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#### Stress-tester

```
#!/bin/bash
# Call as stresstester ITERATIONS [--count]
g++ gen.cpp -o gen
g++ sol.cpp -o sol
g++ brute.cpp -o brute

for i in $(seq 1 "$1") ; do
    echo "Attempt $i/$1"
    ./gen > in.txt
```

```
./sol < in.txt > out1.txt
./brute < in.txt > out2.txt
diff -y out1.txt out2.txt > diff.txt
if [ $? -ne 0 ] ; then
        echo "Differing Testcase Found:"; cat in.txt
        echo -e "\nOutputs:"; cat diff.txt
        break
fi
done
```

# All Macros

/\*--- DEBUG TEMPLATE STARTS HERE ---\*/

```
void show(int x) {cerr << x;}</pre>
void show(long long x) {cerr << x;}</pre>
void show(double x) {cerr << x;}</pre>
void show(char x) {cerr << '\'' << x << '\'';}</pre>
void show(const string &x) {cerr << '\"' << x << '\"';}</pre>
void show(bool x) {cerr << (x ? "true" : "false");}</pre>
template<typename T, typename V>
void show(pair<T, V> x) { cerr << '\{'; show(x.first);</pre>
    cerr << ", "; show(x.second); cerr << '}'; }</pre>
template<typename T>
void show(T x) {int f = 0; cerr << "{"; for (auto &i: x)</pre>
    cerr << (f++ ? ", " : ""), show(i); cerr << "}";}</pre>
void debug_out(string s) {
  s.clear();
  cerr << s << '\n';
template <typename T, typename... V>
void debug_out(string s, T t, V... v) {
  s.erase(remove(s.begin(), s.end(), ''), s.end());
  cerr << "
                  "; // 8 spaces
  cerr << s.substr(0, s.find(','));</pre>
  s = s.substr(s.find(',') + 1);
  cerr << " = ":
  show(t);
  cerr << endl;
  if(sizeof...(v)) debug_out(s, v...);
#define dbg(x...) cerr << "LINE: " << _LINE_ << endl;</pre>
    debug_out(#x, x); cerr << endl;</pre>
/*--- DEBUG TEMPLATE ENDS HERE ---*/
//#pragma GCC optimize("Ofast")
//#pragma GCC optimization ("03")
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("unroll-loops")
//#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm
    ,mmx,avx,tune=native")
```

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
 //find_by_order(k) --> returns iterator to the kth
      largest element counting from 0
 //order_of_key(val) --> returns the number of items in
      a set that are strictly smaller than our item
template <typename DT>
using ordered_set = tree <DT, null_type, less<DT>,
    rb_tree_tag,tree_order_statistics_node_update>;
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
    ().count());
#ifdef LOCAL
#include "dbg.h"
#else
#define dbg(x...)
#endif
int main() {
 cin.tie(0) -> sync_with_stdio(0);
```

#### 2 Data Structure

#### 2.1 Sparse Table

# 2.2 BIT template <typename

```
template <typename T> class BIT
 public:
   int n; vector<T> tree;
   BIT(int size) // 1-indexed
     n = size; tree.assign(n+1, 0);
   BIT(const vector<T> &a) : BIT(a.size())
     for(int i = 1; i <= n; i++) update(i, a[i-1]);</pre>
   T query(int i)
     T ans = 0;
     for( ; i >= 1; i-= (i & -i)) ans+= tree[i];
     return ans;
   }
   T query(int 1, int r)
     return query(r) - query(1-1);
   void update(int i, T delta)
     for( ; i <= n; i+= (i & -i)) tree[i]+= delta;</pre>
};
```

# 2.3 Lazy SegmentTree

```
build(a, b, mid, 2*v);
 build(a, mid+1, e,2*v+1);
 tree[v] = tree[2*v] + tree[2*v+1];
11 query(int 1, int r, int b, int e, int v=1, ll carry =
 if(b > r || e < 1) return 0;
 if (b \ge 1 \&\& e \le r) return tree [v] + carry * (e-b+1);
 int mid = (b+e)/2;
 11 lseg = query(1, r, b, mid, 2*v, carry+lazy[v]);
 ll rseg = query(1, r, mid+1, e, 2*v+1, carry+lazy[v]);
 return lseg + rseg;
void update(int 1, int r, 11 val, int b, int e, int v =
    1)
 if(b > r || e < 1) return;</pre>
 if(b >= 1 \&\& e <= r)
   tree[v]+= (e-b+1)*val;
   lazy[v]+= val;
   return;
 }
 int mid = (b+e)/2;
 update(1, r, val, b, mid, 2*v);
 update(1, r, val, mid+1, e, 2*v+1);
 tree[v] = tree[2*v] + tree[2*v+1] + (e-b+1)*lazy[v];
```

### 2.4 Generic SegmentTree

```
template<typename ST, typename LZ>
class SegmentTree {
private:
   int n;
   ST *tree, identity;
   ST (*merge) (ST, ST);

   LZ *lazy, unmark;
   void (*mergeLazy)(int, int, LZ&, LZ);
   void (*applyLazy)(int, int, ST&, LZ);

void build(vector<ST> &arr, int lo, int hi, int cur=1) {
   if(lo == hi)
   {
```

```
tree[cur] = arr[lo-1]:
   return:
 int mid = (hi+lo)/2, left = 2*cur, right = 2*cur+1;
 build(arr, lo, mid, left);
 build(arr, mid+1, hi, right);
  tree[cur] = merge(tree[left], tree[right]);
void propagate(int lo, int hi, int cur)
  applyLazy(lo, hi, tree[cur], lazy[cur]);
 if(lo < hi)
   int mid = (lo+hi)/2, left = 2*cur, right = 2*cur+1;
   mergeLazy(lo, mid, lazy[left], lazy[cur]);
   mergeLazy(mid+1, hi, lazy[right], lazy[cur]);
 lazv[cur] = unmark:
}
void update(int from, int upto, LZ delta, int lo, int
    hi, int cur=1)
 if(lo>hi) return;
 propagate(lo, hi, cur);
 if(from > hi or upto < lo) return;</pre>
 if(from<= lo and upto >= hi)
 {
   mergeLazy(lo, hi, lazy[cur], delta);
   propagate(lo, hi, cur);
   return;
  int mid = (lo+hi)/2, left = 2*cur, right = 2*cur+1;
  update(from, upto, delta, lo, mid, left);
  update(from, upto, delta, mid+1, hi, right);
 tree[cur] = merge(tree[left], tree[right]);
ST query(int from, int upto, int lo, int hi, int cur=1)
 if(lo>hi) return identity;
 propagate(lo, hi, cur);
  if(from > hi or upto < lo) return identity;</pre>
  if(from<= lo and upto >= hi) return tree[cur];
  int mid = (lo+hi)/2, left = 2*cur, right = 2*cur+1;
  ST lseg = query(from, upto, lo, mid, left);
  ST rseg = query(from, upto, mid+1, hi, right);
```

```
return merge(lseg, rseg);
public:
  SegmentTree(
   vector<ST> arr, ST (*merge) (ST, ST), ST identity,
   void (*mergeLazy)(int, int, LZ&, LZ),
   void (*applyLazy)(int, int, ST&, LZ), LZ unmark
    n(arr.size()), merge(merge), identity(identity),
   mergeLazy(mergeLazy), applyLazy(applyLazy), unmark(
        unmark)
    tree = new ST[n*4];
   lazv = new LZ[n*4];
   build(arr, 1, n);
   fill(lazy, lazy+n*4, unmark);
  void update(int from, int upto, LZ delta)
   update(from, upto, delta, 1, n);
  ST query(int from, int upto)
   return query(from, upto, 1, n);
  ~SegmentTree()
   delete[] tree;
   delete[] lazy;
 }
};
11 add(l1 1, l1 r) { return 1+r;}
void mergeAdd(int lo, int hi, ll &cur, ll pending) { cur
    += pending:}
void applyAdd(int lo, int hi, ll &cur, ll pending) { cur
    += pending*(hi-lo+1);}
void solve(int tcase)
  vector<ll> a(n):
  SegmentTree<11, 11> st(a, add, 0, mergeAdd, applyAdd,
      0);
```

```
2.5 MO
struct node {
 LL 1, r, idx;
bool cmp(const node &x, const node &y) {
 return x.r < y.r;</pre>
void add(LL x) {
 if(mp[x] % 2) curr++;
 mp[x]++;
void diminish(LL x) {
 if(mp[x] \% 2 == 0) curr--;
 mp[x]--;
void solve()
 BLOCK_SIZE = sqrt(n) + 1;
 rep(i, 0, q-1) {
   LL x, y; cin >> x >> y;
   x--; y--;
   query[x / BLOCK_SIZE].pb({x, y, i});
   m = max(m, x / BLOCK_SIZE);
 rep(i, 0, m) sort(all(query[i]), cmp);
 LL mo_left = 0, mo_right = -1;
 rep(i, 0, m) {
   for(auto [left, right, id] : query[i]) {
     while(mo_right < right) add(v[++mo_right]);</pre>
     while(mo_right > right) diminish(v[mo_right--]);
     while(mo_left < left) diminish(v[mo_left++]);</pre>
     while(mo left > left) add(v[--mo left]);
     answer[id] = curr:
 rep(i, 0, q-1) cout << answer[i] << endl;</pre>
```

## 2.6 MergeSort Tree

```
vector<LL> tree[5*MAXN];
LL A[N];
void build_tree(LL now , LL curLeft, LL curRight) {
        if(curLeft == curRight) {
            tree[now].push_back(A[curLeft]);
            return;
        }
        LL mid = (curLeft + curRight) / 2;
        build_tree(2 * now, curLeft, mid);
        build_tree(2 * now + 1, mid + 1 , curRight);
```

# 3 Graph

### 3.1 Graph Template

```
struct edge {
 int u, v;
 edge(int u = 0, int v = 0) : u(u), v(v) {}
 int to(int node) { return u ^ v ^ node; }
struct graph {
 int n;
 vector<vector<int>> adj;
 vector<edge> edges;
 graph(int n = 0) : n(n), adj(n) {}
 void addEdge(int u, int v, bool dir = 1) {
   adj[u].push_back(edges.size());
   if (dir) adj[v].push_back(edges.size());
   edges.emplace_back(u, v);
 }
 int addNode() {
   adj.emplace_back();
   return n++;
 edge &operator()(int idx) { return edges[idx]; }
 vector<int> &operator[](int u) { return adj[u]; }
```

### 3.2 Lifting, LCA, HLD

```
using Tree = vector<vector<int>>;
int anc[B][N], sz[N], lvl[N], st[N], en[N], nxt[N], t =
    0;

void initLifting(int n) {
  for (int b = 1; b < B; b++) {
    for (int i = 0; i < n; i++) {
        anc[b][i] = anc[b - 1][anc[b - 1][i]];
    }
}</pre>
```

```
}
int kthAncestor(int u, int k) {
 for (int b = 0; b < B; b++) {</pre>
   if (k >> b & 1) u = anc[b][u];
 return u;
int lca(int u, int v) {
 if (lvl[u] > lvl[v]) swap(u, v);
 v = kthAncestor(v, lvl[v] - lvl[u]);
 if (u == v) return u;
 for (int b = B - 1; b >= 0; b--) {
   if (anc[b][u] != anc[b][v]) u = anc[b][u], v = anc[b
       ][v]:
 }
 return anc[0][u];
int dis(int u, int v) {
 int g = lca(u, v);
 return lvl[u] + lvl[v] - 2 * lvl[g];
bool isAncestor(int u, int v) { return st[v] <= st[u] and
     en[u] <= en[v]; }
void tour(int u, int p, Tree &T) {
 st[u] = t++:
 int idx = 0;
 for (int v : T[u]) {
   if (v == p) continue;
   nxt[v] = (idx++ ? v : nxt[u]); // only for hld
   anc[0][v] = u, lvl[v] = lvl[u] + 1;
   tour(v, u, T);
 en[u] = t; // [st, en] contains subtree range
void hld(int u, int p, Tree &T) {
 sz[u] = 1;
 int eld = 0, mx = 0, idx = 0;
 for (int i = 0; i < T[u].size(); i++) {</pre>
   int v = T[u][i];
   if (v == p) continue;
   hld(v, u, T);
```

```
if (sz[v] > mx) mx = sz[v], eld = i;
   sz[u] += sz[v];
  swap(T[u][0], T[u][eld]);
LL climbQuery(int u, int g) {
 LL ans = -INF:
 while (1) {
   int _u = nxt[u];
   if (isAncestor(g, _u)) _u = g;
   ans = max(ans, rmq ::query(st[_u], st[u]));
   if (_u == g) break;
   u = anc[0][u];
  return ans;
LL pathQuery(int u, int v) {
 int g = lca(u, v);
 return max(climbQuery(u, g), climbQuery(v, g));
void init(int u, Tree &T) {
 int n = T.size();
 anc[0][u] = nxt[u] = u;
 lvl[u] = 0;
 hld(u, u, T);
 tour(u, u, T);
 initLifting(n);
3.3 SCC
vector<int> order, comp, idx;
vector<bool> vis;
vector<vector<int>> comps;
Graph dag;
void dfs1(int u, Graph &G, string s = "") {
 vis[u] = 1;
 for (int e : G[u]) {
   int v = G(e).to(u);
   if (!vis[v]) dfs1(v, G, s);
 order.push_back(u);
void dfs2(int u, Graph &R) {
```

comp.push\_back(u);

```
idx[u] = comps.size();
 for (int e : R[u]) {
   int v = R(e).to(u):
   if (idx[v] == -1) dfs2(v, R);
void init(Graph &G) {
 int n = G.n;
 vis.assign(n, 0);
 idx.assign(n, -1);
 for (int i = 0; i < n; i++) {
   if (!vis[i]) dfs1(i, G);
 reverse(order.begin(), order.end());
 Graph R(n):
 for (auto &e : G.edges) R.addEdge(e.v, e.u, 0);
 for (int u : order) {
   if (idx[u] != -1) continue;
   comp.clear();
   dfs2(u, R);
   comps.push_back(comp);
Graph &createDAG(Graph &G) {
 int sz = comps.size();
 dag = Graph(sz);
 vector<bool> taken(sz);
 vector<int> cur;
 for (int i = 0; i < sz; i++) {</pre>
   cur.clear():
   taken[i] = 1;
   for (int u : comps[i]) {
     for (int e : G[u]) {
       int v = G(e).to(u);
       int j = idx[v];
       if (taken[j]) continue; // rejects multi-edge
       dag.addEdge(i, j, 0);
       taken[j] = 1;
       cur.push_back(j);
   }
   for (int j : cur) taken[j] = 0;
```

```
}
return dag;
}
```

# 3.4 Centroid Decompose

```
namespace ct {
int par[N], cnt[N], cntp[N];
LL sum[N], sump[N];
void activate(int u) {
 int v = u, u = u;
 ans += sum[u];
  cnt[u]++;
  while (par[u] != -1) {
   u = par[u];
   LL d = ta :: dis(_u, u);
   ans += sum[u] - sump[v];
   ans += d * (cnt[u] - cntp[v]);
   sum[u] += d;
   cnt[u]++;
   sump[v] += d;
   cntp[v]++;
   v = u;
namespace ctrd {
int sz[N];
bool blk[N];
int szCalc(Tree &T, int u, int p = -1) {
 sz[u] = 1:
 for (int v : T[u]) {
   if (v == p or blk[v]) continue;
   sz[u] += szCalc(T, v, u);
 return sz[u];
int getCentroid(Tree &T, int u, int s, int p = -1) {
 for (int v : T[u]) {
   if (v == p or blk[v]) continue;
   if (sz[v] * 2 >= s) return getCentroid(T, v, s, u);
 return u;
void decompose(Tree &T, int u, int p = -1) {
```

```
szCalc(T, u);
u = getCentroid(T, u, sz[u]);
ct ::par[u] = p;

blk[u] = 1;
for (int v : T[u]) {
   if (!blk[v]) decompose(T, v, u);
}
}
```

# 3.5 Euler Tour on Edge

```
// for simplicity, G[idx] contains the adjacency list of
// while G(e) is a reference to the e-th edge.
const int N = 2e5 + 5;
int in[N], out[N], fwd[N], bck[N];
int t = 0;
void dfs(graph &G, int node, int par) {
 out[node] = t;
 for (int e : G[node]) {
   int v = G(e).to(node);
   if (v == par) continue;
   fwd[e] = t++;
   dfs(G, v, node):
   bck[e] = t++;
 in[node] = t - 1;
void init(graph &G, int node) {
 t = 0:
 dfs(G, node, node);
```

#### 3.6 Virtual Tree

```
namespace lca1 {
int st[N], lvl[N];
int tbl[B][2 * N];
int t = 0;

void dfs(int u, int p, Tree &T) {
  st[u] = t;
  tbl[0][t++] = u;
  for(int v: T[u]) {
    if(v == p) continue;
    lvl[v] = lvl[u] + 1;
    dfs(v, u, T);
    tbl[0][t++] = u;
}
```

```
int low(int u, int v) {
 return make_pair(lvl[u], u) < make_pair(lvl[v], v) ? u</pre>
void makeTable(int n) {
 int. m = 2 * n - 1:
 for(int b = 1; b < B; b++) {
   for(int i = 0: i < m: i++) {</pre>
     tbl[b][i] = low(tbl[b - 1][i], tbl[b - 1][i + (1 <<
          b - 1)]);
   }
 }
int lca(int u, int v) {
 int 1 = st[u], r = st[v];
 if(1 > r) swap(1, r);
 int k = _{-}lg(r - 1 + 1);
 return low(tbl[k][1], tbl[k][r - (1 << k) + 1]);</pre>
void init(int root, Tree &T) {
 lvl[root] = 0;
 t = 0:
 dfs(root, root, T);
 makeTable(T.size());
namespace vt {
int st[N], en[N], t;
vector <int> adj[N];
void dfs(int u, int p, Tree &T) {
 st[u] = t++:
 for(int v: T[u]) if(v != p) dfs(v, u, T);
 en[u] = t++;
bool comp(int u, int v) {
 return st[u] < st[v];</pre>
bool isAncestor(int u, int p) {
 return st[p] <= st[u] and en[u] <= en[p];</pre>
void construct(vector <int> &nodes) {
 sort(nodes.begin(), nodes.end(), comp);
 int n = nodes.size();
 for(int i = 0; i + 1 < n; i++) {
   nodes.push_back(lca1 :: lca(nodes[i], nodes[i + 1]));
```

```
sort(nodes.begin(), nodes.end(), comp);
  nodes.erase(unique(nodes.begin(), nodes.end()), nodes.
      end());
  n = nodes.size():
  stack <int> s;
  s.push(nodes[0]);
  for(int i = 1; i < n; i++) {</pre>
    int u = nodes[i];
    while(not isAncestor(u, s.top())) s.pop();
    adj[s.top()].push_back(u);
    s.push(u);
  }
}
void clear(vector <int> &nodes) {
 for(int u: nodes) {
     adi[u].clear();
 }
}
void init(int root, Tree &T) {
 lca1 :: init(root, T);
  t = 0:
  dfs(root, root, T);
```

#### 3.7 Dinic Max Flow

```
/// flow with demand(lower bound) only for DAG
// create new src and sink
// add_edge(new src, u, sum(in_demand[u]))
// add edge(u, new sink, sum(out demand[u]))
// add_edge(old sink, old src, inf)
// if (sum of lower bound == flow) then demand satisfied
// flow in every edge i = demand[i] + e.flow
using Ti = long long;
const Ti INF = 1LL << 60;</pre>
struct edge {
 int v, u;
 Ti cap, flow = 0;
  edge(int v, int u, Ti cap) : v(v), u(u), cap(cap) {}
const int N = 1e5 + 50;
vector<edge> edges;
vector<int> adj[N];
int m = 0, n;
int level[N], ptr[N];
queue<int> q;
bool bfs(int s, int t) {
 for (q.push(s), level[s] = 0; !q.empty(); q.pop()) {
```

```
for (int id : adj[q.front()]) {
     auto &ed = edges[id];
     if (ed.cap - ed.flow > 0 and level[ed.u] == -1)
       level[ed.u] = level[ed.v] + 1, q.push(ed.u);
   }
 }
 return level[t] != -1;
Ti dfs(int v, Ti pushed, int t) {
 if (pushed == 0) return 0;
 if (v == t) return pushed;
 for (int &cid = ptr[v]; cid < adj[v].size(); cid++) {</pre>
   int id = adj[v][cid];
   auto &ed = edges[id];
   if (level[v] + 1 != level[ed.u] || ed.cap - ed.flow <</pre>
   Ti tr = dfs(ed.u, min(pushed, ed.cap - ed.flow), t);
   if (tr == 0) continue;
   ed.flow += tr;
   edges[id ^ 1].flow -= tr;
   return tr;
 return 0;
void init(int nodes) {
 m = 0, n = nodes;
 for (int i = 0; i < n; i++) level[i] = -1, ptr[i] = 0,
      adj[i].clear();
void addEdge(int v, int u, Ti cap) {
 edges.emplace back(v, u, cap), adj[v].push back(m++);
  edges.emplace_back(u, v, 0), adj[u].push_back(m++);
Ti maxFlow(int s, int t) {
 Ti f = 0:
 for (auto &ed : edges) ed.flow = 0;
 for (; bfs(s, t); memset(level, -1, n * 4)) {
   for (memset(ptr, 0, n * 4); Ti pushed = dfs(s, INF, t
        ); f += pushed)
 }
 return f;
3.8 Min Cost Max Flow
```

```
mt19937 rnd(chrono::steady_clock::now().time_since_epoch
          ().count());
const LL inf = 1e9;
struct edge {
   int v, rev;
```

```
LL cap, cost, flow;
 edge() {}
 edge(int v, int rev, LL cap, LL cost)
     : v(v), rev(rev), cap(cap), cost(cost), flow(0) {}
struct mcmf {
 int src. sink. n:
 vector<int> par, idx, Q;
 vector<bool> inq;
 vector<LL> dis;
 vector<vector<edge>> g;
 mcmf() {}
 mcmf(int src, int sink, int n)
     : src(src),
       sink(sink),
       n(n),
       par(n),
       idx(n),
       inq(n),
       dis(n),
       g(n),
       Q(10000005) {} // use Q(n) if not using random
 void add_edge(int u, int v, LL cap, LL cost, bool
     directed = true) {
   edge _u = edge(v, g[v].size(), cap, cost);
   edge _v = edge(u, g[u].size(), 0, -cost);
   g[u].pb(_u);
   g[v].pb(v);
   if (!directed) add_edge(v, u, cap, cost, true);
 bool spfa() {
   for (int i = 0; i < n; i++) {</pre>
     dis[i] = inf, inq[i] = false;
   }
   int f = 0, 1 = 0;
   dis[src] = 0, par[src] = -1, Q[l++] = src, inq[src] =
        true;
   while (f < 1) {
     int u = Q[f++];
     for (int i = 0; i < g[u].size(); i++) {</pre>
       edge &e = g[u][i];
       if (e.cap <= e.flow) continue;</pre>
       if (dis[e.v] > dis[u] + e.cost) {
        dis[e.v] = dis[u] + e.cost;
         par[e.v] = u, idx[e.v] = i;
         if (!inq[e.v]) inq[e.v] = true, Q[1++] = e.v;
        // if (!inq[e.v]) {
        // inq[e.v] = true;
         // if (f \&\& rnd() \& 7) Q[--f] = e.v;
             else Q[1++] = e.v;
```

```
// }
       }
     }
     inq[u] = false;
   return (dis[sink] != inf);
  pair<LL, LL> solve() {
   LL mincost = 0, maxflow = 0;
   while (spfa()) {
     LL bottleneck = inf;
     for (int u = par[sink], v = idx[sink]; u != -1; v = |vector<int> id;
           idx[u], u = par[u]) {
       edge &e = g[u][v];
       bottleneck = min(bottleneck, e.cap - e.flow);
     for (int u = par[sink], v = idx[sink]; u != -1; v =
          idx[u], u = par[u]) {
       edge &e = g[u][v];
       e.flow += bottleneck;
       g[e.v][e.rev].flow -= bottleneck;
     mincost += bottleneck * dis[sink], maxflow +=
         bottleneck;
   return make_pair(mincost, maxflow);
};
// want to minimize cost and don't care about flow
// add edge from sink to dummy sink (cap = inf, cost = 0) |}
// add edge from source to sink (cap = inf, cost = 0)
// run mcmf, cost returned is the minimum cost
```

#### Block Cut Tree

```
vector<vector<int> > components;
vector<int> cutpoints, start, low;
vector<bool> is_cutpoint;
stack<int> st;
void find_cutpoints(int node, graph &G, int par = -1, int void dfs(int u, Graph &G, int ed = -1, int d = 0) {
 low[node] = start[node] = d++;
 st.push(node);
 int cnt = 0;
 for (int e : G[node])
   if (int to = G(e).to(node); to != par) {
     if (start[to] == -1) {
      find_cutpoints(to, G, node, d + 1);
       cnt++;
       if (low[to] >= start[node]) {
        is_cutpoint[node] = par != -1 or cnt > 1;
```

```
components.push_back({node}); // starting a new
             block with the point
         while (st.top() != node)
           components.back().push_back(st.top()), st.pop
       }
     low[node] = min(low[node], low[to]);
graph tree;
void init(graph &G) {
 int n = G.n;
 start.assign(n, -1), low.resize(n), is_cutpoint.resize(
      n), id.assign(n, -1);
 find_cutpoints(0, G);
 for (int u = 0; u < n; ++u)
   if (is_cutpoint[u]) id[u] = tree.addNode();
 for (auto &comp : components) {
   int node = tree.addNode();
   for (int u : comp)
     if (!is_cutpoint[u])
       id[u] = node;
     else
       tree.addEdge(node, id[u]);
 }
 if (id[0] == -1) // corner - 1
   id[0] = tree.addNode();
```

#### 3.10 Bridge Tree

```
vector<vector<int>> comps;
vector<int> depth, low, id;
stack<int> st;
vector<Edge> bridges;
Graph tree;
 low[u] = depth[u] = d;
  st.push(u);
 for (int e : G[u]) {
   if (e == ed) continue;
   int v = G(e).to(u);
   if (depth[v] == -1) dfs(v, G, e, d + 1);
   low[u] = min(low[u], low[v]);
   if (low[v] <= depth[u]) continue;</pre>
   bridges.emplace_back(u, v);
   comps.emplace_back();
```

```
do {
     comps.back().push_back(st.top()), st.pop();
   } while (comps.back().back() != v);
 if (ed == -1) {
   comps.emplace_back();
   while (!st.empty()) comps.back().push_back(st.top()),
         st.pop();
Graph &createTree() {
 for (auto &comp : comps) {
   int idx = tree.addNode();
   for (auto &e : comp) id[e] = idx;
 for (auto &[l, r] : bridges) tree.addEdge(id[l], id[r])
 return tree;
void init(Graph &G) {
 int n = G.n;
 depth.assign(n, -1), id.assign(n, -1), low.resize(n);
 for (int i = 0; i < n; i++) {</pre>
   if (depth[i] == -1) dfs(i, G);
```

#### 3.11 Tree Isomorphism

```
mp["01"] = 1;
ind = 1;
int dfs(int u, int p) {
 int cnt = 0;
 vector<int> vs:
 for (auto v : g1[u]) {
   if (v != p) {
     int got = dfs(v, u);
     vs.pb(got);
     cnt++;
   }
 }
 if (!cnt) return 1;
 sort(vs.begin(), vs.end());
 string s = "0";
 for (auto i : vs) s += to_string(i);
 vs.clear();
 s.pb('1');
 if (mp.find(s) == mp.end()) mp[s] = ++ind;
 int ret = mp[s];
```

```
return ret;
4 Math
4.1 Combi
array<int, N + 1> fact, inv, inv fact;
void init() {
 fact[0] = inv fact[0] = 1;
  for (int i = 1; i <= N; i++) {</pre>
   inv[i] = i == 1 ? 1 : (LL)inv[i - mod % i] * (mod / i
         + 1) % mod;
   fact[i] = (LL)fact[i - 1] * i % mod;
   inv_fact[i] = (LL)inv_fact[i - 1] * inv[i] % mod;
}
LL C(int n, int r) {
 return (r < 0 or r > n) ? 0 : (LL)fact[n] * inv_fact[r] | LL bigmod(LL num, LL pow, LL mod) {
      % mod * inv_fact[n - r] % mod;
4.2 Linear Sieve
const int N = 1e7;
```

```
vector<int> primes;
int spf[N + 5], phi[N + 5], NOD[N + 5], cnt[N + 5], POW[N
     + 5];
bool prime[N + 5];
int SOD[N + 5];
void init() {
 fill(prime + 2, prime + N + 1, 1);
 SOD[1] = NOD[1] = phi[1] = spf[1] = 1;
 for (LL i = 2; i <= N; i++) {</pre>
   if (prime[i]) {
     primes.push_back(i), spf[i] = i;
     phi[i] = i - 1;
     NOD[i] = 2, cnt[i] = 1;
     SOD[i] = i + 1, POW[i] = i;
   }
   for (auto p : primes) {
     if (p * i > N or p > spf[i]) break;
     prime[p * i] = false, spf[p * i] = p;
     if (i % p == 0) {
       phi[p * i] = p * phi[i];
       NOD[p * i] = NOD[i] / (cnt[i] + 1) * (cnt[i] + 2)
              cnt[p * i] = cnt[i] + 1;
       SOD[p * i] = SOD[i] / SOD[POW[i]] * (SOD[POW[i]]
           + p * POW[i]),
              POW[p * i] = p * POW[i];
       break;
     } else {
```

```
phi[p * i] = phi[p] * phi[i];
     NOD[p * i] = NOD[p] * NOD[i], cnt[p * i] = 1;
     SOD[p * i] = SOD[p] * SOD[i], POW[p * i] = p;
 }
}
```

#### 4.3 Pollard Rho

LL mul(LL a, LL b, LL mod) {

```
return ( int128)a * b % mod;
   // LL ans = a * b - mod * (LL) (1.L / mod * a * b);
   // return ans + mod * (ans < 0) - mod * (ans >= (LL)
        mod);
   LL ans = 1;
   for (; pow > 0; pow >>= 1, num = mul(num, num, mod))
       if (pow & 1) ans = mul(ans, num, mod);
   return ans;
bool is prime(LL n) {
   if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;
   LL a[] = \{2, 325, 9375, 28178, 450775, 9780504,
       1795265022}:
   LL s = \_builtin\_ctzll(n - 1), d = n >> s;
   for (LL x : a) {
       LL p = bigmod(x \% n, d, n), i = s;
       for (; p != 1 and p != n - 1 and x % n and i--; p
            = mul(p, p, n))
       if (p != n - 1 and i != s) return false;
   return true;
LL get_factor(LL n) {
   auto f = [\&](LL x) \{ return mul(x, x, n) + 1; \};
   LL x = 0, y = 0, t = 0, prod = 2, i = 2, q;
   for (; t++ % 40 or gcd(prod, n) == 1; x = f(x), y = f
       (f(y)) {
       (x == y) ? x = i++, y = f(x) : 0;
       prod = (q = mul(prod, max(x, y) - min(x, y), n))
           ? q : prod;
   return gcd(prod, n);
map<LL, int> factorize(LL n) {
   map<LL, int> res;
   if (n < 2) return res;</pre>
```

```
LL small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23,
    29, 31, 37, 41,
                                      43, 47, 53, 59,
                                            61, 67,
                                           71, 73, 79,
                                            83, 89,
                                           97};
for (LL p : small primes)
   for (; n % p == 0; n /= p, res[p]++)
auto _factor = [&](LL n, auto &_factor) {
   if (n == 1) return;
   if (is_prime(n))
       res[n]++;
   else {
       LL x = get_factor(n);
       _factor(x, _factor);
       _factor(n / x, _factor);
};
_factor(n, _factor);
return res;
```

#### 4.4 Chinese Remainder Theorem

```
// given a, b will find solutions for
// ax + by = 1
tuple<LL, LL, LL> EGCD(LL a, LL b) {
 if (b == 0)
   return {1, 0, a};
 else {
   auto [x, y, g] = EGCD(b, a \% b);
   return \{y, x - a / b * y, g\};
// given modulo equations, will apply CRT
PLL CRT(vector<PLL> &v) {
 LL V = 0, M = 1;
 for (auto &[v, m] : v) { // value % mod
   auto [x, y, g] = EGCD(M, m);
   if ((v - V) % g != 0) return {-1, 0};
   V += x * (v - V) / g % (m / g) * M, M *= m / g;
   V = (V \% M + M) \% M;
 return make_pair(V, M);
```

#### 4.5 Mobius Function

```
const int N = 1e6 + 5;
```

```
int mob[N];
void mobius() {
 memset(mob, -1, sizeof mob);
 mob[1] = 1:
 for (int i = 2; i < N; i++)</pre>
   if (mob[i]) {
    for (int j = i + i; j < N; j += i) mob[j] -= mob[i
         ];
   }
```

#### 4.6 FFT

```
using CD = complex<double>;
typedef long long LL;
const double PI = acos(-1.0L);
int N;
vector<int> perm;
vector<CD> wp[2];
void precalculate(int n) {
 assert((n & (n - 1)) == 0), N = n;
 perm = vector<int>(N, 0);
 for (int k = 1; k < N; k <<= 1) {
   for (int i = 0; i < k; i++) {</pre>
     perm[i] <<= 1;
     perm[i + k] = 1 + perm[i];
   }
 }
  wp[0] = wp[1] = vector < CD > (N);
 for (int i = 0; i < N; i++) {</pre>
   wp[0][i] = CD(cos(2 * PI * i / N), sin(2 * PI * i / N)
       )):
   wp[1][i] = CD(cos(2 * PI * i / N), -sin(2 * PI * i / N))
       N)):
 }
void fft(vector<CD> &v, bool invert = false) {
 if (v.size() != perm.size()) precalculate(v.size());
 for (int i = 0; i < N; i++)</pre>
   if (i < perm[i]) swap(v[i], v[perm[i]]);</pre>
 for (int len = 2; len <= N; len *= 2) {</pre>
   for (int i = 0, d = N / len; i < N; i += len) {</pre>
     for (int j = 0, idx = 0; j < len / 2; j++, idx += d
         ) {
       CD x = v[i + j];
       CD y = wp[invert][idx] * v[i + j + len / 2];
       v[i + j] = x + y;
       v[i + j + len / 2] = x - y;
   }
```

```
if (invert) {
   for (int i = 0; i < N; i++) v[i] /= N;</pre>
void pairfft(vector<CD> &a, vector<CD> &b, bool invert =
    false) {
 int N = a.size();
 vector<CD> p(N);
 for (int i = 0; i < N; i++) p[i] = a[i] + b[i] * CD(0,
      1);
 fft(p, invert);
 p.push_back(p[0]);
 for (int i = 0; i < N; i++) {</pre>
   if (invert) {
     a[i] = CD(p[i].real(), 0);
     b[i] = CD(p[i].imag(), 0);
   } else {
     a[i] = (p[i] + conj(p[N - i])) * CD(0.5, 0);
     b[i] = (p[i] - conj(p[N - i])) * CD(0, -0.5);
   }
 }
vector<LL> multiply(const vector<LL> &a, const vector<LL>
     &b) {
 int n = 1;
 while (n < a.size() + b.size()) n <<= 1;</pre>
  vector<CD> fa(a.begin(), a.end()), fb(b.begin(), b.end
      ());
 fa.resize(n):
 fb.resize(n);
          fft(fa); fft(fb);
 pairfft(fa, fb);
 for (int i = 0; i < n; i++) fa[i] = fa[i] * fb[i];</pre>
 fft(fa, true);
 vector<LL> ans(n);
 for (int i = 0; i < n; i++) ans[i] = round(fa[i].real()</pre>
      );
 return ans;
const int M = 1e9 + 7, B = sqrt(M) + 1;
vector<LL> anyMod(const vector<LL> &a, const vector<LL> &
    b) {
 int n = 1;
  while (n < a.size() + b.size()) n <<= 1;</pre>
  vector<CD> al(n), ar(n), bl(n), br(n);
 for (int i = 0; i < a.size(); i++) al[i] = a[i] % M / B</pre>
      , ar[i] = a[i] % M % B;
 for (int i = 0; i < b.size(); i++) bl[i] = b[i] % M / B</pre>
      , br[i] = b[i] % M % B;
```

```
pairfft(al, ar);
pairfft(bl, br);
         fft(al); fft(ar); fft(bl); fft(br);
for (int i = 0: i < n: i++) {</pre>
 CD ll = (al[i] * bl[i]), lr = (al[i] * br[i]);
 CD rl = (ar[i] * bl[i]), rr = (ar[i] * br[i]);
  al[i] = 11:
 ar[i] = lr;
 bl[i] = rl:
 br[i] = rr;
pairfft(al, ar, true);
pairfft(bl, br, true);
         fft(al, true); fft(ar, true); fft(bl, true);
    fft(br, true);
vector<LL> ans(n);
for (int i = 0; i < n; i++) {</pre>
 LL right = round(br[i].real()), left = round(al[i].
      real()):
 LL mid = round(round(bl[i].real()) + round(ar[i].real
  ans[i] = ((left \% M) * B * B + (mid \% M) * B + right)
       % M;
return ans;
```

#### 4.7 NTT

```
const LL N = 1 << 18;</pre>
const LL MOD = 786433;
vector<LL> P[N];
LL rev[N], w[N | 1], a[N], b[N], inv_n, g;
LL Pow(LL b, LL p) {
 LL ret = 1;
 while (p) {
   if (p & 1) ret = (ret * b) % MOD;
   b = (b * b) \% MOD;
   p >>= 1;
 return ret;
LL primitive_root(LL p) {
 vector<LL> factor;
 LL phi = p - 1, n = phi;
 for (LL i = 2; i * i <= n; i++) {
   if (n % i) continue;
   factor.emplace_back(i);
   while (n \% i == 0) n /= i;
```

```
if (n > 1) factor.emplace_back(n);
 for (LL res = 2; res <= p; res++) {</pre>
   bool ok = true:
   for (LL i = 0; i < factor.size() && ok; i++)</pre>
     ok &= Pow(res, phi / factor[i]) != 1;
   if (ok) return res;
 return -1;
void prepare(LL n) {
 LL sz = abs(31 - _builtin_clz(n));
 LL r = Pow(g, (MOD - 1) / n);
 inv_n = Pow(n, MOD - 2);
  w[0] = w[n] = 1;
 for (LL i = 1; i < n; i++) w[i] = (w[i-1] * r) % MOD;
 for (LL i = 1: i < n: i++)
   rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (sz - 1));
void NTT(LL *a, LL n, LL dir = 0) {
 for (LL i = 1; i < n - 1; i++)
   if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
 for (LL m = 2; m <= n; m <<= 1) {
   for (LL i = 0; i < n; i += m) {
     for (LL j = 0; j < (m >> 1); j++) {
       LL &u = a[i + j], &v = a[i + j + (m >> 1)];
       LL t = v * w[dir ? n - n / m * j : n / m * j] %
       v = u - t < 0 ? u - t + MOD : u - t;
       u = u + t >= MOD ? u + t - MOD : u + t:
     }
   }
   for (LL i = 0; i < n; i++) a[i] = (inv_n * a[i]) %</pre>
vector<LL> mul(vector<LL> p, vector<LL> q) {
 LL n = p.size(), m = q.size();
 LL t = n + m - 1, sz = 1;
 while (sz < t) sz <<= 1;
 prepare(sz);
 for (LL i = 0; i < n; i++) a[i] = p[i];</pre>
 for (LL i = 0; i < m; i++) b[i] = q[i];</pre>
 for (LL i = n; i < sz; i++) a[i] = 0;
 for (LL i = m; i < sz; i++) b[i] = 0;</pre>
 NTT(a, sz);
 NTT(b, sz);
```

```
for (LL i = 0; i < sz; i++) a[i] = (a[i] * b[i]) % MOD;
NTT(a, sz, 1);

vector<LL> c(a, a + sz);
while (c.size() && c.back() == 0) c.pop_back();
return c;
}
```

```
4.8 WalshHadamard
#include <bits/stdc++.h>
using namespace std;
typedef long long LL;
#define bitwiseXOR 1
// #define bitwiseAND 2
// #define bitwiseOR 3
const LL MOD = 30011;
LL BigMod(LL b, LL p) {
 LL ret = 1:
 while (p > 0) {
   if (p % 2 == 1) {
     ret = (ret * b) % MOD;
   p = p / 2;
   b = (b * b) \% MOD;
 return ret % MOD;
void FWHT(vector<LL>& p, bool inverse) {
 LL n = p.size();
 assert((n & (n - 1)) == 0);
 for (LL len = 1; 2 * len <= n; len <<= 1) {
   for (LL i = 0; i < n; i += len + len) {</pre>
     for (LL j = 0; j < len; j++) {
       LL u = p[i + j];
       LL v = p[i + len + j];
#ifdef bitwiseXOR
       p[i + j] = (u + v) \% MOD;
       p[i + len + j] = (u - v + MOD) \% MOD;
#endif // bitwiseXOR
#ifdef bitwiseAND
       if (!inverse) {
        p[i + j] = v \% MOD;
         p[i + len + j] = (u + v) \% MOD;
       } else {
         p[i + j] = (-u + v) \% MOD;
```

```
p[i + len + j] = u \% MOD;
#endif // bitwiseAND
#ifdef bitwiseOR
       if (!inverse) {
         p[i + j] = u + v;
         p[i + len + j] = u;
       } else {
         p[i + j] = v;
         p[i + len + j] = u - v;
#endif // bitwiseOR
  }
 }
#ifdef bitwiseXOR
 if (inverse) {
   LL val = BigMod(n, MOD - 2); // Option 2: Exclude
   for (LL i = 0; i < n; i++) {</pre>
     // assert(p[i]%n==0); //Option 2: Include
     p[i] = (p[i] * val) % MOD; // Option 2: p[i]/=n;
#endif // bitwiseXOR
```

#### 4.9 Berlekamp Massey

```
struct berlekamp_massey { // for linear recursion
 typedef long long LL;
 static const int SZ = 2e5 + 5;
 static const int MOD = 1e9 + 7; /// mod must be a prime
LL m, a[SZ], h[SZ], t_[SZ], s[SZ], t[SZ];
 // bigmod goes here
 inline vector <LL> BM( vector <LL> &x ) {
  LL lf , ld;
  vector <LL> ls , cur;
  for ( int i = 0; i < int(x.size()); ++i ) {</pre>
    for ( int j = 0; j < int(cur.size()); ++j ) t = (t</pre>
        + x[i - j - 1] * cur[j]) % MOD;
    if ((t - x[i]) \% MOD == 0) continue;
    if (!cur.size()) {
      cur.resize( i + 1 );
      lf = i; ld = (t - x[i]) \% MOD;
      continue;
    LL k = -(x[i] - t) * bigmod(ld, MOD - 2, MOD) %
```

```
vector <LL> c(i - lf - 1);
   c.push back( k );
   for ( int j = 0; j < int(ls.size()); ++j ) c.</pre>
        push_back(-ls[j] * k % MOD);
   if ( c.size() < cur.size() ) c.resize( cur.size() )</pre>
   for ( int j = 0; j < int(cur.size()); ++j ) c[j] =</pre>
        (c[i] + cur[i]) % MOD;
   if (i - lf + (int)ls.size() >= (int)cur.size() ) ls
         = cur, lf = i, ld = (t - x[i]) \% MOD;
   cur = c;
 for ( int i = 0; i < int(cur.size()); ++i ) cur[i] =</pre>
      (cur[i] % MOD + MOD) % MOD;
 return cur;
inline void mull( LL *p , LL *q ) {
 for ( int i = 0; i < m + m; ++i ) t [i] = 0;
 for ( int i = 0; i < m; ++i ) if ( p[i] )
     for ( int j = 0; j < m; ++j ) t_[i + j] = (t_[i +
          j] + p[i] * q[j]) % MOD;
 for ( int i = m + m - 1; i >= m; --i ) if ( t_[i] )
     for ( int j = m - 1; \tilde{j}; --j ) t_{i} = 1 = (
         t_{i} = i - j - 1 + t_{i} * h_{i} % MOD;
 for ( int i = 0; i < m; ++i ) p[i] = t_[i];</pre>
inline LL calc( LL K ) {
 for ( int i = m; ~i; --i ) s[i] = t[i] = 0;
 s[0] = 1; if (m!=1) t[1] = 1; else t[0] = h[0];
 while (K) {
  if (K & 1 ) mull(s, t);
   mull( t , t ); K >>= 1;
 }
 LL su = 0;
 for ( int i = 0; i < m; ++i ) su = (su + s[i] * a[i])
 return (su % MOD + MOD) % MOD;
/// already calculated upto k , now calculate upto n.
inline vector <LL> process( vector <LL> &x , int n ,
    int k ) {
 auto re = BM( x );
 x.resize(n+1):
 for ( int i = k + 1; i <= n; i++ ) {
  for ( int j = 0; j < re.size(); j++ ) {</pre>
     x[i] += 1LL * x[i - j - 1] % MOD * re[j] % MOD; x
         [i] %= MOD;
   }
 }
 return x;
```

```
inline LL work( vector <LL> &x , LL n ) {
   if ( n < int(x.size()) ) return x[n] % MOD;
   vector <LL> v = BM( x ); m = v.size(); if ( !m )
        return 0;
   for ( int i = 0; i < m; ++i ) h[i] = v[i], a[i] = x[i
        ];
   return calc( n ) % MOD;
}
return calc( n ) % MOD;
}
rec;
vector <LL> v;
void solve() {
   int n;
   cin >> n;
   cout << rec.work(v, n - 1) << endl;
}
</pre>
```

#### 4.10 Lagrange

```
// p is a polynomial with n points.
// p(0), p(1), p(2), ... p(n-1) are given.
// Find p(x).
LL Lagrange(vector<LL> &p, LL x) {
 LL n = p.size(), L, i, ret;
 if (x < n) return p[x];</pre>
 L = 1:
 for (i = 1; i < n; i++) {
   L = (L * (x - i)) \% MOD;
   L = (L * bigmod(MOD - i, MOD - 2)) \% MOD;
 ret = (L * p[0]) % MOD;
 for (i = 1; i < n; i++) {
   L = (L * (x - i + 1)) \% MOD;
   L = (L * bigmod(x - i, MOD - 2)) % MOD;
   L = (L * bigmod(i, MOD - 2)) % MOD;
   L = (L * (MOD + i - n)) % MOD;
   ret = (ret + L * p[i]) % MOD;
 return ret;
```

#### 4.11 Shanks' Baby Step, Giant Step

```
// Finds a^x = b (mod p)

LL bigmod(LL b, LL p, LL m) {}

LL babyStepGiantStep(LL a, LL b, LL p) {
   LL i, j, c, sq = sqrt(p);
   map<LL, LL> babyTable;
```

#### 4.12 Xor Basis

```
struct XorBasis {
 static const int sz = 64:
 array<ULL, sz> base = {0}, back;
 array<int, sz> pos;
 void insert(ULL x, int p) {
  ULL cur = 0;
  for (int i = sz - 1; ~i; i--)
    if (x >> i & 1) {
      if (!base[i]) {
        base[i] = x, back[i] = cur, pos[i] = p;
      } else x ^= base[i], cur |= 1ULL << i;</pre>
 pair<ULL, vector<int>> construct(ULL mask) {
  ULL ok = 0, x = mask;
  for (int i = sz - 1; ~i; i--)
     if (mask >> i & 1 and base[i]) mask ^= base[i], ok
         |= 1ULL << i;
   vector<int> ans;
   for (int i = 0; i < sz; i++)</pre>
    if (ok >> i & 1) {
      ans.push_back(pos[i]);
      ok ^= back[i];
   return {x ^ mask, ans};
```

# 5 String

#### 5.1 Aho Corasick

```
struct AC {
int N, P;
const int A = 26;
vector<vector<int>> next;
vector<int> link, out link;
```

```
vector<vector<int>> out:
AC(): N(0), P(0) { node(); }
int node() {
 next.emplace_back(A, 0);
 link.emplace_back(0);
 out_link.emplace_back(0);
 out.emplace_back(0);
 return N++;
inline int get(char c) { return c - 'a'; }
int add_pattern(const string T) {
 int u = 0:
 for (auto c : T) {
   if (!next[u][get(c)]) next[u][get(c)] = node();
   u = next[u][get(c)];
  out[u].push_back(P);
 return P++;
void compute() {
 queue<int> q;
 for (q.push(0); !q.empty();) {
   int u = q.front(); q.pop();
   for (int c = 0; c < A; ++c) {</pre>
     int v = next[u][c]:
     if (!v) next[u][c] = next[link[u]][c];
     else {
       link[v] = u ? next[link[u]][c] : 0;
       out_link[v] = out[link[v]].empty() ? out_link[
           link[v]] : link[v];
       q.push(v);
     }
   }
 }
int advance(int u, char c) {
  while (u && !next[u][get(c)]) u = link[u];
 u = next[u][get(c)];
 return u;
void match(const string S) {
 int u = 0;
 for (auto c : S) {
   u = advance(u, c);
   for (int v = u; v; v = out_link[v]) {
     for (auto p : out[v]) cout << "match " << p << endl</pre>
   }
 }
```

```
};
int main() {
    AC aho; int n; cin >> n;
    while (n--) {
        string s; cin >> s;
        aho.add_pattern(s);
    }
    aho.compute(); string text;
    cin >> text; aho.match(text);
    return 0;
}
```

```
5.2 Double hash
// define +, -, * for (PLL, LL) and (PLL, PLL), % for (
    PLL, PLL);
PLL base(1949313259, 1997293877);
PLL mod(2091573227, 2117566807);
PLL power(PLL a, LL p) {
 PLL ans = PLL(1, 1);
 for(; p; p >>= 1, a = a * a % mod) {
     if(p \& 1) ans = ans * a % mod;
 return ans:
PLL inverse(PLL a) { return power(a, (mod.ff - 1) * (mod.
    ss - 1) - 1); }
PLL inv_base = inverse(base);
PLL val;
vector<PLL> P;
void hash_init(int n) {
 P.resize(n + 1):
 P[0] = PLL(1, 1);
  for (int i = 1; i <= n; i++) P[i] = (P[i - 1] * base) %
PLL append(PLL cur, char c) { return (cur * base + c) %
/// prepends c to string with size k
PLL prepend(PLL cur, int k, char c) { return (P[k] * c +
    cur) % mod: }
/// replaces the i-th (0-indexed) character from right
    from a to b:
PLL replace(PLL cur, int i, char a, char b) {
 cur = (cur + P[i] * (b - a)) \% mod;
 return (cur + mod) % mod;
/// Erases c from the back of the string
```

```
PLL pop_back(PLL hash, char c) {
 return (((hash - c) * inv base) % mod + mod) % mod;
/// Erases c from front of the string with size len
PLL pop_front(PLL hash, int len, char c) {
 return ((hash - P[len - 1] * c) % mod + mod) % mod;
/// concatenates two strings where length of the right is
PLL concat(PLL left, PLL right, int k) { return (left * P
    [k] + right) % mod; }
/// Calculates hash of string with size len repeated cnt
/// This is O(log n). For O(1), pre-calculate inverses
PLL repeat(PLL hash, int len, LL cnt) {
 PLL mul = (P[len * cnt] - 1) * inverse(P[len] - 1);
 mul = (mul % mod + mod) % mod;
 PLL ret = (hash * mul) % mod;
 if (P[len].ff == 1) ret.ff = hash.ff * cnt:
 if (P[len].ss == 1) ret.ss = hash.ss * cnt;
 return ret;
LL get(PLL hash) { return ((hash.ff << 32) ^ hash.ss); }
struct hashlist {
 int len:
 vector<PLL> H, R;
 hashlist() {}
 hashlist(string& s) {
   len = (int)s.size();
   hash_init(len);
   H.resize(len + 1, PLL(0, 0)), R.resize(len + 2, PLL
   for (int i = 1; i <= len; i++) H[i] = append(H[i -</pre>
        1], s[i - 1]);
   for (int i = len; i >= 1; i--) R[i] = append(R[i +
        1], s[i - 1]);
 /// 1-indexed
 PLL range_hash(int 1, int r) {
   return ((H[r] - H[l - 1] * P[r - l + 1]) % mod + mod)
        % mod:
 PLL reverse_hash(int 1, int r) {
   return ((R[1] - R[r + 1] * P[r - 1 + 1]) \% mod + mod)
         % mod;
 PLL concat_range_hash(int 11, int r1, int 12, int r2) {
   return concat(range_hash(11, r1), range_hash(12, r2),
         r2 - 12 + 1);
 }
```

#### 5.3 Manacher's

```
vector<int> d1(n):
// d[i] = number of palindromes taking s[i] as center
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
  while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k])
      k++:
 d1[i] = k--;
 if (i + k > r) l = i - k, r = i + k;
vector<int> d2(n);
// d[i] = number of palindromes taking s[i-1] and s[i] as
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1)
 while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s
      [i + k]) k++;
 d2[i] = k--;
 if (i + k > r) l = i - k - 1, r = i + k;
```

### 5.4 Suffix Array

```
c[0][p[i]] = classes - 1;
 VI pn(n), cn(n);
 cnt.resize(n);
 for (int h = 0; (1 << h) < n; h++) {
   for (int i = 0; i < n; i++) {</pre>
     pn[i] = p[i] - (1 << h);
     if (pn[i] < 0) pn[i] += n;</pre>
   }
   fill(cnt.begin(), cnt.end(), 0);
   /// radix sort
   for (int i = 0; i < n; i++) cnt[c[h][pn[i]]]++;</pre>
   for (int i = 1; i < classes; i++) cnt[i] += cnt[i -</pre>
   for (int i = n - 1; i >= 0; i--) p[--cnt[c[h][pn[i
       ]]]] = pn[i];
   cn[p[0]] = 0;
   classes = 1;
   for (int i = 1; i < n; i++) {</pre>
     PII cur = \{c[h][p[i]], c[h][(p[i] + (1 << h)) \% n\}
     PII prev = \{c[h][p[i-1]], c[h][(p[i-1] + (1 <<
         h)) % n]};
     if (cur != prev) ++classes;
     cn[p[i]] = classes - 1;
   c.push_back(cn);
 return p;
VI suffix_array_construction(string s) {
 VI sorted_shifts = sort_cyclic_shifts(s);
 sorted shifts.erase(sorted shifts.begin());
 return sorted_shifts;
/// LCP between the ith and jth (i != j) suffix of the
    STRING
int suffixLCP(int i, int j) {
 assert(i != j);
 int log_n = c.size() - 1;
 int ans = 0:
 for (int k = log_n; k >= 0; k--) {
   if (c[k][i] == c[k][j]) {
     ans += 1 << k;
     i += 1 << k;
```

```
j += 1 << k;
 }
 return ans;
VI lcp_construction(const string &s, const VI &sa) {
 int n = s.size();
 VI rank(n. 0):
 VI lcp(n-1, 0);
 for (int i = 0; i < n; i++) rank[sa[i]] = i;</pre>
 for (int i = 0, k = 0; i < n; i++, k -= (k != 0)) {
   if (rank[i] == n - 1) {
    k = 0;
     continue;
   int j = sa[rank[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
       ]) k++;
   lcp[rank[i]] = k;
 return lcp;
```

#### 5.5 Z Algo

```
vector<int> calcz(string s) {
int n = s.size();
 vector<int> z(n);
 int 1 = 0, r = 0;
 for (int i = 1; i < n; i++) {
  if (i > r) {
    l = r = i:
    while (r < n \&\& s[r] == s[r - 1]) r++;
    z[i] = r - 1, r--;
  } else {
    int k = i - 1;
    if (z[k] < r - i + 1) z[i] = z[k];
    else {
      while (r < n \&\& s[r] == s[r - 1]) r++;
      z[i] = r - 1, r--;
  }
}
return z;
```

# 6 DP

#### 6.1 1D-1D

```
/// Author: anachor
#include <bits/stdc++.h>
using namespace std;
/// Solves dp[i] = min(dp[j] + cost(j+1, i)) given that
    cost() is QF
long long solve1D(int n, long long cost(int, int)) {
  vector<long long> dp(n + 1), opt(n + 1);
 deque<pair<int, int>> dq;
  dq.push_back({0, 1});
 dp[0] = 0;
  for (int i = 1; i <= n; i++) {</pre>
   opt[i] = dq.front().first;
   dp[i] = dp[opt[i]] + cost(opt[i] + 1, i);
   if (i == n) break;
   dq[0].second++;
   if (dq.size() > 1 \&\& dq[0].second == dq[1].second) dq
        .pop_front();
   int en = n;
   while (dq.size()) {
     int o = dq.back().first, st = dq.back().second;
     if (dp[o] + cost(o + 1, st) >= dp[i] + cost(i + 1,
          st))
       dq.pop back();
     else {
       int lo = st, hi = en;
       while (lo < hi) {</pre>
         int mid = (lo + hi + 1) / 2;
         if (dp[o] + cost(o + 1, mid) < dp[i] + cost(i +</pre>
             1, mid))
          lo = mid;
         else
           hi = mid - 1:
       if (lo < n) dq.push_back({i, lo + 1});</pre>
       break;
     }
     en = st - 1;
   if (dq.empty()) dq.push_back({i, i + 1});
 return dp[n];
```

#### 6.2 Convex Hull Trick

struct line {

```
11 m, c;
    line() {}
     line(ll m, ll c) : m(m), c(c) {}
struct convex_hull_trick {
      vector<line> lines;
       int ptr = 0;
      convex_hull_trick() {}
      bool bad(line a, line b, line c) {
             return 1.0 * (c.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (b.c - a.c) * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (b.c - a.c) * (b.c - a.c) * (a.m - b.m) < 1.0 * (b.c - a.c) * (b.c) * (b.c - a.c) * (b.c - a.c) * (b.c - a.c) * (b.c - a.c) * (b.c) * (b.c - a.c) * (b.c) *
                                   a.c) * (a.m - c.m):
      void add(line L) {
             int sz = lines.size();
              while (sz >= 2 && bad(lines[sz - 2], lines[sz - 1], L
                              )) {
                    lines.pop_back();
                     sz--;
             }
             lines.pb(L);
      ll get(int idx, int x) { return (1ll * lines[idx].m * x
                           + lines[idx].c); }
      11 query(int x) {
             if (lines.empty()) return 0;
              if (ptr >= lines.size()) ptr = lines.size() - 1;
              while (ptr < lines.size() - 1 && get(ptr, x) > get(
                               ptr + 1, x)) ptr++;
             return get(ptr, x);
```

```
}
};
ll sum[MAX];
ll dp[MAX];
int arr[MAX];
int main() {
 fastio:
  int t;
  cin >> t:
  while (t--) {
   int n, a, b, c;
   cin >> n >> a >> b >> c;
   for (int i = 1; i <= n; i++) cin >> sum[i];
   for (int i = 1; i <= n; i++) dp[i] = 0, sum[i] += sum</pre>
        [i - 1];
   convex hull trick cht;
   cht.add(line(0, 0));
   for (int pos = 1; pos <= n; pos++) {</pre>
     dp[pos] = cht.query(sum[pos]) - 111 * a * sqr(sum[
          pos]) - c;
     cht.add(line(211 * a * sum[pos], dp[pos] - a * sqr(
          sum[pos])));
   11 \text{ ans} = (-111 * dp[n]);
   ans += (111 * sum[n] * b);
   cout << ans << "\n";
 }
```

#### 6.3 Divide and Conquer dp

```
const int K = 805, N = 4005;
LL dp[2][N], _cost[N][N];
// 1-indexed for convenience
LL cost(int 1, int r) {
 return _cost[r][r] - _cost[l - 1][r] - _cost[r][l - 1]
      + _cost[1 - 1][1 - 1] >> 1;
void compute(int cnt, int 1, int r, int optl, int optr) {
 if (1 > r) return;
 int mid = 1 + r >> 1:
 LL best = INT MAX;
 int opt = -1;
 for (int i = optl; i <= min(mid, optr); i++) {</pre>
   LL cur = dp[cnt ^1][i - 1] + cost(i, mid);
   if (cur < best) best = cur, opt = i;</pre>
 dp[cnt][mid] = best;
 compute(cnt, 1, mid - 1, optl, opt);
 compute(cnt, mid + 1, r, opt, optr);
```

```
LL dnc_dp(int k, int n) {
 fill(dp[0] + 1, dp[0] + n + 1, INT_MAX);
 for (int cnt = 1; cnt <= k; cnt++) {</pre>
   compute(cnt & 1, 1, n, 1, n);
 return dp[k & 1][n];
```

# 6.4 Dynamic CHT typedef long long LL;

```
const LL IS QUERY = -(1LL << 62);</pre>
struct line {
  LL m, b;
  mutable function <const line*()> succ;
  bool operator < (const line &rhs) const {</pre>
    if (rhs.b != IS_QUERY) return m < rhs.m;</pre>
    const line *s = succ();
    if (!s) return 0;
    LL x = rhs.m:
    return b - s -> b < (s -> m - m) * x;
};
struct HullDynamic : public multiset <line> {
  bool bad (iterator y) {
    auto z = next(y);
    if (y == begin()) {
     if (z == end()) return 0;
      return y -> m == z -> m && y -> b <= z -> b;
    auto x = prev(y);
    if (z == end()) return y \rightarrow m == x \rightarrow m &  y \rightarrow b <=
    return 1.0 * (x \rightarrow b - y \rightarrow b) * (z \rightarrow m - y \rightarrow m) >=
          1.0 * (y \rightarrow b - z \rightarrow b) * (y \rightarrow m - x \rightarrow m);
  }
  void insert line (LL m, LL b) {
    auto y = insert({m, b});
    y -> succ = [=] {return next(y) == end() ? 0 : &*next
         (y);};
    if (bad(y)) {erase(y); return;}
    while (next(y) != end() && bad(next(y))) erase(next(y))
    while (y != begin() && bad(prev(y))) erase(prev(y));
```

```
LL eval (LL x) {
   auto 1 = *lower_bound((line) {x, IS_QUERY});
   return 1.m * x + 1.b;
};
```

#### 6.5 FFT Online

```
void fftOnline(vector <LL> &a, vector <LL> b) {
 int n = a.size();
 auto call = [&](int 1, int r, auto &call){
   if(1 \ge r) return:
   int mid = 1 + r >> 1;
   call(1, mid, call);
   vector <LL> _a(a.begin() + 1, a.begin() + mid + 1);
   vector \langle LL \rangle _b(b.begin(), b.begin() + (r - 1 + 1));
   auto c = fft :: anyMod(_a, _b);
   for(int i = mid + 1; i <= r; i++) {</pre>
    a[i] += c[i - 1];
     a[i] -= (a[i] >= mod) * mod;
   call(mid + 1, r, call);
 call(0, n - 1, call);
```

#### 6.6 Knuth optimization

```
const int N = 1005;
LL dp[N][N], a[N];
int opt[N][N];
LL cost(int i, int j) { return a[j + 1] - a[i]; }
LL knuth_optimization(int n) {
 for (int i = 0; i < n; i++) {</pre>
   dp[i][i] = 0;
   opt[i][i] = i;
 for (int i = n - 2; i \ge 0; i--) {
   for (int j = i + 1; j < n; j++) {
     LL mn = LLONG MAX:
     LL c = cost(i, j);
     for (int k = opt[i][j - 1]; k <= min(j - 1, opt[i +</pre>
           1][i]): k++) {
       if (mn > dp[i][k] + dp[k + 1][j] + c) {
         mn = dp[i][k] + dp[k + 1][j] + c;
         opt[i][i] = k;
       }
     dp[i][j] = mn;
```

```
return dp[0][n - 1];
```

```
6.7 Li Chao Tree
struct line {
 LL m. c:
 line(LL m = 0, LL c = 0) : m(m), c(c) {}
LL calc(line L, LL x) { return 1LL * L.m * x + L.c; }
struct node {
 LL m, c;
 line L:
 node *lft, *rt;
 node(LL m = 0, LL c = 0, node *lft = NULL, node *rt =
     : L(line(m, c)), lft(lft), rt(rt) {}
};
struct LiChao {
 node *root:
 LiChao() { root = new node(); }
 void update(node *now, int L, int R, line newline) {
   int mid = L + (R - L) / 2;
   line lo = now->L, hi = newline;
   if (calc(lo, L) > calc(hi, L)) swap(lo, hi);
   if (calc(lo, R) <= calc(hi, R)) {</pre>
     now->L = hi;
     return:
   if (calc(lo, mid) < calc(hi, mid)) {</pre>
     now->L = hi:
     if (now->rt == NULL) now->rt = new node();
     update(now->rt, mid + 1, R, lo);
   } else {
     now->L = lo:
     if (now->lft == NULL) now->lft = new node();
     update(now->lft, L, mid, hi);
   }
 }
 LL query(node *now, int L, int R, LL x) {
   if (now == NULL) return -inf:
   int mid = L + (R - L) / 2;
   if (x \le mid)
     return max(calc(now->L, x), query(now->lft, L, mid,
          x));
   else
     return max(calc(now->L, x), query(now->rt, mid + 1,
          R, x));
 }
```

### 7 Geometry

#### 7.1 Point

```
typedef double Tf;
typedef double Ti; /// use long long for exactness
const Tf PI = acos(-1), EPS = 1e-9;
int dcmp(Tf x) \{ return abs(x) < EPS ? 0 : (x < 0 ? -1 : 
    1); }
struct Point {
   Ti x, y;
   Point(Ti x = 0, Ti y = 0) : x(x), y(y) {}
   Point operator+(const Point& u) const { return Point(
        x + u.x, y + u.y; }
   Point operator-(const Point& u) const { return Point(
       x - u.x, y - u.y); }
   Point operator*(const LL u) const { return Point(x *
       u, y * u); }
   Point operator*(const Tf u) const { return Point(x *
       u, y * u); }
   Point operator/(const Tf u) const { return Point(x /
       u, y / u); }
   bool operator==(const Point& u) const {
       return dcmp(x - u.x) == 0 && dcmp(y - u.y) == 0;
   bool operator!=(const Point& u) const { return !(*
        this == u); }
   bool operator<(const Point& u) const {</pre>
       return dcmp(x - u.x) < 0 \mid \mid (dcmp(x - u.x) == 0
           && dcmp(y - u.y) < 0);
   }
};
Ti dot(Point a, Point b) { return a.x * b.x + a.y * b.y;
Ti cross(Point a, Point b) { return a.x * b.y - a.y * b.x
Tf length(Point a) { return sqrt(dot(a, a)); }
Ti sqLength(Point a) { return dot(a, a); }
Tf distance(Point a, Point b) { return length(a - b); }
Tf angle(Point u) { return atan2(u.y, u.x); }
// returns angle between oa, ob in (-PI, PI]
Tf angleBetween(Point a, Point b) {
   Tf ans = angle(b) - angle(a);
   return ans <= -PI ? ans + 2 * PI : (ans > PI ? ans -
        2 * PI : ans);
// Rotate a ccw by rad radians, Tf Ti same
```

```
Point rotate(Point a, Tf rad) {
    return Point(a.x * cos(rad) - a.y * sin(rad),
                           a.x * sin(rad) + a.y * cos(rad)
                               )):
// rotate a ccw by angle th with cos(th) = co && sin(th)
    = si. tf ti same
Point rotatePrecise(Point a, Tf co, Tf si) {
    return Point(a.x * co - a.y * si, a.y * co + a.x * si
Point rotate90(Point a) { return Point(-a.y, a.x); }
// scales vector a by s such that length of a becomes s,
    Tf Ti same
Point scale(Point a, Tf s) { return a / length(a) * s; }
// returns an unit vector perpendicular to vector a, Tf
    Ti same
Point normal(Point a) {
   Tf 1 = length(a):
   return Point(-a.y / 1, a.x / 1);
// returns 1 if c is left of ab, 0 if on ab && -1 if
    right of ab
int orient(Point a, Point b, Point c) { return dcmp(cross
    (b - a, c - a): }
/// Use as sort(v.begin(), v.end(), polarComp(0, dir))
/// Polar comparator around O starting at direction dir
struct polarComp {
   Point O, dir;
    polarComp(Point 0 = Point(0, 0), Point dir = Point(1,
         0)) : O(O), dir(dir) {}
    bool half(Point p) {
       return dcmp(cross(dir, p)) < 0 ||</pre>
                    (dcmp(cross(dir, p)) == 0 && dcmp(dot
                        (dir, p)) > 0);
    bool operator()(Point p, Point q) {
       return make_tuple(half(p), 0) < make_tuple(half(q))</pre>
           ), cross(p, q));
   }
struct Segment {
   Point a, b;
   Segment(Point aa, Point bb) : a(aa), b(bb) {}
typedef Segment Line;
struct Circle {
   Point o;
   Tf r;
```

```
Circle(Point o = Point(0, 0), Tf r = 0) : o(o), r(r)
   // returns true if point p is in || on the circle
   bool contains(Point p) { return dcmp(sqLength(p - o)
        - r * r) <= 0; }
   // returns a point on the circle rad radians away
       from +X CCW
   Point point(Tf rad) {
       static_assert(is_same<Tf, Ti>::value);
       return Point(o.x + cos(rad) * r, o.y + sin(rad) *
            r);
   // area of a circular sector with central angle rad
   Tf area(Tf rad = PI + PI) { return rad * r * r / 2; }
   // area of the circular sector cut by a chord with
        central angle alpha
   Tf sector(Tf alpha) { return r * r * 0.5 * (alpha - return)
       sin(alpha)); }
};
```

#### 7.2 Linear

```
// **** LINE LINE INTERSECTION START ****
// returns true if point p is on segment s
bool onSegment(Point p, Segment s) {
 return dcmp(cross(s.a - p, s.b - p)) == 0 && dcmp(dot(s
      .a - p, s.b - p)) <= 0;
// returns true if segment p && q touch or intersect
bool segmentsIntersect(Segment p, Segment q) {
 if (onSegment(p.a, q) || onSegment(p.b, q)) return true
 if (onSegment(q.a, p) || onSegment(q.b, p)) return true
 Ti c1 = cross(p.b - p.a, q.a - p.a);
 Ti c2 = cross(p.b - p.a, q.b - p.a);
 Ti c3 = cross(q.b - q.a, p.a - q.a);
 Ti c4 = cross(q.b - q.a, p.b - q.a);
 return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) <</pre>
bool linesParallel(Line p, Line q) {
 return dcmp(cross(p.b - p.a, q.b - q.a)) == 0;
// lines are represented as a ray from a point: (point,
// returns false if two lines (p, v) && (q, w) are
    parallel or collinear
// true otherwise, intersection point is stored at o via
    reference, Tf Ti Same
```

```
bool lineLineIntersection(Point p, Point v, Point q,
    Point w, Point& o) {
 if (dcmp(cross(v, w)) == 0) return false;
 Point u = p - q;
  o = p + v * (cross(w, u) / cross(v, w));
 return true;
// returns false if two lines p && g are parallel or
// true otherwise, intersection point is stored at o via
    reference
bool lineLineIntersection(Line p, Line q, Point& o) {
 return lineLineIntersection(p.a, p.b - p.a, q.a, q.b -
      q.a, o);
// returns the distance from point a to line 1
// **** LINE LINE INTERSECTION FINISH ****
Tf distancePointLine(Point p, Line 1) {
 return abs(cross(l.b - l.a, p - l.a) / length(l.b - l.a
     ));
// returns the shortest distance from point a to segment
Tf distancePointSegment(Point p, Segment s) {
  if (s.a == s.b) return length(p - s.a);
 Point v1 = s.b - s.a, v2 = p - s.a, v3 = p - s.b;
  if (dcmp(dot(v1, v2)) < 0)
   return length(v2);
  else if (dcmp(dot(v1, v3)) > 0)
   return length(v3);
  else
   return abs(cross(v1, v2) / length(v1));
// returns the shortest distance from segment p to
Tf distanceSegmentSegment(Segment p, Segment q) {
 if (segmentsIntersect(p, q)) return 0;
 Tf ans = distancePointSegment(p.a, q);
  ans = min(ans, distancePointSegment(p.b, q));
  ans = min(ans, distancePointSegment(q.a, p));
  ans = min(ans, distancePointSegment(q.b, p));
 return ans;
// returns the projection of point p on line 1, Tf Ti
Point projectPointLine(Point p, Line 1) {
 Point v = 1.b - 1.a;
 return 1.a + v * ((Tf)dot(v, p - 1.a) / dot(v, v));
```

#### 7.3 Circular

```
// Extremely inaccurate for finding near touches
// compute intersection of line 1 with circle c
// The intersections are given in order of the ray (1.a,
    1.b). Tf Ti same
vector<Point> circleLineIntersection(Circle c, Line 1) {
   vector<Point> ret:
   Point b = 1.b - 1.a, a = 1.a - c.o;
   Tf A = dot(b, b), B = dot(a, b);
   Tf C = dot(a, a) - c.r * c.r, D = B * B - A * C;
   if (D < -EPS) return ret;</pre>
   ret.push_back(1.a + b * (-B - sqrt(D + EPS)) / A);
   if (D > EPS) ret.push back(1.a + b * (-B + sqrt(D)))
   return ret;
// signed area of intersection of circle(c.o, c.r) &&
// triangle(c.o, s.a, s.b) [cross(a-o, b-o)/2]
Tf circleTriangleIntersectionArea(Circle c, Segment s) {
   using Linear::distancePointSegment;
   Tf OA = length(c.o - s.a);
   Tf OB = length(c.o - s.b);
   // sector
   if (dcmp(distancePointSegment(c.o, s) - c.r) >= 0)
       return angleBetween(s.a - c.o, s.b - c.o) * (c.r
           * c.r) / 2.0:
   // triangle
   if (dcmp(OA - c.r) \le 0 \&\& dcmp(OB - c.r) \le 0)
       return cross(c.o - s.b, s.a - s.b) / 2.0;
   // three part: (A, a) (a, b) (b, B)
   vector<Point> Sect = circleLineIntersection(c, s);
   return circleTriangleIntersectionArea(c, Segment(s.a,
         Sect[0])) +
                circleTriangleIntersectionArea(c,
                     Segment(Sect[0], Sect[1])) +
                circleTriangleIntersectionArea(c,
                    Segment(Sect[1], s.b));
// area of intersecion of circle(c.o, c.r) && simple
    polyson(p[])
Tf circlePolyIntersectionArea(Circle c, Polygon p) {
   Tf res = 0;
   int n = p.size();
   for (int i = 0; i < n; ++i)
       res += circleTriangleIntersectionArea(c, Segment(
           p[i], p[(i + 1) % n]);
   return abs(res);
// locates circle c2 relative to c1
// interior
                      (d < R - r)
                                        ---> -2
```

```
// interior tangents (d = R - r)
// concentric
                   (d = 0)
// secants
                     (R - r < d < R + r) \longrightarrow 0
// exterior tangents (d = R + r)
// exterior
                      (d > R + r)
int circleCirclePosition(Circle c1, Circle c2) {
   Tf d = length(c1.o - c2.o);
   int in = dcmp(d - abs(c1.r - c2.r)), ex = dcmp(d - (abs(c1.r - c2.r)))
        c1.r + c2.r):
   return in < 0 ? -2 : in == 0 ? -1 : ex == 0 ? 1 : ex
       > 0 ? 2 : 0;
// compute the intersection points between two circles c1
     && c2, Tf Ti same
vector<Point> circleCircleIntersection(Circle c1, Circle
    c2) {
   vector<Point> ret:
   Tf d = length(c1.o - c2.o);
   if (dcmp(d) == 0) return ret;
   if (dcmp(c1.r + c2.r - d) < 0) return ret;
   if (dcmp(abs(c1.r - c2.r) - d) > 0) return ret;
   Point v = c2.0 - c1.0;
   Tf co = (c1.r * c1.r + sqLength(v) - c2.r * c2.r) /
        (2 * c1.r * length(v));
   Tf si = sqrt(abs(1.0 - co * co));
   Point p1 = scale(rotatePrecise(v, co, -si), c1.r) +
   Point p2 = scale(rotatePrecise(v, co, si), c1.r) + c1
   ret.push_back(p1);
   if (p1 != p2) ret.push_back(p2);
   return ret;
// intersection area between two circles c1, c2
Tf circleCircleIntersectionArea(Circle c1, Circle c2) {
   Point AB = c2.0 - c1.0:
   Tf d = length(AB);
   if (d \ge c1.r + c2.r) return 0:
   if (d + c1.r <= c2.r) return PI * c1.r * c1.r;</pre>
   if (d + c2.r <= c1.r) return PI * c2.r * c2.r;</pre>
   Tf alpha1 = acos((c1.r * c1.r + d * d - c2.r * c2.r))
        /(2.0 * c1.r * d)):
   Tf alpha2 = acos((c2.r * c2.r + d * d - c1.r * c1.r))
        /(2.0 * c2.r * d));
   return c1.sector(2 * alpha1) + c2.sector(2 * alpha2);
```

```
// returns tangents from a point p to circle c, Tf Ti
vector<Point> pointCircleTangents(Point p, Circle c) {
    vector<Point> ret:
   Point u = c.o - p;
   Tf d = length(u);
   if (d < c.r)
    else if (dcmp(d - c.r) == 0) {
       ret = {rotate(u, PI / 2)};
   } else {
       Tf ang = asin(c.r / d);
       ret = {rotate(u, -ang), rotate(u, ang)};
   }
   return ret;
}
// returns the points on tangents that touches the circle
    , Tf Ti Same
vector<Point> pointCircleTangencyPoints(Point p, Circle c
    ) {
   Point u = p - c.o;
   Tf d = length(u);
   if (d < c.r)
       return {};
   else if (dcmp(d - c.r) == 0)
       return {c.o + u};
   else {
       Tf ang = acos(c.r / d);
       u = u / length(u) * c.r;
       return {c.o + rotate(u, -ang), c.o + rotate(u,
           ang)};
   }
}
// for two circles c1 && c2, returns two list of points a
// such that a[i] is on c1 && b[i] is c2 && for every i
// Line(a[i], b[i]) is a tangent to both circles
// CAUTION: a[i] = b[i] in case they touch | -1 for c1 =
int circleCircleTangencyPoints(Circle c1, Circle c2,
    vector<Point> &a.
```

```
a.clear(), b.clear();
      int cnt = 0:
      if (dcmp(c1.r - c2.r) < 0) {
         swap(c1, c2);
         swap(a, b);
     Tf d2 = sqLength(c1.o - c2.o);
     Tf rdif = c1.r - c2.r, rsum = c1.r + c2.r;
      if (dcmp(d2 - rdif * rdif) < 0) return 0;</pre>
      if (dcmp(d2) == 0 \&\& dcmp(c1.r - c2.r) == 0) return
      Tf base = angle(c2.o - c1.o);
      if (dcmp(d2 - rdif * rdif) == 0) {
         a.push_back(c1.point(base));
         b.push_back(c2.point(base));
         cnt++:
         return cnt;
     }
      Tf ang = acos((c1.r - c2.r) / sqrt(d2));
      a.push_back(c1.point(base + ang));
      b.push_back(c2.point(base + ang));
      a.push_back(c1.point(base - ang));
      b.push_back(c2.point(base - ang));
      cnt++;
      if (dcmp(d2 - rsum * rsum) == 0) {
         a.push_back(c1.point(base));
         b.push_back(c2.point(PI + base));
         cnt++:
      else if (dcmp(d2 - rsum * rsum) > 0) {
         Tf ang = acos((c1.r + c2.r) / sqrt(d2));
         a.push back(c1.point(base + ang));
         b.push_back(c2.point(PI + base + ang));
         a.push_back(c1.point(base - ang));
         b.push_back(c2.point(PI + base - ang));
         cnt++;
vector
    < }
    Porietturn cnt;
  7.4 Convex
      minkowski sum of two polygons in O(n)
  Polygon minkowskiSum(Polygon A, Polygon B) {
```

int n = A.size(), m = B.size();

```
{rotate(A.begin(), min_element(A.begin(), A.end()), A.
        end());
   rotate(B.begin(), min_element(B.begin(), B.end()), B.
        end()):
   A.push_back(A[0]);
   B.push_back(B[0]);
   for (int i = 0; i < n; i++) A[i] = A[i + 1] - A[i];
   for (int i = 0; i < m; i++) B[i] = B[i + 1] - B[i];</pre>
   Polygon C(n + m + 1);
   C[0] = A.back() + B.back();
   merge(A.begin(), A.end() - 1, B.begin(), B.end() - 1,
         C.begin() + 1,
              polarComp(Point(0, 0), Point(0, -1)));
   for (int i = 1; i < C.size(); i++) C[i] = C[i] + C[i</pre>
        - 1]:
   C.pop_back();
   return C:
// finds the rectangle with minimum area enclosing a
    convex polygon and
// the rectangle with minimum perimeter enclosing a
    convex polygon
// Tf Ti Same
pair<Tf, Tf> rotatingCalipersBoundingBox(const Polygon &p
    ) {
   using Linear::distancePointLine;
   int n = p.size();
   int l = 1, r = 1, j = 1;
   Tf area = 1e100;
   Tf perimeter = 1e100;
   for (int i = 0; i < n; i++) {</pre>
       Point v = (p[(i + 1) \% n] - p[i]) / length(p[(i +
            1) % n] - p[i]);
       while (dcmp(dot(v, p[r % n] - p[i]) - dot(v, p[(r
            + 1) % n] - p[i])) < 0)
           r++:
       while (j < r \mid | dcmp(cross(v, p[j % n] - p[i]) -
                                              cross(v, p[(
                                                  j + 1)
                                                  % n] -
                                                   p[i]))
                                                   < 0)
           j++;
       while (1 < j ||
                    dcmp(dot(v, p[1 \% n] - p[i]) - dot(v,
                         p[(1 + 1) \% n] - p[i])) > 0)
           1++;
```

```
Tf w = dot(v, p[r \% n] - p[i]) - dot(v, p[1 \% n]
           - p[i]);
       Tf h = distancePointLine(p[j % n], Line(p[i], p[(
           i + 1) % n]));
       area = min(area, w * h);
       perimeter = min(perimeter, 2 * w + 2 * h);
   }
   return make_pair(area, perimeter);
// returns the left side of polygon u after cutting it by
     ray a->b
Polygon cutPolygon(Polygon u, Point a, Point b) {
   using Linear::lineLineIntersection;
   using Linear::onSegment;
   Polygon ret;
   int n = u.size();
   for (int i = 0; i < n; i++) {</pre>
       Point c = u[i], d = u[(i + 1) \% n];
       if (dcmp(cross(b - a, c - a)) >= 0) ret.push_back
       if (dcmp(cross(b - a, d - c)) != 0) {
          Point t;
           lineLineIntersection(a, b - a, c, d - c, t);
           if (onSegment(t, Segment(c, d))) ret.push_back
      }
   }
   return ret;
// returns true if point p is in or on triangle abc
bool pointInTriangle(Point a, Point b, Point c, Point p)
   return dcmp(cross(b - a, p - a)) >= 0 && dcmp(cross(c
        -b, p-b)) >= 0 &&
                dcmp(cross(a - c, p - c)) >= 0;
// pt must be in ccw order with no three collinear points
// returns inside = -1, on = 0, outside = 1
int pointInConvexPolygon(const Polygon &pt, Point p) {
   int n = pt.size();
   assert(n >= 3);
   int lo = 1, hi = n - 1;
   while (hi - lo > 1) {
       int mid = (lo + hi) / 2:
       if (dcmp(cross(pt[mid] - pt[0], p - pt[0])) > 0)
          lo = mid:
       else
          hi = mid;
```

```
bool in = pointInTriangle(pt[0], pt[lo], pt[hi], p);
   if (!in) return 1;
   if (dcmp(cross(pt[lo] - pt[lo - 1], p - pt[lo - 1]))
        == 0) return 0:
   if (dcmp(cross(pt[hi] - pt[lo], p - pt[lo])) == 0)
   if (dcmp(cross(pt[hi] - pt[(hi + 1) % n], p - pt[(hi
       + 1) % n])) == 0)
       return 0;
   return -1;
// Extreme Point for a direction is the farthest point in // #1 If a segment is collinear with the line, only that
     that direction
// u is the direction for extremeness
int extremePoint(const Polygon &poly, Point u) {
   int n = (int)poly.size();
   int a = 0, b = n;
   while (b - a > 1) {
       int c = (a + b) / 2;
       if (dcmp(dot(poly[c] - poly[(c + 1) % n], u)) >=
           0 &&
              dcmp(dot(poly[c] - poly[(c - 1 + n) % n],
                  u)) >= 0) {
           return c;
       }
       bool a_up = dcmp(dot(poly[(a + 1) % n] - poly[a],
            u)) >= 0;
       bool c_up = dcmp(dot(poly[(c + 1) % n] - poly[c],
            u)) >= 0:
       bool a_above_c = dcmp(dot(poly[a] - poly[c], u))
           > 0;
       if (a_up && !c_up)
           b = c:
       else if (!a_up && c_up)
           a = c:
       else if (a_up && c_up) {
           if (a_above_c)
              b = c:
           else
              a = c;
       } else {
           if (!a_above_c)
              b = c;
           else
              a = c;
```

```
}
   }
   if (dcmp(dot(poly[a] - poly[(a + 1) % n], u)) > 0 &&
           dcmp(dot(poly[a] - poly[(a - 1 + n) % n], u))
       return a:
   return b % n;
// For a convex polygon p and a line 1, returns a list of
     segments
// of p that touch or intersect line 1.
// the i'th segment is considered (p[i], p[(i + 1) modulo
     [|q|
    is returned
// #2 Else if 1 goes through i'th point, the i'th segment
     is added
// Complexity: O(lg |p|)
vector<int> lineConvexPolyIntersection(const Polygon &p,
    Line 1) {
   assert((int)p.size() >= 3);
   assert(1.a != 1.b);
   int n = p.size();
   vector<int> ret;
   Point v = 1.b - 1.a;
   int lf = extremePoint(p, rotate90(v));
   int rt = extremePoint(p, rotate90(v) * Ti(-1));
   int olf = orient(l.a, l.b, p[lf]);
   int ort = orient(l.a, l.b, p[rt]);
   if (!olf || !ort) {
       int idx = (!olf ? lf : rt);
       if (orient(1.a, 1.b, p[(idx - 1 + n) \% n]) == 0)
           ret.push back((idx - 1 + n) \% n);
           ret.push_back(idx);
       return ret;
   if (olf == ort) return ret;
   for (int i = 0; i < 2; ++i) {</pre>
       int lo = i ? rt : lf;
       int hi = i ? lf : rt;
       int olo = i ? ort : olf;
       while (true) {
           int gap = (hi - lo + n) \% n;
```

```
if (gap < 2) break;
           int mid = (lo + gap / 2) % n;
           int omid = orient(l.a, l.b, p[mid]);
           if (!omid) {
              lo = mid;
              break;
           if (omid == olo)
              lo = mid;
           else
              hi = mid;
       }
       ret.push_back(lo);
   }
   return ret;
}
// Calculate [ACW, CW] tangent pair from an external
constexpr int CW = -1, ACW = 1;
bool isGood(Point u, Point v, Point Q, int dir) {
    return orient(Q, u, v) != -dir;
Point better(Point u, Point v, Point Q, int dir) {
    return orient(Q, u, v) == dir ? u : v;
Point pointPolyTangent(const Polygon &pt, Point Q, int
    dir, int lo, int hi) {
    while (hi - lo > 1) {
       int mid = (lo + hi) / 2:
       bool pvs = isGood(pt[mid], pt[mid - 1], Q, dir);
       bool nxt = isGood(pt[mid], pt[mid + 1], Q, dir);
       if (pvs && nxt) return pt[mid];
       if (!(pvs || nxt)) {
           Point p1 = pointPolyTangent(pt, Q, dir, mid +
               1, hi);
           Point p2 = pointPolyTangent(pt, Q, dir, lo,
               mid - 1);
           return better(p1, p2, Q, dir);
       }
       if (!pvs) {
           if (orient(Q, pt[mid], pt[lo]) == dir)
              hi = mid - 1;
           else if (better(pt[lo], pt[hi], Q, dir) == pt[
               lo])
              hi = mid - 1;
           else
              lo = mid + 1;
```

```
}
       if (!nxt) {
           if (orient(Q, pt[mid], pt[lo]) == dir)
              lo = mid + 1:
           else if (better(pt[lo], pt[hi], Q, dir) == pt[
               101)
              hi = mid - 1;
           else
              lo = mid + 1:
       }
   }
   Point ret = pt[lo];
   for (int i = lo + 1; i <= hi; i++) ret = better(ret,</pre>
        pt[i], Q, dir);
   return ret;
// [ACW, CW] Tangent
pair<Point, Point> pointPolyTangents(const Polygon &pt,
    Point Q) {
   int n = pt.size();
   Point acw_tan = pointPolyTangent(pt, Q, ACW, 0, n -
   Point cw_tan = pointPolyTangent(pt, Q, CW, 0, n - 1);
   return make_pair(acw_tan, cw_tan);
```

#### 7.5 Polygon

```
typedef vector<Point> Polygon;
// removes redundant colinear points
// polygon can't be all colinear points
Polygon RemoveCollinear(const Polygon &poly) {
   Polygon ret;
   int n = poly.size();
   for (int i = 0; i < n; i++) {</pre>
       Point a = poly[i];
       Point b = poly[(i + 1) \% n];
       Point c = poly[(i + 2) \% n];
       if (dcmp(cross(b - a, c - b)) != 0 && (ret.empty
            () || b != ret.back()))
           ret.push_back(b);
   }
   return ret;
  returns the signed area of polygon p of n vertices
Tf signedPolygonArea(const Polygon &p) {
   Tf ret = 0:
   for (int i = 0; i < (int)p.size() - 1; i++)</pre>
       ret += cross(p[i] - p[0], p[i + 1] - p[0]);
   return ret / 2;
```

```
// given a polygon p of n vertices, generates the convex
   hull in in CCW
// Tested on https://acm.timus.ru/problem.aspx?space=1&
    num=1185
// Caution: when all points are colinear AND
    removeRedundant == false
// output will be contain duplicate points (from upper
    hull) at back
Polygon convexHull(Polygon p, bool removeRedundant) {
   int check = removeRedundant ? 0 : -1;
   sort(p.begin(), p.end());
   p.erase(unique(p.begin(), p.end()), p.end());
   int n = p.size();
   Polygon ch(n + n);
   int m = 0; // preparing lower hull
   for (int i = 0; i < n; i++) {</pre>
       while (m > 1 &&
                    dcmp(cross(ch[m-1]-ch[m-2], p[i
                        ] - ch[m - 1])) <= check)
          m--;
       ch[m++] = p[i];
   int k = m; // preparing upper hull
   for (int i = n - 2; i >= 0; i--) {
       while (m > k &&
                    dcmp(cross(ch[m-1]-ch[m-2], p[i
                        ] - ch[m - 2])) <= check)
          m--;
       ch[m++] = p[i];
   }
   if (n > 1) m--;
   ch.resize(m):
   return ch;
// returns inside = -1, on = 0, outside = 1
int pointInPolygon(const Polygon &p, Point o) {
   using Linear::onSegment;
   int wn = 0, n = p.size();
   for (int i = 0; i < n; i++) {</pre>
       int j = (i + 1) \% n;
       if (onSegment(o, Segment(p[i], p[j])) || o == p[i
           ]) return 0;
       int k = dcmp(cross(p[j] - p[i], o - p[i]));
       int d1 = dcmp(p[i].y - o.y);
       int d2 = dcmp(p[j].y - o.y);
       if (k > 0 && d1 <= 0 && d2 > 0) wn++;
       if (k < 0 \&\& d2 \le 0 \&\& d1 > 0) wn--;
```

```
return wn ? -1 : 1;
// Given a simple polygon p, and a line l, returns (x, y)
// x = longest segment of 1 in p, y = total length of 1
pair<Tf, Tf> linePolygonIntersection(Line 1, const
    Polygon &p) {
   using Linear::lineLineIntersection;
   int n = p.size();
   vector<pair<Tf, int>> ev;
   for (int i = 0; i < n; ++i) {</pre>
       Point a = p[i], b = p[(i + 1) \% n], z = p[(i - 1) \% n]
           + n) % n];
       int ora = orient(l.a, l.b, a), orb = orient(l.a,
           1.b, b),
              orz = orient(l.a, l.b, z);
       if (!ora) {
           Tf d = dot(a - 1.a, 1.b - 1.a);
           if (orz && orb) {
              if (orz != orb) ev.emplace_back(d, 0);
              // else // Point Touch
           } else if (orz)
              ev.emplace_back(d, orz);
           else if (orb)
              ev.emplace_back(d, orb);
       } else if (ora == -orb) {
           Point ins;
           lineLineIntersection(1, Line(a, b), ins);
           ev.emplace back(dot(ins - 1.a, 1.b - 1.a), 0);
       }
   }
   sort(ev.begin(), ev.end());
   If ans = 0, len = 0, last = 0, tot = 0;
   bool active = false;
   int sign = 0;
   for (auto &qq : ev) {
       int tp = qq.second;
       Tf d = qq.first; /// current Segment is (last, d)
       if (sign) {
                       /// On Border
           len += d - last;
           tot += d - last;
           ans = max(ans, len);
           if (tp != sign) active = !active;
           sign = 0;
       } else {
           if (active) { /// Strictly Inside
              len += d - last;
              tot += d - last;
              ans = max(ans, len);
```

```
    if (tp == 0)
        active = !active;
    else
        sign = tp;
}
last = d;
if (!active) len = 0;
}
ans /= length(l.b - l.a);
tot /= length(l.b - l.a);
return {ans, tot};
```

#### 7.6 Half Plane

```
using Linear::lineLineIntersection;
struct DirLine {
   Point p, v;
   Tf ang;
   DirLine() {}
   /// Directed line containing point P in the direction
   DirLine(Point p, Point v) : p(p), v(v) { ang = atan2(
        v.v, v.x); }
   bool operator<(const DirLine& u) const { return ang <</pre>
         u.ang; }
// returns true if point p is on the ccw-left side of ray
bool onLeft(DirLine 1, Point p) { return dcmp(cross(1.v,
    p - 1.p)) >= 0; }
// Given a set of directed lines returns a polygon such
// the polygon is the intersection by halfplanes created
// left side of the directed lines. MAY CONTAIN DUPLICATE
int halfPlaneIntersection(vector<DirLine>& li, Polygon&
    poly) {
   int n = li.size();
   sort(li.begin(), li.end());
   int first, last;
   Point* p = new Point[n];
   DirLine* q = new DirLine[n];
   q[first = last = 0] = li[0];
   for (int i = 1; i < n; i++) {</pre>
```

```
while (first < last && !onLeft(li[i], p[last -</pre>
        1])) last--;
   while (first < last && !onLeft(li[i], p[first]))</pre>
        first++:
   q[++last] = li[i];
   if (dcmp(cross(q[last].v, q[last - 1].v)) == 0) {
       last--;
       if (onLeft(q[last], li[i].p)) q[last] = li[i];
   if (first < last)</pre>
       lineLineIntersection(q[last - 1].p, q[last -
            1].v, q[last].p, q[last].v,
                                                p[last -
                                                     1])
}
while (first < last && !onLeft(q[first], p[last - 1])</pre>
    ) last--;
if (last - first <= 1) {</pre>
   delete[] p;
   delete[] q;
   poly.clear();
   return 0;
lineLineIntersection(q[last].p, q[last].v, q[first].p
    , q[first].v, p[last]);
int m = 0;
poly.resize(last - first + 1);
for (int i = first; i <= last; i++) poly[m++] = p[i];</pre>
delete[] p;
delete[] q;
return m;
```

# Equations and Formulas

#### Catalan Numbers

$$C_n = \frac{1}{n+1} {2n \choose n} C_0 = 1, C_1 = 1 \text{ and } C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k}$$

The number of ways to completely parenthesize n+1 factors. The number of triangulations of a convex polygon with  $n+2|S^d(n,k)$ , to be the number of ways to partition the integers sides (i.e. the number of partitions of polygon into disjoint [1,2,.,n] into k nonempty subsets such that all elements in  $\sum [\gcd(i,n)=k] = \phi\left(\frac{n}{i}\right)$ triangles by using the diagonals).

form n disjoint i.e. non-intersecting chords.

The number of rooted full binary trees with n+1 leaves (ver- $S^d(n,k) = S(n-d+1,k-d+1), n \ge k \ge d$ tices are not numbered). A rooted binary tree is full if every 8.4 Other Combinatorial Identities vertex has either two children or no children.

Number of permutations of 1, n that avoid the pattern 123  $\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$ (or any of the other patterns of length 3); that is, the number of permutations with no three-term increasing sub-sequence. For n = 3, these permutations are 132, 213, 231, 312 and 321

### 8.2 Stirling Numbers First Kind

The Stirling numbers of the first kind count permutations according to their number of cycles (counting fixed points as cycles of length one).

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1), where, S(0,0) = \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S(n,k) + \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S(n,k) + \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S(n,k) + \sum_{n=0}^{\infty} S(n,k) = \sum_{n=0}^{\infty} S($$

$$1, S(n,0) = S(0,n) = 0 \sum_{k=0}^{n} S(n,k) = n!$$

The unsigned Stirling numbers may also be defined algebraically, as the coefficient of the rising factorial:

$$x^{\bar{n}} = x(x+1)...(x+n-1) = \sum_{k=0}^{n} S(n,k)x^{k}$$

Lets [n, k] be the stirling number of the first kind, then

$$[n - k] = \sum_{0 < i_1 < i_2 < i_k < n} i_1 i_2 \dots i_k.$$

# 8.3 Stirling Numbers Second Kind

Stirling number of the second kind is the number of ways to partition a set of n objects into k non-empty subsets.

$$S(n,k) = k \cdot S(n-1,k) + S(n-1,k-1)$$
, where  $S(0,0) = 1$ ,  $S(n,0) = S(0,n) = 0$   $S(n,2) = 2^{n-1} - 1$   $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k \text{ such that if } m \text{ is any integer, then } \gcd(a+m\cdot b,b) = \gcd(a,b)$  The gcd is a multiplicative function in the follow

ber of ways to partition a set of n objects into k subsets, with  $\gcd(a_1,b) \cdot \gcd(a_2,b)$ .

each subset containing at least r elements. It is denoted by  $\gcd(a, \operatorname{lcm}(b, c)) = \operatorname{lcm}(\gcd(a, b), \gcd(a, c))$ .  $S_r(n,k)$  and obeys the recurrence relation.  $S_r(n+1,k) = |\operatorname{lcm}(a,\operatorname{gcd}(b,c))| = \operatorname{gcd}(\operatorname{lcm}(a,b),\operatorname{lcm}(a,c)).$ 

$$kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$

Denote the n objects to partition by the integers 1, 2, ..., n. De- $gcd(a, b) = \sum \phi(k)$ fine the reduced Stirling numbers of the second kind, denoted each subset have pairwise distance at least d. That is, for  $\overline{i=1}$ The number of ways to connect the 2n points on a circle to any integers i and j in a given subset, it is required that  $|i-j| \geq d$ . It has been shown that these numbers satisfy,

$$\binom{n}{k} = \frac{n}{k} \binom{n-1}{k-1}$$

$$\sum_{i=0}^{k} \binom{n+i}{i} = \sum_{i=0}^{k} \binom{n+i}{n} = \binom{n+k+1}{k}$$

$$n, r \in N, n > r, \sum_{i=r}^{n} \binom{i}{r} = \binom{n+1}{r+1}$$

If 
$$P(n) = \sum_{k=0}^{n} {n \choose k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^{n-k} \binom{n}{k} \cdot P(k)$$

If 
$$P(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot Q(k)$$
, then,

$$Q(n) = \sum_{k=0}^{n} (-1)^k \binom{n}{k} \cdot P(k)$$

#### 8.5 Different Math Formulas

Picks Theorem : A = i + b/2 - 1

**Deragements:**  $d(i) = (i-1) \times (d(i-1) + d(i-2))$ 

$$\frac{n}{ab}$$
 -  $\left\{\frac{b\prime n}{a}\right\}$  -  $\left\{\frac{a\prime n}{b}\right\}$  +

The gcd is a multiplicative function in the following sense: An r-associated Stirling number of the second kind is the num- if  $a_1$  and  $a_2$  are relatively prime, then  $gcd(a_1 \cdot a_2, b) =$ 

For non-negative integers a and b, where a and b are not both zero,  $gcd(n^a - 1, n^b - 1) = n^{gcd(a,b)} - 1$  $\sum_{k=1}^{n} \gcd(k, n) = \sum_{d \mid n} d \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} x^{\gcd(k,n)} = \sum_{l=1}^{n} x^{d} \cdot \phi\left(\frac{n}{d}\right)$  $\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{\text{all}} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{\text{all}} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$  $\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for } n > 1$  $\left| \sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2} \right|$  $\left|\sum_{i=1}^{n}\sum_{j=1}^{n}\gcd(i,j)=\sum_{d=1}^{n}\phi(d)\left\lfloor\frac{n}{d}\right\rfloor^{2}\right|$ 

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$

$$F(n) = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left(\frac{\left(1 + \lfloor \frac{n}{l} \rfloor\right)\left(\lfloor \frac{n}{l} \rfloor\right)}{2}\right)^{2} \sum_{l=1}^{n} \mu(d)ld$$