);
 for (int i = 1; i <= n; i++) P[i] = (P[i - 1] * base) %
       mod;
PLL append(PLL cur, char c) { return (cur * base + c) %
    mod; }
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) { return (P[k] * c +
    cur) % mod; }
/// replaces the i-th (0-indexed) character from right
    from a to b;
PLL replace(PLL cur, int i, char a, char b) {
 cur = (cur + P[i] * (b - a)) \% mod:
 return (cur + mod) % mod;
/// Erases c from the back of the string
PLL pop back(PLL hash, char c) {
 return (((hash - c) * inv_base) % mod + mod) % mod;
/// Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod + mod) % mod;
/// concatenates two strings where length of the right is
PLL concat(PLL left, PLL right, int k) { return (left * P
    [k] + right) % mod; }
/// Calculates hash of string with size len repeated cnt
/// This is O(\log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
 PLL ret = (hash * mul) % mod:
 if (P[len].ff == 1) ret.ff = hash.ff * cnt;
 if (P[len].ss == 1) ret.ss = hash.ss * cnt;
 return ret:
```

```
struct hashlist {
 int len:
 vector<PLL> H. R:
 hashlist() {}
  hashlist(string& s) {
   len = (int)s.size();
   hash init(len);
   H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL
   for (int i = 1; i <= len; i++) H[i] = append(H[i -</pre>
       1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i] = append(R[i +
       1], s[i - 1]);
 /// 1-indexed
  PLL range_hash(int 1, int r) {
   return ((H[r] - H[1 - 1] * P[r - 1 + 1]) % mod + mod)
        % mod:
 }
  PLL reverse hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 + 1]) \% mod + mod)
        % mod;
  PLL concat_range_hash(int l1, int r1, int l2, int r2) {
   return concat(range_hash(l1, r1), range_hash(l2, r2),
        r2 - 12 + 1);
 PLL concat reverse hash(int 11, int r1, int 12, int r2)
   return concat(reverse hash(12, r2), reverse hash(11,
       r1), r1 - 11 + 1);
 }
};
```

# 4.3 Aho Corasick

```
struct AC {
int N, P;
const int A = 26;
vector<vector<int>> next;
vector<int> link, out_link;
vector<vector<int>> out;
AC() : N(0), P(0) \{ node(); \}
int node() {
 next.emplace_back(A, 0);
 link.emplace back(0);
 out_link.emplace_back(0);
 out.emplace_back(0);
 return N++;
```

```
LL get(PLL hash) { return ((hash.ff << 32) ^ hash.ss); } inline int get(char c) { return c - 'a'; }
                                                          int add pattern(const string T) {
                                                            int u = 0:
                                                            for (auto c : T) {
                                                              if (!next[u][get(c)]) next[u][get(c)] = node();
                                                              u = next[u][get(c)];
                                                            out[u].push back(P);
                                                            return P++:
                                                           void compute() {
                                                            queue<int> q;
                                                            for (q.push(0); !q.empty();) {
                                                              int u = q.front(); q.pop();
                                                              for (int c = 0; c < A; ++c) {
                                                                int v = next[u][c];
                                                                if (!v) next[u][c] = next[link[u]][c];
                                                                else {
                                                                 link[v] = u ? next[link[u]][c] : 0;
                                                                 out link[v] = out[link[v]].empty() ? out_link[
                                                                      link[v]] : link[v];
                                                                 q.push(v);
                                                                }
                                                          int advance(int u, char c) {
                                                            while (u && !next[u][get(c)]) u = link[u];
                                                            u = next[u][get(c)];
                                                            return u:
                                                           void match(const string S) {
                                                            int u = 0:
                                                            for (auto c : S) {
                                                              u = advance(u, c);
                                                              for (int v = u; v; v = out_link[v]) {
                                                                for (auto p : out[v]) cout << "match " << p << endl</pre>
                                                              }
                                                            }
                                                          int main() {
                                                            AC aho; int n; cin >> n;
                                                            while (n--) {
                                                             string s; cin >> s;
                                                              aho.add_pattern(s);
                                                            aho.compute(); string text;
                                                            cin >> text; aho.match(text);
```

```
return 0;
```

#### 4.4 KMP

```
vector<int> prefix_function(string s) {
 int n = (int)s.length();
 vector<int> pi(n);
 for (int i = 1; i < n; i++) {</pre>
   int j = pi[i - 1];
   while (j > 0 \&\& s[i] != s[j])
    i = pi[i - 1];
   if (s[i] == s[i])
     j++;
  pi[i] = j;
 return pi;
```

#### 4.5 Manacher's

```
vector<int> d1(n):
// d[i] = number of palindromes taking s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
 while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k])
      k++:
 d1[i] = k--:
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
// d[i] = number of palindromes taking s[i-1] and s[i] as
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1)
 while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s
      [i + k]) k++;
 d2[i] = k--;
 if (i + k > r) l = i - k - 1, r = i + k;
```

#### 4.6 Suffix Match FFT

```
// Find occurrences of t in s where '?'s are
    automatically matched with any character
// \text{ res}[i + m - 1] = \text{sum}_j = 0 \text{ to } m - 1 \{ s[i + j] * t[j] * \}
    (s[i + j] - t[j]) 
vector<int> string_matching(string &s, string &t) {
 int n = s.size(), m = t.size();
 vector<int> s1(n), s2(n), s3(n);
```

```
for(int i = 0; i < n; i++)</pre>
  s1[i] = s[i] == ?????0 : s[i] - a? + 1; // assign
      any non zero number for non '?'s
for(int i = 0: i < n: i++)
  s2[i] = s1[i] * s1[i];
for(int i = 0; i < n; i++)</pre>
  s3[i] = s1[i] * s2[i]:
vector<int> t1(m), t2(m), t3(m);
for(int i = 0; i < m; i++)</pre>
 t1[i] = t[i] == '?' ? 0 : t[i] - 'a' + 1;
for(int i = 0; i < m; i++)</pre>
 t2[i] = t1[i] * t1[i];
for(int i = 0; i < m; i++)</pre>
  t3[i] = t1[i] * t2[i];
reverse(t1.begin(), t1.end());
reverse(t2.begin(), t2.end());
reverse(t3.begin(), t3.end());
vector<int> s1t3 = multiply(s1, t3);
vector<int> s2t2 = multiply(s2, t2);
vector<int> s3t1 = multiply(s3, t1);
vector<int> res(n):
for(int i = 0; i < n; i++)</pre>
 res[i] = s1t3[i] - s2t2[i] * 2 + s3t1[i];
vector<int> oc;
for(int i = m - 1; i < n; i++)</pre>
 if(res[i] == 0)
   oc.push_back(i - m + 1);
return oc;
```

## 4.7 Suffix Array

```
for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;</pre>
 c[0][p[0]] = 0;
 int classes = 1:
 for (int i = 1; i < n; i++) {
   if (s[p[i]] != s[p[i - 1]]) classes++;
   c[0][p[i]] = classes - 1;
 VI pn(n), cn(n);
 cnt.resize(n);
 for (int h = 0; (1 << h) < n; h++) {
   for (int i = 0; i < n; i++) {
     pn[i] = p[i] - (1 << h);
     if (pn[i] < 0) pn[i] += n;
   fill(cnt.begin(), cnt.end(), 0);
   /// radix sort
   for (int i = 0; i < n; i++) cnt[c[h][pn[i]]]++;</pre>
   for (int i = 1; i < classes; i++) cnt[i] += cnt[i -</pre>
   for (int i = n - 1; i >= 0; i--) p[--cnt[c[h][pn[i
       ]]]] = pn[i];
   cn[p[0]] = 0;
   classes = 1;
   for (int i = 1; i < n; i++) {</pre>
     PII cur = \{c[h][p[i]], c[h][(p[i] + (1 << h)) \% n\}
     PII prev = \{c[h][p[i-1]], c[h][(p[i-1] + (1 << \})\}
         h)) % n]};
     if (cur != prev) ++classes;
     cn[p[i]] = classes - 1;
   }
   c.push back(cn);
 return p;
VI suffix_array_construction(string s) {
 VI sorted_shifts = sort_cyclic_shifts(s);
 sorted_shifts.erase(sorted_shifts.begin());
 return sorted shifts;
/// LCP between the ith and jth (i != j) suffix of the
int suffixLCP(int i, int j) {
 assert(i != j);
```

```
int log_n = c.size() - 1;
 int ans = 0;
 for (int k = log_n; k >= 0; k--) {
   if (c[k][i] == c[k][j]) {
     ans += 1 << k;
    i += 1 << k:
     i += 1 << k;
 }
 return ans;
VI lcp_construction(const string &s, const VI &sa) {
 int n = s.size();
 VI rank(n, 0);
 VI lcp(n - 1, 0);
 for (int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
 for (int i = 0, k = 0; i < n; i++, k -= (k != 0)) {
   if (rank[i] == n - 1) {
    k = 0;
     continue;
   int j = sa[rank[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
       1) k++;
   lcp[rank[i]] = k;
 return lcp;
```

#### 4.8 Trie

```
endMark[cur] = true;
  bool search(const string &s) {
   int cur = 0;
   for (auto c : s) {
     int val = c - 'a':
     if (!trie[cur][val])
       return false:
     cur = trie[cur][val];
   return endMark[cur];
private:
  int id, endMark[100005];
  vector<int> trie[100005];
};
```

#### 4.9 Z Algo

```
vector<int> calcz(string s) {
 int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {</pre>
   if (i > r) {
    l = r = i;
     while (r < n \&\& s[r] == s[r - 1]) r++;
     z[i] = r - 1, r--;
   } else {
     int k = i - 1:
     if (z[k] < r - i + 1) z[i] = z[k];
     else {
      1 = i:
       while (r < n \&\& s[r] == s[r - 1]) r++;
      z[i] = r - 1, r--;
     }
   }
 }
 return z;
```

# DP

#### 5.1 Bitmask

```
for(int mask= 0; mask < (1 << 4); mask++){</pre>
 ll sum of set = 0;
 for(int i = 0; (111 << i) <= mask; i++) if(mask& (111</pre>
      << i)) sum_of_set += v[i];
  if(sum of set == S){
    cout << "Yes\n";</pre>
```

```
flg = true;
   break;
 }
if(!flg) cout << "No\n";</pre>
```

### 5.2 LIS

```
vector<pair<11, 11>> LIS(vector<11> &v){
 11 n=v.size():
 vector<pair<11, 11>> seq(n); //{size, last element}
 set<ll> s; //multiset for non_dcrs
 for(int i=0; i<n; ++i){</pre>
   auto it=s.lower bound(v[i]);
   if(it==s.end()) s.insert(v[i]);
   else{
     s.erase(it);
     s.insert(v[i]):
   seq[i]={s.size(), *(s.rbegin())};
 return seq;
} //seq[i] = {size of LIS in v[0, i], largest element in
    that sequence}
```

## 5.3 Divide and Conquer DP

```
const int K = 805, N = 4005;
LL dp[2][N], _cost[N][N];
// 1-indexed for convenience
LL cost(int 1, int r) {
 return _cost[r][r] - _cost[1 - 1][r] - _cost[r][1 - 1]
      + cost[1 - 1][1 - 1] >> 1;
void compute(int cnt, int 1, int r, int optl, int optr) { | 11 bigmod(11 a, 11 b, 11 m) {
 if (1 > r) return:
 int mid = 1 + r >> 1;
 LL best = INT MAX;
 int opt = -1;
 for (int i = optl; i <= min(mid, optr); i++) {</pre>
   LL cur = dp[cnt ^1][i - 1] + cost(i, mid);
   if (cur < best) best = cur, opt = i;</pre>
  dp[cnt][mid] = best;
  compute(cnt, 1, mid - 1, optl, opt);
  compute(cnt, mid + 1, r, opt, optr);
LL dnc dp(int k, int n) {
 fill(dp[0] + 1, dp[0] + n + 1, INT_MAX);
 for (int cnt = 1; cnt <= k; cnt++) {</pre>
   compute(cnt & 1, 1, n, 1, n);
```

```
return dp[k & 1][n];
```

# 5.4 Knuth Optimization

```
const int N = 1005;
LL dp[N][N], a[N];
int opt[N][N];
LL cost(int i, int j) { return a[j + 1] - a[i]; }
LL knuth_optimization(int n) {
 for (int i = 0; i < n; i++) {
   dp[i][i] = 0;
   opt[i][i] = i;
 for (int i = n - 2; i \ge 0; i--) {
   for (int j = i + 1; j < n; j++) {
    LL mn = LLONG MAX:
    LL c = cost(i, j);
     for (int k = opt[i][j-1]; k \le min(j-1, opt[i+1])
          1][i]); k++) {
       if (mn > dp[i][k] + dp[k + 1][j] + c) {
         mn = dp[i][k] + dp[k + 1][j] + c;
         opt[i][i] = k;
      }
     dp[i][j] = mn;
 return dp[0][n - 1];
```

# 6 Math

# 6.1 BigMod

```
if(b == 0) return 1:
11 x = bigmod(a, b/2, m);
x = (x * x) % m;
if(b \% 2) x = (x * a) \% m;
return x;
```

#### 6.2 Combi

```
array<int, N + 1> fact, inv, inv_fact;
void init() {
 fact[0] = inv fact[0] = 1;
 for (int i = 1; i <= N; i++) {</pre>
   inv[i] = i == 1 ? 1 : (LL)inv[i - mod % i] * (mod / i
        + 1) % mod:
   fact[i] = (LL)fact[i - 1] * i % mod;
   inv_fact[i] = (LL)inv_fact[i - 1] * inv[i] % mod;
```

#### 6.3 Sieve

```
const ll m = 10e6;
vector<ll> lp(m+1);
vector<ll> prime;
void ln_sieve() {
  for(ll i = 2; i <= m; i++){
    if(!lp[i]){
     lp[i] = i;
     prime.push_back(i);
    }

  for(ll j = 0; i * prime[j] <= m; j++){
     lp[i * prime[j]] = prime[j];
     if(prime[j] == lp[i]) break;
    }
}</pre>
```

#### 6.4 Linear Sieve

```
const int N = 1e7;
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5], cnt[N + 5], POW[N
     + 5];
bool prime[N + 5];
int SOD[N + 5];
void init() {
 fill(prime + 2, prime + N + 1, 1);
 SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
 for (LL i = 2; i <= N; i++) {</pre>
   if (prime[i]) {
     primes.push_back(i), spf[i] = i;
     phi[i] = i - 1;
     NOD[i] = 2, cnt[i] = 1;
     SOD[i] = i + 1, POW[i] = i;
   }
   for (auto p : primes) {
     if (p * i > N or p > spf[i]) break;
     prime[p * i] = false, spf[p * i] = p;
     if (i % p == 0) {
      phi[p * i] = p * phi[i];
       NOD[p * i] = NOD[i] / (cnt[i] + 1) * (cnt[i] + 2)
              cnt[p * i] = cnt[i] + 1;
```

// LL ans = a \* b - mod \* (LL) (1.L / mod \* a \* b);

# 6.5 Pollard Rho

LL mul(LL a, LL b, LL mod) {

return (\_\_int128)a \* b % mod;

```
// return ans + mod * (ans < 0) - mod * (ans >= (LL)
       mod):
LL bigmod(LL num, LL pow, LL mod) {
   LL ans = 1:
   for (; pow > 0; pow >>= 1, num = mul(num, num, mod))
       if (pow & 1) ans = mul(ans, num, mod);
   return ans;
bool is_prime(LL n) {
   if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;
   LL a[] = \{2, 325, 9375, 28178, 450775, 9780504,
       1795265022};
   LL s = builtin ctzll(n - 1), d = n \gg s;
   for (LL x : a) {
      LL p = bigmod(x \% n, d, n), i = s;
       for (; p != 1 and p != n - 1 and x % n and i--; p
            = mul(p, p, n))
       if (p != n - 1 and i != s) return false;
   return true;
LL get_factor(LL n) {
   auto f = [\&](LL x) \{ return mul(x, x, n) + 1; \};
   LL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
   for (; t++ % 40 or gcd(prod, n) == 1; x = f(x), y = f(x)
       (f(y))) {
       (x == y) ? x = i++, y = f(x) : 0;
       prod = (q = mul(prod, max(x, y) - min(x, y), n))
           ? q : prod;
   return gcd(prod, n);
```

```
map<LL, int> factorize(LL n) {
   map<LL, int> res;
   if (n < 2) return res;</pre>
   LL small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23,
        29, 31, 37, 41,
                                           43, 47, 53, 59,
                                                61, 67,
                                               71, 73, 79,
                                                83, 89,
                                               97};
   for (LL p : small_primes)
       for (; n % p == 0; n /= p, res[p]++)
   auto _factor = [&](LL n, auto &_factor) {
       if (n == 1) return;
       if (is_prime(n))
           res[n]++;
       else {
           LL x = get_factor(n);
           _factor(x, _factor);
           _factor(n / x, _factor);
   }:
   _factor(n, _factor);
   return res;
```

### 6.6 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
if (b == 0)
   return {1, 0, a};
 else {
   auto [x, y, g] = EGCD(b, a \% b);
   return \{y, x - a / b * y, g\};
// given modulo equations, will apply CRT
PLL CRT(vector<PLL> &v) {
 LL V = 0, M = 1:
 for (auto &[v, m] : v) { // value % mod
   auto [x, y, g] = EGCD(M, m);
   if ((v - V) % g != 0) return {-1, 0};
   V += x * (v - V) / g % (m / g) * M, M *= m / g;
   V = (V \% M + M) \% M;
 return make_pair(V, M);
```

#### 6.7 Mobius Function

```
const int N = 1e6 + 5;
int mob[N];
void mobius() {
  memset(mob, -1, sizeof mob);
  mob[1] = 1;
  for (int i = 2; i < N; i++)
    if (mob[i]) {
      for (int j = i + i; j < N; j += i) mob[j] -= mob[i
        ];
    }
}</pre>
```

#### 6.8 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
  assert((n & (n - 1)) == 0), N = n;
  perm = vector<int>(N, 0);
  for (int k = 1; k < N; k <<= 1) {
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
   }
  wp[0] = wp[1] = vector < CD > (N);
  for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N)
   wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
        N));
 }
void fft(vector<CD> &v, bool invert = false) {
  if (v.size() != perm.size()) precalculate(v.size());
  for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
  for (int len = 2; len <= N; len *= 2) {</pre>
   for (int i = 0, d = N / len; i < N; i += len) {</pre>
     for (int j = 0, idx = 0; j < len / 2; j++, idx += d
         ) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j + len / 2];
```

```
v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
 if (invert) {
   for (int i = 0; i < N; i++) v[i] /= N;
void pairfft(vector<CD> &a, vector<CD> &b, bool invert =
    false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0,
     1);
 fft(p, invert);
 p.push_back(p[0]);
 for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
 }
vector<LL> multiply(const vector<LL> &a, const vector<LL>
     &b) {
 int n = 1:
 while (n < a.size() + b.size()) n <<= 1;</pre>
 vector <CD > fa(a.begin(), a.end()), fb(b.begin(), b.end
      ());
 fa.resize(n);
 fb.resize(n);
          fft(fa); fft(fb);
 pairfft(fa, fb);
 for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];</pre>
 fft(fa, true);
 vector<LL> ans(n):
 for (int i = 0; i < n; i++) ans[i] = round(fa[i].real()</pre>
     );
 return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const vector<LL> &
    b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<= 1;</pre>
 vector<CD> al(n), ar(n), bl(n), br(n);
```

```
for (int i = 0; i < a.size(); i++) al[i] = a[i] % M / B</pre>
    , ar[i] = a[i] % M % B;
for (int i = 0; i < b.size(); i++) bl[i] = b[i] % M / B</pre>
    , br[i] = b[i] % M % B;
pairfft(al, ar);
pairfft(bl, br);
         fft(al); fft(ar); fft(bl); fft(br);
for (int i = 0; i < n; i++) {</pre>
  CD ll = (al[i] * bl[i]), lr = (al[i] * br[i]);
  CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
  al[i] = 11;
  ar[i] = lr;
 bl[i] = rl;
  br[i] = rr;
pairfft(al, ar, true);
pairfft(bl, br, true);
         fft(al, true); fft(ar, true); fft(bl, true);
    fft(br. true):
vector<LL> ans(n):
for (int i = 0; i < n; i++) {</pre>
 LL right = round(br[i].real()), left = round(al[i].
      real());
 LL mid = round(round(bl[i].real()) + round(ar[i].real
  ans[i] = ((left \% M) * B * B + (mid \% M) * B + right)
       % M;
return ans;
```

#### 6.9 NTT

```
const LL N = 1 << 18;
const LL MOD = 786433;

vector<LL> P[N];
LL rev[N], w[N | 1], a[N], b[N], inv_n, g;
LL Pow(LL b, LL p) {
    LL ret = 1;
    while (p) {
        if (p & 1) ret = (ret * b) % MOD;
        b = (b * b) % MOD;
        p >>= 1;
    }
    return ret;
}
LL primitive_root(LL p) {
    vector<LL> factor;
    LL phi = p - 1, n = phi;
```

```
for (LL i = 2; i * i <= n; i++) {</pre>
   if (n % i) continue;
   factor.emplace_back(i);
   while (n % i == 0) n /= i:
  if (n > 1) factor.emplace_back(n);
  for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true;
   for (LL i = 0; i < factor.size() && ok; i++)</pre>
     ok &= Pow(res, phi / factor[i]) != 1;
   if (ok) return res;
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - builtin clz(n));
 LL r = Pow(g, (MOD - 1) / n);
 inv_n = Pow(n, MOD - 2);
  w[0] = w[n] = 1:
  for (LL i = 1; i < n; i++) w[i] = (w[i-1] * r) % MOD;
 for (LL i = 1; i < n; i++)</pre>
   rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
 for (LL i = 1: i < n - 1: i++)
   if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (LL m = 2; m <= n; m <<= 1) {</pre>
   for (LL i = 0; i < n; i += m) {</pre>
     for (LL j = 0; j < (m >> 1); j++) {
       LL &u = a[i + j], &v = a[i + j + (m >> 1)];
       LL t = v * w[dir ? n - n / m * j : n / m * j] %
       v = u - t < 0 ? u - t + MOD : u - t;
       u = u + t >= MOD ? u + t - MOD : u + t;
     }
   }
 }
   for (LL i = 0; i < n; i++) a[i] = (inv_n * a[i]) %</pre>
        MOD:
vector<LL> mul(vector<LL> p, vector<LL> q) {
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
  while (sz < t) sz <<= 1;
  prepare(sz);
 for (LL i = 0; i < n; i++) a[i] = p[i];</pre>
 for (LL i = 0; i < m; i++) b[i] = q[i];
 for (LL i = n; i < sz; i++) a[i] = 0;</pre>
```

```
for (LL i = m; i < sz; i++) b[i] = 0;</pre>
 NTT(a, sz);
 NTT(b. sz):
 for (LL i = 0; i < sz; i++) a[i] = (a[i] * b[i]) % MOD;</pre>
 NTT(a, sz, 1);
 vector<LL> c(a, a + sz);
 while (c.size() && c.back() == 0) c.pop_back();
 return c;
6.10 ModInverse
//solves ax+by=gcd(a, b) i guess
int gcd(int a, int b, int& x, int& y) {
 x = 1, y = 0;
 int x1 = 0, y1 = 1, a1 = a, b1 = b;
 while (b1) {
   int q = a1 / b1;
   tie(x, x1) = make_tuple(x1, x - q * x1);
   tie(y, y1) = make_tuple(y1, y - q * y1);
   tie(a1, b1) = make tuple(b1, a1 - q * b1);
 return a1:
//finds mod inverse?
int x, v;
int g = gcd(a, m, x, y);
if (g != 1) {
 cout << "No solution!":</pre>
else {
 x = (x \% m + m) \% m;
 cout << x << endl;</pre>
6.11 Diophantine
```

```
int gcd(int a, int b, int& x, int& y) {
   if (b == 0) {
      x = 1;
      y = 0;
      return a;
   }
   int x1, y1;
   int d = gcd(b, a % b, x1, y1);
   x = y1;
   y = x1 - y1 * (a / b);
   return d;
```

```
bool find_any_solution(int a, int b, int c, int &x0, int
    &y0, int &g) {
 g = gcd(abs(a), abs(b), x0, y0);
 if (c % g) {
  return false;
 x0 *= c / g;
 v0 *= c / g;
 if (a < 0) x0 = -x0;
 if (b < 0) y0 = -y0;
 return true;
void shift_solution(int & x, int & y, int a, int b, int
    cnt) {
 x += cnt * b:
 y -= cnt * a;
int find_all_solutions(int a, int b, int c, int minx, int
     maxx, int miny, int maxy) {
 int x, y, g;
 if (!find_any_solution(a, b, c, x, y, g))
   return 0;
 a /= g;
 b /= g;
 int sign a = a > 0 ? +1 : -1;
 int sign_b = b > 0 ? +1 : -1;
 shift_solution(x, y, a, b, (minx - x) / b);
 if (x < minx)
   shift_solution(x, y, a, b, sign_b);
 if (x > maxx)
   return 0:
 int lx1 = x;
 shift_solution(x, y, a, b, (maxx - x) / b);
 if (x > maxx)
   shift_solution(x, y, a, b, -sign_b);
 int rx1 = x;
 shift_solution(x, y, a, b, -(miny - y) / a);
 if (y < miny)</pre>
   shift_solution(x, y, a, b, -sign_a);
 if (y > maxy)
   return 0;
```

```
int 1x2 = x;

shift_solution(x, y, a, b, -(maxy - y) / a);
if (y > maxy)
    shift_solution(x, y, a, b, sign_a);
int rx2 = x;

if (1x2 > rx2)
    swap(1x2, rx2);
int 1x = max(1x1, 1x2);
int rx = min(rx1, rx2);

if (1x > rx)
    return 0;
return (rx - 1x) / abs(b) + 1;
}
```

# 7 Geometry

#### 7.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) \{ return abs(x) < EPS ? 0 : (x < 0 ? -1 : 
struct Point {
   Ti x, y;
   Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
   Point operator+(const Point& u) const { return Point(
        x + u.x, y + u.y; }
   Point operator-(const Point& u) const { return Point(
        x - u.x, y - u.y); }
   Point operator*(const LL u) const { return Point(x *
        u, y * u); }
   Point operator*(const Tf u) const { return Point(x *
        u, y * u); }
   Point operator/(const Tf u) const { return Point(x /
        u, y / u); }
   bool operator==(const Point& u) const {
       return dcmp(x - u.x) == 0 && dcmp(y - u.y) == 0;
   bool operator!=(const Point& u) const { return !(*
        this == u): }
   bool operator<(const Point& u) const {</pre>
       return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x) == 0
           && dcmp(y - u.y) < 0);
};
```

```
Ti dot(Point a, Point b) { return a.x * b.x + a.y * b.y;
Ti cross(Point a, Point b) { return a.x * b.y - a.y * b.x
Tf length(Point a) { return sqrt(dot(a, a)); }
Ti sqLength(Point a) { return dot(a, a); }
Tf distance(Point a, Point b) { return length(a - b); }
Tf angle(Point u) { return atan2(u.v, u.x); }
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Point a, Point b) {
   Tf ans = angle(b) - angle(a);
   return ans <= -PI ? ans + 2 * PI : (ans > PI ? ans -
        2 * PI : ans):
// Rotate a ccw by rad radians, Tf Ti same
Point rotate(Point a, Tf rad) {
    return Point(a.x * cos(rad) - a.y * sin(rad),
                           a.x * sin(rad) + a.y * cos(rad)
                               )):
// rotate a ccw by angle th with cos(th) = co && sin(th)
    = si, tf ti same
Point rotatePrecise(Point a, Tf co, Tf si) {
   return Point(a.x * co - a.y * si, a.y * co + a.x * si
Point rotate90(Point a) { return Point(-a.y, a.x); }
// scales vector a by s such that length of a becomes s,
    Tf Ti same
Point scale(Point a, Tf s) { return a / length(a) * s; }
// returns an unit vector perpendicular to vector a, Tf
    Ti same
Point normal(Point a) {
   Tf 1 = length(a);
   return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if on ab && -1 if
    right of ab
int orient(Point a, Point b, Point c) { return dcmp(cross
    (b - a, c - a): }
/// Use as sort(v.begin(), v.end(), polarComp(0, dir))
/// Polar comparator around O starting at direction dir
struct polarComp {
   Point O, dir;
    polarComp(Point 0 = Point(0, 0), Point dir = Point(1,
         0)) : O(O), dir(dir) {}
    bool half(Point p) {
       return dcmp(cross(dir, p)) < 0 ||</pre>
```

```
(dcmp(cross(dir, p)) == 0 && dcmp(dot
                        (dir, p)) > 0);
   bool operator()(Point p, Point q) {
       return make_tuple(half(p), 0) < make_tuple(half(q))</pre>
           ), cross(p, q));
struct Segment {
   Point a, b;
   Segment(Point aa, Point bb) : a(aa), b(bb) {}
typedef Segment Line;
struct Circle {
   Point o;
   Tf r;
   Circle(Point o = Point(0, 0), Tf r = 0) : o(o), r(r)
   // returns true if point p is in || on the circle
   bool contains(Point p) { return dcmp(sqLength(p - o)
       - r * r) <= 0; }
   // returns a point on the circle rad radians away
       from +X CCW
   Point point(Tf rad) {
       static_assert(is_same<Tf, Ti>::value);
       return Point(o.x + cos(rad) * r, o.y + sin(rad) *
            r);
   // area of a circular sector with central angle rad
   Tf area(Tf rad = PI + PI) { return rad * r * r / 2: }
   // area of the circular sector cut by a chord with
       central angle alpha
   Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - 1)
       sin(alpha)); }
```

#### 7.2 Linear

```
Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) <</pre>
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.a)) == 0;
// lines are represented as a ray from a point: (point,
// returns false if two lines (p, v) && (q, w) are
    parallel or collinear
// true otherwise, intersection point is stored at o via
    reference, Tf Ti Same
bool lineLineIntersection(Point p, Point v, Point q,
    Point w, Point& o) {
 if (dcmp(cross(v, w)) == 0) return false;
 Point u = p - q;
  o = p + v * (cross(w, u) / cross(v, w));
 return true:
// returns false if two lines p && q are parallel or
    collinear
// true otherwise, intersection point is stored at o via
    reference
bool lineLineIntersection(Line p, Line q, Point& o) {
 return lineLineIntersection(p.a, p.b - p.a, q.a, q.b -
      q.a, o);
// returns the distance from point a to line 1
// **** LINE LINE INTERSECTION FINISH ****
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(l.b - l.a, p - l.a) / length(l.b - l.a
// returns the shortest distance from point a to segment
Tf distancePointSegment(Point p, Segment s) {
  if (s.a == s.b) return length(p - s.a);
  Point v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
  if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
  else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
  else
   return abs(cross(v1, v2) / length(v1));
```

```
// returns the shortest distance from segment p to
    segment q

Tf distanceSegmentSegment(Segment p, Segment q) {
    if (segmentsIntersect(p, q)) return 0;
    Tf ans = distancePointSegment(p.a, q);
    ans = min(ans, distancePointSegment(p.b, q));
    ans = min(ans, distancePointSegment(q.a, p));
    ans = min(ans, distancePointSegment(q.b, p));
    return ans;
}

// returns the projection of point p on line 1, Tf Ti
    Same
Point projectPointLine(Point p, Line 1) {
    Point v = 1.b - 1.a;
    return 1.a + v * ((Tf)dot(v, p - 1.a) / dot(v, v));
}

7.3 Circular
```

```
// Extremely inaccurate for finding near touches
// compute intersection of line 1 with circle c
// The intersections are given in order of the ray (1.a,
    1.b), Tf Ti same
vector<Point> circleLineIntersection(Circle c, Line 1) {
   vector<Point> ret:
   Point b = 1.b - 1.a, a = 1.a - c.o;
   Tf A = dot(b, b), B = dot(a, b);
   Tf C = dot(a, a) - c.r * c.r, D = B * B - A * C;
   if (D < -EPS) return ret;</pre>
   ret.push_back(1.a + b * (-B - sqrt(D + EPS)) / A);
   if (D > EPS) ret.push_back(l.a + b * (-B + sqrt(D)) /
         A);
   return ret:
// signed area of intersection of circle(c.o, c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)/2]
Tf circleTriangleIntersectionArea(Circle c, Segment s) {
   using Linear::distancePointSegment;
   Tf OA = length(c.o - s.a);
   Tf OB = length(c.o - s.b);
   // sector
   if (dcmp(distancePointSegment(c.o, s) - c.r) >= 0)
       return angleBetween(s.a - c.o, s.b - c.o) * (c.r
           * c.r) / 2.0;
   // triangle
   if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r) \le 0)
       return cross(c.o - s.b, s.a - s.b) / 2.0;
   // three part: (A, a) (a, b) (b, B)
   vector<Point> Sect = circleLineIntersection(c, s);
   return circleTriangleIntersectionArea(c, Segment(s.a,
         Sect[0])) +
```

```
circleTriangleIntersectionArea(c,
                     Segment(Sect[0], Sect[1])) +
                circleTriangleIntersectionArea(c,
                     Segment(Sect[1], s.b));
// area of intersecion of circle(c.o, c.r) && simple
    polyson(p[])
Tf circlePolyIntersectionArea(Circle c, Polygon p) {
   Tf res = 0:
   int n = p.size();
   for (int i = 0; i < n; ++i)</pre>
       res += circleTriangleIntersectionArea(c, Segment(
           p[i], p[(i + 1) % n]));
   return abs(res);
// locates circle c2 relative to c1
// interior
                      (d < R - r)
                                         ---> -2
// interior tangents (d = R - r)
                                      ---> -1
// concentric
                   (d = 0)
// secants
                     (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r)
// exterior
                      (d > R + r)
int circleCirclePosition(Circle c1, Circle c2) {
   Tf d = length(c1.o - c2.o);
   int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (abs(c1.r - c2.r)))
        c1.r + c2.r));
   return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex
       > 0 ? 2 : 0;
// compute the intersection points between two circles c1
     && c2, Tf Ti same
vector<Point> circleCircleIntersection(Circle c1. Circle
    c2) {
   vector<Point> ret;
   Tf d = length(c1.o - c2.o);
   if (dcmp(d) == 0) return ret;
   if (dcmp(c1.r + c2.r - d) < 0) return ret;
   if (dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
   Point v = c2.0 - c1.0:
   Tf co = (c1.r * c1.r + sqLength(v) - c2.r * c2.r) /
        (2 * c1.r * length(v));
   Tf si = sqrt(abs(1.0 - co * co));
   Point p1 = scale(rotatePrecise(v, co, -si), c1.r) +
   Point p2 = scale(rotatePrecise(v, co, si), c1.r) + c1
        .0;
   ret.push_back(p1);
   if (p1 != p2) ret.push back(p2);
```

```
return ret;
// intersection area between two circles c1, c2
Tf circleCircleIntersectionArea(Circle c1, Circle c2) {
    Point AB = c2.o - c1.o;
   Tf d = length(AB);
    if (d >= c1.r + c2.r) return 0;
    if (d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
    if (d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
   Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r))
        /(2.0 * c1.r * d));
   Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r))
        /(2.0 * c2.r * d));
    return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
// returns tangents from a point p to circle c, Tf Ti
vector<Point> pointCircleTangents(Point p, Circle c) {
    vector<Point> ret;
    Point u = c.o - p;
    Tf d = length(u);
    if (d < c.r)
    else if (dcmp(d - c.r) == 0) {
       ret = {rotate(u, PI / 2)};
   } else {
       Tf ang = asin(c.r / d);
       ret = {rotate(u, -ang), rotate(u, ang)};
   }
    return ret;
// returns the points on tangents that touches the circle
    , Tf Ti Same
vector<Point> pointCircleTangencyPoints(Point p, Circle c
    ) {
    Point u = p - c.o;
    Tf d = length(u);
    if (d < c.r)
       return {}:
    else if (dcmp(d - c.r) == 0)
       return {c.o + u};
    else {
       Tf ang = acos(c.r / d);
       u = u / length(u) * c.r;
       return {c.o + rotate(u, -ang), c.o + rotate(u,
            ang)};
   }
}
```

```
// for two circles c1 && c2, returns two list of points a
// such that a[i] is on c1 && b[i] is c2 && for every i
// Line(a[i], b[i]) is a tangent to both circles
// CAUTION: a[i] = b[i] in case they touch | -1 for c1 =
int circleCircleTangencyPoints(Circle c1, Circle c2,
    vector<Point> &a, vector<Point> &b) {
       a.clear(), b.clear();
       int cnt = 0;
       if (dcmp(c1.r - c2.r) < 0) {
              swap(c1, c2);
              swap(a, b);
       }
       Tf d2 = sqLength(c1.o - c2.o);
       Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r;
       if (dcmp(d2 - rdif * rdif) < 0)
              return 0;
       if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r) == 0)
              return -1;
       Tf base = angle(c2.o - c1.o);
       if (dcmp(d2 - rdif * rdif) == 0) {
              a.push_back(c1.point(base));
              b.push_back(c2.point(base));
              cnt++;
              return cnt;
       }
       Tf ang = acos((c1.r - c2.r) / sqrt(d2));
       a.push_back(c1.point(base + ang));
       b.push_back(c2.point(base + ang));
       cnt++;
       a.push_back(c1.point(base - ang));
       b.push_back(c2.point(base - ang));
       cnt++:
       if (dcmp(d2 - rsum * rsum) == 0) {
              a.push_back(c1.point(base));
              b.push_back(c2.point(PI + base));
       } else if (dcmp(d2 - rsum * rsum) > 0) {
              Tf ang = acos((c1.r + c2.r) / sqrt(d2));
              a.push_back(c1.point(base + ang));
              b.push_back(c2.point(PI + base + ang));
              cnt++;
```

```
a.push_back(c1.point(base - ang));
b.push_back(c2.point(PI + base - ang));
cnt++;
}
return cnt;
```

```
7.4 Convex
 /// minkowski sum of two polygons in O(n)
Polygon minkowskiSum(Polygon A, Polygon B) {
          int n = A.size(), m = B.size();
          rotate(A.begin(), min_element(A.begin(), A.end()), A.
                     end());
         rotate(B.begin(), min_element(B.begin(), B.end()), B.
                     end());
          A.push_back(A[0]);
          B.push back(B[0]);
          for (int i = 0; i < n; i++) A[i] = A[i + 1] - A[i];
          for (int i = 0; i < m; i++) B[i] = B[i + 1] - B[i];
          Polygon C(n + m + 1);
          C[0] = A.back() + B.back();
          merge(A.begin(), A.end() - 1, B.begin(), B.end() - 1,
                        C.begin() + 1,
                                       polarComp(Point(0, 0), Point(0, -1)));
         for (int i = 1; i < C.size(); i++) C[i] = C[i] + C[i</pre>
                     - 1]:
         C.pop_back();
         return C;
// finds the rectangle with minimum area enclosing a
            convex polygon and
// the rectangle with minimum perimeter enclosing a
            convex polygon
// Tf Ti Same
pair<Tf, Tf> rotatingCalipersBoundingBox(const Polygon &p
           ) {
          using Linear::distancePointLine;
          int n = p.size();
          int 1 = 1, r = 1, j = 1;
         Tf area = 1e100;
          Tf perimeter = 1e100;
         for (int i = 0; i < n; i++) {</pre>
                    Point v = (p[(i + 1) \% n] - p[i]) / length(p[(i + 1) \% n]) / length(p
                                  1) % n] - p[i]);
                    while (dcmp(dot(v, p[r % n] - p[i]) - dot(v, p[(r
                                 + 1) % n] - p[i])) < 0)
                    while (j < r \mid | dcmp(cross(v, p[j % n] - p[i]) -
```

```
cross(v, p[(
                                                 j + 1)
                                                 % n] -
                                                  p[i]))
                                                  < 0)
           j++;
       while (1 < j ||
                    dcmp(dot(v, p[1 % n] - p[i]) - dot(v,
                         p[(1 + 1) \% n] - p[i])) > 0)
           1++;
       Tf w = dot(v, p[r \% n] - p[i]) - dot(v, p[1 \% n])
       Tf h = distancePointLine(p[j % n], Line(p[i], p[(
           i + 1) % n]));
       area = min(area, w * h);
       perimeter = min(perimeter, 2 * w + 2 * h);
   }
   return make_pair(area, perimeter);
// returns the left side of polygon u after cutting it by
     ray a->b
Polygon cutPolygon(Polygon u, Point a, Point b) {
   using Linear::lineLineIntersection;
   using Linear::onSegment;
   Polygon ret;
   int n = u.size();
   for (int i = 0; i < n; i++) {</pre>
       Point c = u[i], d = u[(i + 1) \% n];
       if (dcmp(cross(b - a, c - a)) >= 0) ret.push_back
       if (dcmp(cross(b - a, d - c)) != 0) {
           Point t;
           lineLineIntersection(a, b - a, c, d - c, t);
           if (onSegment(t, Segment(c, d))) ret.push_back
               (t);
       }
   return ret;
// returns true if point p is in or on triangle abc
bool pointInTriangle(Point a, Point b, Point c, Point p)
   return dcmp(cross(b - a, p - a)) >= 0 && dcmp(cross(c
         -b, p-b)) >= 0 &&
                dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no three collinear points
// returns inside = -1, on = 0, outside = 1
int pointInConvexPolygon(const Polygon &pt, Point p) {
```

```
int n = pt.size();
   assert(n >= 3);
   int lo = 1, hi = n - 1;
   while (hi - lo > 1) {
       int mid = (lo + hi) / 2;
       if (dcmp(cross(pt[mid] - pt[0], p - pt[0])) > 0)
           lo = mid;
       else
           hi = mid;
   }
   bool in = pointInTriangle(pt[0], pt[lo], pt[hi], p);
   if (!in) return 1;
   if (dcmp(cross(pt[lo] - pt[lo - 1], p - pt[lo - 1]))
        == 0) return 0;
   if (dcmp(cross(pt[hi] - pt[lo], p - pt[lo])) == 0)
       return 0:
   if (dcmp(cross(pt[hi] - pt[(hi + 1) % n], p - pt[(hi
       + 1) % n])) == 0)
       return 0;
   return -1;
// Extreme Point for a direction is the farthest point in // #1 If a segment is collinear with the line, only that
     that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly, Point u) {
   int n = (int)poly.size();
   int a = 0, b = n;
   while (b - a > 1) {
       int c = (a + b) / 2;
       if (dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >=
              dcmp(dot(poly[c] - poly[(c - 1 + n) % n],
                  u)) >= 0) {
           return c;
       }
       bool a_up = dcmp(dot(poly[(a + 1) % n] - poly[a],
            u)) >= 0:
       bool c_up = dcmp(dot(poly[(c + 1) % n] - poly[c],
            u)) >= 0:
       bool a_above_c = dcmp(dot(poly[a] - poly[c], u))
           > 0;
       if (a_up && !c_up)
           b = c;
       else if (!a_up && c_up)
           a = c;
```

```
else if (a_up && c_up) {
          if (a_above_c)
              b = c;
          else
              a = c;
       } else {
          if (!a_above_c)
              b = c;
          else
              a = c;
       }
   }
   if (dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 &&
          dcmp(dot(poly[a] - poly[(a - 1 + n) % n], u))
       return a:
   return b % n;
// For a convex polygon p and a line 1, returns a list of
     segments
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i], p[(i + 1) modulo
// #2 Else if 1 goes through i'th point, the i'th segment
     is added
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const Polygon &p,
    Line 1) {
   assert((int)p.size() >= 3);
   assert(1.a != 1.b);
   int n = p.size();
   vector<int> ret;
   Point v = 1.b - 1.a:
   int lf = extremePoint(p, rotate90(v));
   int rt = extremePoint(p, rotate90(v) * Ti(-1));
   int olf = orient(l.a, l.b, p[lf]);
   int ort = orient(l.a, l.b, p[rt]);
   if (!olf || !ort) {
       int idx = (!olf ? lf : rt);
       if (orient(1.a, 1.b, p[(idx - 1 + n) \% n]) == 0)
          ret.push_back((idx - 1 + n) \% n);
          ret.push_back(idx);
       return ret;
```

```
if (olf == ort) return ret;
   for (int i = 0; i < 2; ++i) {</pre>
       int lo = i ? rt : lf;
       int hi = i ? lf : rt;
       int olo = i ? ort : olf;
       while (true) {
           int gap = (hi - lo + n) \% n;
           if (gap < 2) break;</pre>
           int mid = (lo + gap / 2) % n;
           int omid = orient(l.a, l.b, p[mid]);
           if (!omid) {
              lo = mid;
              break;
           if (omid == olo)
              lo = mid:
           else
              hi = mid;
       }
       ret.push_back(lo);
   }
   return ret;
// Calculate [ACW, CW] tangent pair from an external
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int dir) {
    return orient(Q, u, v) != -dir;
}
Point better(Point u, Point v, Point Q, int dir) {
    return orient(Q, u, v) == dir ? u : v;
Point pointPolyTangent(const Polygon &pt, Point Q, int
    dir, int lo, int hi) {
    while (hi - lo > 1) {
       int mid = (lo + hi) / 2:
       bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
       bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
       if (pvs && nxt) return pt[mid];
       if (!(pvs || nxt)) {
           Point p1 = pointPolyTangent(pt, Q, dir, mid +
               1, hi);
           Point p2 = pointPolyTangent(pt, Q, dir, lo,
               mid - 1);
           return better(p1, p2, Q, dir);
```

```
}
       if (!pvs) {
           if (orient(Q, pt[mid], pt[lo]) == dir)
              hi = mid - 1;
           else if (better(pt[lo], pt[hi], Q, dir) == pt[ Tf signedPolygonArea(const Polygon &p) {
              hi = mid - 1;
           else
              lo = mid + 1;
       }
       if (!nxt) {
           if (orient(Q, pt[mid], pt[lo]) == dir)
              lo = mid + 1;
           else if (better(pt[lo], pt[hi], Q, dir) == pt[
              hi = mid - 1;
           else
              lo = mid + 1:
       }
   }
   Point ret = pt[lo];
   for (int i = lo + 1; i <= hi; i++) ret = better(ret,</pre>
        pt[i], Q, dir);
   return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(const Polygon &pt,
    Point Q) {
   int n = pt.size();
   Point acw_tan = pointPolyTangent(pt, Q, ACW, 0, n -
   Point cw_tan = pointPolyTangent(pt, Q, CW, 0, n - 1);
   return make_pair(acw_tan, cw_tan);
7.5 Polygon
typedef vector<Point> Polygon;
```

```
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &poly) {
   Polygon ret;
   int n = poly.size();
   for (int i = 0; i < n; i++) {</pre>
       Point a = poly[i];
       Point b = poly[(i + 1) \% n];
       Point c = poly[(i + 2) \% n];
       if (dcmp(cross(b - a, c - b)) != 0 && (ret.empty
            () || b != ret.back()))
```

```
ret.push_back(b);
   }
   return ret;
// returns the signed area of polygon p of n vertices
   Tf ret = 0:
   for (int i = 0; i < (int)p.size() - 1; i++)</pre>
       ret += cross(p[i] - p[0], p[i + 1] - p[0]);
   return ret / 2;
// given a polygon p of n vertices, generates the convex
    hull in in CCW
// Tested on https://acm.timus.ru/problem.aspx?space=1&
// Caution: when all points are colinear AND
    removeRedundant == false
// output will be contain duplicate points (from upper
    hull) at back
Polygon convexHull(Polygon p, bool removeRedundant) {
   int check = removeRedundant ? 0 : -1;
   sort(p.begin(), p.end());
   p.erase(unique(p.begin(), p.end()), p.end());
   int n = p.size();
   Polygon ch(n + n);
   int m = 0; // preparing lower hull
   for (int i = 0; i < n; i++) {</pre>
       while (m > 1 &&
                    dcmp(cross(ch[m-1]-ch[m-2], p[i
                        ] - ch[m - 1])) <= check)
          m--;
       ch[m++] = p[i];
   int k = m; // preparing upper hull
   for (int i = n - 2; i >= 0; i--) {
       while (m > k &&
                    dcmp(cross(ch[m-1]-ch[m-2], p[i
                        ] - ch[m - 2])) \leftarrow check)
          m--;
       ch[m++] = p[i];
   if (n > 1) m--;
   ch.resize(m);
   return ch;
// returns inside = -1, on = 0, outside = 1
int pointInPolygon(const Polygon &p, Point o) {
   using Linear::onSegment;
   int wn = 0, n = p.size();
```

```
for (int i = 0; i < n; i++) {</pre>
       int j = (i + 1) \% n;
       if (onSegment(o, Segment(p[i], p[j])) || o == p[i
           1) return 0:
       int k = dcmp(cross(p[j] - p[i], o - p[i]));
       int d1 = dcmp(p[i].y - o.y);
       int d2 = dcmp(p[j].y - o.y);
       if (k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
       if (k < 0 \&\& d2 \le 0 \&\& d1 > 0) wn--;
   }
   return wn ? -1 : 1;
// Given a simple polygon p, and a line 1, returns (x, y)
// x = longest segment of 1 in p, y = total length of 1
pair<Tf, Tf> linePolygonIntersection(Line 1, const
    Polygon &p) {
   using Linear::lineLineIntersection;
   int n = p.size();
   vector<pair<Tf, int>> ev;
   for (int i = 0; i < n; ++i) {</pre>
       Point a = p[i], b = p[(i + 1) \% n], z = p[(i - 1)
           + n) % n];
       int ora = orient(l.a, l.b, a), orb = orient(l.a,
           1.b, b),
              orz = orient(l.a, l.b, z);
       if (!ora) {
           Tf d = dot(a - 1.a, 1.b - 1.a);
           if (orz && orb) {
              if (orz != orb) ev.emplace_back(d, 0);
              // else // Point Touch
           } else if (orz)
              ev.emplace_back(d, orz);
           else if (orb)
              ev.emplace_back(d, orb);
       } else if (ora == -orb) {
           Point ins;
           lineLineIntersection(1, Line(a, b), ins);
           ev.emplace_back(dot(ins - 1.a, 1.b - 1.a), 0);
       }
   sort(ev.begin(), ev.end());
   Tf ans = 0, len = 0, last = 0, tot = 0;
   bool active = false;
   int sign = 0;
   for (auto &qq : ev) {
       int tp = qq.second;
       Tf d = qq.first; /// current Segment is (last, d)
       if (sign) {
                       /// On Border
```

```
len += d - last:
           tot += d - last;
           ans = max(ans, len);
           if (tp != sign) active = !active;
           sign = 0;
       } else {
           if (active) { /// Strictly Inside
              len += d - last;
              tot += d - last;
              ans = max(ans, len);
           }
           if (tp == 0)
              active = !active;
           else
              sign = tp;
       }
       last = d:
       if (!active) len = 0;
   ans /= length(l.b - l.a);
   tot /= length(l.b - l.a);
   return {ans, tot};
7.6 Half Plane
```

```
using Linear::lineLineIntersection;
struct DirLine {
   Point p, v;
   Tf ang;
   DirLine() {}
   /// Directed line containing point P in the direction
   DirLine(Point p, Point v) : p(p), v(v) { ang = atan2(
       v.y, v.x); }
   bool operator<(const DirLine& u) const { return ang <</pre>
        u.ang; }
// returns true if point p is on the ccw-left side of ray
bool onLeft(DirLine 1, Point p) { return dcmp(cross(1.v,
    p - 1.p)) >= 0; }
// Given a set of directed lines returns a polygon such
// the polygon is the intersection by halfplanes created
// left side of the directed lines. MAY CONTAIN DUPLICATE
     POINTS
int halfPlaneIntersection(vector<DirLine>& li, Polygon&
    poly) {
```

```
int n = li.size();
sort(li.begin(), li.end());
int first, last;
Point* p = new Point[n];
DirLine* q = new DirLine[n];
q[first = last = 0] = li[0];
for (int i = 1; i < n; i++) {</pre>
   while (first < last && !onLeft(li[i], p[last -</pre>
        1])) last--;
   while (first < last && !onLeft(li[i], p[first]))</pre>
        first++;
   q[++last] = li[i];
   if (dcmp(cross(q[last].v, q[last - 1].v)) == 0) {
       last--:
       if (onLeft(q[last], li[i].p)) q[last] = li[i];
   if (first < last)</pre>
       lineLineIntersection(q[last - 1].p, q[last -
            1].v, q[last].p, q[last].v,
                                                p[last -
                                                     1])
}
while (first < last && !onLeft(q[first], p[last - 1])</pre>
    ) last--:
if (last - first <= 1) {</pre>
   delete[] p;
   delete[] q;
   poly.clear();
   return 0;
lineLineIntersection(q[last].p, q[last].v, q[first].p
    , q[first].v, p[last]);
int m = 0:
poly.resize(last - first + 1);
for (int i = first; i <= last; i++) poly[m++] = p[i];</pre>
delete[] p;
delete[] q;
return m;
```

# Equations and Formulas

## Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1 = 1 \text{ and } C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

The number of ways to completely parenthesize n+1 factors. The number of triangulations of a convex polygon with  $n+2|S^d(n,k)$ , to be the number of ways to partition the integers sides (i.e. the number of partitions of polygon into disjoint [1,2,.,n] into k nonempty subsets such that all elements in  $\sum [\gcd(i,n)=k] = \phi\left(\frac{n}{i}\right)$ triangles by using the diagonals).

form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (ver- $S^d(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$ tices are not numbered). A rooted binary tree is full if every 8.4 Other Combinatorial Identities vertex has either two children or no children.

Number of permutations of 1, n that avoid the pattern 123  $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$ (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321

# 8.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) =$$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$[n - k] = \sum_{0 \le i_1 < i_2 < i_k < n} i_1 i_2 \dots i_k.$$

# 8.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

$$S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$$
, where  $S(0,0) = 1$ ,  $S(n,0) = S(0,n) = 0$   $S(n,2) = 2^{n-1} - 1$   $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k \text{ such that if } m \text{ is any integer, then } \gcd(a+m\cdot b,b) = \gcd(a,b)$  The gcd is a multiplicative function in the follow

ber of ways to partition a set of n objects into k subsets, with  $\gcd(a_1,b) \cdot \gcd(a_2,b)$ .

each subset containing at least r elements. It is denoted by  $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c))$ .  $S_r(n,k)$  and obeys the recurrence relation.  $S_r(n+1,k) = |\operatorname{lcm}(a,\operatorname{gcd}(b,c))| = \operatorname{gcd}(\operatorname{lcm}(a,b),\operatorname{lcm}(a,c))$ .  $\left| kS_r(n,k) + \binom{n}{r-1} S_r(n-r+1,k-1) \right|$ 

Denote the n objects to partition by the integers 1, 2, ..., n. De  $gcd(a, b) = \sum_{a} \phi(k)$ fine the reduced Stirling numbers of the second kind, denoted each subset have pairwise distance at least d. That is, for  $\overline{i=1}$ The number of ways to connect the 2n points on a circle to any integers i and j in a given subset, it is required that  $|i-j| \geq d$ . It has been shown that these numbers satisfy,

$$\begin{pmatrix} n \\ k \end{pmatrix} = \frac{n}{k} \binom{n-1}{k-1} 
\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n} = \binom{n+k+1}{k} 
n, r \in N, n > r, \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1} 
If  $P(n) = \sum_{k=0}^{n} \binom{n}{k} \cdot Q(k)$ , then,$$

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If 
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

# 8.5 Different Math Formulas

Picks Theorem: A = i + b/2 - 1

**Deragements:**  $d(i) = (i-1) \times (d(i-1) + d(i-2))$ 

$$\frac{n}{ab}$$
 -  $\left\{\frac{b'n}{a}\right\}$  -  $\left\{\frac{a'n}{b}\right\}$  +

The gcd is a multiplicative function in the following sense: An r-associated Stirling number of the second kind is the num- if  $a_1$  and  $a_2$  are relatively prime, then  $gcd(a_1 \cdot a_2, b) =$ 

For non-negative integers a and b, where a and b are not both zero,  $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$ k|a and k|b $\sum_{k=1}^{n} \gcd(k, n) = \sum_{d \mid n} d \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{l=1}^{n} x^{d} \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{\text{all}} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{\text{all}} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d \mid n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d \mid n} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$  $\left| \sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2} \right|$  $\left| \sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2} \right|$  $\sum \sum i \cdot j[\gcd(i,j) = 1] = \sum \phi(i)i^2$  $F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left( \frac{\left(1 + \lfloor \frac{n}{l} \rfloor\right) \left(\lfloor \frac{n}{l} \rfloor\right)}{2} \right)^{2} \sum_{i} \mu(d) l d$ 

# 8.7 Geometry

Cone:  $V = \frac{1}{3}\pi r^2 h$ ,  $A = \pi r(r + \sqrt{h^2 + r^2})$ 

**Pyramid:**  $V = \frac{1}{3} \times \text{base} \times \text{height}, A = \text{base area} + \frac{1}{2} \times \text{base}$ perimeter × slant height

**Triangular Prism:**  $V = \frac{1}{2} \times \text{base} \times \text{height} \times \text{depth}, A =$ base × height +  $3 \times (\frac{1}{2} \times \text{side} \times \text{perimeter})$ 

**Torus:**  $V = 2\pi^2 R r^2$ ,  $A = 4\pi^2 R r$ 

Ellipsoid: 
$$V = \frac{4}{3}\pi abc$$
,  $A = 4\pi \left(\frac{(ab)^{1.6} + (bc)^{1.6} + (ca)^{1.6}}{3}\right)^{1/1.6}$