

# CS 208 HW 3 QUESTION 3

---

SAKSHAM RATHI

22B1003



### Question 3

1.  $L = \{ w \in \Sigma^* \mid \exists u, v \in \Sigma^*, w = u \cdot v, n_0(u) = n_1(v) \}$

We will claim that  $L$  is  $\Sigma^*$  i.e. every string of 1s and 0s belongs to  $L$ .

To prove this fact, we need to provide a possible split of  $u \cdot v$  for every string  $w \in \Sigma^*$ .

Consider  $w = a_1 a_2 \dots a_n$ , where  $a_1, a_2 \dots a_n$  are letters of the string  $w$ ; clearly  $a_i = 0$  or  $1$ .

I will try to provide an algorithmic proof for every such  $w$ .

Now, there are two extreme cases (i)  $u = \epsilon, v = w$  or  
(ii)  $u = w, v = \epsilon$

In the first case (i)  $n_0(u) = 0$   
 $n_1(v) = n_1(w)$

In case (ii)  $n_0(u) = n_0(w)$   
 $n_1(v) = 0$

If  $w$  contained only 0s, then we can take case (i) as the possible split. ( $n_0(u) = n_1(v) = 0$ )

If  $w$  contained only 1s, then we can take case (ii) as the possible split. ( $n_0(u) = n_1(v) = 0$ )

Thus, we can safely assume that  $w$  contains both 1s and 0s. Therefore  $n_1(v)$  in case (i)  $\neq 0$  and  $n_0(u)$  in case (ii)  $\neq 0$ .

$$w = \underset{\substack{\uparrow \\ \text{case (i)}}}{a_1} a_2 a_3 \dots a_{n-1} a_n \underset{\substack{\uparrow \\ \text{case (ii)}}}{a_n}$$

Now, if we add  $a_1$  to  $u$  and remove it from  $v$ , there are two possible cases:

$$a_1 = 1 \rightarrow \begin{aligned} n_0(u) &= 0 \\ n_1(v) &= n_1(w) - 1 \end{aligned} \rightarrow \text{This decreased}$$

$$a_1 = 0 \rightarrow \begin{aligned} n_0(u) &< 1 \rightarrow \text{This increased} \\ n_1(v) &= n_1(w) \end{aligned}$$

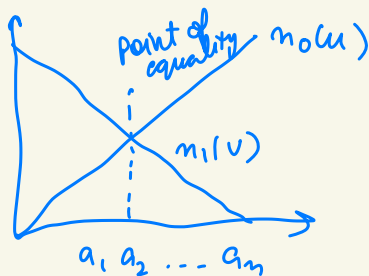
So, in both the cases, one of the sides changed by 1.

In the start (case (i)), we had  $n_0(u) = 0$  and  $n_1(v)$  non-zero.

In the end (case (ii)), we had  $n_0(u)$  non-zero and  $n_1(v) = 0$ .

And in each step, only one of them changed.

So, clearly there must exist a point where both of them are the same.



Therefore, we are successful in finding  $u, v$  for every such  $w$ .

Hence, every string  $w \in \Sigma^*$  also belongs to  $L$ .

Since  $\Sigma^*$  is regular,  $L$  is also regular. Here is the DFA:



2. Yes, there is a unique way to split  $w$  as  $u \cdot v$  for every  $w \in L$ .

Although the previous proof and the graph thereafter shows this fact, but we write the proof more formally here.

Let us assume, we have found a single way  $w = u \cdot v$  such that

$$n_0(u) = n_1(v) \quad \dots \quad (i)$$

(we can always do so, since  $w \in L$ )

We assume to the contrary that there exists another way  $w = u_1 \cdot v_1$ .

Then there are two cases

- (i)  $\text{len}(u_1) > \text{len}(u)$  and  $\text{len}(v_1) < \text{len}(v)$
- (ii)  $\text{len}(u_1) < \text{len}(u)$  and  $\text{len}(v_1) > \text{len}(v)$

There can't be any other possibility because

$$\text{len}(u_1) + \text{len}(v_1) = \text{len}(u) + \text{len}(v) = \text{len}(w)$$

[If both correspondingly equal, then the split will be same.]

(i) In this case we can split  $v$  as  $v_2 \cdot v_1$

$$\text{So } w = u \cdot v = \underbrace{(u \cdot v_2)}_{= u_1} (v_1)$$

$$n_0(u_1) = n_0(u) + n_0(v_2)$$

$$n_0(u_1) = n_1(v_1) \quad (\text{Because, this is a valid split})$$

$$n_0(u) + n_0(v_2) = n_1(v_1) \quad \dots \quad (ii)$$

Putting (i) in (ii), we get  $n_1(v) + n_0(v_2) = n_1(v_1)$

$$n_1(v) = n_1(v_1) + n_1(v_2)$$

$$n_1(v_1) + n_1(v_2) + n_0(v_2) = n_1(v_1)$$

$$\Rightarrow n_1(v_2) + n_0(v_2) = 0 \Rightarrow \text{len}(v_2) = 0 \Rightarrow v_2 = \epsilon$$

This implies, we have achieved the same splitting.

Hence, this case is rejected.

$$(ii) \quad u = u_1 \cdot u_2$$

$$w = u \cdot v = \underbrace{u_1 \cdot u_2}_{u_1} \cdot v$$

$$n_0(u_1) = n_1(u_1) \quad [\text{valid splitting}]$$

$$n_0(u_1) = n_1(u_2) + \underbrace{n_1(v)}_{n_0(u) \text{ [from (i)]}}$$

$$n_0(u_1) = n_1(u_2) + n_0(u_1) + n_0(u_2)$$

$$\Rightarrow n_0(u_2) + n_1(u_2) = 0 \Rightarrow \text{len}(u_2) = 0 \Rightarrow u_2 = \epsilon$$

$\Rightarrow$  Same, splitting

$\Rightarrow$  This case also rejected.

$\Rightarrow$  Both cases rejected  $\rightarrow$  Contradiction!

Therefore, we have proved that for every  $w \in L$ , the splitting will be unique.