

EPT-TEST- 54(CD, Lexical, Syntax)
Total Questions: 15
Time: 60 Minutes

Q1.[MCQ]

Consider the following given grammar:

$S \rightarrow Aa$

$A \rightarrow BD$

$B \rightarrow b|\epsilon$

$D \rightarrow d|\epsilon$

Let a, b, d and \$ be indexed as

a	b	d	\$
3	2	1	0

Compute the FOLLOW set of the non-terminal B and write the index values for the symbols in the FOLLOW set in the descending order. (For example, if the FOLLOW set is {a, b, d, \$}, then the answer should be 3210).

- (A) 31**
- (B) 310**
- (C) 230**
- (D) 23**

Q2. [MCQ]

Consider the following grammar G.

$$S \rightarrow F \mid H$$
$$F \rightarrow p \mid c$$
$$H \rightarrow d \mid c$$

Where S, F and H are non-terminal symbols, p, d and c are terminal symbols. Which of the following statement(s) is/are correct

S1: LL(1) can parse all strings that are generated using grammar G.

S2: LR(1) can parse all strings that are generated using grammar G.

(A) Only S1

(B) Only S2

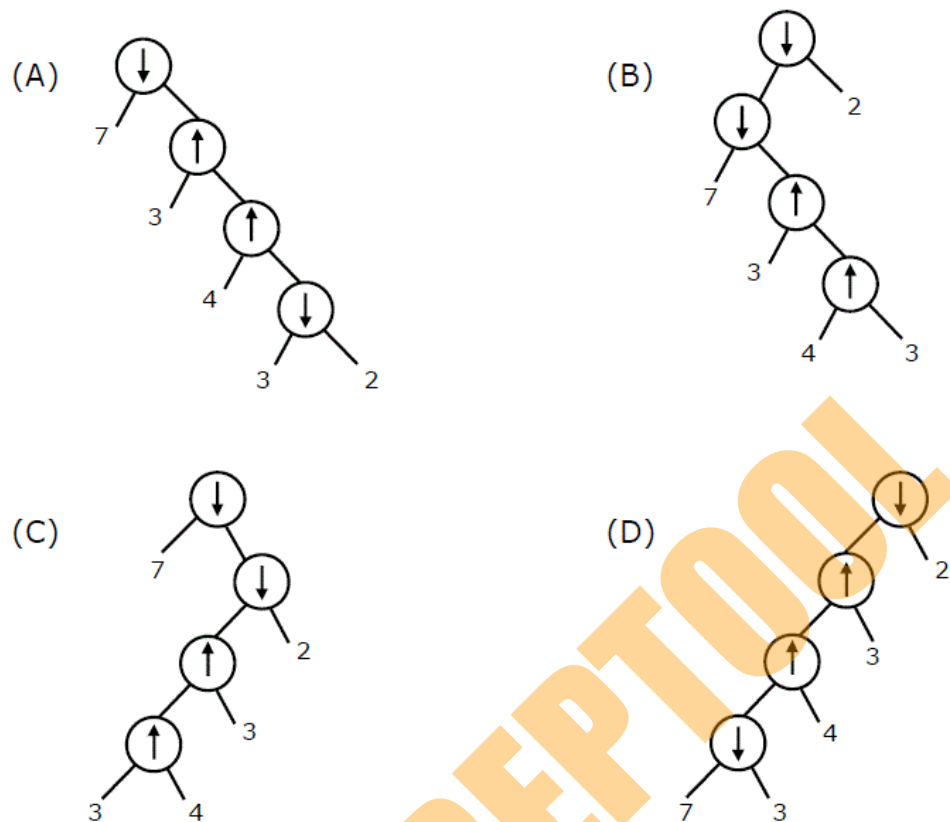
(C) Both S1 and S2

(D) Neither S1 and S2

Q3.[MCQ]

Consider two binary operators ' \uparrow ' and ' \downarrow ' with the precedence of operator \downarrow being lower than that of the operator \uparrow . Operator \uparrow is right associative while operator \downarrow is left associative.

Which one of the following represents the parse tree for expression $(7\downarrow 3\uparrow 4\uparrow 3\downarrow 2)$



Q4.[MCQ]

Which of the following suffices to convert an arbitrary CFG to an LL(1) grammar?

- (A) Removing left recursion alone
- (B) Factoring the grammar alone
- (C) Removing left recursion and factoring the grammar
- (D) None of these

Q5. [MCQ]

The grammar $A \rightarrow AA \mid (A) \mid \epsilon$ is not suitable for predictive-parsing because the grammar is

- (A) ambiguous
- (B) left-recursive
- (C) right-recursive
- (D) an operator-grammar

Q6. [MSQ]

For context-free grammar, FOLLOW(A) is the set of terminals that can appear immediately to the right of non-terminal A in some "sentential" form. We define two sets LFOLLOW(A) and RFOLLOW(A) by replacing the word "sentential" by "left sentential" and "right most sentential" respectively in the definition of FOLLOW (A).

- A. FOLLOW(A) and LFOLLOW(A) may be different.
- B. FOLLOW(A) and RFOLLOW(A) are always the same.
- C. All the three sets are identical.
- D. All the three sets are different.

Q7. [MCQ]

Match the following according to input(from the left column) to the compiler phase(in the right column) that process it:

(P) Syntax tree	(i) Code generator
(Q) Character stream	(ii) Syntax analyzer
(R) Intermediate representation	(iii) Semantic analyzer
(S) Token stream	(iv) Lexical analyzer

- (A) P -> (ii), Q -> (iii), R -> (iv), S -> (i)
 (B) P -> (ii), Q -> (i), R -> (iii), S -> (iv)
 (C) P -> (iii), Q -> (iv), R -> (i), S -> (ii)
 (D) P -> (i), Q -> (iv), R -> (ii), S -> (iii)

Q8.[MCQ]

The number of tokens in the following C statement is
`printf("i = %d, &i = %x", i, &i);`

- (A) 3
 (B) 26
 (C) 10
 (D) 21

Q9. [MCQ]

Consider the following two statements:

- P: Every regular grammar is LL(1)
- Q: Every regular set has a LR(1) grammar

Which of the following is TRUE?

- A.Both P and Q are true
 B.P is true and Q is false

- C. P is false and Q is true
D. Both P and Q are false

Q10.[MCQ]

Consider the CFG with $\{S, A, B\}$ as the non-terminal alphabet, $\{a, b\}$ as the terminal alphabet, S as the start symbol and the following set of production rules:

$S \rightarrow aB$ $S \rightarrow bA$
 $B \rightarrow b$ $A \rightarrow a$
 $B \rightarrow bS$ $A \rightarrow aS$
 $B \rightarrow aBB$ $S \rightarrow bAA$

For the string aabbab, how many derivation trees are there?

- A. 1
B. 2
C. 3
D. 4

Q11.[MCQ]

A lexical analyzer uses the following patterns to recognize three tokens T_1 , T_2 , and T_3 over the alphabet $\{a, b, c\}$.

$T_1: a?(b|c)^*a$
 $T_2: b?(a|c)^*b$
 $T_3: c?(b|a)^*c$

Note that 'x?' means 0 or 1 occurrence of the symbol x. Note also that the analyzer outputs the token that

matches the longest possible prefix.

If the string *bbaacabc* is processed by the analyzer, which one of the following is the sequence of tokens it outputs?

(A) $T_1 T_2 T_3$

(B) $T_1 T_1 T_3$

(C) $T_2 T_1 T_3$

(D) $T_3 T_3$

Q12. [MCQ]

Consider the grammar shown below

$S \rightarrow i E t S S' \mid a$

$S' \rightarrow e S \mid \epsilon$

$E \rightarrow b$

In the predictive parse table M , of this grammar, the entries $M[S', e]$ and $M[S', \$]$ respectively are

A. $\{S' \rightarrow e S\}$ and $\{S' \rightarrow e\}$

B. $\{S' \rightarrow e S\}$ and $\{\}$

C. $\{S' \rightarrow \epsilon\}$ and $\{S' \rightarrow \epsilon\}$

D. $\{S' \rightarrow e S, S' \rightarrow \epsilon\}$ and $\{S' \rightarrow \epsilon\}$

Q13.[MCQ]

Given the following expression grammar:

$E \rightarrow E * F \mid F + E \mid F$

$F \rightarrow F - F \mid id$

Which of the following is true?

- A. $*$ has higher precedence than $+$**
- B. $-$ has higher precedence than $*$**
- C. $+$ and $-$ have same precedence**
- D. $+$ has higher precedence than $*$**

Q14. [MCQ]

In the following grammar

$X ::= X \oplus Y / Y$

$Y ::= Z * Y / Z$

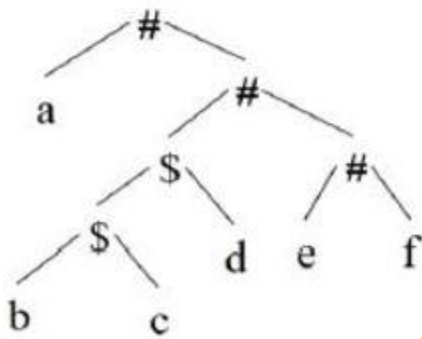
$Z ::= id$

Which of the following is true?

- A. \oplus is left associative while $*$ is right associative**
- B. Both \oplus and $*$ are left associative**
- C. \oplus is right associative while $*$ is left associative**
- D. None of the above**

Q15. [MCQ]

Consider the following parse tree for the expression $a\#b\$c\$d\#e\#f$, involving two binary operators $\$$ and $\#$.



Which one of the following is correct for the given parse tree?

- A. $\$$ has higher precedence and is left associative; $\#$ is right associative**
- B. $\#$ has higher precedence and is left associative; $\$$ is right associative**
- C. $\$$ has higher precedence and is left associative; $\#$ is left associative**
- D. $\#$ has higher precedence and is right associative; $\$$ is left associative**

ANSWERS

A1. A

A2. D

The given grammar is ambiguous as there are two possible leftmost derivations for string "c"

A3. B

A4. D

A5. A

A6. A, B

Ans - A,B.

LFOLLOW may be different from FOLLOW but RFOLLOW and FOLLOW will be the same,

In FOLLOW(A), we add all terminals which appear on the immediate right of A in some sentential form. (LMD/RMD)

In RFOLLOW(A), we add all terminals which appear on the immediate right of A in some right sentential form. (RMD)

In LFOLLOW(A), we add all terminals which appear on the immediate right of A in some left sentential form. (LMD)

Consider the following grammar

$S \rightarrow AB$
 $A \rightarrow a \quad B \rightarrow b$
 $\text{LMD} \mid \text{RMD}$

$S \rightarrow AB \rightarrow A\underline{b} \rightarrow ab$
 $\text{Follow}(A) = b$

LMD
 $S \rightarrow AB \rightarrow aB \rightarrow ab$
 $\text{LFOLLOW}(A) = \emptyset$

$S \rightarrow AB \rightarrow A\underline{b} \rightarrow ab$
 $\text{RFOLLOW}(A) = b$
 RMD

The above example proves that

LFOLLOW may not always be the same as FOLLOW but d prove that RFOLLOW and FOLLOW will always be the same.

A7. C

A8. C

A9. C

P is false

$S \rightarrow Sa/b$

$S \rightarrow aS/a$

These are regular grammars but not LL(1)

Q is true

Every Regular Language is DCFL

Every DCFL has LR(1) Grammar, hence all regular languages are parsed by LR(1) parser.

Hence, Statement Q is True.

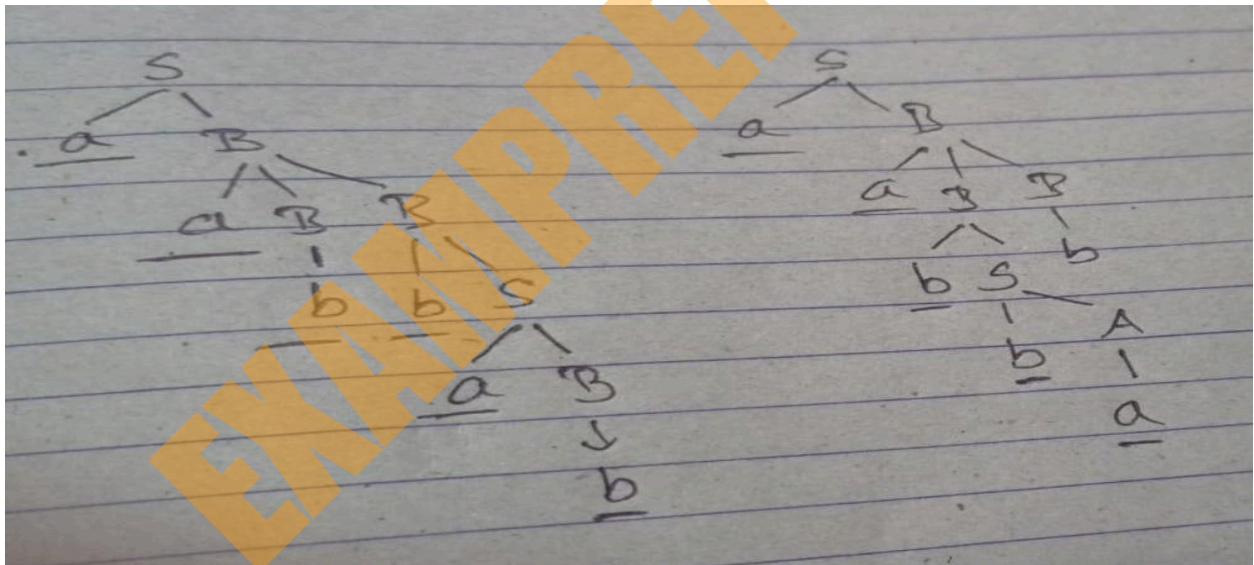
Note also that

Every LR(1) grammar is Unambiguous

So Every DCFL is unambiguous

So Every regular language is also Unambiguous

A10. B



A11. D

A12. D

A13. B

A14. A

A15. A

EXAMPLEPREP1001