

# CS378 Lab-8

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## Result Plots

Here are the plots:

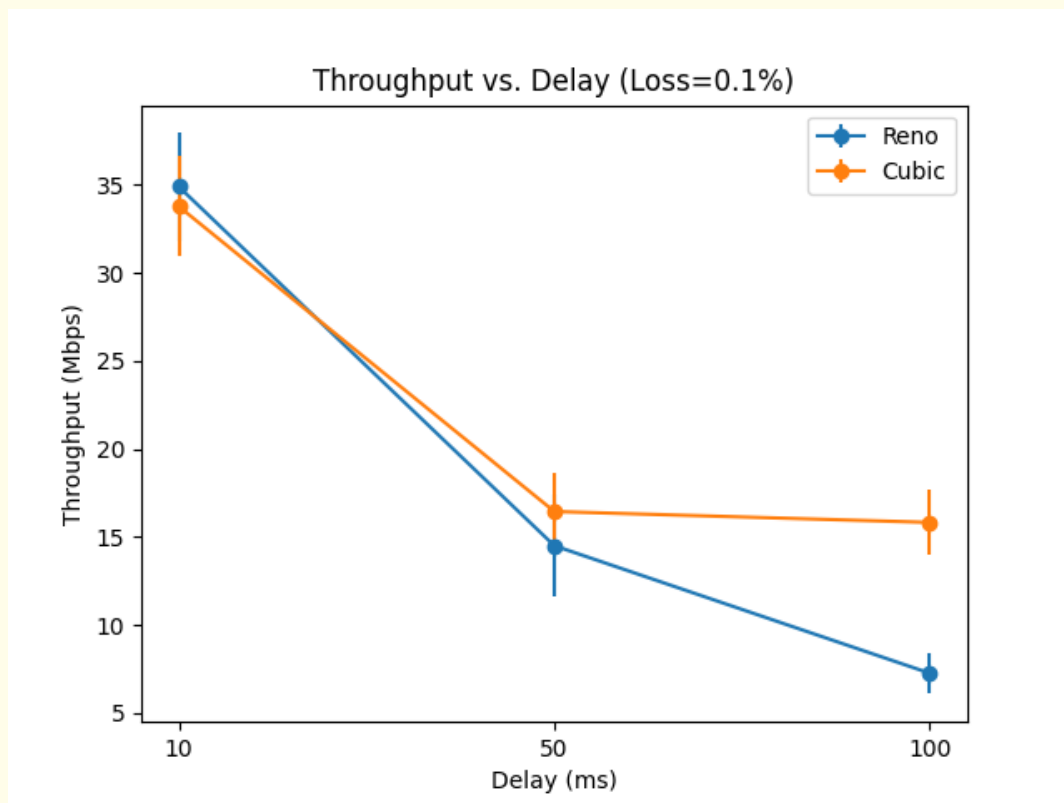


Figure 1: Throughput vs Delay for Loss 0.1%

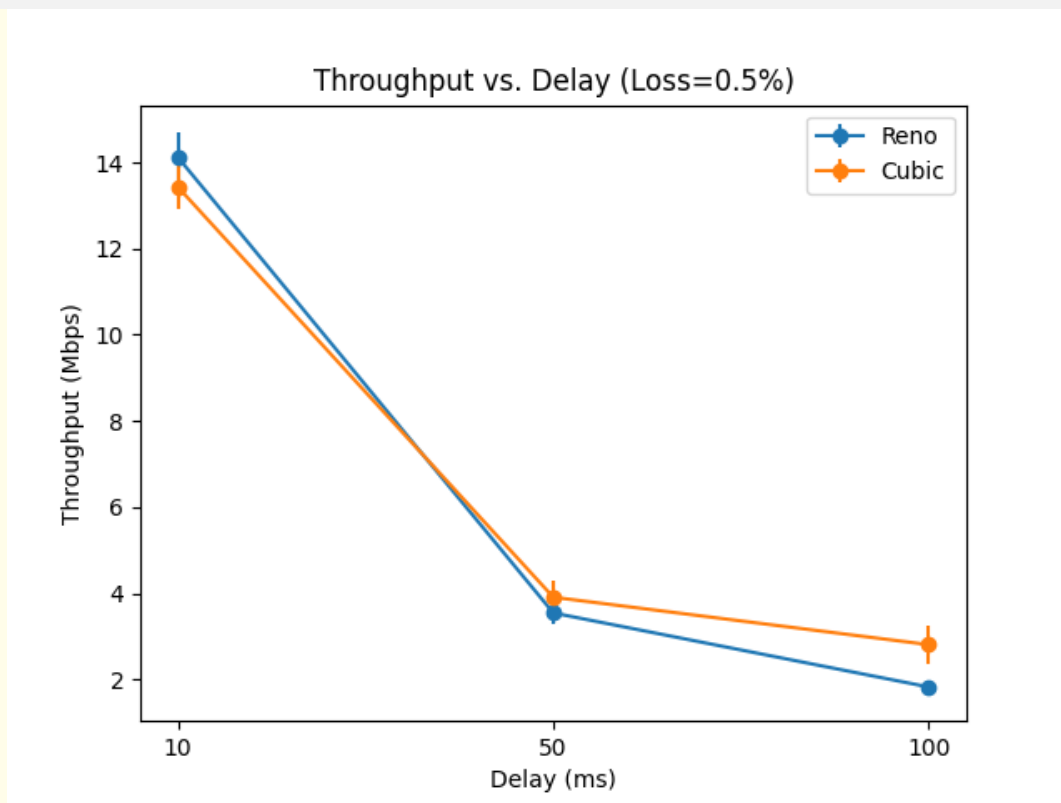


Figure 2: Throughput vs Delay for Loss 0.5%

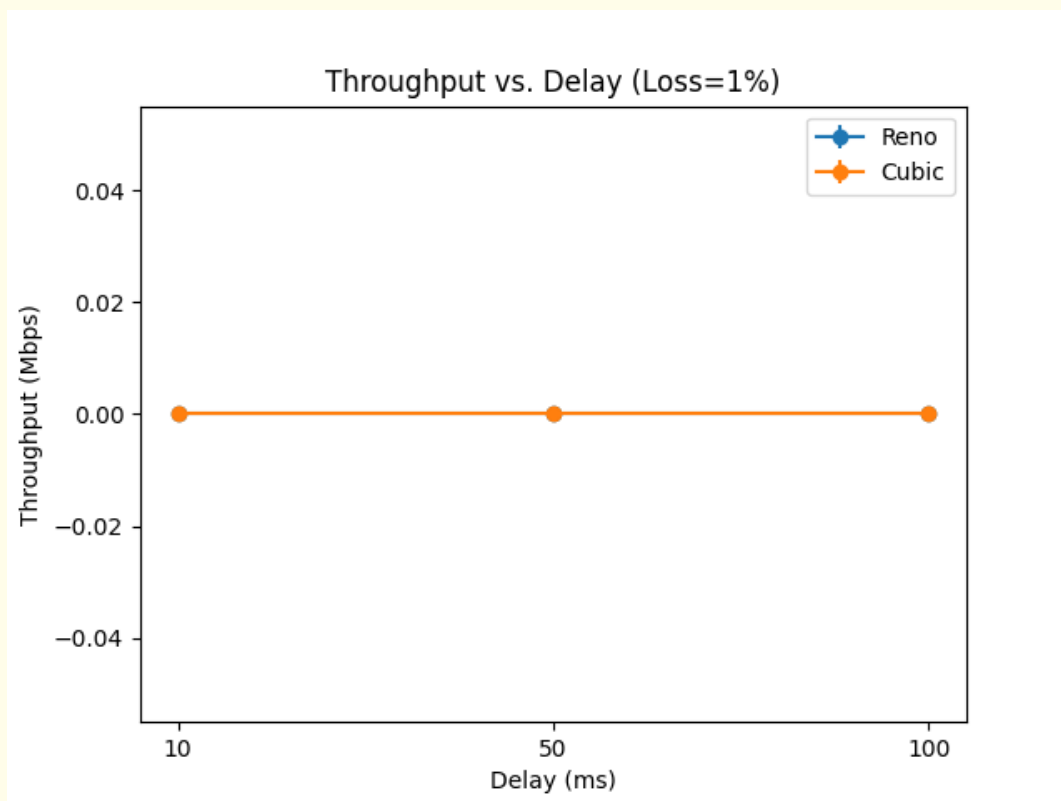


Figure 3: Throughput vs Delay for Loss 1.0%

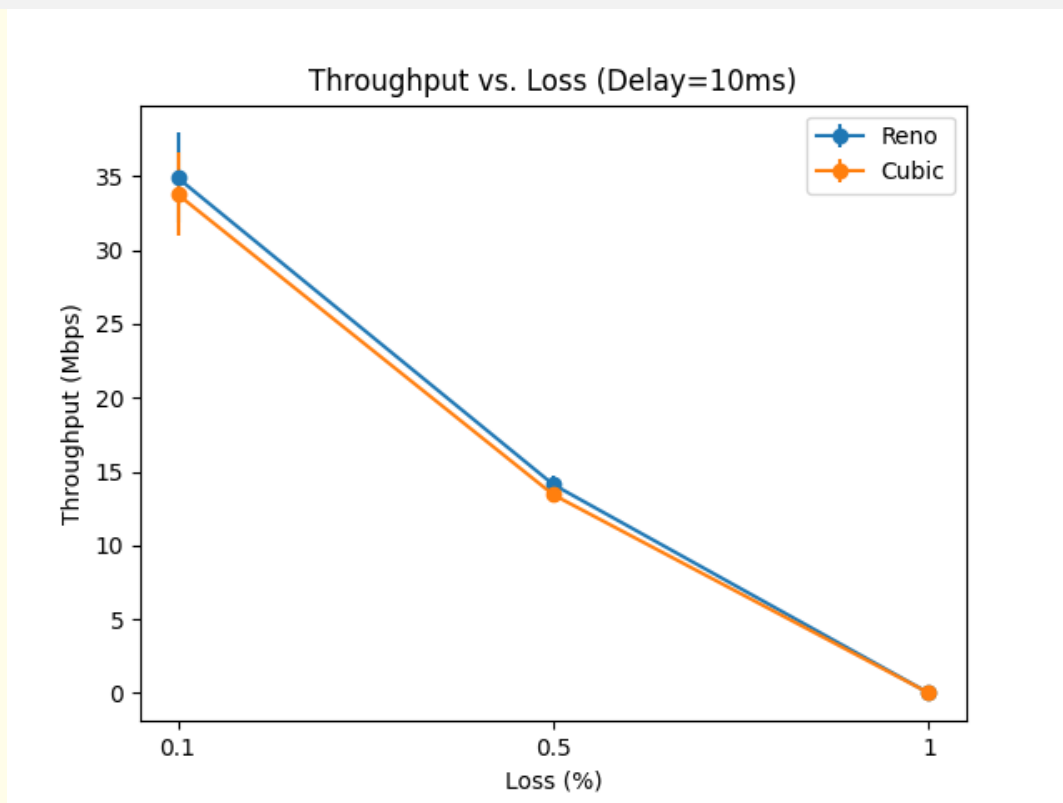


Figure 4: Throughput vs Loss for Delay 10ms

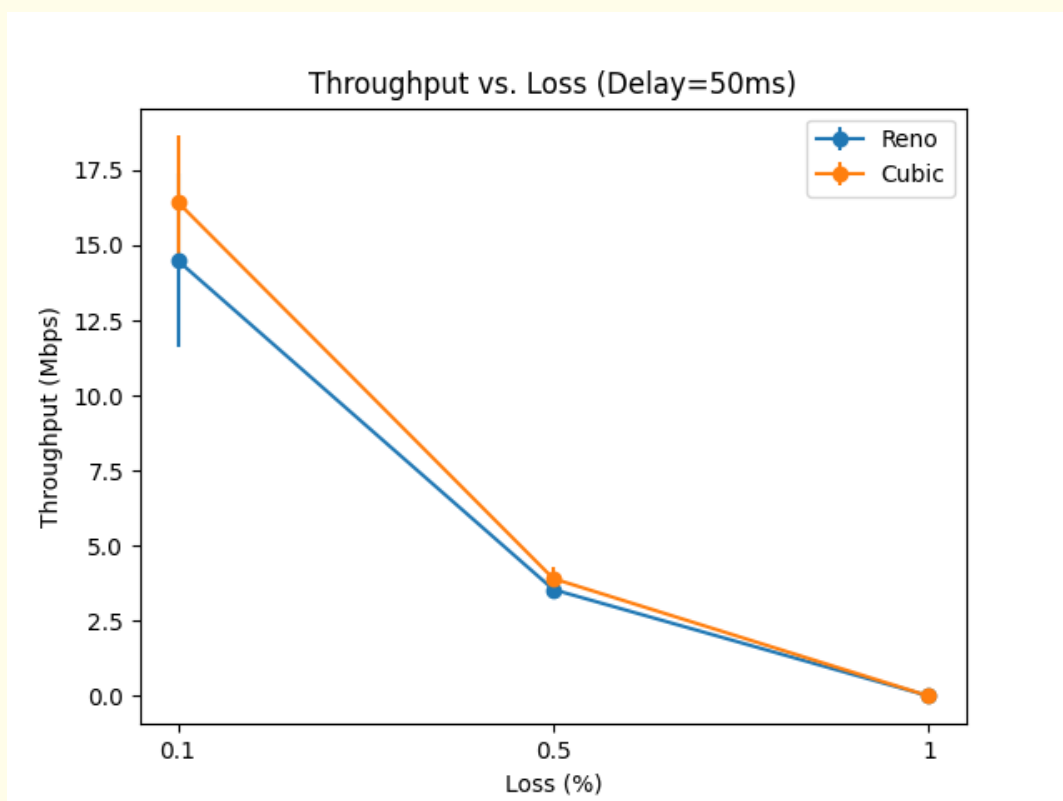


Figure 5: Throughput vs Loss for Delay 50ms

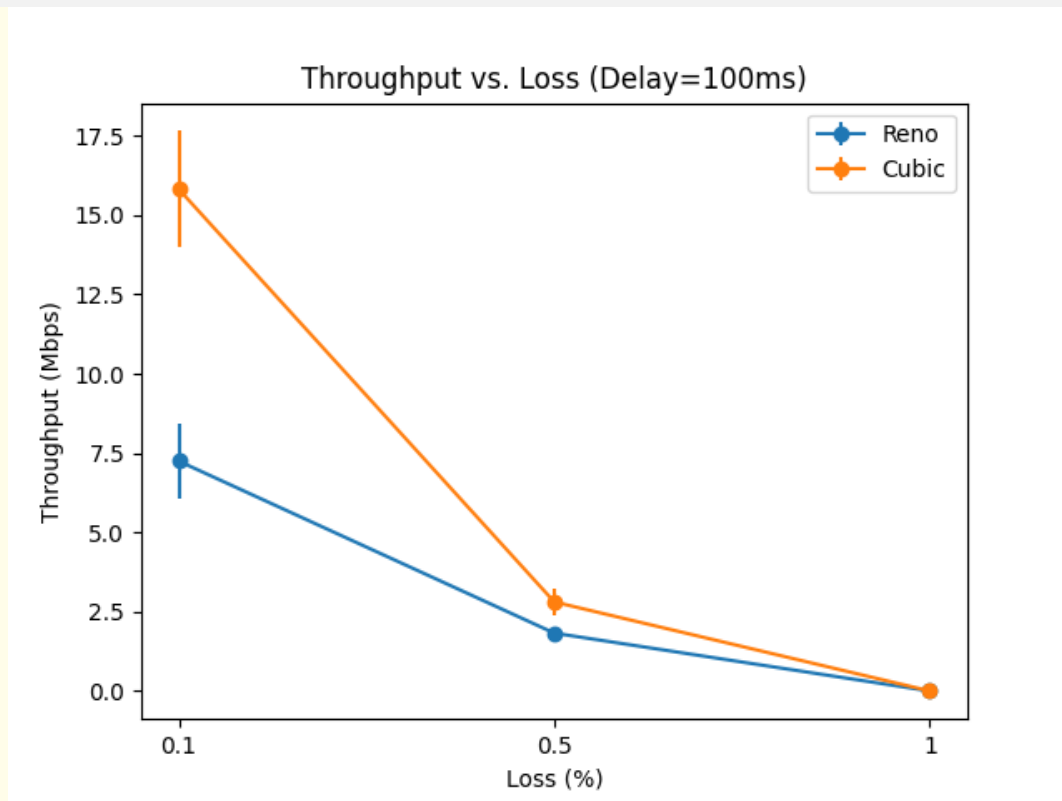


Figure 6: Throughput vs Loss for Delay 100ms

## Observations

- In all of the plots (and for all of the datapoints), the throughput from cubic is higher than that from Reno. This is because cubic is more aggressive in increasing the congestion window size. Therefore it is able to send more packets in the same time (and a higher throughput is observed).
- The throughput decreases as the delay increases. This is because the packets take more time to reach the destination and the sender has to wait for the ACKs. This increases the RTT and hence the throughput decreases. This again confirms the theoretical result that throughput is inversely proportional to the RTT.
- The loss percentage field represents the rate at which packets are dropped. As the loss percentage increases, the throughput decreases. This is because the sender has to retransmit the packets that are lost. This increases the RTT and hence the throughput decreases. When the packet loss rate is 1%, the throughput is almost zero. This is because the sender has to retransmit almost all the packets and the throughput is very low.

## Wireshark Plots

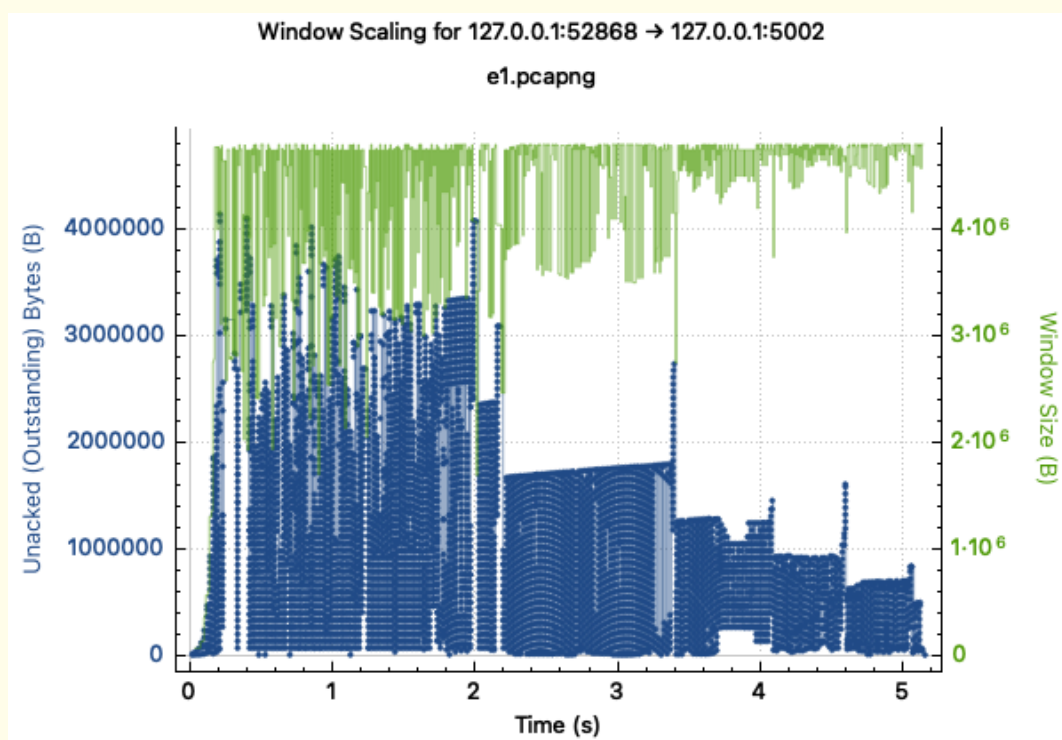


Figure 7: Window Scaling Graph for Delay=10ms and Loss=0.1%

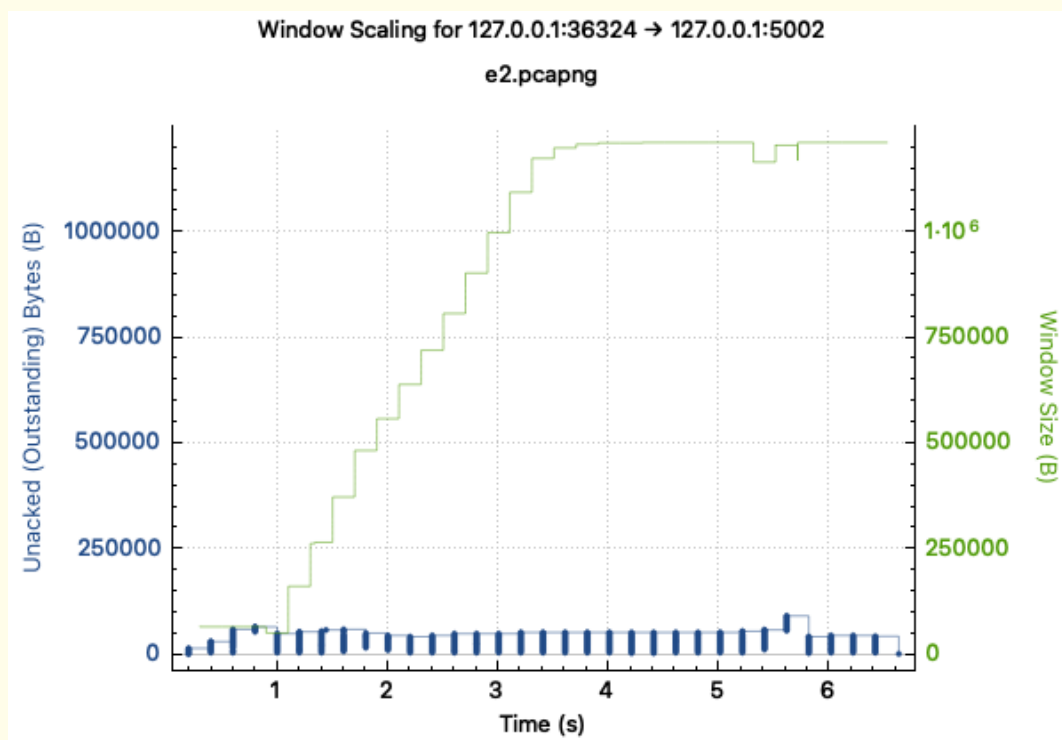


Figure 8: Window Scaling Graph for Delay=100ms and Loss=1%