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**University**

LAB TASK

COUSE CODE : CSE214

COURSE TITLE : ALGORITHM LAB

TOPIC :

BUBBLE SORT , LINEAR SORT , INSERTION SORT , SELECTION SORT IMPLEMENTATION &

ANALYSIS OF TIME COMPLEXITY .

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## Bubble Sort

```
#include<stdio.h>
```

```
int main()
{
    int x1[] = {2,3,5,7,11,13,17,19,23,29};
    int x[]={5,2,8,14,10,9};
    int t,i,j;
    for(i=0;i<6;i++)
    {
        for(j=0;j<6-i-1;j++)
        {
            if(x[j]>x[j+1])
            {
                t=x[j];
                x[j]=x[j+1];
                x[j+1]=t;
            }
        }
    }
    for(i=0;i<6;i++)
    {
        printf("%d ",x[i]);
    }
}
```

We have to do  $n-1$  comparisons in 1st iteration and  $n-2$  in 2nd one or  $n-3$  in 3rd and so on..

So,

$(n-1) + (n-2) + (n-3) + \dots + 3 + 2 + 1$

Sum =  $n(n-1)/2$

Hence,  $O(n^2)$

## Linear Search

```
#include<stdio.h>
```

```
int main()
{
    int x[] = {23,2,11,5,19,3,7,17,23,29};
    int i,data=5;
    for(i=0; i<10; i++)
    {
        if(data==x[i])printf("%d found at %d",data,i);

    }
}
```

Time complexity:

In best case the time complexity would be  $O(1)$

The worst case would be  $O(n)$  because there could be case where we would iterate through the full array but didn't find the data or the data is at the last.

## Insertion Sort

```
#include<stdio.h>
int main()
{
    int x1[] = {2,3,5,7,11,13,17,19,23,29};
    int x[]={5,2,8,14,10,9};
    int t,i,j;
    for(i=0;i<6;i++)
    {
        for(j=0;j<6-i-1;j++)
        {
            if(x[j]>x[j+1])
            {
                t=x[j];
                x[j]=x[j+1];
                x[j+1]=t;
            }
        }
    }
    for(i=0;i<6;i++)
    {
        printf("%d ",x[i]);
    }
}
```

1st iteration => | 4 | 3 | 2 | 1 | No. of comparisons = 1 | No. of movements = 1

2nd iteration => | 3 | 4 | 2 | 1 | No. of comparisons = 2 | No. of movements = 2

3rd iteration => | 2 | 3 | 4 | 1 | No. of comparisons = 3 | No. of movements = 3

4th iteration => | 1 | 2 | 3 | 4 | No. of comparisons = 4 | No. of movements = 4

$T(n) = 2 + 4 + 6 + 8 + \text{-----} + 2(n-1)$

$T(n) = 2 * (1 + 2 + 3 + 4 + \text{-----} + (n-1))$

$T(n) = 2 * (n(n-1))/2$

$T(n) = O(n^2)$

## Selection Sort

```
#include<stdio.h>

int main()

{
    int x[]={5,4,3,2,1};
    int i,j,min,t;
    for(i=0;i<5;i++)
    {
        min=i;
        for(j=i+1;j<5;j++)
        {
            if(x[min]>x[j])
            {
                min=j;
            }
            if(min!=i)
            {
                t=x[i];
                x[i]=x[min];
                x[min]=t;
            }
        }
    }
    for(i=0;i<5;i++)
    {
        printf("%d ",x[i]);
    }
}
```

iteration 1 -> 1 to n

iteration 2 -> 2 to n

iteration 3 -> 3 to n

iteration 4 -> 4 to n

So we can say,

$n+n-1+n-2+\dots+1$

$n(n-1)/2$  if we apply big oh notion  $O(n^2-n)$  or  $O(n^2)$

Hence  $\theta(n^2)$  is the time complexity.

