

Feature	Dijkstra's Algorithm	Bellman-Ford Algorithm
Approach	Greedy	Dynamic Programming
Edge Weights Supported	Only non-negative weights	Handles negative weights (except negative cycles in undirected graphs)
Negative Cycle Detection	Cannot detect negative cycles	Can detect negative cycles in directed graphs
Graph Types Supported	Works on directed and undirected graphs with non-negative weights	Works on directed graphs with negative weights; undirected only if all weights non-negative
Time Complexity	$O(V^2)$ (simple), $O((E+V) \log V)$ with binary heap, $O(V \log V + E)$ with Fibonacci heap	$O(VE)$
Space Complexity	$O(V)$ + priority queue	$O(V)$ (simple arrays for distances)
Path Update Mechanism	Selects the node with the minimum tentative distance, relaxes its outgoing edges	Relaxes all edges $V-1$ times, updating distances if a shorter path is found
Algorithm Steps	<ul style="list-style-type: none"> - Initialize distances, use priority queue - Pick closest unvisited node - Update neighbors 	<ul style="list-style-type: none"> - Initialize distances - Repeat $V-1$ times: relax all edges
Distributed Implementation	Difficult (needs global knowledge of graph)	Easier (can use only local neighbor info, suitable for distributed routing)
Typical Use Cases	Link-state routing protocols (OSPF, IS-IS), GPS navigation	Distance-vector routing protocols (RIP, IGRP), graphs with negative weights

Handling of Undirected Graphs	Works for non-negative weights	Negative weights in undirected graphs lead to negative cycles (problematic)
Optimality	Always finds the shortest path if all weights are non-negative	Finds shortest path even with negative weights, unless negative cycle exists
Implementation Complexity	Needs priority queue, more complex data structures	Simpler: just loops and arrays
Limitations	Fails with negative weights; cannot detect negative cycles	Slower on large graphs; fails on undirected graphs with negative weights

Summary of Key Differences:

- Dijkstra is faster and more efficient for graphs with non-negative weights, but cannot handle negative weights or cycles.
- Bellman-Ford is slower but can handle negative weights and detect negative cycles in directed graphs, making it essential for certain network and optimization problems.
- Dijkstra relies on a global view and priority queue, while Bellman-Ford is amenable to distributed and asynchronous computation.
- Bellman-Ford is more robust in edge cases (negative weights/cycles), while Dijkstra is preferred for performance when the graph is well-behaved (no negative weights).